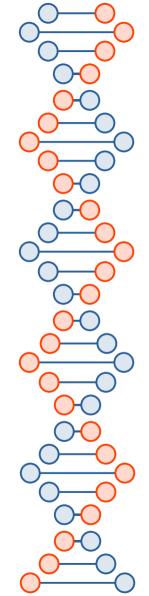


Hash functions and certificates CSE 548 Fall 2024

jedimaestro@asu.edu

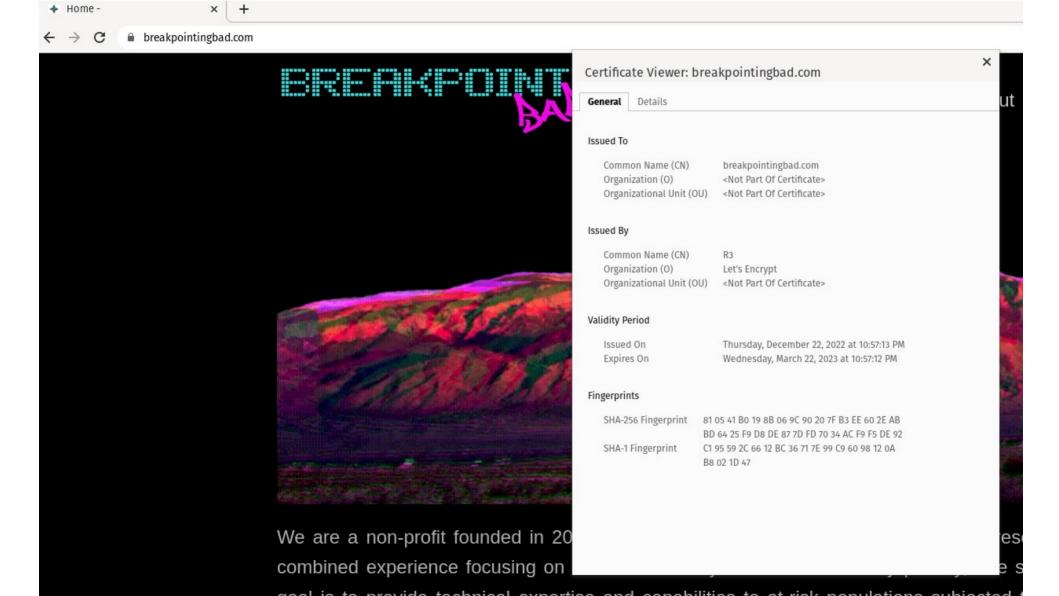


Interesting...

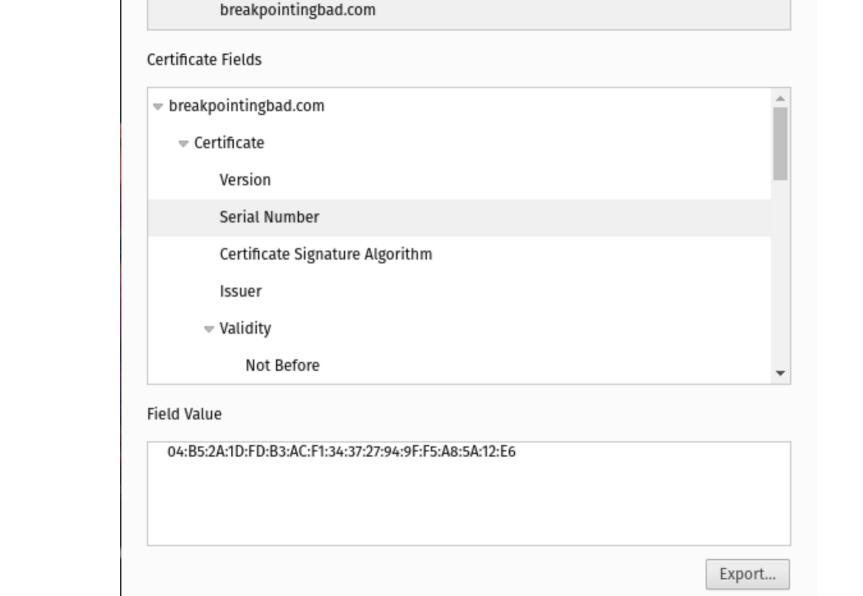
https://media.ccc.de/v/25c3-3023-en-making_the_theoretical_possible

Also check out:

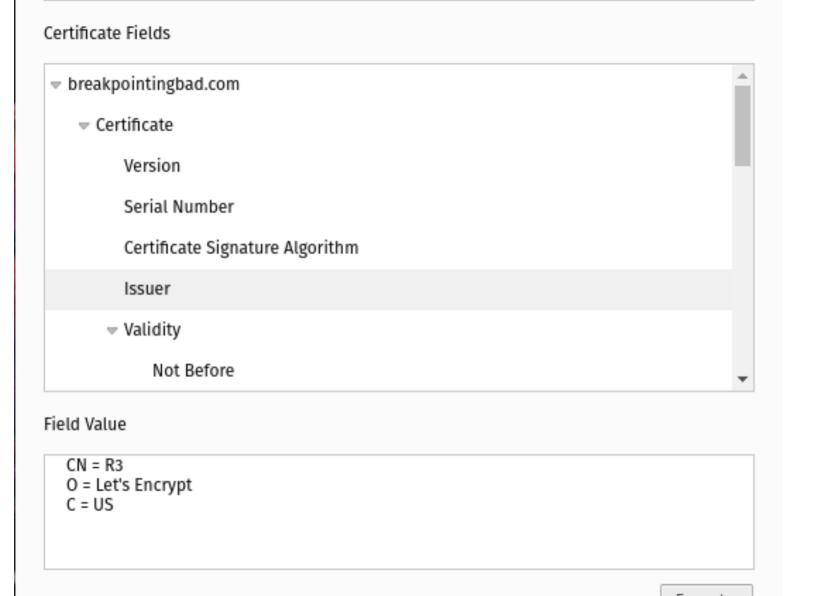
https://www.win.tue.nl/hashclash/rogue-ca/

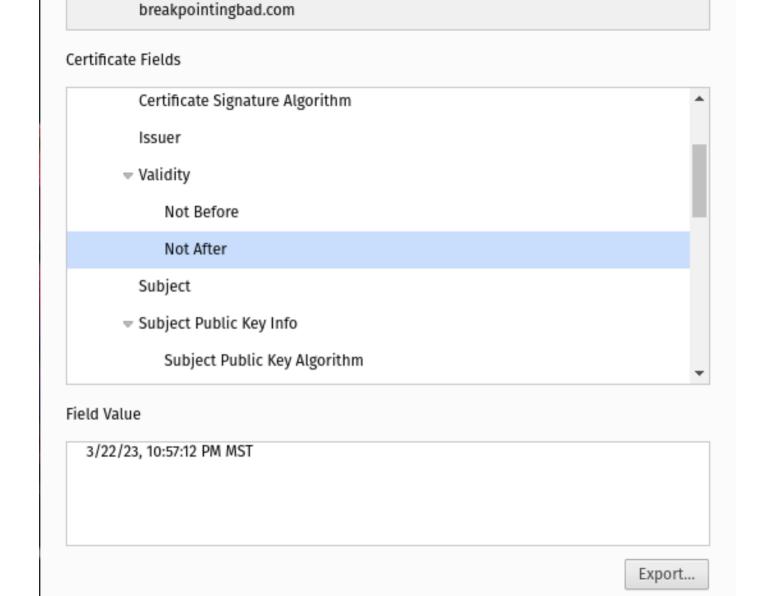


Common Name (CN) breakpointingbad.com Organization (O) <not certificate="" of="" part=""> Organizational Unit (OU) <not certificate="" of="" part=""> ssued By Common Name (CN) R3 Organization (O) Let's Encrypt Organizational Unit (OU) <not certificate="" of="" part=""> /alidity Period Issued On Thursday, December 22, 2022 at 10:57:13 PM Expires On Wednesday, March 22, 2023 at 10:57:12 PM</not></not></not>	ertificate Viewer: I	breakpointingbad.com
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	SHA-1 Fingerprint	C1 95 59 2C 66 12 BC 36 71 7E 99 C9 60 98 12 0A



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Certificate	Fields	
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	ificate	
1	Version	
9	Serial Number	
(Certificate Signature Algorithm	
ı	Issuer	
₩ <i>1</i>	Validity	
	Not Before	,
Field Value		
	SHA-256 With RSA Encryption	



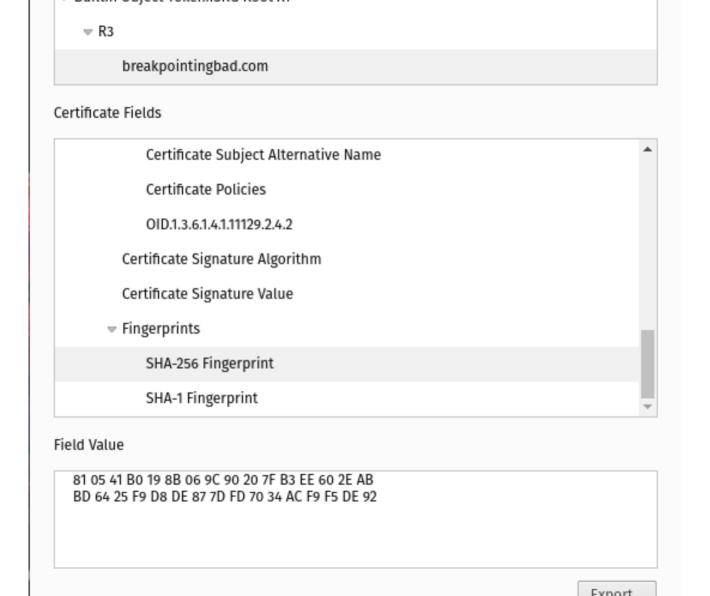


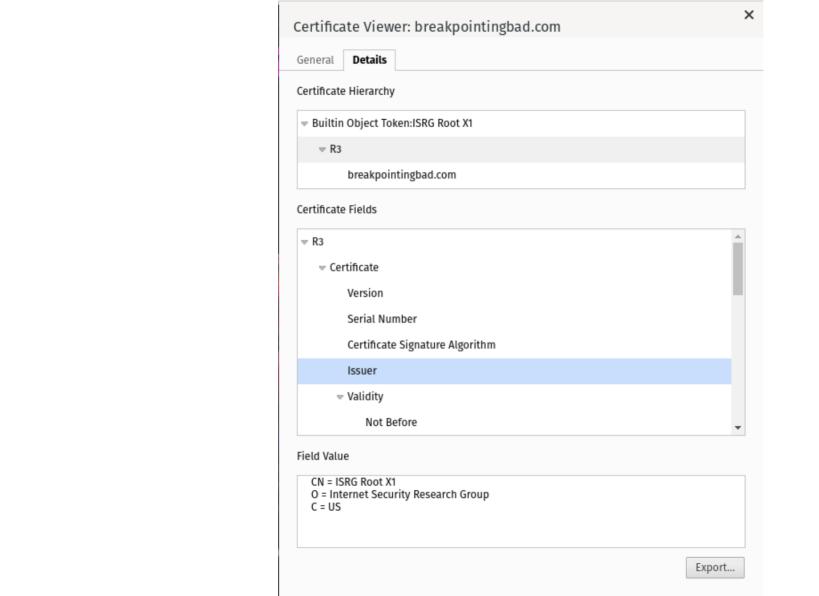
Certificat	e Signature Algorithm	
Issuer		
Not E	efore	
Not A	fter	
Subject		
	Public Key Info	
Subje	ct Public Key Algorithm	
Field Value		
CN = breakpoint	ngbad.com	

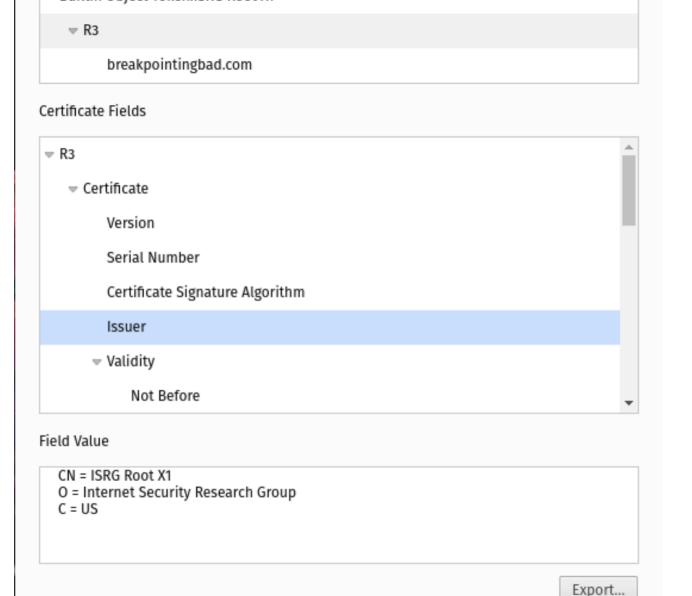
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ertificate Fields		
Not After		
Subject		
)	
Subject Public Key	Algorithm	
Subject's Public Ke	y	
Certificate Key Usa	ge	
Extended Key Usag	ge .	
eld Value		
Modulus (2048 bits): DC DA F0 96 47 5C 62 91 27 27	AD B2 95 EE 3D 51	í
CF 26 EB EC 27 EE ED 2E 9F DA 6 EA CC 96 5B 8C C1 3E A1 C6 46 9		
5C 9B 62 BB 82 3A 7F 77 7C 85 C		

Certificate Fields breakpointingbad.com Certificate Version Serial Number Certificate Signature Algorithm Issuer Validity Not Before Field Value 98 2F 52 F3 68 5E 6D BC 18 2C 93 42 8B C5 41 DT 40 B4 0F 53 D9 BD BA 22 F9 52 90 76 37 F0 C4 56 31 F8 8D C7 B8 21 3E FB 0F 83 B8 A7 CF F3 B4 A1 Public Exponent (17 bits): 01 00 01 Export...

 ▼ R3 breakpointingbad.com Certificate Fields Subject's Public Key ▼ Extensions Certificate Key Usage Extended Key Usage Certificate Basic Constraints 	
Subject's Public Key Extensions Certificate Key Usage Extended Key Usage	
✓ ExtensionsCertificate Key UsageExtended Key Usage	
Certificate Key Usage Extended Key Usage	
Extended Key Usage	
Certificate Basic Constraints	
Certificate Subject Key ID	
Certification Authority Key ID	
Authority Information Access	
Field Value	



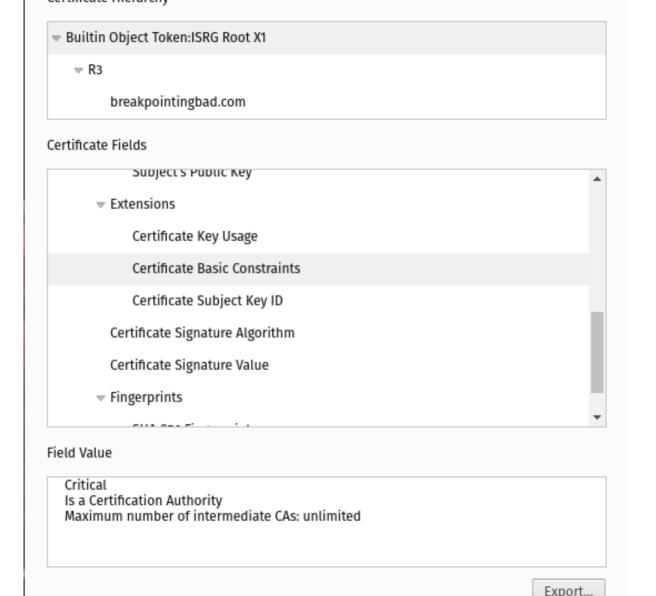


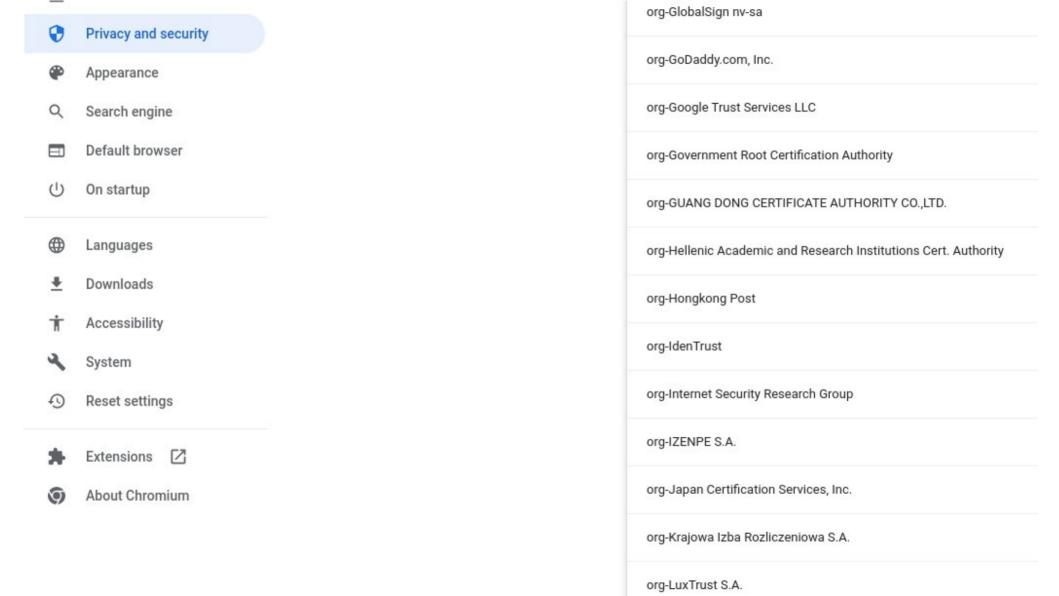


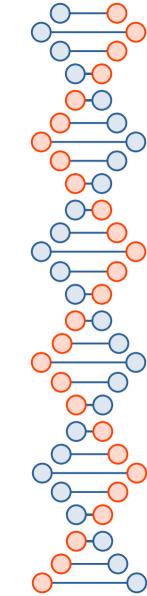
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~	Extensions	
	Certificate Key Usage	
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	Certificate Basic Constraints	
	Certificate Subject Key ID	
	Certification Authority Key ID	
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▼ R3	
breakpointingbad.com	
Certificate Fields	
■ Builtin Object Token:ISRG Root X1	
✓ Certificate	
Version	
Serial Number	
Certificate Signature Algorithm	
Issuer	
Not Before	
Field Value	
CN = ISRG Root X1 O = Internet Security Research Group C = US	

Certificate Hierarchy

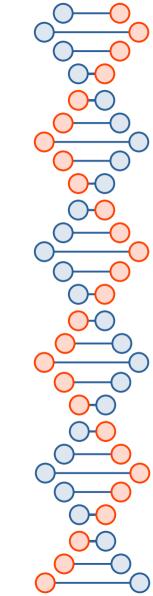






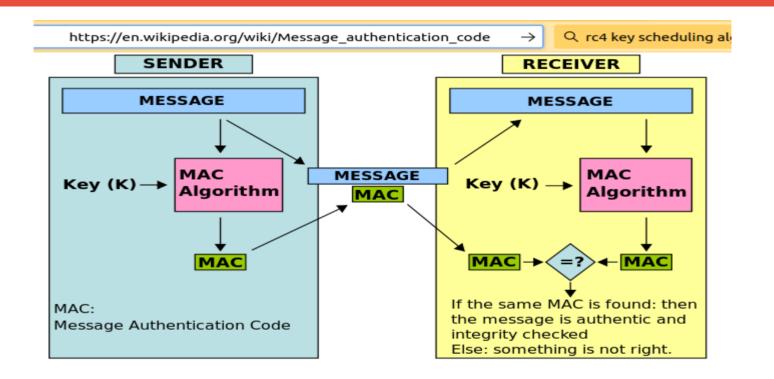
Why hash functions?

- Speed
 - Symmetric crypto is generally faster than asymmetric
 - Hashes are generally faster than either
- Error detection (e.g., checksum)
- Security and privacy

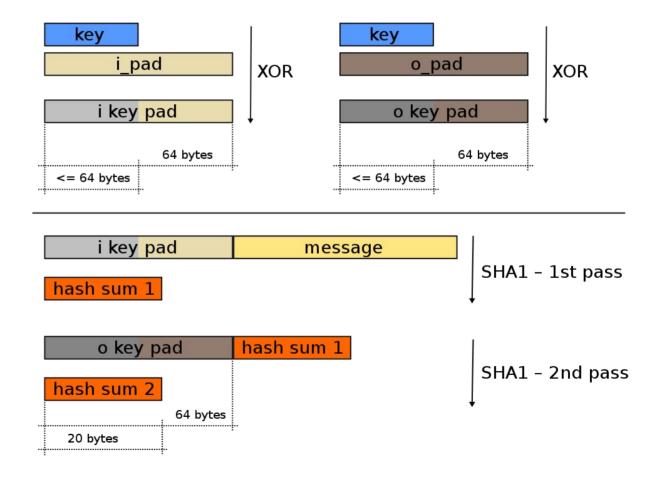


Why cryptographic hash functions?

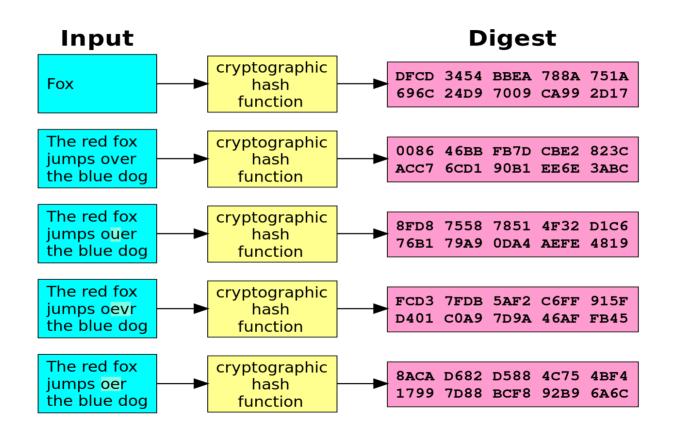
- Unique identifier for an object
- Integrity of an object
 - *E.g.*, message authentication codes
- Digital signatures
 - Sign the digest
 - E.g., 1024-bit RSA, 100MB message, 256-bit digest
- Passwords
- Proof of work



https://en.wikipedia.org/wiki/HMAC



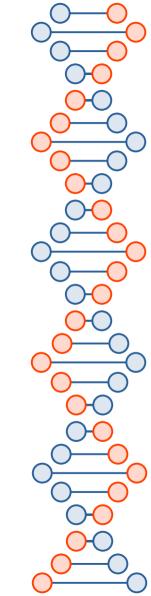
Hash function example



By User:Jorge Stolfi based on Image:Hash_function.svg by Helix84 - Original work for Wikipedia, Public Domain, https://commons.wikimedia.org/w/index.php?curid=5290240

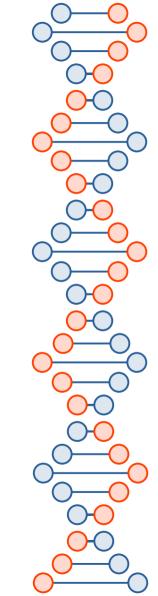


- One-way function
- Deterministic (same input, same output)
- Infeasible to find message that digests to specific hash value
- Infeasible to find two messages that digest to the same hash
- Avalanche effect (small change in message leads to big changes in digest---digests seemingly uncorrelated)
- Still want it to be quick



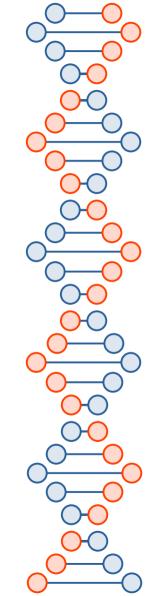
Example algorithms

- MD5: 128-bit digest
 - seriously broken
- SHA-1: 160-bit digest
 - not secure against well-funded adversaries
- SHA-2: 224 to 512 bit digest
 - Merkle–Damgård construction
- SHA-3: 224 to 512 bit digest
 - Sponge construction
 - adopted in August of 2015
- CRC32: not cryptographic, very poor choice



Example algorithms

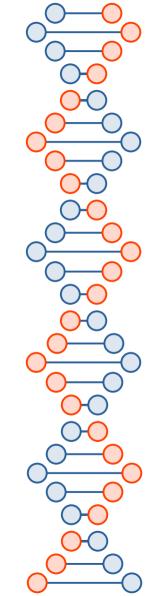
- MD5: 128-bit digest, seriously broken
- SHA-1: 160-bit digest, not secure against well-funded adversaries
- SHA-3: 224 to 512 bit digest, adopted in August of 2015
- CRC32: not cryptographic, very poor choice



Property #1

- Pre-image resistance
- Given h, it should be infeasible to find m such that h = hash(m)

Neither MD5 nor SHA-3 are broken in this way, but MD5 digests are small.

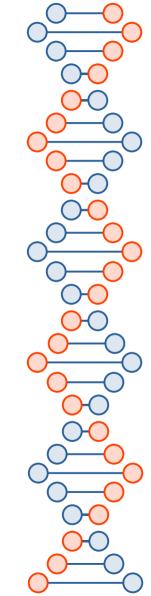


Property #2

- Second pre-image resistance
- Given a message m_1 , it should be infeasible to find another message m_2 such that...

$$hash(m_1) = hash(m_2)$$

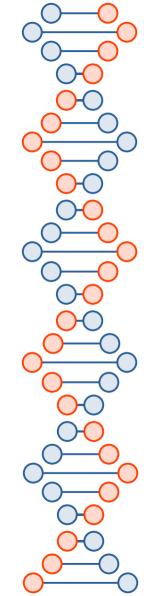
Neither MD5 nor SHA-3 are broken in this way, but MD5 digests are small.



Property #3

- Collision resistance
- It should be infeasible to find two messages, m_1 and m_2 such that... $hash(m_1) = hash(m_2)$

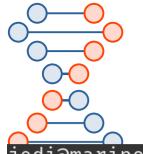
SHA-3 is not broken in this way, MD5 broken in seconds on your laptop, SHA-1 with \$100K or so.



Wang Xiaoyun



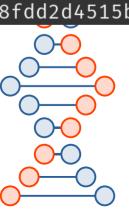
- Tsinghua University
- Contributed a lot of ideas to cracking MD5, SHA-0, and SHA-1



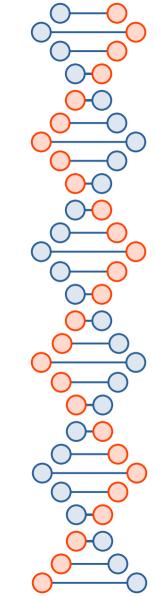
Length extension attack

```
jedi@mariposa:~$ echo "password='lDEnr45#d3'&donut=choc&quantity=1" | md5sum
91a9fc74a98997dba291a26a91c9648e -
jedi@mariposa:~$ echo "password='lDEnr45#d3'&donut=choc&quantity=100" | md5sum
8fdd2d4515bcba887b1b80a653f21e0c -
```

```
jedi@mariposa:~$ echo "password= 'Sdonut=choc&quantity=1" | md5sum
91a9fc74a98997dba291a26a91c9648e -
jedi@mariposa:~$ echo "password= 'Sdonut=choc&quantity=100" | md5sum
8fdd2d4515bcba887b1b80a653f21e0c -
```



MD5 and SHA-1 vulnerable, SHA-2 basically is, SHA-3 is not

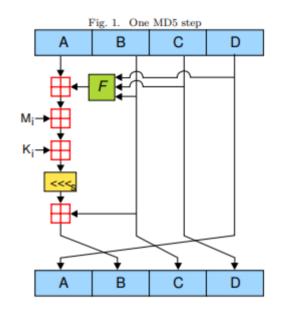


Length extension attack

- One issue is if the attacker doesn't know the password
- Another issue is if the password is different but the attacker finds a collision later on

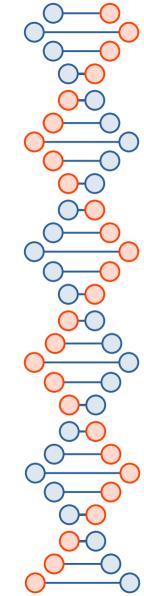
MD5

- Pad to multiple of 512 bits
- 4 rounds
- 4 32-bit words at a time
- Concatenate them at the end for a 128-bit digest
- F is non-linear, varies by round



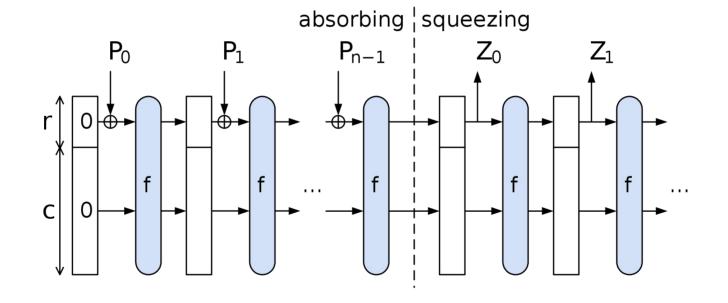
Round (i)	F(X,Y,Z)	g
0	$(X \wedge Y) \vee (\neg X \wedge Z)$	i
1	$(X \wedge Z) \vee (Y \wedge \neg Z)$	$(5 \times i + 1) \mod 16$
2	$(X \oplus Y \oplus Z)$	$i(3 \times i + 5) \mod 16$
3	$(Y \oplus (X \vee \neg Z))$	$(7 \times i) \mod 16$

http://koclab.cs.ucsb.edu/teaching/cren/project/2008/savage.pdf

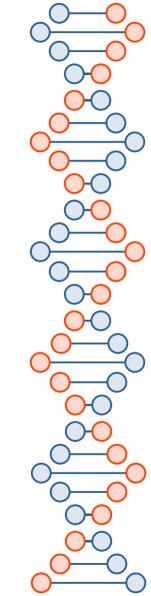


SHA-3

Sponge construction, 1600 bits of internal state



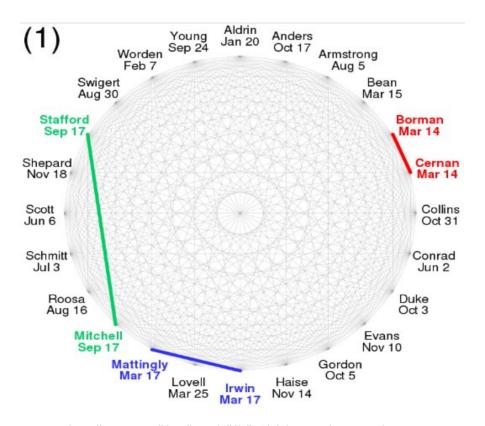
https://en.wikipedia.org/wiki/SHA-3



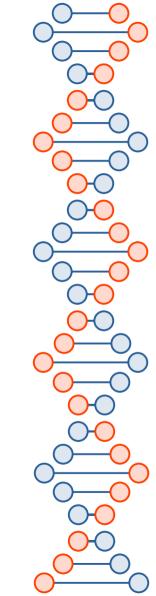
Birthday attack

- Probability of collision is 1 in 2^n , but the expected number of hashes until two of them collide is $sqrt(2^n)=2^{n/2}$
 - Why? Third try has two opportunities to collide, fourth has three opportunities, fifth has six, and so on...

24 people, same birthday?

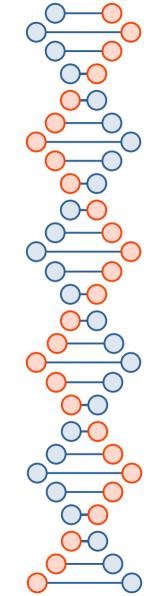


https://commons.wikimedia.org/wiki/File:Birthday_attack_vs_paradox.svg



Chosen-prefix collision attack

- Given two prefixes p_1 and p_2 , find m_1 and m_2 such that $hash(p_1||m_1)=hash(p_2||m_2)$
- p1 and p2 could be domain names in a certificate, images, PDFs, etc. ... any digital image.

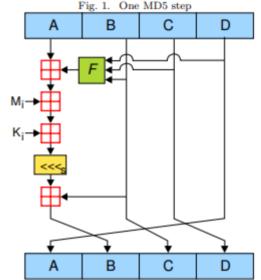


Ingredients for a practical chosen prefix attack on MD5

- Collision attack on MD5
 - That works for any initialization vector (so you can put bits in front)
- Length extension attack
 - So you can put identical bits on the end
 - Birthday attack
 - So you can bridge the prefix to a block that meets the requirements of the collision attack

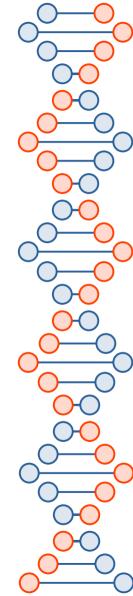
MD5 collision attack by Wang and Yu

 $C_0 = (0, 0, 0, 0, 2^{31}, 0, 0, 0, 0, 0, 0, 2^{15}, 0, 0, 2^{31}, 0)$ and $C_1 = (0, 0, 0, 0, 2^{31}, 0, 0, 0, 0, 0, 0, -2^{15}, 0, 0, 2^{31}, 0)$



Round (i)	F(X,Y,Z)	g
0	$(X \wedge Y) \vee (\neg X \wedge Z)$	i
1	$(X \wedge Z) \vee (Y \wedge \neg Z)$	$(5 \times i + 1) \mod 16$
2	$(X \oplus Y \oplus Z)$	$i(3 \times i + 5) \mod 16$
3	$(Y \oplus (X \vee \neg Z))$	$(7 \times i) \mod 16$

http://koclab.cs.ucsb.edu/teaching/cren/project/2008/savage.pdf



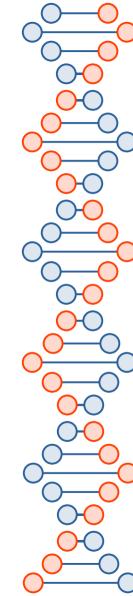
An example

Both have digest 79054025255fb1a26e4bc422aef54eb4

d131dd02c5e6eec4693d9a0698aff95c2fcab58712467eab4004583eb8fb7f89 55ad340609f4b30283e488832571415a085125e8f7cdc99fd91dbdf280373c5b d8823e3156348f5bae6dacd436c919c6dd53e2b487da03fd02396306d248cda0 e99f33420f577ee8ce54b67080a80d1ec69821bcb6a8839396f9652b6ff72a70

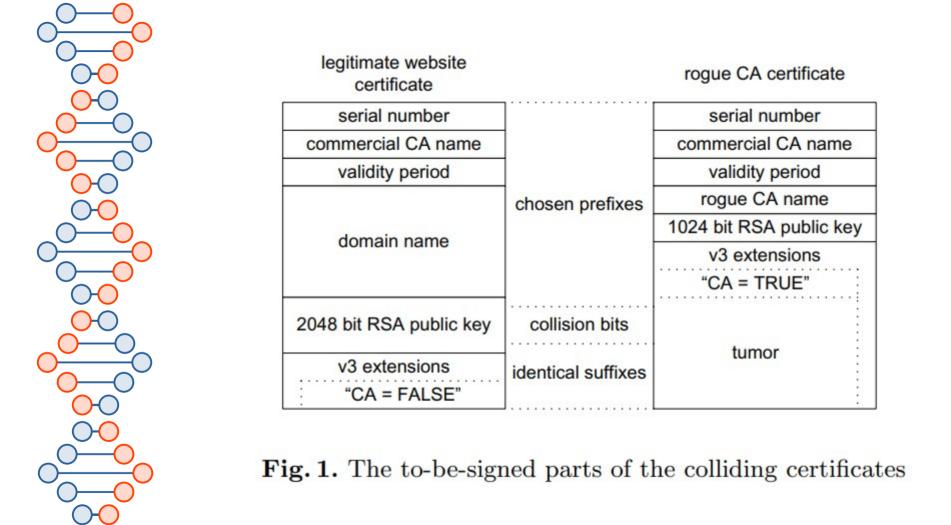
d131dd02c5e6eec4693d9a0698aff95c2fcab50712467eab4004583eb8fb7f89 55ad340609f4b30283e4888325f1415a085125e8f7cdc99fd91dbd7280373c5b d8823e3156348f5bae6dacd436c919c6dd53e23487da03fd02396306d248cda0 e99f33420f577ee8ce54b67080280d1ec69821bcb6a8839396f965ab6ff72a70

https://www.mscs.dal.ca/~selinger/md5collision/

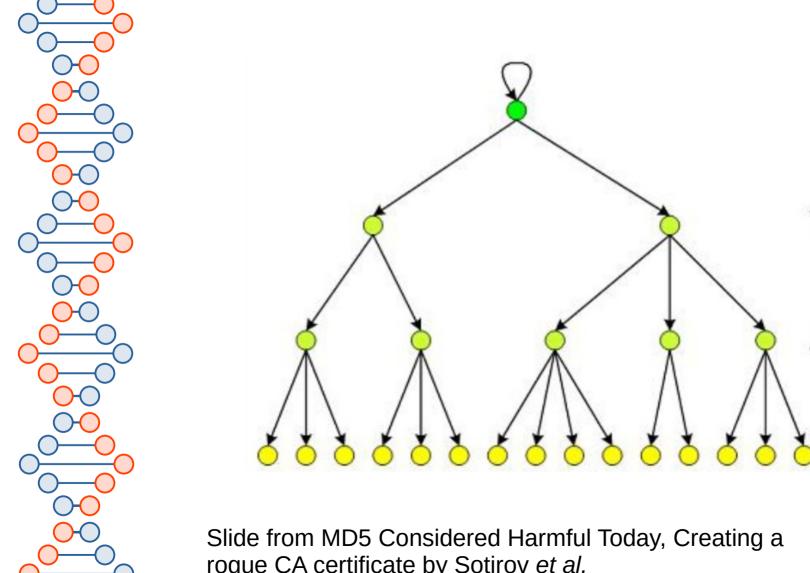


Short Chosen-Prefix Collisions for MD5 and the Creation of a Rogue CA Certificate

Marc Stevens¹, Alexander Sotirov², Jacob Appelbaum³, Arjen Lenstra^{4,5}, David Molnar⁶, Dag Arne Osvik⁴, and Benne de Weger⁷



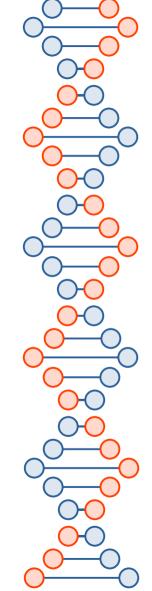




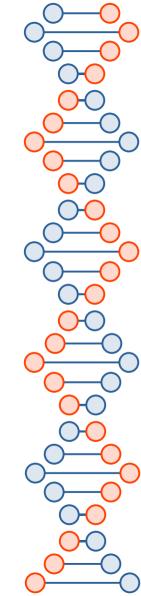
rogue CA certificate by Sotirov et al.

root CA

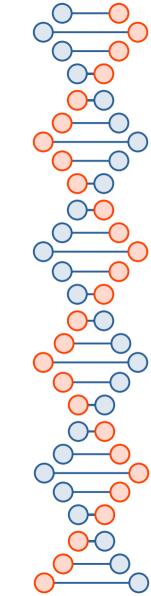
intermediate CA's



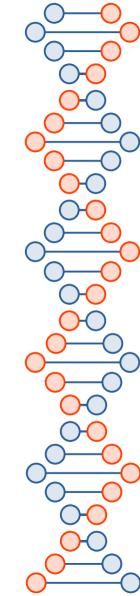
CN=Starfield Services Root Certificate Authority
- G2,O=Starfield Technologies,
Inc.,L=Scottsdale,ST=Arizona,C=US



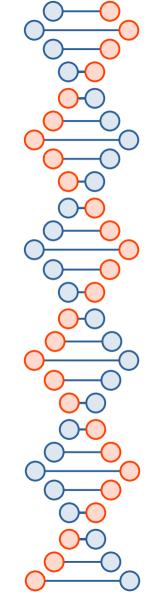
CN=COMODO RSA Certification Authority,O=COMODO CA Limited,L=Salford,ST=Greater Manchester,C=GB



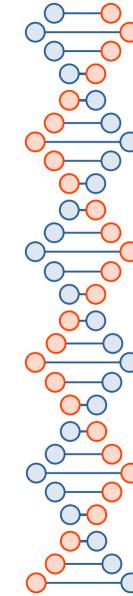
CN=Amazon Root CA 3,O=Amazon,C=US



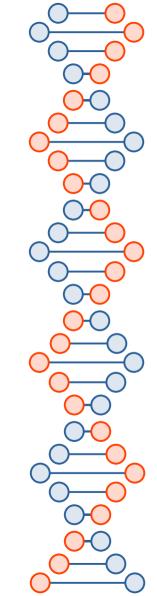
CN=Microsoft RSA Root Certificate Authority 2017,O=Microsoft Corporation,C=US



CN=Autoridad de Certificacion Firmaprofesional CIF A62634068,C=ES



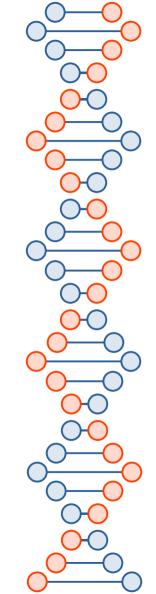
OU=certSIGN ROOT CA G2,O=CERTSIGN SA,C=RO



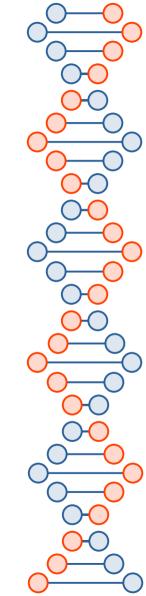
CN=TUBITAK Kamu SM SSL Kok Sertifikasi - Surum 1,OU=Kamu Sertifikasyon Merkezi - Kamu SM,O=Turkiye Bilimsel ve Teknolojik Arastirma Kurumu - TUBITAK,L=Gebze - Kocaeli,C=TR



CN=CFCA EV ROOT,O=China Financial Certification Authority,C=CN

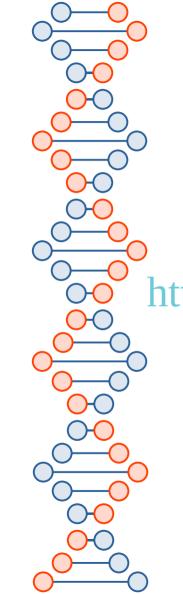


CN=GlobalSign,OU=GlobalSign ECC Root CA - R5,O=GlobalSign



References

- [Cryptography Engineering] *Cryptography Engineering: Design Principles and Applications,* by Niels Ferguson, Bruce Schneier, and Tadayoshi Kohno. Wiley Publishing, 2010.
- Lots of images and info plagiarized from Wikipedia



Some preview of what's to come...

https://www.youtube.com/watch?v=fZ1R9RliM1w