Denote  $\mathcal{H}_{\circ}$  as the hypothesis class defined in the problem. Since we assume that  $\mathcal{H}_{\circ}$  is realizable, there is a perfect  $h_*$  that with  $\operatorname{err}_{\mathcal{D}}(h_*) = 0$ . Denote  $h_{ERM}$  as the hypothesis which achieves 0 error-rate on  $\mathcal{S}$ . First assume there is at least one positive label. We choose  $h_{ERM}$  by finding the smallest radius such that every positive label is contained inside (the boundary is allowed as well). We must have that  $h_*$  contains  $h_{ERM}$ , otherwise  $h_*$  has non-zero error. Let the annulus B between  $h_{ERM}$  and  $h_*$  have probability mass  $\epsilon$ . This area is where  $h_*$  and  $h_{ERM}$  disagree: for  $b \in B$ ,  $h_*(b) = 1$ ,  $h_{ERM}(b) = 0$ , and is therefore the generalization error. Then note that if there were  $x_i \in B$ ,  $h_{ERM}$  would be larger. Therefore,  $\mathbf{P}\{x_i \notin B\} \leq 1 - \epsilon$  since B has probability mass  $\epsilon$ . Thus the probability that no  $x_i \in \mathcal{S}$  is  $\epsilon$  is  $\epsilon$  is  $\epsilon$  is  $\epsilon$  is  $\epsilon$  in  $\epsilon$  in  $\epsilon$ .

Then  $\mathbf{P}\{\operatorname{err}_{\mathcal{D}}(h_{ERM}) > \epsilon\} \leq (1 - \epsilon)^{|\mathcal{S}|} \leq e^{-\epsilon|\mathcal{S}|} = \delta$ , and we choose  $|S| = m_{\mathcal{H}_{\circ}}(\epsilon, \delta) = \frac{1}{\epsilon} \ln\left(\frac{1}{\delta}\right)$ .

If there are no positive labels, then choose  $h_{ERM}$  so that it has a radius  $r_{ERM}$  not containing a negative label such that  $r_{ERM}+\epsilon_0=R_{\rm smallest\ negative\ label}$ . Then we have a two sided bound for  $h_*$ , since  $h_*$  can be inside  $h_{ERM}$  if there are no positive labels in the sample. Denote the annulus of radii  $r_{ERM} < r < r_{ERM} + \epsilon_0$  as  $B_+$ , and the annulus of radius  $r_{h_*} < r < r_{ERM}$  as  $B_-$ . Let the area of  $B_+$  be  $\epsilon_+$ , and the area of  $B_-$  be  $\epsilon_-$ , and let  $\epsilon = \max(\epsilon_+, \epsilon_-)$ .

Then  $\mathbf{P}\{\operatorname{err}_{\mathcal{D}}(h_{ERM}) > \epsilon\} \leq 2(1-\epsilon)^{|\mathcal{S}|} \leq 2e^{-\epsilon|\mathcal{S}|} = \delta$ , and we choose  $|S| = \frac{1}{\epsilon}\ln\left(\frac{2}{\delta}\right)$ .

Then we choose the maximum between the two cases as our lower bound,  $|S| = \frac{1}{\epsilon} \ln \left( \frac{2}{\delta} \right)$ .

Since this is not necessarily a positive integer, and the number of samples is an integer, we know that the minimum number of required samples is bounded above by  $\lceil \frac{1}{\epsilon} \ln \left( \frac{2}{\delta} \right) \rceil$ , and thus  $\mathcal{H}_{\circ}$  is PAC-learnable since  $m_{\mathcal{H}_{\circ}}(\epsilon, \delta)$  is polynomial in all required quantities and obtains the desired error bound.

Let  $\mathcal{D}$  be the original distribution over concept  $h \in \mathcal{C}$  ( $(x,y) \in \mathcal{X} \times \mathcal{Y}$  have zero mass in  $\mathcal{D}$  unless they satisfy the concept h), and let  $\mathcal{D}'$  be the distribution under the uniform noise. We are given that concept class  $\mathcal{C}$  is agnostically learnable. Therefore, with sample complexity poly  $\left(\frac{1}{\epsilon}, \ln\left(\frac{1}{\delta}\right)\right)$ , some algorithm  $\mathcal{A}$  returns a hypothesis  $h_A$  such that

$$\operatorname{err}_{\mathcal{D}'}(h_A) \leq \min_{h^* \in \mathcal{C}} \operatorname{err}_{\mathcal{D}'}(h^*) + \epsilon$$

Now consider that

$$\min_{h^* \in \mathcal{C}} \operatorname{err}_{\mathcal{D}'}(h^*) = \min_{h^* \in \mathcal{C}} \left\{ (1 - \epsilon_0) * \operatorname{err}_{\mathcal{D}}(h^*) + \frac{\epsilon_0}{2} * 1 + \frac{\epsilon_0}{2} * 0 \right\} 
= \frac{\epsilon_0}{2} + (1 - \epsilon_0) * \min_{h^* \in \mathcal{C}} \operatorname{err}_{\mathcal{D}}(h^*)$$
(1)

since any hypothesis has a  $\epsilon_0/2$  chance of getting the wrong answer in the noisy case and thus takes on  $\epsilon_0/2$  error in expectation. Then, since we know that our original concept  $h \in \mathcal{C}$ , we know that

$$\min_{h^* \in \mathcal{C}} \operatorname{err}_{\mathcal{D}}(h^*) = 0$$

and thus

$$\min_{h^* \in \mathcal{C}} \operatorname{err}_{\mathcal{D}'}(h^*) = \frac{\epsilon_0}{2}$$

Thus, we conclude

$$\operatorname{err}_{\mathcal{D}'}(h_A) \le \frac{\epsilon_0}{2} + \epsilon$$

as desired.

We have  $x_i$  are i.i.d Bernoulli random variables. We have  $X = \sum_{i=1}^k x_i$ . First we have that  $X \geq t$  implies  $e^{\lambda X} \geq e^{\lambda t}$  for  $\lambda \geq 0$ , which holds since  $e^z$  is monotone increasing over its whole domain and since  $\lambda$  does not change the signs of the exponent terms since  $\lambda \geq 0$ . The other direction also holds for the same reason. Therefore,  $\mathbf{P}\{X \geq t\} = \mathbf{P}\{e^{\lambda X} \geq e^{\lambda t}\} = \mathbf{P}\{e^{-\lambda t} \prod_{i=1}^k e^{\lambda x_i} \geq 1\}$ . The probability that this event occurs is bounded above by  $\mathbf{E}[e^{-\lambda t} \prod_{i=1}^k e^{\lambda x_i}]$  by Markov's inequality, since by taking exponents we made all values nonnegative. Then since  $x_i$  are independent, and therefore  $e^{\lambda x_i}$  are independent,  $\mathbf{E}[e^{-\lambda t} \prod_{i=1}^k e^{\lambda x_i}] = e^{-\lambda t} \prod_{i=1}^k \mathbf{E}[e^{\lambda x_i}]$ . Then  $x_i = 1$  or -1, each with probability  $\frac{1}{2}$ , so  $\mathbf{E}[e^{\lambda x_i}] = \frac{1}{2}(e^{\lambda} + e^{-\lambda})$ . Now we bound this quantity using Taylor Expansion  $e^z = \sum_{i=0}^\infty \frac{z^i}{i!}$ . Summing the Taylor expansions of  $e^{\lambda}$  and  $e^{-\lambda}$ , the odd terms cancel and the even terms are doubled. Diving by half, we get  $\frac{1}{2}(e^{\lambda} + e^{-\lambda}) = \sum_{i=0}^\infty \frac{\lambda^{2i}}{(2i)!}$ . Then we note that by Taylor expansion,  $e^{\frac{\lambda^2}{2}} = \sum_{i=0}^\infty \frac{\lambda^{2i}}{2^i i!}$ . Therefore we only need to show  $(2i)! \geq 2^i i! \to \frac{(2i)!}{i!} \geq 2^i$  to get our upper bound. Consider that  $(2i) * (2i-1) * \cdots * (i+1)$  has i terms, and each term is greater than or equal to 2. Therefore the inequality follows, and we have  $\mathbf{P}\{X \geq t\} \leq e^{-\lambda t} e^{\frac{k\lambda^2}{2}}$ . Then since t > 0, we can set  $\lambda = \frac{t}{k}$  to get  $\mathbf{P}\{X \geq t\} \leq e^{-\frac{t^2}{k}} e^{\frac{k\lambda^2}{2k^2}} = e^{-\frac{t^2}{2k}}$ , as desired.

(a) We prove that if a concept f is PAC-learnable by  $\mathcal{H}$ , then  $f \in \mathcal{H}$  for finite domain  $\mathcal{X}$ .

Since f is PAC-learnable by  $\mathcal{H}$ , there is an algorithm  $\mathcal{A}$  that returns a hypothesis h such that  $\operatorname{err}_{\mathcal{D}}(h) < \epsilon$ , where sample complexity  $m_{\mathcal{H}}(\epsilon, \delta) = \operatorname{poly}\left(\frac{1}{\epsilon}, \ln\left(\frac{1}{\delta}\right), \ln\left(|\mathcal{H}|\right)\right)$ . Since  $\mathcal{X}$  is finite, and the number of samples cannot exceed the domain, we have that  $|\mathcal{X}|$  is  $\operatorname{poly}\left(\frac{1}{\epsilon}\right)$ .

Therefore  $|\mathcal{X}| \sim \left(\frac{1}{\epsilon}\right)^c$  and  $\epsilon \leq \left(\frac{1}{|\mathcal{X}|}\right)^{\frac{1}{c}} \leq \frac{1}{|\mathcal{X}|}$  for  $c \geq 1$ . Since  $\operatorname{err}_{\mathcal{D}}(h) < \epsilon \leq \frac{1}{|\mathcal{X}|}$ ,  $\operatorname{err}_{\mathcal{D}}(h)$  must be 0, since the smallest non-zero error probability for finite domain  $|\mathcal{X}|$  is  $\frac{1}{|\mathcal{X}|}$ . Therefore, h = f and we have  $f \in \mathcal{H}$  as desired.

(b) We provide a counter-example  $(f, \mathcal{H}, \mathcal{X})$  to show that this theorem does not hold in general if  $\mathcal{X}$  has infinite cardinality. Let  $\mathcal{X} = \mathbb{R}$ , and let  $\mathcal{H} = \mathcal{H}_+$  be the hypothesis class of positive half-lines as defined in Lecture 2 (Definition 3.2), and define concept  $f: \mathcal{X} \to \{0,1\}$  as the 0 constant function; i.e.  $f(x) = 0 \ \forall x \in \mathbb{R}$ .

We have that  $f \notin \mathcal{H}_+$  since the constant function does not change its value and the positive half-line function necessarily does, therefore implying non-zero generalization error for all  $h \in \mathcal{H}$ . We now show that this concept is in fact PAC-learnable by  $\mathcal{H}_+$ . Let us choose the hypothesis h for a given sample set S as

$$h_r(x) = \begin{cases} 1 : x \ge r \\ 0 : x < r \end{cases} \text{ where } r = 1 + \max_{(x_i, y_i) \in \mathcal{S}}(x_i)$$

Then let the probability mass of  $\{x > r\}$  be  $\epsilon > 0$ . Note that this is the generalization error, since we should like to choose r at  $\infty$  (but there is no such perfect classifier). The probability that a sample lands in this bad set is  $\epsilon$ , therefore we have  $\mathbf{P}\{\operatorname{err}_{\mathcal{D}}(h_r) > \epsilon\} \leq (1-\epsilon)^{|\mathcal{S}|} \leq e^{-\epsilon|\mathcal{S}|} = \delta$ , and we get that for  $m_{\mathcal{H}}(\epsilon, \delta) \geq \frac{1}{\epsilon} \ln\left(\frac{1}{\delta}\right)$ , we have  $\operatorname{err}_{\mathcal{D}}(h_r) < \epsilon$  with probability  $1 - \delta$ , demonstrating that f is PAC-learnable even though  $f \notin \mathcal{H}$ .

We essentially repeat the proof of No Free Lunch from the book.

We will choose  $|\mathcal{X}| \geq \frac{m}{2\epsilon}$  for some  $\epsilon > 0$  to prove the statement. Then we can let m be any number smaller than  $|\mathcal{X}| * 2\epsilon$  representing a training set size. Let C be the subset of  $\mathcal{X}$  of size  $m/2\epsilon$ . Note there are  $T = 2^{\frac{m}{2\epsilon}}$  possible functions from  $C \to \{0,1\}$ . Denote these functions by  $f_1, \dots, f_T$ . For each function, let distribution  $\mathcal{D}_i$  over  $\mathcal{X} \times \{0,1\}$  be uniform (probability 1/|C|) for point-labels (x,y) such that  $f_i(x) = y$ , and 0 everywhere else. By this definition,  $\operatorname{err}_{\mathcal{D}_i}(f_i) = 0$ .

Now, we will prove that for every algorithm A that receives m examples from  $C \times \{0,1\}$  returns a function  $h_A: C \to \{0,1\}$  such that

$$\max_{i \in [T]} \mathbf{E}_{S \sim \mathcal{D}_i^m} [\operatorname{err}_{\mathcal{D}_i}(h_A)] \ge \frac{1}{2} - \epsilon$$

This result demonstrates that for every algorithm receiving m samples there is a  $(f_{i^*}, \mathcal{D}_{i^*})$  over which  $\operatorname{err}_{\mathcal{D}_{i^*}}(f_{i^*}) = 0$  and additionally for which

$$\mathbf{E}_{S \sim \mathcal{D}_{i^*}^m}[\operatorname{err}_{\mathcal{D}_{i^*}}(h_A)] \ge \frac{1}{2} - \epsilon$$

as desired.

Now we prove it. Note there are  $k = (m/2\epsilon)^m$  possible sequences of m examples from C. Denote these sequences by  $S_1, \dots, S_k$ . We use the index j to denote which sequence, and the index i to denote the fact that we drew from  $(f_i, \mathcal{D}_i)$ . Indexing over both we get all possible training sets  $\{S_i^i\}_{i\in[T],j\in[k]}$ . We have

$$\max_{i \in [T]} \mathbf{E}_{S \sim \mathcal{D}_i^m} [\operatorname{err}_{\mathcal{D}_i}(h_A)] = \max_{i \in [T]} \frac{1}{k} \sum_{j=1}^k \operatorname{err}_{\mathcal{D}_i}(h_A(S_j^i))$$

$$\geq \frac{1}{T} \sum_{i=1}^T \frac{1}{k} \sum_{j=1}^k \operatorname{err}_{\mathcal{D}_i}(h_A(S_j^i))$$

$$= \frac{1}{k} \sum_{j=1}^k \frac{1}{T} \sum_{i=1}^T \operatorname{err}_{\mathcal{D}_i}(h_A(S_j^i))$$

$$\geq \min_{j \in [k]} \frac{1}{T} \sum_{i=1}^T \operatorname{err}_{\mathcal{D}_i}(h_A(S_j^i))$$

$$(2)$$

Now fix  $j \in [k]$  as the minimum in the above, denote  $S_j = (x_1, \dots, x_m)$  and let us consider the samples  $v_1, \dots, v_p$  which do not appear in  $S_j$ . Since  $|S_j| = m/2\epsilon$ , we have that  $p = m/2\epsilon - m = (\frac{1}{2\epsilon} - 1) m \ge m$ . Therefore, for every  $h: C \to \{0, 1\}$ , since the error on a subset is  $\le$  the error on the whole set,

$$\operatorname{err}_{\mathcal{D}_{i}}(h) = \frac{2\epsilon}{m} \sum_{x \in C} \mathbf{1}\{h(x) \neq f_{i}(x)\}$$

$$\geq \frac{2\epsilon}{m} \sum_{r=1}^{p} \mathbf{1}\{h(v_{r}) \neq f_{i}(v_{r})\}$$

$$= \frac{2\epsilon \left(\frac{1}{2\epsilon} - 1\right)}{p} \sum_{r=1}^{p} \mathbf{1}\{h(v_{r}) \neq f_{i}(v_{r})\}$$

$$= \frac{1 - 2\epsilon}{p} \sum_{r=1}^{p} \mathbf{1}\{h(v_{r}) \neq f_{i}(v_{r})\}$$
(3)

Plugging this back in to what we had before, we get

$$\frac{1}{T} \sum_{i=1}^{T} \operatorname{err}_{\mathcal{D}_{i}}(h_{A}(S_{j}^{i})) \geq \frac{1}{T} \sum_{i=1}^{T} \frac{1 - 2\epsilon}{p} \sum_{r=1}^{p} \mathbf{1}\{h(v_{r}) \neq f_{i}(v_{r})\} 
= (1 - 2\epsilon) * \frac{1}{p} \sum_{r=1}^{p} \frac{1}{T} \sum_{i=1}^{T} \mathbf{1}\{h_{A}(S_{j}^{i})(v_{r}) \neq f_{i}(v_{r})\} 
\geq (1 - 2\epsilon) * \min_{r \in [p]} \frac{1}{T} \sum_{i=1}^{T} \mathbf{1}\{h_{A}(S_{j}^{i})(v_{r}) \neq f_{i}(v_{r})\}$$
(4)

Fix  $r \in [p]$  as the minimum. Now, we can partition the functions  $f_1, \dots, f_T$  into T/2 disjoint pairs where for pair of concepts  $(f_a, f_b)$  we have that they ONLY differ on the point  $v_r$ . That is  $f_a(v_r) \neq f_b(v_r)$ , and only at that point. Since  $v_r \notin S_j^a$  and  $v_r \notin S_j^b$  by definition, we have that  $S_j^a = S_j^b$ ; i.e. the training sets will be the same. Furthermore,  $h_A(S_j^a) = h_A(S_j^b)$  and if we denote pairs by  $(f_a, f_b)_i$ , we can just call this  $h_A(S_j^i)$ . Thus, we can rewrite

$$\frac{1}{T} \sum_{i=1}^{T} \mathbf{1} \{ h_A(S_j^i)(v_r) \neq f_i(v_r) \} = \frac{1}{T} \sum_{i=1}^{T/2} \mathbf{1} \{ h_A(S_j^i)(v_r) \neq f_{a_i}(v_r) \} + \mathbf{1} \{ h_A(S_j^i)(v_r) \neq f_{b_i}(v_r) \}$$

$$= \frac{1}{T} \sum_{i=1}^{T/2} 1$$

$$= \frac{1}{T} * \frac{T}{2} = \frac{1}{2}$$
(5)

since exactly one of the indicator functions in each sum term was 1. Thus, we can conclude that

$$(1 - 2\epsilon) * \min_{r \in [p]} \frac{1}{T} \sum_{i=1}^{T} \mathbf{1} \{ h_A(S_j^i)(v_r) \neq f_i(v_r) \} = (1 - 2\epsilon) * (1/2)$$

$$= \frac{1}{2} - \epsilon$$
(6)

as desired.

(a) We demonstrate that we cannot shatter the labelling (0,0,1,0).

Consider a sin curve with half-period  $k \in \mathbb{R}^+$  (this corresponds to taking  $\omega = \pi/k$ ) (WLOG consider  $\omega > 0$ , since  $-\sin(x) = \sin(-x)$  and we'll just flip all labels).

Now consider that we're trying to use this sin curve to classify

(x,0); (2x,0); (3x,1); (4x,0). Also note that x will only be below the sin curve on intervals of the form [2m\*k,(2m+1)\*k) for  $m \in \mathbb{Z}$ .

Suppose that  $x = (n + \delta) * k$ , where n is an even integer and  $\delta \in [0, 1)$  is the fractional part. Note that n must be even so that x is below the sin curve.

Now consider  $2x = 2(n+\delta)*k = (2n+2\delta)*k$ . We have 2n is even. Since  $2\delta < 2$ ,  $\lfloor 2\delta \rfloor = 0$  or 1. We require it to be 0 to maintain evenness so that 2x is under the sin curve as well (we want to get the classification (2x,0)). Therefore,  $2\delta < 1$  and  $\delta < 1/2$ .

Now consider  $3x = (3n + 3\delta) * k$ . Again since we assumed n is even, 3n is also even. This time,  $3\delta < 3/2$ , we want  $\lfloor 3\delta \rfloor = 1$  instead of 0 since even + odd = odd and since we want 3x to be above the sin curve. Therefore, we must have  $3\delta > 1$ , or  $\delta > 1/3$ .

Thus, we have established that to classify (x, 0); (2x, 0); (3x, 1), we need

$$1/3 < \delta < 1/2$$

Finally, to shatter (0,0,1,0), we need  $\lfloor 4n+4\delta \rfloor$  to be even. 4n is even, thus  $4\delta$  must be even as well. We have from before that

$$4/3 < 4\delta < 4/2 = 2$$

implying that  $4\delta$  is necessarily odd, a contradiction. Thus, we cannot shatter ((x,0),(2x,0),(3x,1),(4x,0)) with the family  $\{t \to \sin(\omega t)\}$ .

(b) We prove that we can shatter any configuration of points  $S_m = \{2^{-i} : i \in [1, ..., m]\}$  with  $\sin(\omega x)$  for any integer m > 1. This directly demonstrates that the VC-dimension of the class is infinite.

Let  $x_i = 2^{-i}$ , where i goes from 1 to some positive integer m > 1. Choose an arbitrary classification of these m points  $y \in \{0, 1\}^m$ , denoting the classification of  $x_i$  by  $y_i$ . Then, define

$$\omega = \pi \left( \sum_{i=1}^{m} y_i 2^{i-1} \right)$$

Note that our classifier is the sign of the sin function. We will show that  $sgn(sin(\omega x_i)) = 1 - 2y_i$  in all cases. Since y was arbitrary, we have demonstrated that we can shatter  $S_m$  for all m > 1.

Let us calculate the classification of  $x_j$ ,  $1 \le j \le m$ . Then,

$$\sin(\omega x_{j}) = \sin(\omega 2^{-j})$$

$$= \sin\left(\pi \left(\sum_{i=1}^{m} y_{i} 2^{i-1} 2^{-j}\right)\right)$$

$$= \sin\left(\pi y_{i} + \pi \sum_{i>j+1} y_{i} 2^{i-(j+1)} + \pi \sum_{i< j+1} y_{i} 2^{i-(j+1)}\right)$$
(7)

Note that the last sum in the sin function is necessarily an even factor of  $\pi$ . Using the identity  $\sin(x+2\pi)=\sin(x)$ , we can throw that term out. Thus we get

$$\sin(\omega x_j) = \sin\left(\pi y_i + \pi \sum_{i>j+1} y_i 2^{i-(j+1)}\right)$$

$$= \sin\left(\alpha \pi + \pi y_i\right)$$
(8)

defining the remaining sum term coefficient of  $\pi$  to be  $\alpha$ .

Now consider that

$$\alpha = \sum_{i>j+1} y_i 2^{i-(j+1)} \le \sum_{i=j+2}^m \frac{1}{2^{(j+1)-i}}$$

$$= \sum_{i=1}^{m-(j+1)} \left(\frac{1}{2}\right)^i$$

$$< \sum_{i=1}^{\infty} \left(\frac{1}{2}\right)^i = 1$$
(9)

Now we consider both cases  $y_j = 0$  and  $y_j = 1$ . If  $y_j = 0$ , then since  $0 < \alpha < 1$  we have that the argument is in the first or second quadrant, meaning sin is positive which corresponds with 1 - 2 \* 1 = -1. If  $y_j = 1$ , then we add an extra term of  $\pi$  to the argument of sin, and since  $\sin(x + \pi) = -\sin(x)$ , the sign changes to positive which corresponds to 1 - 2 \* 0 = 1, as desired.

Thus we have demonstrated that we can completely shatter any  $S_m$  for all m > 1.