

Prove or disprove every complete graph (2) of odd order is Hamilton cycle decomposible, Also determine the humber of Hamilton cycles is that decomposition.

Are lareides a complete graph Kn; where n is sold. Thus the degree of each vertex in the complete graph is even.

three, for a cycle - every vertex has degree 2 and thus; in seder to decompose a graph into transition cycles, the graph should have even degree. Therefore, every complete graph of odd seder is a transition cycle. Set 'n' he the vertices of complete graph & let 'n' he odd. Thus, the ro. of bousition cycles 'y n(n-1).

DIS a complete graph of even order hamilton cycle decomposible? Explain.

Be It is not decomposible.

Loreides a a complete graph Kn; where 'n' is even.

The degree of each vertex in the much is

the degree of each vertex in the graph is (n-1) which is odd. Now for a cycle to every vertex should have degree 2, which is even. But in this cope, degree of each vertex is odd. Herce, we cannot decompose a graph of even order intre Hamilton cycles.

3 Prove or disprove the following: (i) An Eulerian graph does not cut edges.

At set us offence that Eulerian graph has
cut edges on semonal of a particular
edge 'e' from the Euler graph G; the
graph becomes disconnected. But the end
westices of the graph 'G-e' have
add degree. Hence it could be

decomposed into eycler.

Therefore, it violates the aff conditions that is it is not a Eulerian graph. Hence on ossumption is wrong. Thus the Eulerian graph does not have cut edges.

cii) A hamilton graph does not ensure or have cut vertices.

And let us ossume that a transition graph boy out westices.

the removal of vertices from the graph, the graph becomes disconnected. Hence it does not cover all vertices of the graph.

Herce, it is not a Hamiltonian graph.

There is a contradiction be our assumption is wrong.

. Hamiltonian geaph does not have cut weter.