

ANT Algorithms

Ant Algorithms, (ACO: Ant Colony Optimization) is a meta heuristics for solving combinatorial ^{optimization} problems (hard ones)

(How Ants control where to go?)

→ Initialization: (need an admissible solution)

The Pheromone values are initialized to $\tau_{ij}(0)$, which maybe chosen to the user's discretion.

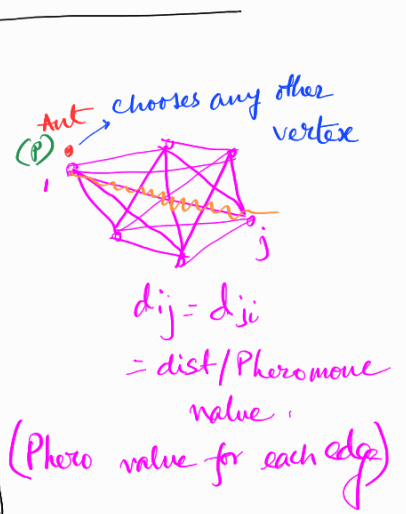
Ant P has already visited the nodes $N_p(t)$ (at time t)

(It is not allowed to visit same node again)

$$prob_{ij}^{(p)} = \begin{cases} \frac{(\tau_{ij})^\alpha (\eta_{ij})^\beta}{\sum_{s \notin N_p(t)} (\tau_{is})^\alpha (\eta_{is})^\beta} & \text{if } j \notin N_p(t) \text{ where } \alpha, \beta \in [0, 1] \\ 0 & \text{otherwise} \end{cases}$$

Prefer smaller edges

"Double Bridge" Experiment.



η_{ij} is called Visibility, or heuristic info.

(long edge is less visible, prefer short edges)

Here, we could use $\eta_{ij} = \frac{1}{d_{ij}}$

- Generate Solution candidates by ants, starting for ant P, with $N_p(0) = \emptyset$ (or $N_p(0) = \{\text{starting position}\}$)

At time $t+1$:
Ant P at position i goes to $j \notin N_p(t)$ with prob $prob_{ij}^{(p)}$
 $N_p(t+1) = N_p(t) \cup \{j\}$

- α, β control the influence of Pheromone values and heuristic info.

Fixed number Param.

• Update Pheromone:
 $\tau_{ij} = \tau_{ij} \cdot \rho$ if ant P goes from i to j between time $(t, t+1)$

• Ant Density : $g_{ij}^{(p)}(t, t+1) = \begin{cases} 0 & \text{otherwise} \end{cases}$

↳ Influence on edge pheromones (independent of length) spread out over whole edge.

• Ant Quantity : $g_{ij}^{(p)}(t, t+1) = \begin{cases} \frac{\alpha_2}{\text{dist}(ij)} & \text{if ant } p \text{ goes } \dots \\ 0 & \end{cases}$

• Ant Cycle : $g_{ij}^{(p)}(t, t+1) = \begin{cases} \frac{\alpha_3}{L^{(p)}} & \text{if ant } p \text{ goes } \dots \\ 0 & \end{cases}$

$L^{(p)}$ is length of ant p 's best solution found so far.

→ Pheromone update now,

$$\tau_{ij}(t+1) = \underline{(1-\rho)} \tau_{ij}^{(t)} + \Delta \tau_{ij}(t, t+1)$$

parameter ρ denotes the decay, or evaporation

$$\Delta \tau_{ij}(t, t+1) = \sum_{p \in P} g_{ij}^{(p)}(t, t+1)$$

Goal: To find Hamiltonian Path at back at So.

(Travel Salesman Problem)

✓
decay of already seen pheromone value

- Cannot find soln after one walk through Graph and returning
So need → Ant Cycle → after completion, update Phero values, and
restart Cycle while remembering old solution
(Ant always reach So after N steps where $N = \text{no of nodes}$)
- solves Min Spanning Tree prob, without being
aware of it, if we come close (in running time) to
special MST problem, then the solution is more general
than the algo that is not even invented.

