





## C

## **Q1** Team Name

0 Points

Goldfish

### **Q2** Commands

10 Points

List the commands used in the game to reach the ciphertext

go,enter,pick,c,back,give,back,back,thrnxxtzy,read

# **Q3** Analysis

50 Points

Give a detailed analysis of how you figured out the password? (Explain in less than 500 words)

After giving above commands , we got hints in form of group theory equations, which we solved in following ways to get the password.

prime modulus p = 455470209427676832372575348833

we had three equation in form of  $\,q^*g^{\wedge}a$  , where q is password, g is a element of the group and a is an integer

Given

 $q*g^{429}=431955503618234519808008749742 mod(p)$  -----eqn(1)

 $q*g^{1973}=176325509039323911968355873643 mod(p)$  -----eqn(2)

 $q * q^{7596} = 98486971404861992487294722613 mod(p)$  -----eqn(3)

# Assignment 3



GROUP

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View or edit group

TOTAL POINTS

70 / 70 pts

QUESTION 1

Team Name **0** / 0 pts

**QUESTION 2** 

Commands 10 / 10 pts

QUESTION 3

Analysis 50 / 50 pts

QUESTION 4

Password 10 / 10 pts

**QUESTION 5** 

we will refer  $q*g^a$  term in each equation as  $x_1,x_2,x_3$  respectively Dividing eqn(2) by eqn(1):  $g^{1544}=x_2*x_1^{-1}(\bmod p) ----- \mathrm{eqn(4)}$  Dividing eqn(3) by eqn(2):  $g^{5623}=x_3*x_2^{-1}(\bmod p) ----- \mathrm{eqn(5)}$  Dividing eqn(3) by eqn(1):  $g^{7167}=x_3*x_1^{-1}(\bmod p) ----- \mathrm{eqn(6)}$ 

To perform above operation we need to get the multiplicative inverses of  $x_1,x_2$ . as g is member of the multiplicative group, so  $x_1,x_2$  must also be members of the group and since p is prime  $x_1,x_2$  are coprime to p. Hence their multiplicative inverse also exist.

let  $y_1$  be the inverse of  $x_1$ 

Then,  $(x_1 * y_1) \bmod p = 1$ 

since p is a prime number , we can use Fermat little theorem to get the inverse  $x_1^{p-1} \equiv 1 (\bmod p)$ 

multiplying both side by  $x_1^{-1}$ 

$$x_1^{p-2}\equiv x_1^{-1}(\bmod p)=70749996790223471732904681640(\bmod p)$$
 from eqn(4) 
$$g^{1544}=70749996790223471732904681640(\bmod p)$$
 -----eqn(7)

similarly by solving eqn(5) and eqn(6) we get

$$g^{5623} = 420413074251022028027270785553 \pmod{p}$$
 ----eqn(8)

$$g^{7167}=110411376670918912626907526185 \pmod{p}$$
 ----eqn(9) Multiplying both side of eqn(9) by inverse of  $(g^{1544})^3$ 

$$g^{7167}*((g^{1544})^4)^{-1}(\bmod p)\equiv g^{991}=\\161798558270556961732424822635\bmod p$$

Similarly

$$(g^{991})^2 * ((g^{1544})^4)^{-1} (\bmod p) \equiv g^{438} = 327597482298082119695568192760 (\bmod p)$$

$$(g^{991}) * ((g^{438})^2)^{-1} (\text{mod } p) \equiv g^{115} = 212427760325417336316893638262 (\text{mod } p)$$

 $(g^{115})^4 * ((g^{438})^1)^{-1} (\bmod p) \equiv g^{22} = 62875864560156876567783127811 (\bmod p)$ 

 $(g^{22})^2 1 * ((g^{115})^4)^{-1} (\bmod p) \equiv g^2 = 108044907665466013935627786069 (\bmod p)$ 

 $(g^{115})^1*((g^{22})^5)^{-1}(\bmod p) \equiv g^5 = 254662155980870723273334022569(\bmod p)$ 

 $(g^5)^1*((g^2)^2)^{-1}(\bmod p)\equiv g=52565085417963311027694339(\bmod p)$  It is mentioned in the hints that g is 5\_\_\_50\_\_4\_\_\_31\_\_\_94\_\_9 , which actually matches with our answer.

Now we determine q (password) by putting value of g in eqn(1) which is  $q*g^{429}=431955503618234519808008749742 mod(p)$  We multiply both side of eqn(1) by  $(g^{429})^{-1}$ ,  $q=((g^{429})^{-1}*431955503618234519808008749742) mod(p)$  The multiplicative inverse of  $(g^{429})$  is 442956820316148690889301696615

That gives us value of q = password as 134721542097659029845273957

Hence the password is 134721542097659029845273957

## **Q4** Password

10 Points

What was the final command used to clear this level?

134721542097659029845273957

### **Q5** Codes

0 Points

Upload any code that you have used to solve this level

```
In [25]:
             def gcd(x,y):
                 if(x==0):
                     return y
                 return gcd(y%x,x)
In [26]:
             def po(a,b,m):
                 if(b==0):
                     return 1
                 p = po(a,b // 2,m)%m
                 p = (p*p)%m
                 if(b % 2 == 0):
                     return p
                 else:
                     return ((a*p) % m)
In [32]:
             def mI(a , m):
                 g = gcd(a, m)
                 if(g == 1):
                     y = po(a, m-2, m)
                     print("The modular inverse is ",y)
                     return y
                     print("Does not exist")
In [33]:
             y1 = 431955503618234519808008749742
             y2 = 176325509039323911968355873643
             y3 = 98486971404861992487294722613
             p = 455470209427676832372575348833
In [34]:
            y1_{inverse} = mI(y1,p)
             y2_{inverse} = mI(y2,p)
             The modular inverse is 70749996790223471732904681640
             The modular inverse is 228947149478752602606353685125
In [35]:
             g_{1544} = (y_{2} * y_{inverse})%p
             g_1544
Out [35]:
            111590994894663139264552154672
```

```
In [36]: g_7167 = (y3 * y1_inverse)\%p
             g_7167
Out [36]:
             110411376670918912626907526185
In [37]:
             g_{5623} = (y_3 * y_{inverse})%p
             g_5623
Out [37]:
             420413074251022028027270785553
In [38]:
             g_{991} = (g_{7167} * mI(po(g_{1544,4,p}),p))%p
             g_991
             The modular inverse is 304296090672599420401986286302
Out [38]:
            161798558270556961732424822635
In [39]:
             g_{438} = (po(g_{991,2},p) * mI(g_{1544,p}))%p
             g_438
             The modular inverse is 218173384175465437436454958180
Out [39]:
             327597482298082119695568192760
In [40]:
             g_{115} = (mI(po(g_{438,2,p}),p) * g_{991})%p
             g_115
             The modular inverse is 297374246948059676278983181681
Out [40]:
             212427760325417336316893638262
In [41]:
             g_22 = (mI(g_438,p) * po(g_115,4,p))%p
             g_22
             The modular inverse is 453530656410176241507046342872
Out [41]:
            62875864560156876567783127811
In [42]:
             g_2 = (mI(po(g_115,4,p),p) * po(g_22,21,p))%p
             g_2
```

```
The modular inverse is 105600931401330644523752113862
```

Out [42]: 108044907665466013935627786069

In [43]: 
$$g_5 = (mI(po(g_22,5,p),p)*g_115)%p$$
  
 $g_5$ 

The modular inverse is 100551735247242729663164535176

Out [43]: 254662155980870723273334022569

In [44]: 
$$g = (mI(po(g_2,2,p),p) * g_5)%p$$

The modular inverse is 254950689434017345885415339945

Out [44]: 52565085417963311027694339

In [45]: password = 
$$(y1*mI(po(g,429,p),p))%p$$
  
password

The modular inverse is 442956820316148690889301696615

Out [45]: 134721542097659029845273957