

AIFA

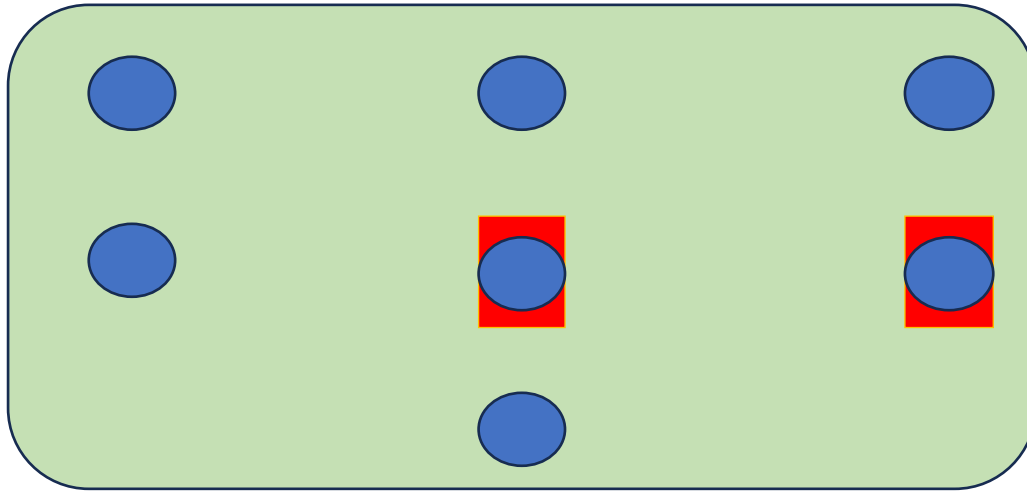
First Order Logic

30/01/2024

Koustav Rudra

Quantifiers

- There are two basic quantifiers in FOL
- \forall “for all” Universal Quantifier
- \exists “there exists” Existential Quantifier




Days in a week

$\forall_x \text{sunrise}(x)$

$\exists_x \text{holiday}(x)$

Universal Quantifiers

- All dogs are faithful
 - faithful(x): x is faithful
 - dog(x): x is a dog
 - $\forall_x(\text{dog}(x) \rightarrow \text{faithful}(x))$
- All birds cannot fly
 - fly(x): x can fly
 - bird(x): x is a bird
 -  $(\forall_x(\text{bird}(x) \rightarrow \text{fly}(x)))$

Existential Quantifiers

- At least one planet has life on it
 - planet(x): x is a planet
 - haslife(x): x has life on it
 - $\exists_x(\text{planet}(x) \wedge \text{haslife}(x))$
- There exists a bird that can't fly
 - fly(x): x can fly
 - bird(x): x is a bird
 - $\exists_x(\text{bird}(x) \wedge \sim \text{fly}(x))$

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Duality of Quantifiers

- All men are mortal
 - No man is immortal
- $\forall_x(\text{man}(x) \rightarrow \text{mortal}(x))$
- $\neg(\exists_x(\text{man}(x) \wedge \neg \text{mortal}(x)))$

Universal quantifiers could be expressed in a different way with existential quantifiers

Sentences

- A predicate is a sentence
- If sen , sen' are sentences and x is a variable then following are sentences
 - (sen) , $\sim\text{sen}$, $\exists_x \text{sen}$, $\forall_x \text{sen}$
 - $\text{sen} \wedge \text{sen}'$
 - $\text{sen} \vee \text{sen}'$
 - $\text{sen} \rightarrow \text{sen}'$


First-order Logic

- Sentence \rightarrow AtomicSentence
 - | Sentence Connective Sentence
 - Quantifier Variable, ... Sentence
 - \sim Sentence | (Sentence)
- AtomicSentence \rightarrow Predicate(Term, ...)
 - | Term = Term
- Term \rightarrow Function(Term, ...) | Constant | Variable
- Connective $\rightarrow \Rightarrow | \vee | \wedge | \Leftrightarrow$
- Quantifier $\rightarrow \forall | \exists$

Difference from second order logic

- In FOL we quantify only variables
- In second order logic we can quantify predicates
 - $\exists_P \forall_x \forall_y P(x, y)$

Examples

- Not all students take history and biology
 - Student(x): x is a student
 - Takes(x,y): subject x is taken by y
 -  $[\forall_x \text{Student}(x) \rightarrow \text{Takes}(\text{History}, x) \wedge \text{Takes}(\text{Biology}, x)]$
 - $\exists_x \text{Student}(x) \wedge [\sim \text{Takes}(\text{History}, x) \vee \sim \text{Takes}(\text{Biology}, x)]$
- Only one student failed biology
 - Failed(x,y): student y failed in subject x
 - $\exists_x [\text{Student}(x) \wedge \text{Failed}(\text{biology}, x) \wedge$
 - $\forall_y [(\sim(x=y) \wedge \text{Student}(y)) \rightarrow \sim \text{Failed}(\text{biology}, y)]]$

Examples

- Only one student failed both history and biology
 - Failed(x,y): student y failed in subject x
 - $\exists_x [\text{Student}(x) \wedge \text{Failed}(\text{biology},x) \wedge \text{Failed}(\text{history},x) \wedge$
 - $\forall_y [(\sim(x=y) \wedge \text{Student}(y)) \rightarrow \sim\text{Failed}(\text{biology},y) \vee \sim\text{Failed}(\text{history},y)]]$
- The best score in history is better than the best score in biology
 - Function: score(subject, student)
 - Greater(x,y): $x > y$
 - $\forall_x [\text{Student}(x) \wedge \text{Takes}(\text{biology},x) \rightarrow \exists_y [\text{Student}(y) \wedge \text{Takes}(\text{history},y) \wedge \text{Greater}(\text{score}(\text{history},y), \text{score}(\text{biology},x))]]$
 - $\exists_x [\text{Student}(x) \wedge \text{Takes}(\text{biology},x) \wedge \exists_y [\text{Student}(y) \wedge \text{Takes}(\text{history},y) \wedge \text{Greater}(\text{score}(\text{history},y), \text{score}(\text{biology},x))]]$

Example of sentences

- $\text{Brother}(x,y)$: x is y's brother
- $\text{Loves}(x,y)$: x loves y
- $\forall_x \forall_y \text{Brother}(x,y) \rightarrow \text{Loves}(x,y)$
 - Everyone loves his/her brothers
- A sentence is evaluated to be true in a particular domain, the sentence is said to be satisfied
- A sentence is valid in terms of FOL, for all possible domains this sentence will evaluate to be true

Takeaway

- Predicate logic is a more powerful version of proposition logic
- Provide support for variables and quantifiers
- Can capture sentences more naturally

How we can perform inferencing with predicates?

Inference in FOPL

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Inference Rules

- Universal Elimination
 - $\forall_x \text{Likes}(x, \text{flower})$
 - Substituting x by Akash $\{x/\text{Akash}\}$
 - $\text{Likes}(\text{Akash}, \text{flower})$
- The substitution should be done by a constant term

Inference Rules

- Existential elimination (Skolemization)
 - $\exists_x \text{Likes}(x, \text{flower}) \rightarrow \text{Likes}(\text{Person}, \text{flower})$
 - as long as Person is not in the knowledge base
- This method of finding out a particular constant that satisfies the predicate there exists x that likes flower is known as Skolemization.
- Existential introduction
 - Likes(John, flower)
 - Can be written as
 - $\exists_x \text{Likes}(x, \text{flower})$

Example1: Reasoning in FOL

- The law says that it is a crime for a Gaul to sell potion formulas to hostile nations.
- The country Rome, an enemy of Gaul, has acquired some potion formulas, and all of its formulas were sold to it by Druid Traitorix.
- Traitorix is a Gaul
- Is Traitorix a criminal?

Example1: Reasoning in FOL

- The law says that it is a crime for a Gaul to sell potion formulas to hostile nations.
- Gaul(x): x is a Gaul
- Hostile(z): z is hostile nation
- Potion(y): y is a potion
- Criminal(x): x is criminal
- Sells(x,y,z): x sells y to z
- Owns(x,y): x owns y
- $\forall_x \forall_y \forall_z \text{Gaul}(x) \wedge \text{Potion}(y) \wedge \text{Sells}(x,y,z) \wedge \text{hostile}(z) \rightarrow \text{Criminal}(x)$

Example1: Reasoning in FOL

- The country Rome, an enemy of Gaul, has acquired some potion formulas, and all of its formulas were sold to it by Druid Traitorix.
- Gaul(x): x is a Gaul
- Hostile(z): z is hostile nation
- Potion(y): y is a potion
- Criminal(x): x is criminal
- Sells(x,y,z): x sells y to z
- Owns(x,y): x owns y
- Hostile(Rome)
- $\exists_y \text{Potion}(y) \wedge \text{Owns}(\text{Rome}, y)$
- $\forall_y \text{Potion}(y) \wedge \text{Owns}(\text{Rome}, y) \rightarrow \text{Sells}(\text{Traitorix}, y, \text{Rome})$

Example1: Reasoning in FOL

- Traitorix is a Gaul
- Gaul(x): x is a Gaul
- Hostile(z): z is hostile nation
- Potion(y): y is a potion
- Criminal(x): x is criminal
- Sells(x,y,z): x sells y to z
- Owns(x,y): x owns y
- Gaul(Traitorix)
- ?- Criminal(Traitorx) [We have to deduce it]

Example1: Reasoning in FOL

- $\forall_x \forall_y \forall_z \text{Gaul}(x) \wedge \text{Potion}(y) \wedge \text{Sells}(x,y,z) \wedge \text{hostile}(z) \rightarrow \text{Criminal}(x)$
- $\text{Hostile}(\text{Rome})$
- $\exists_y \text{Potion}(y) \wedge \text{Owns}(\text{Rome},y)$
- $\forall_y \text{Potion}(y) \wedge \text{Owns}(\text{Rome},y) \rightarrow \text{Sells}(\text{Traitorix}, y, \text{Rome})$
- $\text{Gaul}(\text{Traitorix})$
- ?- $\text{Criminal}(\text{Traitorix})$

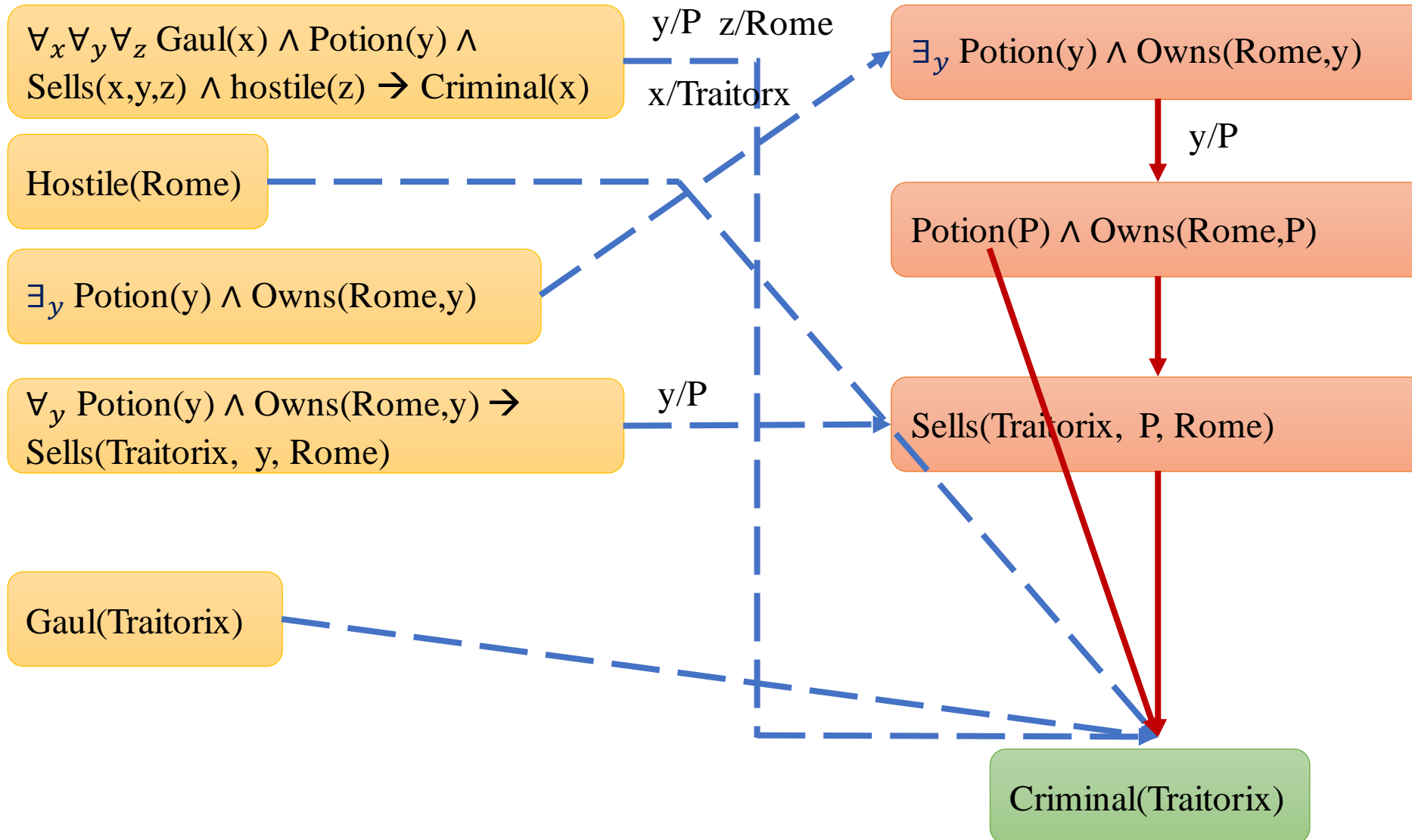
Forward Chaining

Started from existing facts and rules in knowledge base and reach the goal

Backward Chaining

Started from the goal and use existing facts and rules in knowledge base to prove goal

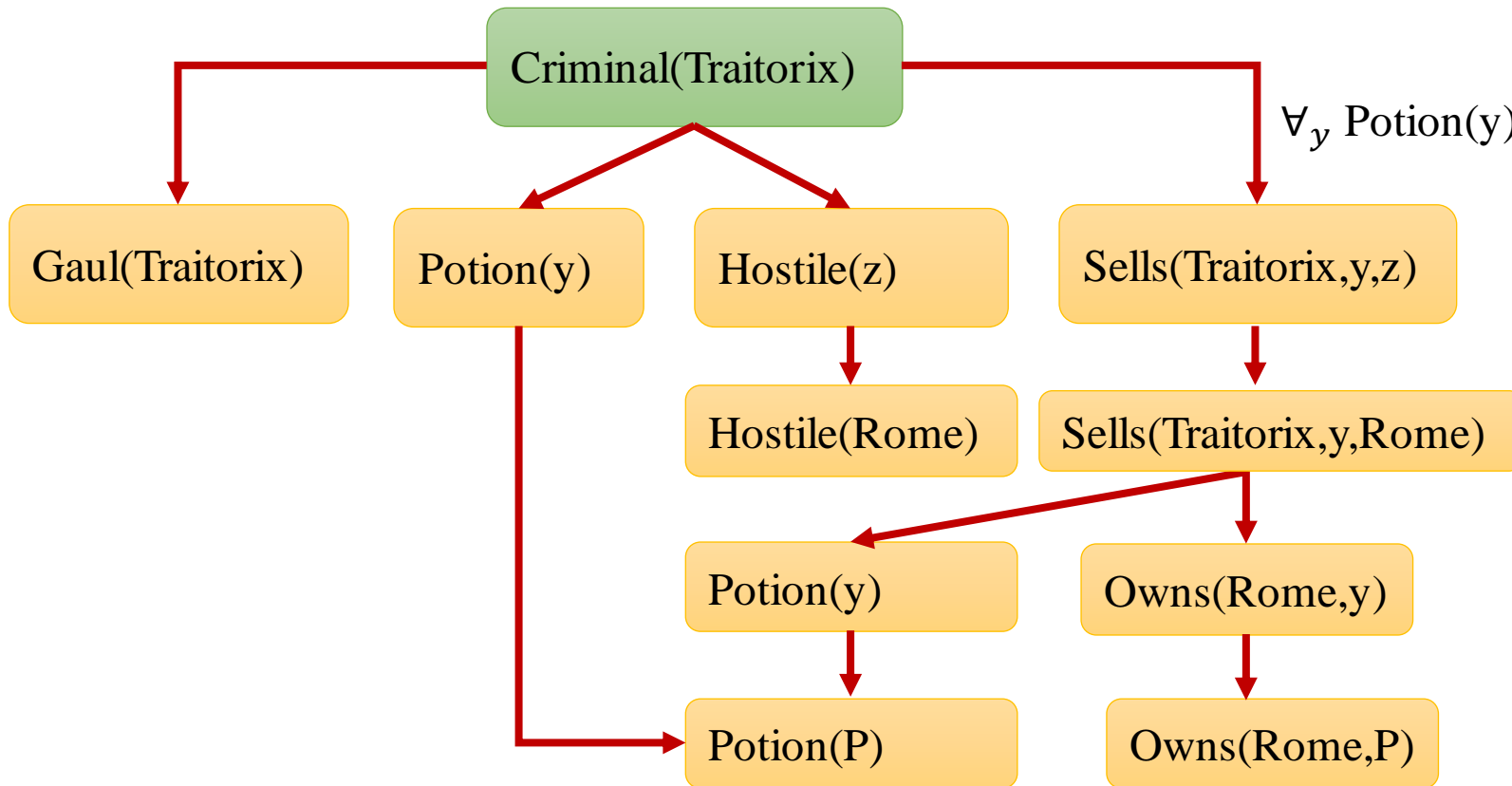
Reasoning in FOL: Forward Chaining



Reasoning in FOL: Backward Chaining

- $\forall_x \forall_y \forall_z \text{Gaul}(x) \wedge \text{Potion}(y) \wedge \text{Sells}(x,y,z) \wedge \text{hostile}(z) \rightarrow \text{Criminal}(x)$
- $\forall_y \forall_z \text{Gaul}(\text{Traitorix}) \wedge \text{Potion}(y) \wedge \text{Sells}(\text{Traitorix},y,z) \wedge \text{hostile}(z) \rightarrow \text{Criminal}(\text{Traitorix})$

Hostile(Rome)



$\forall_y \text{Potion}(y) \wedge \text{Owns}(\text{Rome},y) \rightarrow \text{Sells}(\text{Traitorix}, y, \text{Rome})$

$\exists_y \text{Potion}(y) \wedge \text{Owns}(\text{Rome},y)$

Thank You