

AIFA A^* Analysis

11/01/2024

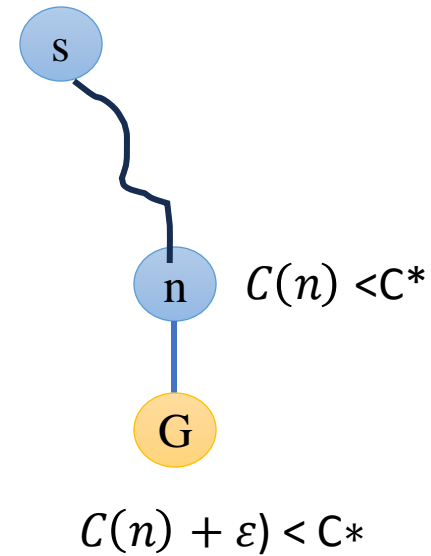
Koustav Rudra

Algorithm A*: Benefit

- Reduces number of expanded nodes
- Performs the lookahead and tells us promising paths
- What about optimality?

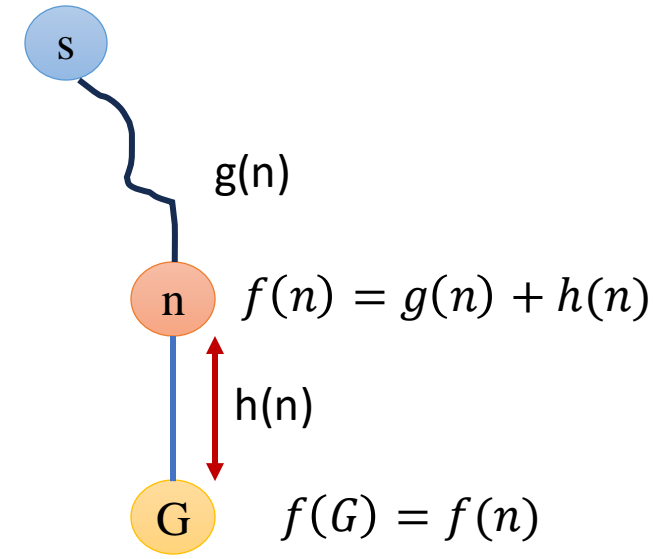
Uniform Cost Search

- **Claim:** If $C(n) < C^*$ (optimal cost) then n must be expanded
- Let algorithm A does not expand n
- For the class of algorithms without any heuristics
 - All states that have cost $< C^*$ will have to be expanded



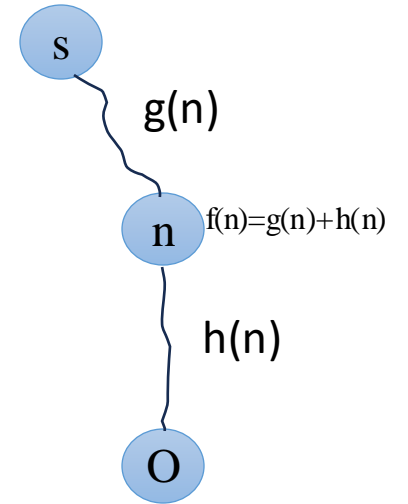
Algorithm A*: Benefit

- **Claim:** $f(n) < C^*$ then n must be expanded
- The heuristic function underestimates
 - $h(n) \leq h^*(n)$
 - Cost of reaching goal from n
 - All costs are +ve
- If we do not expand n, we can't find the goal
- If we have a state whose cost is less than C^*
 - Then every algorithm which guarantees finding optimal solution have to expand it



Algorithm A*

- **Initialize:** Set $OPEN = \{s\}$, $CLOSED = \{\}$, $g(s)=0$, $f(s) = h(s)$
- **Fail:**
 - If $OPEN = \{\}$, Terminate with failure
- **Select:** Select the minimum cost state, n , from $OPEN$ and save in $CLOSED$
- **Terminate:**
 - If $n \in G$, terminate with success

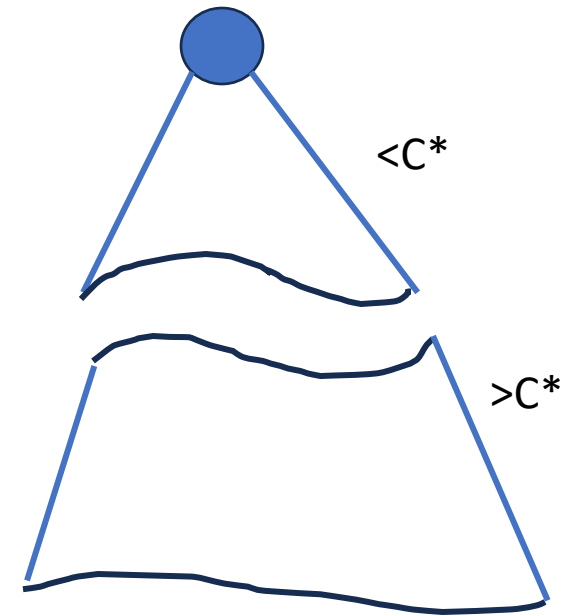


How to break the tie?

A*: Result

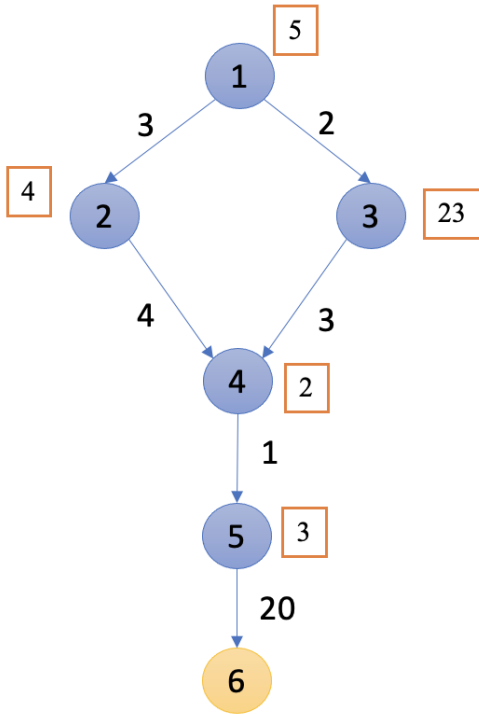
- What is an admissible heuristic?

- If it always underestimates
- We always have $h(n) \leq h^*(n)$, where $h^*(n)$ denotes minimum distance to a goal state from n



A*: Result

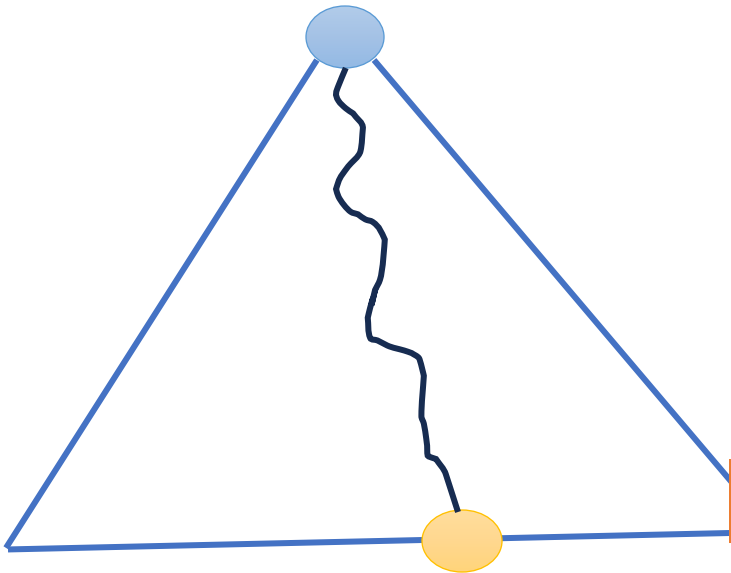
- At any time before A* terminates, there exists in OPEN a state **n**
 - That is on an optimal path from s to a goal state, with
 - $f(n) \leq f^*(s)$



OPEN SET	SELECT	GOAL	EXPANDED	CLOSED
[3(25),5(11)]	5(11)	N	[3(25),6(28)]	[1(5),2(7),4(9),5(11)]
[3(25),6(28)]	3(25)	N	[6(28),4(7)]	[1(5),2(7),4(9),5(11),3(25)]
[6(28),4(7)]	4(7)	N	[6(28),5(9)]	[1(5),2(7),4(9),5(11),3(25),4(7)]
[6(28),5(9)]	5(9)	N	[6(26)]	[1(5),2(7),4(9),5(11),3(25),4(7),5(9)]
[6(26)]	6(26)	Y		

A*: Result

- At any time before A* terminates, there exists in OPEN a state **n**
 - That is on an optimal path from s to a goal state, with
 - $f(n) \leq f^*(s)$



OPEN SET	SELECT	GOAL	EXPANDED	CLOSED
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[3(25),6(28)]	3(25)	N	[6(28),4(7)]	[1(5),2(7), 4(9) ,5(11),3(25)]
[6(28),4(7)]	4(7)	N	[6(28),5(9)]	[1(5),2(7), 4(9) , 5(11) ,3(25),4(7)]
[6(28),5(9)]	5(9)	N	[6(26)]	[1(5),2(7), 4(9) , 5(11) ,3(25),4(7),5(9)]
[6(26)]	6(26)	Y		

A*: Result

- If there is a path from s to a goal state, A* terminates
- What is the worst case scenario?
- Approximation algorithm

A*: Result

- Algorithm A* is admissible
 - If there is a path from s to a goal state,
 - A* terminates by finding an optimal path

A* Optimality

- Suppose some **suboptimal goal path G2** has been generated and is in OPEN
- Let **n** be an **unexpanded node** in OPEN such that **n** is on a **shortest path** to an **optimal goal G**

$$h(G_2) = h(G) = 0$$

$$f(G_2) = g(G_2) \quad f(G) = g(G)$$

$$g(G_2) > g(G) \quad G_2 \text{ is suboptimal}$$

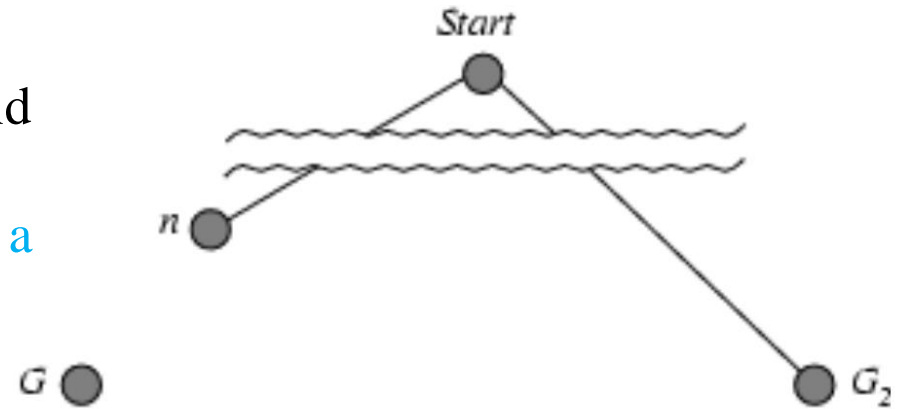
$$f(G_2) > f(G)$$

$$h(n) \leq h^*(n) \quad h \text{ is admissible, } h^* \text{ is minimal distance}$$

$$g(n) + h(n) \leq g(n) + h^*(n)$$

$$f(n) < f^*(G)$$

$$f(G_2) > f(n) \quad \text{A* will never select } G_2 \text{ for expansion}$$



A*: Result

- Algorithm A* is admissible
 - If there is a path from s to a goal state,
 - A* terminates by finding an optimal path
- If A1 and A2 are two versions of A* such that A2 is more informed than A1
 - A1 expands at least as many states as does A2

A*: Result

- We have two good heuristic functions h_1 and h_2 but do not know which one is more informed
- How much effort should we put in computing heuristics?

Thank You