

For this homework assignment, I will be using **LINDO** to calculate the problems.

1). A). Find the distance of the shortest path from G to C in the graph below.

Objective function is max c as to find the shortest path to c from g, and the constraints are the code between ST and END below from the screenshot in which they are the distances from one vertex to another. For example, $h-g \leq 3$ is interpreted as the distance from h to g is 3.

The shortest path from g to c is 16.

Copy of Code

```
max c
ST
    g=0
    h-g<=3
    d-g<=2
    g-e<=7
    e-d<=25
    d-e<=9
    e-b<=10
    b-h<=9
    a-h<=4
    e-f<=2
    b-a<=8
    a-f<=5
    f-a<=10
    b-f<=7
    c-b<=4
    c-f<=3
    f-d<=18
    d-c<=3
END
```

Result (Output)

LP OPTIMUM FOUND AT STEP | 6

OBJECTIVE FUNCTION VALUE

1) 16.000000

LP Code

max c
ST

g=0
 h-g<=3
 d-g<=2
 g-e<=7
 e-d<=25
 d-e<=9
 e-b<=10
 b-h<=9
 a-h<=4
 e-f<=2
 b-a<=8
 a-f<=5
 f-a<=10
 b-f<=7
 c-b<=4
 c-f<=3
 f-d<=18
 d-c<=3

END

VARIABLE	VALUE	REDUCED COST
C	16.000000	0.000000
G	0.000000	0.000000
H	3.000000	0.000000
D	0.000000	0.000000
E	0.000000	0.000000
B	12.000000	0.000000
A	4.000000	0.000000
F	13.000000	0.000000

ROW	SLACK OR SURPLUS	DUAL PRICES
2)	0.000000	1.000000
3)	0.000000	1.000000
4)	2.000000	0.000000
5)	7.000000	0.000000
6)	25.000000	0.000000
7)	9.000000	0.000000
8)	22.000000	0.000000
9)	0.000000	1.000000
10)	3.000000	0.000000
11)	15.000000	0.000000
12)	0.000000	0.000000
13)	14.000000	0.000000
14)	1.000000	0.000000
15)	8.000000	0.000000
16)	0.000000	1.000000
17)	0.000000	0.000000
18)	5.000000	0.000000
19)	19.000000	0.000000

NO. ITERATIONS= 6

B). Find the distances of the shortest paths from G to all other vertices.

The objective function is $\max a+b+c+d+e+f+h$ to find the shortest paths from g to all the vertices. Although we will get the sum of all the shortest paths, but that doesn't matter because we only care about the shortest path in each vertex from g, which is underneath the objective function value in the picture below. The constraints the code between ST and END below from the screenshot in which they are the distances from one vertex to another. For example, $h-g \leq 3$ is interpreted as the distance from h to g is 3.

The shortest path from G to all the other vertices:

A	B	C	D	E	F	H
7	12	16	2	19	17	3

The value "76" is meaningless because that is the sum of all the shortest paths, but the shortest paths have been listed below from A to H.

Copy of Code

$\max a+b+c+d+e+f+h$

ST

$g=0$

$h-g \leq 3$

$d-g \leq 2$

$g-e \leq 7$

$e-d \leq 25$

$d-e \leq 9$

$e-b \leq 10$

$b-h \leq 9$

$a-h \leq 4$

$e-f \leq 2$

$b-a \leq 8$

$a-f \leq 5$

$f-a \leq 10$

$b-f \leq 7$

$c-b \leq 4$

$c-f \leq 3$

$f-d \leq 18$

$d-c \leq 3$

END

Result (Output)

LP OPTIMUM FOUND AT STEP 4

OBJECTIVE FUNCTION VALUE

1) 76.000000

VARIABLE	VALUE	REDUCED COST
A	7.000000	0.000000
B	12.000000	0.000000
C	16.000000	0.000000
D	2.000000	0.000000
E	19.000000	0.000000
F	17.000000	0.000000
H	3.000000	0.000000
G	0.000000	0.000000

LP Code

```
max a+b+c+d+e+f+h
ST
```

```
g=0
h-g<=3
d-g<=2
g-e<=7
e-d<=25
d-e<=9
e-b<=10
b-h<=9
a-h<=4
e-f<=2
b-a<=8
a-f<=5
f-a<=10
b-f<=7
c-b<=4
c-f<=3
f-d<=18
d-c<=3
```

ROW	SLACK OR SURPLUS	DUAL PRICES
2)	0.000000	8.000000
3)	0.000000	6.000000
4)	0.000000	1.000000
5)	26.000000	0.000000
6)	8.000000	0.000000
7)	26.000000	0.000000
8)	3.000000	0.000000
9)	0.000000	2.000000
10)	0.000000	3.000000
11)	0.000000	1.000000
12)	3.000000	0.000000
13)	15.000000	0.000000
14)	0.000000	2.000000
15)	12.000000	0.000000
16)	0.000000	1.000000
17)	4.000000	0.000000
18)	3.000000	0.000000
19)	17.000000	0.000000

END

NO. ITERATIONS= 4

2). Formulate the problem as a linear program with an objective function and all constraints. Determine the optimal solution for the linear program using any software you want. Include a copy of the code and output. What are the optimal numbers of ties of each type to maximize profit?

Type	Selling Price	Labor Cost	Material Cost	Profit
Silk	6.7	0.75	2.5	3.45
Polyester	3.55	0.75	0.48	2.32
Blend 1	4.31	0.75	0.75	2.81
Blend 2	4.81	0.75	0.81	3.25

Formulate the problem as linear program

max $3.45s + 2.32p + 2.81b + 3.25c$

ST

$0.125s \leq 1000$: This is for Silk

$0.08p + 0.05b + 0.03c \leq 2000$: This is for Polyester

$0.05b + 0.07c \leq 1250$: This is for Cotton

$s \leq 7000$; $s \geq 6000$

$p \leq 14000$; $p \geq 10000$

$b \leq 16000$; $b \geq 13000$

$c \leq 8500$; $c \geq 6000$

Copy of Code

max $3.45s + 2.32p + 2.81b + 3.25c$

ST

$0.125s \leq 1000$

$0.08p + 0.05b + 0.03c \leq 2000$

$0.05b + 0.07c \leq 1250$

$s \leq 7000$

$s \geq 6000$

$p \leq 14000$

$p \geq 10000$

$b \leq 16000$

$b \geq 13000$

$c \leq 8500$

$c \geq 6000$

END

Result (output)

LP OPTIMUM FOUND AT STEP 4

OBJECTIVE FUNCTION VALUE

1) 120196.0

LP Code

VARIABLE	VALUE	REDUCED COST
S	7000.000000	0.000000
P	13625.000000	0.000000
B	13100.000000	0.000000
C	8500.000000	0.000000

ROW	SLACK OR SURPLUS	DUAL PRICES
2)	125.000000	0.000000
3)	0.000000	29.000000
4)	0.000000	27.200001
5)	0.000000	3.450000
6)	1000.000000	0.000000
7)	375.000000	0.000000
8)	3625.000000	0.000000
9)	2900.000000	0.000000
10)	100.000000	0.000000
11)	0.000000	0.476000
12)	2500.000000	0.000000

NO. ITERATIONS= 4

```

max 3.45s + 2.32p + 2.81b + 3.25c
ST
    0.125s <= 1000
    0.08p + 0.05b + 0.03c <= 2000
    0.05b + 0.07c <= 1250
s <= 7000
s >= 6000
p <= 14000
p >= 10000
b <= 16000
b >= 13000
c <= 8500
c >= 6000
END

```

The maximum profit is \$120,196.00 from producing 7,000 units of silk ties, 13,625 units of polyester ties, 13,100 units of Blend 1, and 8,500 units of Blend 2.

3). A). Determine the number of refrigerators to be shipped from the plants to the warehouses and then warehouses to retailers to minimize the cost. Formulate the problem as a linear program with an objective function and all constraints. Determine the optimal solution for the linear program using any software you want. What are the optimal shipping routes and minimum cost?

Objective Function:

Minimize transshipment costs (Z) = $10P1W1 + 15P1W2 + 11P2W1 + 8P2W2 + 13P3W1 + 8P3W2 + 9P3W3 + 14P4W2 + 8P4W3 + 5W1R1 + 6W1R2 + 7W1R3 + 10W1R4 + 12W2R3 + 8W2R4 + 10W2R5 + 14W2R6 + 14W3R4 + 12W3R5 + 12W3R6 + 6W3R7$

Constraints:

Constraints on Plants (P):

$P1W1 + P1W2 \leq 150$
 $P2W1 + P2W2 \leq 450$
 $P3W1 + P3W2 + P3W3 \leq 250$
 $P4W2 + P4W3 \leq 150$

Constraints on Retailers (R):

$W1R1 \geq 100$
 $W1R2 \geq 150$
 $W1R3 + W2R3 \geq 100$
 $W1R4 + W2R4 + W3R4 \geq 200$
 $W2R5 + W3R5 \geq 200$
 $W2R6 + W3R6 \geq 150$
 $W3R7 \geq 100$

Constraints on Warehouses (W) since nothing should be stored in the end

$W1R1 + W1R2 + W1R3 + W1R4 - P1W1 - P2W1 - P3W1 = 0$
 $W2R3 + W2R4 + W2R5 + W2R6 - P1W2 - P2W2 - P3W2 - P4W2 = 0$
 $W3R4 + W3R5 + W3R6 + W3R7 - P3W3 - P4W3 = 0$

Non-negativity Constraints:

$P1W1 \geq 0$	$W1R1 \geq 0$	$W3R5 \geq 0$
$P1W2 \geq 0$	$W1R2 \geq 0$	$W3R6 \geq 0$
$P2W1 \geq 0$	$W1R3 \geq 0$	$W3R7 \geq 0$
$P2W2 \geq 0$	$W1R4 \geq 0$	
$P3W1 \geq 0$	$W2R3 \geq 0$	
$P3W2 \geq 0$	$W2R4 \geq 0$	
$P3W3 \geq 0$	$W2R5 \geq 0$	
$P4W2 \geq 0$	$W2R6 \geq 0$	
$P4W3 \geq 0$	$W3R4 \geq 0$	

LP Code

```

min 10P1W1 + 15P1W2 + 11P2W1 + 8P2W2 + 13P3W1 + 8P3W2 + 9P3W3 + 14P4W2 + 8P4W3 + 5W1R1 + 6W1R2 + 7W1R3 + 10W1R4 + 12W2R3 + 8W2R4 + 10W2R5 + 14W2R6 + 14W3R4 + 12W3R5 + 12W3R6 + 6W3R7
ST
    P1W1 + P1W2 <= 150
    P2W1 + P2W2 <= 450
    P3W1 + P3W2 + P3W3 <= 250
    P4W2 + P4W3 <= 150
    W1R1 >= 100
    W1R2 >= 150
    W1R3 + W2R3 >= 100
    W1R4 + W2R4 + W3R4 >= 200
    W2R5 + W3R5 >= 200
    W2R6 + W3R6 >= 150
    W3R7 >= 100
    W1R1 + W1R2 + W1R3 + W1R4 - P1W1 - P2W1 - P3W1 = 0
    W2R3 + W2R4 + W2R5 + W2R6 - P1W2 - P2W2 - P3W2 - P4W2 = 0
    W3R4 + W3R5 + W3R6 + W3R7 - P3W3 - P4W3 = 0
    P1W1 >= 0
    P1W2 >= 0
    P2W1 >= 0
    P2W2 >= 0
    P3W1 >= 0
    P3W2 >= 0
    P3W3 >= 0
    P4W2 >= 0
    P4W3 >= 0
    W1R1 >= 0
    W1R2 >= 0
    W1R3 >= 0
    W1R4 >= 0
    W2R3 >= 0
    W2R4 >= 0
    W2R5 >= 0
    W2R6 >= 0
    W3R4 >= 0
    W3R5 >= 0
    W3R6 >= 0
    W3R7 >= 0
END

```

Copy of Code

$\min 10P1W1 + 15P1W2 + 11P2W1 + 8P2W2 + 13P3W1 + 8P3W2 + 9P3W3 + 14P4W2 + 8P4W3 + 5W1R1$
 $+ 6W1R2 + 7W1R3 + 10W1R4 + 12W2R3 + 8W2R4 + 10W2R5 + 14W2R6 + 14W3R4 + 12W3R5 + 12W3R6$
 $+ 6W3R7$

ST

$P1W1 + P1W2 \leq 150$
 $P2W1 + P2W2 \leq 450$
 $P3W1 + P3W2 + P3W3 \leq 250$
 $P4W2 + P4W3 \leq 150$
 $W1R1 \geq 100$
 $W1R2 \geq 150$
 $W1R3 + W2R3 \geq 100$
 $W1R4 + W2R4 + W3R4 \geq 200$
 $W2R5 + W3R5 \geq 200$
 $W2R6 + W3R6 \geq 150$
 $W3R7 \geq 100$

$W1R1 + W1R2 + W1R3 + W1R4 - P1W1 - P2W1 - P3W1 = 0$
 $W2R3 + W2R4 + W2R5 + W2R6 - P1W2 - P2W2 - P3W2 - P4W2 = 0$
 $W3R4 + W3R5 + W3R6 + W3R7 - P3W3 - P4W3 = 0$

$P1W1 \geq 0$
 $P1W2 \geq 0$
 $P2W1 \geq 0$
 $P2W2 \geq 0$
 $P3W1 \geq 0$
 $P3W2 \geq 0$
 $P3W3 \geq 0$
 $P4W2 \geq 0$
 $P4W3 \geq 0$

W1R1 >= 0
 W1R2 >= 0
 W1R3 >= 0
 W1R4 >= 0
 W2R3 >= 0
 W2R4 >= 0
 W2R5 >= 0
 W2R6 >= 0
 W3R4 >= 0
 W3R5 >= 0
 W3R6 >= 0
 W3R7 >= 0

END

Result (Output)

The minimum cost is **\$17,100**

Route	Quantity	Price per item	Cost
P1W1	150	10	1500
P2W1	200	11	2200
P2W2	250	8	2000
P3W2	150	8	1200
P3W3	100	9	900
P4W3	150	8	1200
W1R1	100	5	500
W1R2	150	6	900
W1R3	100	7	700
W2R4	200	8	1600
W2R5	200	10	2000
W3R6	150	12	1800
W3R7	100	6	600
Total	2000		17,100

LP OPTIMUM FOUND AT STEP 13

OBJECTIVE FUNCTION VALUE

1) 17100.00

VARIABLE	VALUE	REDUCED COST
P1W1	150.000000	0.000000
P1W2	0.000000	8.000000
P2W1	200.000000	0.000000
P2W2	250.000000	0.000000
P3W1	0.000000	2.000000
P3W2	150.000000	0.000000
P3W3	100.000000	0.000000
P4W2	0.000000	7.000000
P4W3	150.000000	0.000000
W1R1	100.000000	0.000000
W1R2	150.000000	0.000000
W1R3	100.000000	0.000000
W1R4	0.000000	5.000000
W2R3	0.000000	2.000000
W2R4	200.000000	0.000000
W2R5	200.000000	0.000000
W2R6	0.000000	1.000000
W3R4	0.000000	7.000000
W3R5	0.000000	3.000000
W3R6	150.000000	0.000000
W3R7	100.000000	0.000000

ROW	SLACK OR SURPLUS	DUAL PRICES
2)	0.000000	1.000000
3)	0.000000	0.000000
4)	0.000000	0.000000
5)	0.000000	1.000000
6)	0.000000	-16.000000
7)	0.000000	-17.000000
8)	0.000000	-18.000000
9)	0.000000	-16.000000
10)	0.000000	-18.000000
11)	0.000000	-21.000000
12)	0.000000	-15.000000
13)	0.000000	11.000000
14)	0.000000	8.000000
15)	0.000000	9.000000
16)	150.000000	0.000000
17)	0.000000	0.000000
18)	200.000000	0.000000
19)	250.000000	0.000000
20)	0.000000	0.000000
21)	150.000000	0.000000
22)	100.000000	0.000000
23)	0.000000	0.000000
24)	150.000000	0.000000
25)	100.000000	0.000000
26)	150.000000	0.000000
27)	100.000000	0.000000
28)	0.000000	0.000000
29)	0.000000	0.000000
30)	200.000000	0.000000
31)	200.000000	0.000000
32)	0.000000	0.000000
33)	0.000000	0.000000
34)	0.000000	0.000000
35)	150.000000	0.000000
36)	100.000000	0.000000

NO. ITERATIONS= 13

3). B). Due to old infrastructure Warehouse 2 is going to close eliminating all of the associated routes. What is the optimal solution for this modified model? Is it feasible to ship all the refrigerators to either warehouse 1 or 3 and then to the retailers without using warehouse 2? Why or why not?

This modified model is not feasible or no solution to that because, according to the graph, there are lots of retailers like R5, R6, and R7 cannot be reached through warehouse 1 but only 3 if warehouse 2 is closed. Since P1 and P2 can only access to warehouse 1, those retailers cannot be reached without warehouse 2. It also breaks the constraint for R5 if the warehouse 2 is closed because if only P3 and P4 are able supply R5, R6, and R7 with their supply capacity of 400 in combined, they will not be enough to supply all those retailers. P3 and P4 have a combined supply capacity of 400 but R5, R6, and R7 have a combined demand of 450. Though P3 and P4 can supply R6 and R7 sufficiently, $(400-250=150)$, that leaves the combined supply capacity with only 150 but R5 demands for 200, therefore without warehouse 2, P3 and P4 will not be able to supply the combined demand of R5, R6, and R7. We can also test this hypothesis by modify the code and run it on LINDO, the program resulted an error.

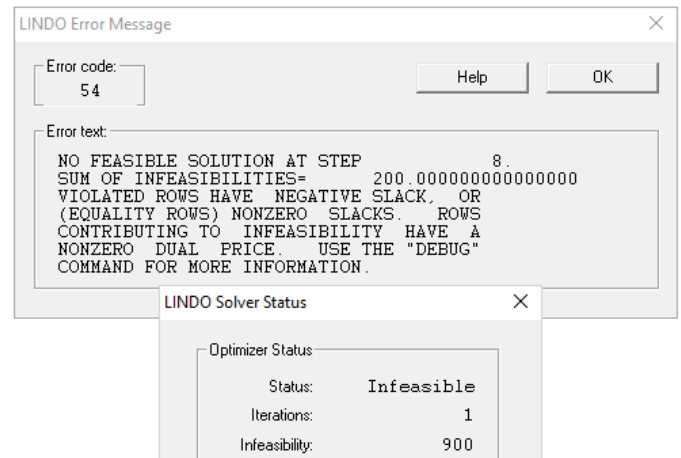
Here's the output result:

```

min 10P1W1 + 11P2W1 + 13P3W1 + 9P3W3 + 8P4W3 + 5W1R1 + 6W1R2 + 7W1R3 + 10W1R4 + 14W3R4 + 12W3R5 + 12W3R6 + 6W3R7
ST
    P1W1 <= 150
    P2W1 <= 450
    P3W1 + P3W3 <= 250
    P4W3 <= 150
    W1R1 >= 100
    W1R2 >= 150
    W1R3 >= 100
    W1R4 + W3R4 >= 200
    W3R5 >= 200
    W3R6 >= 150

    W1R1 + W1R2 + W1R3 + W1R4 - P1W1 - P2W1 - P3W1 = 0
    W3R4 + W3R5 + W3R6 + W3R7 - P4W3 = 0
    P1W1 >= 0
    P2W1 >= 0
    P3W1 >= 0
    P3W3 >= 0
    P4W3 >= 0
    W1R1 >= 0
    W1R2 >= 0
    W1R3 >= 0
    W1R4 >= 0
    W3R4 >= 0
    W3R5 >= 0
    W3R6 >= 0
    W3R7 >= 0
END

```



Copy of Code

```

min 10P1W1 + 11P2W1 + 13P3W1 + 9P3W3 + 8P4W3 + 5W1R1 + 6W1R2 + 7W1R3 + 10W1R4 +
14W3R4 + 12W3R5 + 12W3R6 + 6W3R7
ST
    P1W1 <= 150
    P2W1 <= 450
    P3W1 + P3W3 <= 250
    P4W3 <= 150
    W1R1 >= 100
    W1R2 >= 150
    W1R3 >= 100
    W1R4 + W3R4 >= 200
    W3R5 >= 200
    W3R6 >= 150
    W3R7 >= 100

```

$$W1R1 + W1R2 + W1R3 + W1R4 - P1W1 - P2W1 - P3W1 = 0$$

$$W3R4 + W3R5 + W3R6 + W3R7 - P3W3 - P4W3 = 0$$

$$P1W1 \geq 0$$

$$P2W1 \geq 0$$

$$P3W1 \geq 0$$

$$P3W3 \geq 0$$

$$P4W3 \geq 0$$

$$W1R1 \geq 0$$

$$W1R2 \geq 0$$

$$W1R3 \geq 0$$

$$W1R4 \geq 0$$

$$W3R4 \geq 0$$

$$W3R5 \geq 0$$

$$W3R6 \geq 0$$

$$W3R7 \geq 0$$

END

3). C). Instead of closing Warehouse 2 management has decide to keep a portion of it open but limit shipments to 100 refrigerators per week. Is this feasible? If so what is the optimal solution when warehouse 2 is limited to 100 refrigerators?

The modified model is feasible and the optimal solution for the minimized cost when warehouse 2 is limited to 100 refrigerators is 18,300. Here are the results for the modified model with LINDO:

LP Code

```
min 10P1W1 + 15P1W2 + 11P2W1 + 8P2W2 + 13P3W1 + 8P3W2 + 9P3W3 + 14P4W2 + 8P4W3 + 5W1R1 + 6W1R2 + 7W1R3 + 10W1R4 + 12W2R3 + 8W2R4 + 10W2R5 + 14W2R6 + 14W3R4 + 12W3R5 + 12W3R6 + 6W3R7
ST
P1W1 + P1W2 <= 150
P2W1 + P2W2 <= 450
P3W1 + P3W2 + P3W3 <= 250
P4W2 + P4W3 <= 150
W1R1 >= 100
W1R2 >= 150
W1R3 + W2R3 >= 100
W1R4 + W2R4 + W3R4 >= 200
W2R5 + W3R5 >= 200
W2R6 + W3R6 >= 150
W3R7 >= 100

W1R1 + W1R2 + W1R3 + W1R4 - P1W1 - P2W1 - P3W1 = 0
W2R3 + W2R4 + W2R5 + W2R6 - P1W2 - P2W2 - P3W2 - P4W2 = 0
W3R4 + W3R5 + W3R6 + W3R7 - P3W3 - P4W3 = 0

P1W2 + P2W2 + P3W2 + P4W2 <= 100
W2R3 + W2R4 + W2R5 + W2R6 <= 100

P1W1 >= 0
P1W2 >= 0
P2W1 >= 0
P2W2 >= 0
P3W1 >= 0
P3W2 >= 0
P3W3 >= 0
P4W2 >= 0
P4W3 >= 0
W1R1 >= 0
W1R2 >= 0
W1R3 >= 0
W1R4 >= 0
W2R3 >= 0
W2R4 >= 0
W2R5 >= 0
W2R6 >= 0
W3R4 >= 0
W3R5 >= 0
W3R6 >= 0
W3R7 >= 0

END
```

Copy of Code

min 10P1W1 + 15P1W2 + 11P2W1 + 8P2W2 + 13P3W1 + 8P3W2 + 9P3W3 + 14P4W2 + 8P4W3 + 5W1R1 + 6W1R2 + 7W1R3 + 10W1R4 + 12W2R3 + 8W2R4 + 10W2R5 + 14W2R6 + 14W3R4 + 12W3R5 + 12W3R6 + 6W3R7

ST

P1W1 + P1W2 <= 150
P2W1 + P2W2 <= 450
P3W1 + P3W2 + P3W3 <= 250
P4W2 + P4W3 <= 150
W1R1 >= 100
W1R2 >= 150
W1R3 + W2R3 >= 100
W1R4 + W2R4 + W3R4 >= 200
W2R5 + W3R5 >= 200
W2R6 + W3R6 >= 150
W3R7 >= 100

W1R1 + W1R2 + W1R3 + W1R4 - P1W1 - P2W1 - P3W1 = 0
W2R3 + W2R4 + W2R5 + W2R6 - P1W2 - P2W2 - P3W2 - P4W2 = 0
W3R4 + W3R5 + W3R6 + W3R7 - P3W3 - P4W3 = 0

P1W2 + P2W2 + P3W2 + P4W2 <= 100
W2R3 + W2R4 + W2R5 + W2R6 <= 100

P1W1 >= 0

```
P1W2 >= 0
P2W1 >= 0
P2W2 >= 0
P3W1 >= 0
P3W2 >= 0
P3W3 >= 0
P4W2 >= 0
P4W3 >= 0
W1R1 >= 0
W1R2 >= 0
W1R3 >= 0
W1R4 >= 0
W2R3 >= 0
W2R4 >= 0
W2R5 >= 0
W2R6 >= 0
W3R4 >= 0
W3R5 >= 0
W3R6 >= 0
W3R7 >= 0
END
```

Result (Output)

OBJECTIVE FUNCTION VALUE

1) 18300.00

VARIABLE	VALUE	REDUCED COST
P1W1	150.000000	0.000000
P1W2	0.000000	8.000000
P2W1	350.000000	0.000000
P2W2	100.000000	0.000000
P3W1	0.000000	4.000000
P3W2	0.000000	2.000000
P3W3	250.000000	0.000000
P4W2	0.000000	9.000000
P4W3	150.000000	0.000000
W1R1	100.000000	0.000000
W1R2	150.000000	0.000000
W1R3	100.000000	0.000000
W1R4	150.000000	0.000000
W2R3	0.000000	7.000000
W2R4	50.000000	0.000000
W2R5	50.000000	0.000000
W2R6	0.000000	4.000000
W3R4	0.000000	4.000000
W3R5	150.000000	0.000000
W3R6	150.000000	0.000000
W3R7	100.000000	0.000000

ROW	SLACK OR SURPLUS	DUAL PRICES
2)	0.000000	1.000000
3)	0.000000	0.000000
4)	0.000000	2.000000
5)	0.000000	3.000000
6)	0.000000	-16.000000
7)	0.000000	-17.000000
8)	0.000000	-18.000000
9)	0.000000	-21.000000
10)	0.000000	-23.000000
11)	0.000000	-23.000000
12)	0.000000	-17.000000
13)	0.000000	11.000000
14)	0.000000	8.000000
15)	0.000000	11.000000
16)	0.000000	0.000000
17)	0.000000	5.000000
18)	150.000000	0.000000
19)	0.000000	0.000000
20)	350.000000	0.000000
21)	100.000000	0.000000
22)	0.000000	0.000000
23)	0.000000	0.000000
24)	250.000000	0.000000
25)	0.000000	0.000000
26)	150.000000	0.000000
27)	100.000000	0.000000
28)	150.000000	0.000000
29)	100.000000	0.000000
30)	150.000000	0.000000
31)	0.000000	0.000000
32)	50.000000	0.000000
33)	50.000000	0.000000
34)	0.000000	0.000000
35)	0.000000	0.000000
36)	150.000000	0.000000
37)	150.000000	0.000000
38)	100.000000	0.000000

NO. ITERATIONS= 14

4). A). $V = [1, 5, 10, 25]$ and $A = 202$.

Objective function is the min of $V1 + V2 + V3 + V4$. The constraints are they should be non-negatives and $1V1 + 5V2 + 10V3 + 25V4 = 202$. The minimum number of coins for making a change is 10 with two coins of 1's and eight coins of 25's.

Copy of Code

min $V1 + V2 + V3 + V4$

ST

$1V1 + 5V2 + 10V3 + 25V4 = 202$

$V1 \geq 0$

$V2 \geq 0$

$V3 \geq 0$

$V4 \geq 0$

END

GIN V1

GIN V2

GIN V3

GIN V4

Result (Output)

LP Code		OBJECTIVE FUNCTION VALUE		
		1)	10.00000	
min	$V1 + V2 + V3 + V4$	VARIABLE	VALUE	REDUCED COST
ST		V1	2.000000	1.000000
	$1V1 + 5V2 + 10V3 + 25V4 = 202$	V2	0.000000	1.000000
	$V1 \geq 0$	V3	0.000000	1.000000
	$V2 \geq 0$	V4	8.000000	1.000000
	$V3 \geq 0$	ROW	SLACK OR SURPLUS	DUAL PRICES
	$V4 \geq 0$	2)	0.000000	0.000000
END		3)	2.000000	0.000000
GIN V1		4)	0.000000	0.000000
GIN V2		5)	0.000000	0.000000
GIN V3		6)	8.000000	0.000000
GIN V4		NO. ITERATIONS= 63		
		BRANCHES= 12 DETERM.= 1.000E 0		

B). $V = [1, 3, 7, 12, 27]$ and $A = 293$

Objective function is the min of $V1 + V2 + V3 + V4 + V5$. The constraints are they should be non-negatives and $1V1 + 3V2 + 7V3 + 12V4 + 27V5 = 293$. The minimum number of coins for making a change is 14 with two coins of 7's, three coins of 12's, and nine coins of 27's.

Copy of Code

```
min V1 + V2 + V3 + V4 + V5
ST
    1V1 + 3V2 + 7V3 + 12V4 + 27V5 = 293
    V1>=0
    V2>=0
    V3>=0
    V4>=0
    V5>=0
END
GIN V1
GIN V2
GIN V3
GIN V4
GIN V5
```

Result (Output)

LP Code		OBJECTIVE FUNCTION VALUE		
		1)	14.000000	
min	V1 + V2 + V3 + V4 + V5	VARIABLE	VALUE	REDUCED COST
ST		V1	0.000000	1.000000
	1V1 + 3V2 + 7V3 + 12V4 + 27V5 = 293	V2	0.000000	1.000000
	V1>=0	V3	2.000000	1.000000
	V2>=0	V4	3.000000	1.000000
	V3>=0	V5	9.000000	1.000000
	V4>=0			
	V5>=0	ROW	SLACK OR SURPLUS	DUAL PRICES
END		2)	0.000000	0.000000
GIN V1		3)	0.000000	0.000000
GIN V2		4)	0.000000	0.000000
GIN V3		5)	2.000000	0.000000
GIN V4		6)	3.000000	0.000000
GIN V5		7)	9.000000	0.000000
		NO. ITERATIONS=	98	
		BRANCHES=	34 DETERM.= 1.000E	0