

Transducers

October 16, 2023

1 Transducers

1.1 Data processing inequality

The output Y of a finite state transducer driven by a finite state statistical source (=a stationary and ergodic Markov process) X , is a finite state statistical source with information rate (=entropy per unit time) $R(Y)$ that less than or equal to that of the input entropy $R(X)$

$$R(Y) \leq R(X).$$

If the transducer is non-singular, then the equality holds

$$R(Y) = R(X).$$

1.1.1 Information rate and channel capacity

Assume that the transducer generates a binary sample Y every $T_s = 1ms$. The maximum information rate equals $R_{max} = 1/T_s$ bits per sample or equivalently bits per seconds if all these samples are independent and with equal probability for a ‘0’ and a ‘1’. Hence, the maximal information rate $R_{max} = 1/T_s$ bits\$.

If the transducer is non-singular (= reversible), then the data processing inequality learns us that the channel capacity

$$C = R_{max}(X) = R_{max}(Y) = 1/T_s$$

1.1.2 Fixed code length (=block code)

Consider the bineary encoding with the following fixed code length

i	code	$p(x_i)$
1	00	0.4
2	01	0.3
3	10	0.2
4	11	0.1

Note that this transducer is non-singular.

Compute the symbol length, the average symbol length of the binary source, and the information rate expressed in bit/sample (where sample represents a single bit in the binary stream).

```
p_i = [0.4 0.3 0.2 0.1]
H_X = 1.84644 bit/symbol
N_avg = 2.00000 sample/symbol
R_X = 0.92322 bit/sample
```

Compute the duration of the symbol, the average symbol duration and the information rate expressed in bit/s.

```
tau_avg = 0.00200 s/symbol
R_X = 923.21967 bit/s
```

Compute the absolute redundancy

$$R_X = C - R(X)$$

the relative redundancy

$$r_X = 1 - \frac{R(X)}{C}$$

and the efficiency

$$e_X = \frac{R(X)}{C}.$$

Absolute redundancy: $R_X = 76.78033$ bit/s

Relative redundancy: $r_X = 7.7\%$

Efficiency: $e_X = 92.3\%$

1.1.3 Determine the probability of 0 and 1

First consider N symbols (with N tending to infinity, e.g. $N = 1000$). For these N symbols, the count the number of zeros and the estimated length. The ratio of the number of zeros and the estimated length is then the probability of encounterint zero. (The propability of a one is 1 minus this probability).

$P[0]=0.65$

1.1.4 Variable length code

Consider the bineary encoding with the following fixed code length

i	code	$p(x_i)$
1	0	0.4
2	10	0.3
3	110	0.2
4	111	0.1

This is also a non-singular transducer as this represent a pre-fix code (proof: see later).

Compute the symbol length, the average symbol length of the binary source, and the information

rate expressed in bit/sample.

```
N_avg = 1.90000 sample/symbol  
R_X = 0.97181 bit/sample
```

Compute the duration of the symbol, the average symbol duration and the information rate expressed in bit/s.

```
tau_avg = 0.00190 sample/s  
R_X = 971.81018 bit/s
```

Compute the absolute redundancy, the relative redundancy, and the efficiency.

```
Absolute redundancy: R_X = 28.18982 bit/s  
Relative redundancy: r_X = 2.8%  
Efficiency: e_X = 97.2%
```

1.1.5 Determine the probability of 0 and 1

First consider N symbols (with N tending to infinity, e.g. $N = 1000$). For these N symbols, the count the number of zeros and the estimated length. The ratio of the number of zeros and the estimated length is then the propability of encounterint zero. (The propability of a one is 1 minus this probability).

P[0]=0.4736842105263156