















3. Let us assume H exists

1. Let R(x) be a function that x is a function and $\forall x (H(x,x) = 1 \Leftrightarrow time(R(x)) = \infty) \land (H(x,x) = 0 \Leftrightarrow time(R(x)) < \infty)$ 1. Let R(x) be a function that x is a function and x time(x) $= 0 \Leftrightarrow x$ such that x

3. Let us assume time $(R(R))=\infty$ a. H(R,R)=1 by 3.1
b. time $(R(R))<\infty$ by 2
c. time $(R(R))=\infty$ \(\text{time}(R(R))<00\)
d. This is a contradiction.

4. time(R(R)) \ \neq \infty

5. Let us assume time $(R(R)) < \infty$.

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6. H(R,R) = 0 by 3.1

6. time $(R(R)) = \infty$ by 2

6. time $(R(R)) = \infty$ 1 time $(R(R)) < \infty$ 6. This is a contradiction.

6. time $(R(R)) = \infty$ 7. time $(R(R)) = \infty$ 1 time $(R(R)) \neq \infty$ by 3.4 and 3.6 8. This is a contradiction

4. H does not exist

Q.E.D

BUT I DO NOT UNDERSTAND THE SYNTAX OF YOUR LANGUAGE! Alright here goes ... It seems that you suppose you have a program f(p,i) do not understand these wathomatical that takes another program p and constructs. Let me describe the proof set of input(s) i. in pseudocode then. such that f(p,i) returns 1 if the program halts and O otherwise. Suppose we have the partial function * g(p) that returns 0 if f(p.p) == 0 and executes an infinite loop otherwise... NOTICE THAT procedure compute_g(i) g is partially computable if +(i,i) == 0 then return O Je such that e is else while (i) a program that computes pass. the partical function g(ie. e = = compute g)



