

# FOLLOW THE REGULARIZED LEADER

- To overcome the shortcomings of the FTL algorithm, one can incorporate a regularization function  $\psi : \mathcal{A} \rightarrow \mathbb{R}_+$  into the action choice of FTL, which leads to more stability.
- To be more precise, let for  $t \geq 1$

$$a_t^{\text{FTRL}} \in \arg \min_{a \in \mathcal{A}} \left( \psi(a) + \sum_{s=1}^{t-1} \langle a, z_s \rangle \right),$$

(Technical side note: if there are more than one minimum, then one of them is chosen.)

(Technical side note: if there are more than one minimum, then one of them is chosen.)

then the algorithm choosing  $a_t^{\text{FTRL}}$  in time step  $t$  is called the **Follow the regularized leader (FTRL)** algorithm.



# FOLLOW THE REGULARIZED LEADER

- To overcome the shortcomings of the FTL algorithm, one can incorporate a regularization function  $\psi : \mathcal{A} \rightarrow \mathbb{R}_+$  into the action choice of FTL, which leads to more stability.
- To be more precise, let for  $t \geq 1$

$$a_t^{\text{FTRL}} \in \arg \min_{a \in \mathcal{A}} \left( \psi(a) + \sum_{s=1}^{t-1} \ell(a, z_s) \right),$$

(Technical side note: if there are more than one minimum, then one of them is chosen.)

then the algorithm choosing  $a_t^{\text{FTRL}}$  in time step  $t$  is called the **Follow the regularized leader (FTRL)** algorithm.

- **Interpretation:** The algorithm predicts  $a_t$  as the element in  $\mathcal{A}$ , which minimizes the regularization function plus the cumulative loss so far over the previous  $t - 1$  time periods.



# FOLLOW THE REGULARIZED LEADER

- To overcome the shortcomings of the FTL algorithm, one can incorporate a regularization function  $\psi : \mathcal{A} \rightarrow \mathbb{R}_+$  into the action choice of FTL, which leads to more stability.
- To be more precise, let for  $t \geq 1$

$$a_t^{\text{FTRL}} \in \arg \min_{a \in \mathcal{A}} \left( \psi(a) + \sum_{s=1}^{t-1} \ell(a, z_s) \right),$$

(Technical side note: if there are more than one minimum, then one of them is chosen.)

then the algorithm choosing  $a_t^{\text{FTRL}}$  in time step  $t$  is called the **Follow the regularized leader (FTRL)** algorithm:

- **Interpretation:** The algorithm predicts  $a_t$  as the element in  $\mathcal{A}$ , which minimizes the regularization function plus the cumulative loss so far over the previous  $t - 1$  time periods.
- Obviously, the behavior of the FTRL algorithm is depending heavily on the choice of the regularization function  $\psi$ . If  $\psi \equiv 0$ , then FTRL equals FTL:



# REGRET ANALYSIS OF FTRL: A HELPFUL LEMMA

- Lemma:** Let  $a_1^{\text{FTRL}}, a_2^{\text{FTRL}}, \dots$  be the sequence of actions coming used by the FTRL algorithm for the environmental data sequence,  $z_1, z_2, \dots$ . Then, for all  $\tilde{a} \in \mathcal{A}$  we have

$$\begin{aligned}
 R_T^{\text{FTRL}}(\tilde{a}) &= \sum_{t=1}^T ((a_t^{\text{FTRL}}, z_t) - (\tilde{a}, z_t)) \\
 &\leq \psi(\tilde{a}) - \psi(a_1^{\text{FTRL}}) + \sum_{t=1}^{T-1} ((a_t^{\text{FTRL}}, z_t) - (a_{t+1}^{\text{FTRL}}, z_t)).
 \end{aligned}$$



# REGRET ANALYSIS OF FTRL: A HELPFUL LEMMA

- Lemma:** Let  $a_1^{\text{FTRL}}, a_2^{\text{FTRL}}, \dots$  be the sequence of actions coming used by the FTRL algorithm for the environmental data sequence,  $z_1, z_2, \dots$ . Then, for all  $\tilde{a} \in \mathcal{A}$  we have

$$\begin{aligned}
 R_T^{\text{FTRL}}(\tilde{a}) &= \sum_{t=1}^T ((a_t^{\text{FTRL}}, z_t) - (\tilde{a}, z_t)) \\
 &\leq \psi(\tilde{a}) - \psi(a_1^{\text{FTRL}}) + \sum_{t=1}^T ((a_t^{\text{FTRL}}, z_t) - (a_{t+1}^{\text{FTRL}}, z_t)).
 \end{aligned}$$

- Interpretation:* the regret of the FTRL algorithm is bounded by the difference of cumulated losses of itself compared to its one-step lookahead cheater version and an additional regularization difference term.

$\Rightarrow$  We have seen an analogous result for FTL!

(The proof is similar.)





# FTRL FOR ONLINE LINEAR OPTIMIZATION

- In the following, we analyze the FTRL algorithm for the linear loss  $\ell(a, z) = a^\top z$  for online linear optimization (OLO) problems.
- For this purpose, the squared L2-norm regularization will be used:

$$\psi(a) = \frac{1}{2\eta} \|a\|_2^2 = \frac{a^\top a}{2\eta},$$

where  $\eta$  is some positive scalar, the *regularization magnitude*.

- It is straightforward to compute that if  $\mathcal{A} = \mathbb{R}^d$ , then

$$a_t^{\text{FTRL}} = -\eta \sum_{s=1}^{t-1} z_s.$$



# FTRL FOR ONLINE LINEAR OPTIMIZATION

- In the following, we analyze the FTRL algorithm for the linear loss  $\ell(a, z_t) = a^\top z_t$  for online linear optimization (OLO) problems.
- For this purpose, the squared L2-norm regularization will be used:

$$\psi(a) = \frac{1}{2\eta} \|a\|_2^2 = \frac{a^\top a}{2\eta},$$

where  $\eta$  is some positive scalar, the *regularization magnitude*.

- It is straightforward to compute that if  $\mathcal{A} = \mathbb{R}^d$ , then

$$a_t^{\text{FTRL}} = -\eta \sum_{s=1}^{t-1} z_s.$$

- Hence, in this case we have for the FTRL algorithm the following update rule

$$a_{t+1}^{\text{FTRL}} = a_t^{\text{FTRL}} - \eta z_t, \quad t = 1, \dots, T-1.$$





# FTRL FOR ONLINE LINEAR OPTIMIZATION

- In the following, we analyze the FTRL algorithm for the linear loss  $\ell(a; z) = a^\top z$  for online linear optimization (OLO) problems.
- For this purpose, the squared L2-norm regularization will be used:

$$\psi(a) = \frac{1}{2\eta} \|a\|_2^2 = \frac{a^\top a}{2\eta},$$

where  $\eta$  is some positive scalar, the *regularization magnitude*.

- It is straightforward to compute that if  $\mathcal{A} = \mathbb{R}^d$ , then

$$a_t^{\text{FTRL}} = -\eta \sum_{s=1}^{t-1} z_s.$$

- Hence, in this case we have for the FTRL algorithm the following update rule

$$a_{t+1}^{\text{FTRL}} = a_t^{\text{FTRL}} - \eta z_t, \quad t = 1, \dots, T-1.$$

*Interpretation:*  $-z_t$  is the direction in which the update of  $a_t^{\text{FTRL}}$  to  $a_{t+1}^{\text{FTRL}}$  is conducted with step size  $\eta$  in order to reduce the loss.



# FTRL FOR OLO: THEORETICAL GUARANTEES

- **Proposition:** Using the FTRL algorithm with the squared L2-norm regularization on any online linear optimization (OLO) problem with  $\mathcal{A} \subset \mathbb{R}^d$  leads to a regret of FTRL with respect to any action  $\tilde{a} \in \mathcal{A}$  of

$$R_T^{\text{FTRL}}(\tilde{a}) \leq \frac{1}{2\eta} \|\tilde{a}\|_2^2 + \eta \sum_{t=1}^T \|z_t\|_2^2.$$

- We will show the result only for the case  $\mathcal{A} = \mathbb{R}^d$ .
- For the more general case, where  $\mathcal{A}$  is a strict subset of  $\mathbb{R}^d$ , we need a slight modification of the update formula above:

$$a_t^{\text{FTRL}} = \Pi_{\mathcal{A}} \left( \eta \sum_{i=1}^{t-1} z_i \right) = \arg \min_{a \in \mathcal{A}} \left\| a - \eta \sum_{i=1}^{t-1} z_i \right\|_2^2.$$

In words, the action of the FTRL algorithm has to be projected onto the set  $\mathcal{A}$ . Here,  $\Pi_{\mathcal{A}} : \mathbb{R}^d \rightarrow \mathcal{A}$  is the projection onto  $\mathcal{A}$ .

(The proof is essentially the same, except that the Cauchy-Schwarz inequality is used in between.)



# FTRL FOR OLO: THEORETICAL GUARANTEES

- Proof:

**Reminder (1):** 
$$R_T^{\text{FTRL}}(\tilde{a}) \leq \psi(\tilde{a}) - \psi(a_1^{\text{FTRL}}) + \sum_{t=1}^T \sum_{t=1}^T \left( \ell(a_t^{\text{FTRL}}, z_t) - \ell(a_{t+1}^{\text{FTRL}}, z_t) \right).$$

**Reminder (2):** 
$$a_{t+1}^{\text{FTRL}} = a_t^{\text{FTRL}} - \eta z_t \eta z_t, \quad t = 1, \dots, T-1.$$

- For sake of brevity, we write  $a_1, a_2, \dots$  for  $a_1^{\text{FTRL}}, a_2^{\text{FTRL}}, \dots$



# FTRL FOR OLO: THEORETICAL GUARANTEES

- Proof:

**Reminder (1):**  $R_T^{\text{FTRL}}(\tilde{a}) \leq \psi(\tilde{a}) - \psi(a_1) + \sum_{t=1}^T \sum_{i=1}^T \left( \ell(\tilde{a}_t, z_t) - \ell(a_{t+1}, z_t) \right).$

**Reminder (2):**  $a_{t+1}^{\text{FTRL}} = \tilde{a}_t - \eta z_t \eta z_t, \quad t = 1, \dots, T-1.$

- For sake of brevity, we write  $a_1, a_2, \dots$  for  $a_1^{\text{FTRL}}, a_2^{\text{FTRL}}, \dots$

- With this,

$$\begin{aligned} R_T^{\text{FTRL}}(\tilde{a}) &\leq \psi(\tilde{a}) - \psi(a_1) + \sum_{t=1}^T \left( \ell(\tilde{a}_t, z_t) - \ell(a_{t+1}, z_t) \right) && \text{(Reminder (1))} \\ &\leq \frac{1}{2\eta} \|\tilde{a}\|_2^2 + \sum_{t=1}^T (a_t^\top z_t - a_{t+1}^\top z_t) \quad (\psi(a_1) \geq 0 \text{ and definition of } \psi) \\ &= \frac{1}{2\eta} \|\tilde{a}\|_2^2 + \sum_{t=1}^T (a_t^\top - a_{t+1}^\top) z_t && \text{(Distributivity)} \\ &= \frac{1}{2\eta} \|\tilde{a}\|_2^2 + \eta \sum_{t=1}^T \|z_t\|_2^2. && \text{(Reminder (2))} \end{aligned}$$



# DESIRED RESULTS

- With the FTRL algorithm we can cope with
  - online quadratic optimization (OQO) problems by using no regularity ( $\psi \equiv 0$ ). In this case, we have satisfactory regret guarantees and also a quick update rule for  $a_{t+1}^{\text{FTRL}}$  (It is just the empirical average over all data points seen till  $t$ ),



# DESIRED RESULTS

- With the FTRL algorithm we can cope with
  - online quadratic optimization (OQO) problems by using no regularity ( $\psi \equiv 0$ ). In this case, we have satisfactory regret guarantees and also a quick update rule for  $a_{t+1}^{\text{FTRL}}$  (It is just the empirical average over all data points seen till  $t$ ),
  - online linear optimization (OLO) problems by using a suitable regularization function. In this case, we have quick update formulas and satisfactory regret guarantees as well.



# DESIRED RESULTS

- With the FTRL algorithm we can cope with
  - online quadratic optimization (OQO) problems by using no regularity ( $\psi \equiv 0$ ). In this case, we have satisfactory regret guarantees and also a quick update rule for  $a_{t+1}^{\text{FTRL}}$  (It is just the empirical average over all data points seen till  $t$ ),
  - online linear optimization (OLO) problems by using a suitable regularization function. In this case, we have quick update formulas and satisfactory regret guarantees as well.

⇒ But what about other online learning problems or rather other loss functions?

- What we wish to have is an approach such that we can achieve for a large class of loss functions the advantages of FTRL for OLO and OQO problems:
  - (a) reasonable regret upper bounds;
  - (b) a quick update formula.

