# Some characterizations of reversibility of quantum channels

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# Reversible (sufficient) quantum channels

Let  ${\mathcal S}$  be a set of quantum states,  $\Phi$  a quantum channel.

We say that  $\Phi$  is reversible (sufficient) with respect to  $\mathcal S$  if there exists some channel  $\Psi$  (recovery channel) such that

$$\Psi \circ \Phi(\rho) = \rho, \qquad \rho \in \mathcal{S}.$$

Reference: Denes Petz's papers

# The setting and assumptions

 $B(\mathcal{H})$  - operators on a finite dimensional Hilbert space  $\mathcal{H}$ 

A set of states

$$S \subset {\rho \in B(\mathcal{H}), \ \rho \ge 0, \ \text{Tr} \ \rho = 1}$$

• A channel  $\Phi: B(\mathcal{H}) \to B(\mathcal{K})$ , completely positive and trace preserving

### Assumptions:

There is a faithful (full rank) state  $\sigma \in \mathcal{S}$ , its image  $\Phi(\sigma) \in B(\mathcal{K})$  is also faithful.



### Preservation of the relative entropy

The relative entropy: for states  $\rho, \sigma$ 

$$D(\rho\|\sigma) = \begin{cases} \operatorname{Tr}\left[\rho(\log(\rho) - \log(\sigma))\right], & \operatorname{supp}(\rho) \leq \operatorname{supp}(\sigma) \\ \\ \infty, & \text{otherwise}. \end{cases}$$

ullet Data processing inequality: for a channel  $\Phi$ 

$$D(\Phi(\rho)\|\Phi(\sigma)) \le D(\rho\|\sigma),$$

• If  $D(\rho \| \sigma) < \infty$ , then reversibility is equivalent to

$$D(\Phi(\rho)\|\Phi(\sigma)) = D(\rho\|\sigma), \quad \forall \rho \in \mathcal{S}.$$

Petz

# Universal recovery map

The Petz dual of  $\Phi$  with respect to  $\sigma$ 

$$\Phi_{\sigma}(\cdot) = \sigma^{1/2} \Phi^* (\Phi(\sigma)^{-1/2} \cdot \Phi(\sigma)^{-1/2}) \sigma^{1/2}$$

•  $\Phi_{\sigma}$  is a channel  $B(\mathcal{K}) \to B(\mathcal{H})$  such that

$$\Phi_{\sigma} \circ \Phi(\sigma) = \sigma$$

ullet  $\Phi$  is reversible with respect to  ${\cal S}$  if and only if

$$\Phi_{\sigma} \circ \Phi(\rho) = \rho, \quad \forall \rho \in \mathcal{S}$$

Petz

# Semigroup of channels preserving ${\mathcal S}$

How to describe all channels reversible with respect to  $\mathcal{S}$ ?

Let us consider the set of channels

$$\mathcal{C}_{\mathcal{S}} := \{\Theta : B(\mathcal{H}) \to B(\mathcal{H}), \ \Theta(\rho) = \rho, \ \forall \rho \in \mathcal{S}\}$$

- convex and compact semigroup (closed under composition)
- has a faithful fixed state:  $\sigma \in \mathcal{S}$ .

By the mean ergodic theorem, there is some  $\mathcal{E_S} \in \mathcal{C_S}$  such that

$$\mathcal{E}_{\mathcal{S}} \circ \Theta = \Theta \circ \mathcal{E}_{\mathcal{S}} = \mathcal{E}_{\mathcal{S}}, \quad \forall \Theta \in \mathcal{C}_{\mathcal{S}}.$$

We see that such  $\mathcal{E}_{\mathcal{S}}$  is unique and

$$\mathcal{E}_{\mathcal{S}}^2 = \mathcal{E}_{\mathcal{S}}, \qquad \mathcal{E}_{\mathcal{S}}(\rho) = \rho, \quad \forall \rho \in \mathcal{S}.$$



### The minimal sufficient subalgebra

The adjoint  $\mathcal{E}_{\mathcal{S}}^*$  is a faithful conditional expectation



its range is a subalgebra  $\mathcal{M}_{\mathcal{S}} := \mathcal{E}_{\mathcal{S}}^*(B(\mathcal{H})).$ 

 $\mathcal{M}_{\mathcal{S}}$  is the minimal sufficient subalgebra with respect to  $\mathcal{S}$ :

- $\rho \mapsto \rho|_{\mathcal{M}_{\mathcal{S}}}$  is a sufficient channel
- $\mathcal{M}_{\mathcal{S}}$  is contained in any subalgebra with this property.

# The range of a conditional expectation

Let  $\mathcal{E}: B(\mathcal{H}) \to B(\mathcal{H})$  be such that  $\mathcal{E}^*$  is a conditional expectation.

There is a decomposition  $\mathcal{H} \equiv \oplus_n \mathcal{H}_n^L \otimes \mathcal{H}_n^R$  such that

$$\mathcal{E}^*(B(\mathcal{H})) \equiv \bigoplus_n B(\mathcal{H}_n^L) \otimes I_{\mathcal{H}_n^R}$$
$$\mathcal{E}(B(\mathcal{H})) \equiv \bigoplus_n B(\mathcal{H}_n^L) \otimes \omega_n$$

for some fixed states  $\omega_n \in B(\mathcal{H}_n^R)$ .

### The Koashi-Imoto decomposition

Applying this to  $\mathcal{E}_{\mathcal{S}}$ , we obtain

$$\mathcal{M}_{\mathcal{S}} \equiv \bigoplus_{n} B(\mathcal{H}_{n}^{\mathcal{S},L}) \otimes I_{\mathcal{H}_{n}^{\mathcal{S},R}}$$

$$\rho \equiv \bigoplus_{n} \lambda_{n}(\rho) \rho_{n} \otimes \sigma_{n}, \qquad \rho \in \mathcal{S},$$

- $\lambda_n(\rho)$  is a probability dsitribution (classical part of S)
- $\rho_n \in B(\mathcal{H}_n^{\mathcal{S},L})$  are states (depending on  $\rho$ )
- $\sigma_n \in B(\mathcal{H}_n^{\mathcal{S},R})$  are fixed states.

Koashi-Imoto, Hayden, etc., Luczak, Kuramochi

### Generators of $\mathcal{M}_{\mathcal{S}}$

The minimal sufficient subalgebra is generated by:

• Connes cocycles:

$$\rho^{it}\sigma^{-it}, \qquad \rho \in \mathcal{S}, \ t \in \mathbb{R}.$$

• Radon Nikodym derivatives:

$$\sigma^{it}(\sigma^{-1/2}\rho\sigma^{-1/2})\sigma^{-it}, \qquad \rho \in \mathcal{S}, \ t \in \mathbb{R}.$$



### Reversible channels with respect to ${\cal S}$

Assume that  $\Phi$  is reversible.

• Let  $\Psi$  be a recovery channel, then  $\Psi \circ \Phi \in \mathcal{C}_{\mathcal{S}}$ , so that

$$\mathcal{E}_{\mathcal{S}} \circ (\Psi \circ \Phi) = (\Psi \circ \Phi) \circ \mathcal{E}_{\mathcal{S}} = \mathcal{E}_{\mathcal{S}}.$$

• Note that  $\mathcal{E}_{\mathcal{S}} \circ \Psi$  is also a recovery channel, so we may assume

$$\mathcal{E}_{\mathcal{S}} \circ \Psi = \Psi, \qquad \Psi \circ \Phi = \mathcal{E}_{\mathcal{S}}.$$

• We then have  $\Phi \circ \Psi = \mathcal{E}_{\mathcal{S}_0}$ , where

$$\mathcal{S}_0 := \{ \Phi(\rho), \ \rho \in \mathcal{S} \}.$$



### Reversible channels with respect to ${\cal S}$

A channel  $\Phi: B(\mathcal{H}) \to B(\mathcal{K})$  is reversible with respect to  $\mathcal{S}$  iff

$$\Phi^*|_{\mathcal{M}_{\mathcal{S}_0}}: \mathcal{M}_{\mathcal{S}_0} \xrightarrow{\mathit{iso}} \mathcal{M}_{\mathcal{S}}.$$

### Equivalently, there is

- ullet a decomposition  $\mathcal{K} \equiv \oplus_n \mathcal{K}_n^L \otimes \mathcal{K}_n^R$
- unitaries  $U_n:\mathcal{H}_n^{\mathcal{S},L}\to\mathcal{K}_n^L$
- channels  $\Phi_n: B(\mathcal{H}_n^{\mathcal{S},R}) \to B(\mathcal{K}_n^R)$

#### such that

$$\Phi|_{B(\mathcal{H}_n^{\mathcal{S},L} \otimes \mathcal{H}_n^{\mathcal{S},R})} \equiv U_n^* \cdot U_n \otimes \Phi_n.$$

### Reversible channels with respect to ${\cal S}$

Further conditions for reversibility: preserving the generators

Connes cocycles

$$\Phi^*(\Phi(\rho)^{it}\Phi(\sigma)^{-it}) = \rho^{it}\sigma^{-it}, \qquad \rho \in \mathcal{S}, \ t \in \mathbb{R};$$

Radon-Nikodym derivatives:

$$\Phi^*(\Phi(\sigma)^{-1/2}\Phi(\rho)\Phi(\sigma)^{-1/2}) = \sigma^{1/2}\rho\sigma^{1/2}, \qquad \rho \in \mathcal{S};$$

Petz dual

$$\Phi_{\sigma} \circ \Phi(\rho) = \rho, \qquad \rho \in \mathcal{S}.$$



### Conditions on ${\cal S}$

Given a channel  $\Phi$ , what are the conditions for states in S?

We fix a faithful state  $\sigma \in \mathcal{S}$ . Then we must have

$$\mathcal{S} \subset \operatorname{Fix}(\Phi_{\sigma} \circ \Phi) := \{ \rho, \ \Phi_{\sigma} \circ \Phi(\rho) = \rho \}.$$

Put

$$\mathcal{F} := \lim_{n} \frac{1}{n} \sum_{k=1}^{n} (\Phi_{\sigma} \circ \Phi)^{k},$$

then  $\mathcal{F}^*$  is a conditional expectation and

$$\mathcal{F}(B(\mathcal{H})) = \operatorname{Fix}(\Phi_{\sigma} \circ \Phi).$$

### Conditions on ${\cal S}$

#### There is

- ullet a decomposition  $\mathcal{H}\equiv \oplus_n \mathcal{H}_n^{\Phi,\sigma,L}\otimes \mathcal{H}_n^{\Phi,\sigma,R}$
- and states  $\omega_n \in B(\mathcal{H}_n^{\Phi,\sigma,R})$

such that  $\Phi$  is reversible with respect to  $\mathcal S$  if and only if all  $\rho \in \mathcal S$  have the form

$$\rho \equiv \bigoplus_{n} \mu_n(\rho) \rho_n \otimes \omega_n$$

for some probability distribution  $\mu(\rho)$  and states  $\rho_n \in B(\mathcal{H}^{\Phi,\sigma,L})$ .



# Preservation of standard Rényi relative entropies

The standard (Petz-type) Rényi relative entropies,  $\alpha > 0$ :

$$D_{\alpha}(\rho \| \sigma) = \begin{cases} \frac{1}{\alpha - 1} \log \operatorname{Tr} \, \rho^{\alpha} \sigma^{1 - \alpha} & \alpha \neq 1 \\ \operatorname{Tr} \, \rho(\log \rho - \log \sigma), & \alpha = 1. \end{cases}$$

• satisfy data processing inequality for  $\alpha \in (0, 2]$ .

 $\Phi$  is sufficient with respect to  ${\cal S}$  if and only if

$$D_{\alpha}(\Phi(\rho)\|\Phi(\sigma)) = D_{\alpha}(\rho\|\sigma), \quad \forall \rho \in \mathcal{S}$$

for some  $\alpha \in (0,2)$ . Petz, PetzJA, HMPB, HM,H

# Preservation of sandwiched Rényi relative entropies

The sandwiched (minimal) Rényi relative entropies,  $\alpha > 0$ :

$$\tilde{D}_{\alpha}(\rho \| \sigma) = \begin{cases} \frac{1}{\alpha - 1} \log \operatorname{Tr} \left( \sigma^{\frac{1 - \alpha}{2\alpha}} \rho \sigma^{\frac{1 - \alpha}{2\alpha}} \right)^{\alpha} & \alpha \neq 1 \\ \operatorname{Tr} \rho(\log \rho - \log \sigma), & \alpha = 1. \end{cases}$$

• satisfy data processing inequality for  $\alpha \in [1/2, \infty]$ 

 $\Phi$  is sufficient with respect to  ${\cal S}$  if and only if

$$\tilde{D}_{\alpha}(\Phi(\rho)\|\Phi(\sigma)) = \tilde{D}_{\alpha}(\rho\|\sigma), \quad \forall \rho \in \mathcal{S}$$

for some  $\alpha \in (1/2, \infty)$ . JA, JA

# Sandwiched Rényi relative entropy

We look at

$$\tilde{Q}_{\alpha}(\rho \| \sigma) := \operatorname{Tr} \left( \sigma^{\frac{1-\alpha}{2\alpha}} \rho \sigma^{\frac{1-\alpha}{2\alpha}} \right)^{\alpha},$$

so that

$$\tilde{D}_{\alpha}(\rho \| \sigma) = \frac{1}{\alpha - 1} \log \tilde{Q}_{\alpha}(\rho \| \sigma)$$

- For  $\alpha > 1$ : interpolation  $L_p$ -norms
- For  $\alpha \in (1/2,1)$ : a variational formula, relation to case  $\alpha > 1$
- The case  $\alpha = 1$  (relative entropy): solved by Petz

# An interpolation $L_p$ -norm with respect to a state

Let us define a norm in  $B(\mathcal{H})$ , for  $\alpha \geq 1$ :

$$||X||_{\alpha,\sigma} = \left(\operatorname{Tr} |\sigma^{\frac{1-\alpha}{2\alpha}} X \sigma^{\frac{1-\alpha}{2\alpha}}|^{\alpha}\right)^{\frac{1}{\alpha}}$$

We have for any state  $\rho$ :

$$\tilde{Q}_{\alpha}(\rho \| \sigma) = \| \rho \|_{\alpha, \sigma}^{\alpha}$$

The norm can be obtained by complex interpolation between

$$||X||_{1,\sigma} = \text{Tr } |X| = ||X||_1, \qquad ||X||_{\infty,\sigma} = ||\sigma^{-\frac{1}{2}}X\sigma^{-\frac{1}{2}}||$$

### Hadamard three lines theorem

For any function on 
$$S=\{z\in\mathbb{C},\ {\rm Re}(z)\in[0,1]\},$$
 
$$f:S\to B(\mathcal{H}),\quad \text{continuous, analytic in } {\rm int}(S)$$

• we have for any  $\alpha > 1$ ,

$$||f(1/\alpha)||_{\alpha,\sigma} \le \max_{t \in \mathbb{R}} ||f(it)||_{\infty,\sigma} \max_{t \in \mathbb{R}} ||f(1+it)||_1$$

• If equality holds for some  $\alpha>1$ , then it holds for all

### Hadamard three lines theorem

For any  $\rho \geq 0$  and  $\alpha$ , we define a function

$$f_{\rho,\alpha}(z) = \|\rho\|_{\alpha,\sigma}^{1-z\alpha} \sigma^{\frac{1-z}{2}} \left(\sigma^{\frac{1-\alpha}{2\alpha}} \rho \sigma^{\frac{1-\alpha}{2\alpha}}\right)^{z\alpha} \sigma^{\frac{1-z}{2}}, \qquad z \in S$$

- $f_{\rho,\alpha}(1/\alpha) = \rho$ ,
- The equality in Hadamard three lines theorem is attained:

$$||f_{\rho,\alpha}(1/\alpha)||_{\alpha,\sigma} = \max_{t \in \mathbb{R}} ||f_{\alpha,\sigma}(it)||_{\infty,\sigma} \max_{t \in \mathbb{R}} ||f_{\alpha,\sigma}(1+it)||_1$$

### Positive trace preserving maps are contractions

Let  $\Phi: B(\mathcal{H}) \to B(\mathcal{K})$  be a positive trace preserving linear map:

• For  $\alpha = 1$ .

$$\|\Phi(X)\|_1 \le \|X\|_1, \qquad X \in B(\mathcal{H})$$

• For  $\alpha = \infty$ ,

$$\|\Phi(X)\|_{\infty,\Phi(\sigma)} = \|\Phi_{\sigma}^*(\sigma^{-1/2}X\sigma^{-1/2})\|_{\infty} \le \|X\|_{\infty,\sigma}$$

• For  $\alpha > 1$ , by Riesz-Thorin (complex interpolation)

$$\|\Phi(X)\|_{\alpha,\Phi(\sigma)} \le \|X\|_{\alpha,\sigma}, \qquad X \in B(\mathcal{H}).$$

Beigi

### The case $\alpha = 2$

Let  $\alpha = 2$ .

•  $\|\cdot\|_{s,\sigma}$  is a Hilbert space norm, with the inner product

$$\langle X, Y \rangle_{\sigma} = \operatorname{Tr} X^* \sigma^{1/2} Y \sigma^{1/2}$$

• For a positive trace preserving map  $\Phi: B(\mathcal{H}) \to B(\mathcal{K})$ ,

$$\langle X, \Phi(Y) \rangle_{\Phi(\sigma)} = \langle \Phi_{\sigma}(X), Y \rangle_{\sigma}, \qquad X \in B(\mathcal{K}), \ Y \in B(\mathcal{H})$$

• Since  $\Phi$  is a contraction,

$$\|\Phi(Y)\|_{2,\Phi(\sigma)} = \|Y\|_{2,\sigma} \iff \Phi_{\sigma} \circ \Phi(Y) = Y.$$



### Preservation and reversibility

Let  $\Phi$  be a channel and assume that for some  $\alpha > 1$ ,

$$\|\Phi(\rho)\|_{\alpha,\Phi(\sigma)} = \|\rho\|_{\alpha,\sigma} \bigg( \iff \tilde{D}_{\alpha}(\Phi(\rho)\|\Phi(\sigma)) = \tilde{D}_{\alpha}(\rho\|\sigma) \bigg)$$

• For  $\alpha=2$ , we get

$$\Phi_{\sigma} \circ \Phi(\rho) = \rho, \Longrightarrow \Phi$$
 is reversible.