Notes on asymptotics of quantum hypothesis testing

Anna Jenčová*

1 Preliminaries

Let \mathcal{H} be a finite dimensional Hilbert space.

1.1 Pinching

Let $A \in B(\mathcal{H})$ be self-djoint, with spectral decomposition $A = \sum_i \lambda_i P_i$. We will need the pinching map $B(\mathcal{H}) \to B(\mathcal{H})$, defined as

$$\mathcal{E}_A(X) = \sum_i P_i X P_i.$$

Then A is a cp unital map. Moreover, $\mathcal{E}_A(X)$ commutes with X and we have the pinching inequality [?]

$$\mathcal{E}_A(X) \le |\operatorname{spec}(A)|X, \qquad X \ge 0.$$
 (1)

1.2 Relative entropies

Let ρ and σ be density operators. The (Umegaki) relative entropy is defined as

$$D(\rho \| \sigma) := \begin{cases} \operatorname{Tr} \left[\rho(\log \rho - \log \sigma) \right], & \operatorname{supp}(\rho) \leq \operatorname{supp}(\sigma) \\ \infty, & \text{otherwise.} \end{cases}$$

The standard Rényi relative entropy for $\alpha \in [0,1] \setminus \{1\}$ is defined as

$$D_{\alpha}(\rho \| \sigma) := \begin{cases} \frac{1}{\alpha - 1} \log \operatorname{Tr} \left[\rho^{\alpha} \sigma^{1 - \alpha} \right], & \operatorname{supp}(\rho) \leq \operatorname{supp}(\sigma) \text{ or } \alpha \in (0, 1) \\ \infty, & \text{otherwise.} \end{cases}$$

The sandwiched Rényi relative entropy for $\alpha \in [1/2, \infty] \setminus \{1\}$ is defined as

$$\hat{D}_{\alpha}(\rho \| \sigma) := \begin{cases} \frac{1}{\alpha - 1} \log \operatorname{Tr} \left[\sigma^{\frac{1 - \alpha}{2\alpha}} \rho \sigma^{\frac{1 - \alpha}{2\alpha}} \right], & \operatorname{supp}(\rho) \leq \operatorname{supp}(\sigma) \text{ or } \alpha \in [1/2, 1) \\ \infty, & \text{otherwise.} \end{cases}$$

^{*}Mathematical Institute, Slovak Academy of Sciences, Štefánikova 49, 814 73 Bratislava, Slovakia, jenca@mat.savba.sk

1.3 The function ϕ

We assume $\operatorname{supp}(\rho) \leq \operatorname{supp}(\sigma)$, so that $D(\rho \| \sigma) < \infty$.

Let us define

$$\phi(s) = \log \operatorname{Tr} \left[\rho^{1-s} \sigma^s \right], \quad s \in \mathbb{R}.$$

Then ϕ is a convex and smooth function, with derivative

$$\phi'(s) = (\operatorname{Tr}\left[\rho^{1-s}\sigma^{s}\right])^{-1}\operatorname{Tr}\left[\rho^{1-s}\sigma^{s}(\log\sigma - \log\rho),\right]$$

[?, Exercise 3.5] In particular, $\phi'(0) = -D(\rho \| \sigma)$ and $\phi'(1) = D(\sigma \| \rho)$.

Lemma 1. Put

$$\psi(\lambda) = \inf_{s \in [0,1]} \lambda s + \phi(s), \qquad \lambda \in \mathbb{R}.$$

Then

$$\psi(\lambda) = \begin{cases} 0 & \lambda \ge -\phi'(0) = D(\rho \| \sigma) \\ < 0 & \lambda < D(\rho \| \sigma) \\ \lambda & \lambda \le -\phi'(1) = -D(\sigma \| \rho). \end{cases}$$

Proof. By convexity, the derivative $\phi'(s)$ is nondecreasing. It follows that if $\lambda \geq -\phi'(0)$, then

$$\frac{d}{ds}(\lambda s + \phi(s)) = \lambda + \phi'(s) \ge -\phi'(0) + \phi'(s) \ge 0,$$

so that the function $s \mapsto \lambda s + \phi(s)$ is nondecreasing, so that the infimum is attained at s = 0. Similarly, if $\lambda \leq -\phi'(1)$, then the infimum is attained at s = 1 and hence $\psi(\lambda) = \lambda$. Let now $\lambda < -\phi'(0)$, then we see that the function $s \mapsto \lambda s + \phi(s)$ is strictly decreasing at s = 0, so that we must have $\psi(\lambda) < 0$.

We define

$$\psi^*(\lambda) = \inf_{t \in [-1,0]} t\lambda + \phi(t).$$

Again, if $\lambda > D(\rho \| \sigma) = -\phi'(0)$, then $t \mapsto t\lambda + \phi(t)$ is strictly increasing at t = 0, which implies that $\phi^*(\lambda) < 0$.

1.4 Inequalities

We have two basic inequalities. For $A, B \ge 0$, let $\{A \ge B\}$ be the sum of eigenprojections of A - B corresponding to nonnegative eigenvalues, similarly $\{A \le B\}$, $\{A > B\}$ etc. Then

Lemma 2 (Quantum Neyman-Pearson). We have

$$\min_{0 \le T \le I} \text{Tr} [A(I - T)] + \text{Tr} [BT] = \text{Tr} [A\{A \le B\}] + \text{Tr} [B\{A > B\}].$$

Lemma 3 (Audenaert et al). We have for any $s \in [0, 1]$,

$$\operatorname{Tr}\left[A\{A\leq B\}\right]+\operatorname{Tr}\left[B\{A>B\}\right]\leq \operatorname{Tr}\left[A^{1-s}B^{s}\right].$$

These statements hold in the von Neumann algebra case as well.

2 QHT

Let ρ, σ be a pair of density matrices. We test the hypothesis $H_0 = \rho$ against the alternative $H_1 = \sigma$. A test is given by an operator $0 \le T \le I$, corresponding to accepring H_0 . The two error probabilities are

$$\alpha(T) = \text{Tr}[(I - T)\rho], \qquad \beta(T) = \text{Tr}[T\sigma].$$

We will consider the asymptotic behaviour of the error probabilities

$$\alpha_n(T_n) = \text{Tr}\left[(I - T_n)\rho_n\right], \qquad \beta_n(T_n) = \text{Tr}\left[T_n\sigma_n\right]$$

in testing $H_0 = \rho_n := \rho^{\otimes n}$ against $H_1 = \sigma_n := \sigma^{\otimes n}$.

2.1 Quantum Stein's lemma

We assume $\operatorname{supp}(\rho) \leq \operatorname{supp}(\sigma)$, so that $D(\rho \| \sigma) < \infty$.

Let $\lambda \in \mathbb{R}$ and let $S_n := \{\rho^{\otimes n} > e^{n\lambda}\sigma^{\otimes n}\}$. Then using Lemma 3 (Audenaert) with $A = \rho^{\otimes n}$ and $B = e^{\lambda n}\sigma^{\otimes n}$, we get for any $s \in [0, 1]$

$$\alpha_n(S_n) + e^{n\lambda}\beta_n(S_n) \le \operatorname{Tr} e^{n\lambda s} [(\rho^{\otimes n})^{1-s}(\sigma^{\otimes n})^s] = e^{n\lambda s} (\operatorname{Tr} [\rho^{1-s}\sigma^s])^n = e^{n(\lambda s + \phi(s))}.$$
 (2)

Hence by taking the infimum over $s \in [0, 1]$,

$$\alpha_n(S_n) \le e^{n\psi(\lambda)}, \qquad \beta_n(S_n) \le e^{n(-\lambda + \psi(\lambda))}$$
 (3)

On the other hand, put $p_n = \text{Tr}\left[\rho^{\otimes n}S_n\right]$ and $q_n = \text{Tr}\left[\sigma^{\otimes n}S_n\right]$. Then $p_n \geq e^{n\lambda}q_n$ and therefore $p_n^t \leq e^{n\lambda t}q_n^t$ for for any $t \in [-1,0]$. We get

$$1 - \alpha_n(S_n) = p_n \le e^{n\lambda t} p_n^{1-t} q_n^t \le e^{n\lambda t} (p_n^{1-t} q_n^t + (1 - p_n)^{1-t} (1 - q_n)^t) \le e^{n\lambda t} \operatorname{Tr} \left[(\rho^{\otimes n})^{1-t} (\sigma^{\otimes n})^t \right]$$
$$= e^{n(\lambda t + \phi(t))}$$

for all $t \in [-1,0]$. It follows that for any test T_n , we have

$$1 - \alpha_n(T_n) = \operatorname{Tr}\left[\rho^{\otimes n} T_n\right] = \operatorname{Tr}\left[\left(\rho^{\otimes n} - e^{\lambda n} \sigma^{\otimes n}\right) T_n\right] + e^{\lambda n} \beta_n(T_n) \le \operatorname{Tr}\left[\left(\rho^{\otimes n} - e^{\lambda n} \sigma^{\otimes n}\right) S_n\right] + e^{\lambda n} \beta_n(T_n)$$
$$< 1 - \alpha_n(S_n) + e^{\lambda n} \beta_n(T_n) < e^{n(\lambda t + \phi(t))} + e^{\lambda n} \beta_n(T_n)$$

and hence

$$\beta_n(T_n) \ge e^{-n\lambda} (1 - \alpha_n(T_n) - e^{n\psi^*(\lambda)}) \tag{4}$$

Lemma 4 (Quantum Stein's lemma). [? ?] For all $\epsilon \in (0,1)$, we have

$$\lim_{n \to \infty} -\frac{1}{n} \log \beta_n(\epsilon) = D(\rho \| \sigma).$$

Proof. Let $\lambda < D(\rho \| \sigma)$, then by Lemma 1, $\psi(\lambda) < 0$, so that in this case (3), $\alpha_n(S_n) \to 0$ and

$$-\frac{1}{n}\log\beta_n(S_n) \ge \lambda - \psi(\lambda) > \lambda.$$

For $\epsilon \in (0,1)$ we have $\alpha_n(S_n) \leq \epsilon$ for large enough n, so that $\beta_n(\epsilon) \leq \beta_n(S_n)$. It follows that

$$\liminf_{n} -\frac{1}{n} \log \beta_n(\epsilon) \ge -\frac{1}{n} \log \beta_n(S_n) \ge \lambda.$$

Conversely, by (4) we have for any sequence of tests such that $\alpha_n(T_n) \leq \epsilon$ that

$$\beta_n(T_n) \ge e^{-n\lambda} (1 - \epsilon - e^{n\psi^*(\lambda)}).$$

Since $\psi^*(\lambda) < 0$ for $\lambda > D(\rho \| \sigma)$, this implies that, for such λ ,

$$\limsup_{n} -\frac{1}{n}\beta_{n}(\epsilon) \leq \lambda.$$

Choosing any $\delta > 0$, we obtain

$$D(\rho\|\sigma) - \delta \le \liminf_{n} -\frac{1}{n}\beta_n(\epsilon) \le \limsup_{n} -\frac{1}{n}\beta_n(\epsilon) \le D(\rho\|\sigma) + \delta.$$

Since δ was arbitrary, this implies the statement.

The following quantum Stein's lemma describes the situation when the error of the first kind is constrained by some $\epsilon > 0$. In this case, the optimal value of the second kind error

$$\beta_n(\epsilon) := \min_{0 \le T_n \le I} \{ \beta_n(T_n) \mid \alpha_n(T_n) \le \epsilon \}, \qquad \epsilon > 0$$

converges to zero exponetially fast, that is, $\beta_n(\epsilon) \sim e^{-nr}$. The lemma shows that for any $\epsilon > 0$, the error exponent r is given by the relative entropy