

# Various definitions

Anna Jenčová

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## 1 Symmetric monoidal categories (SMC)

**Monoidal category:** A category  $C$  equipped with

- A functor  $\otimes : C \times C \rightarrow C$ ;
- unit object  $I \in C$ ;
- associator: natural iso  $(A \otimes B) \otimes C \xrightarrow{\alpha_{A,B,C}} A \otimes (B \otimes C)$ ;
- left unitor: natural iso  $I \otimes A \xrightarrow{\lambda_A} A$ ;
- right unitor: natural iso  $A \otimes I \xrightarrow{\rho_A} A$ ;
- **symmetric** if there is a symmetry: natural iso  $A \otimes B \xrightarrow{\sigma_{A,B}} B \otimes A$  such that  $\sigma_{B,A} = \sigma_{A,B}^{-1}$ ,

satisfying triangle, pentagon (+ symmetry) diagrams.

We will always assume that  $C$  is a SMC.

### 1.1 Closed SMC

A SMC  $C$  is **closed** if:

for every  $b \in C$ , the endofunctor  $- \otimes b$  has a right adjoint  $[b, -]$  (internal hom).

What does this mean?

- (1) For all  $a, c \in C$ ,  $C(a \otimes b, c) \simeq C(a, [b, c])$ , naturally in  $a, c$ .
- (2) unit  $\eta_a^b : a \rightarrow [b, a \otimes b]$ , counit:  $\epsilon_a^b : [b, a] \otimes b \rightarrow a$ , natural transformations, triangle identities

‘ Relation of the two:

- Let  $i$  be the iso of (1):

$$\begin{aligned} \eta_a^b &\in C(a, [b, a \otimes b]) \simeq C(a \otimes b, a \otimes b), & \eta_a^b &= i(id_{a \otimes b}) \\ \epsilon_a^b &\in C([b, a] \otimes b, a) \simeq C([b, a], [b, a]), & \epsilon_a^b &= i^{-1}(id_{[b, a]}). \end{aligned}$$

- Conversely, from  $\eta^b, \epsilon^b$  of (2), we define  $i$  as

$$g \in C(a \otimes b, c), \quad i(g) = [b, g] \circ \eta_a^b$$

with inverse

$$h \in C(a, [b, c]), \quad i^{-1}(h) = \epsilon_c^b \circ (h \otimes b)$$

Informally, we may interpret  $\eta_a^b$  as 'labeling of  $b$  by  $a$ ' and  $\epsilon_a^b$  as 'evaluation of  $[b, a]$ '.

## 1.2 Compact SMC

A SMC is **compact** if each object  $a \in C$  has a dual  $a^* \in C$  such that there are maps  $\cup_a : I \rightarrow a^* \otimes a$  and  $\cap_a : a \otimes a^* \rightarrow I$  satisfying the snake identities

$$(\cap_a \otimes id_a) \circ (id_a \otimes \cup_a) = id_a, \quad (id_{a^*} \otimes \cap_a) \circ (\cup_a \otimes id_{a^*}) = id_{a^*}.$$

The following are easily seen by the snake identities (pictures):

- (1)  $a^*$  is determined up to iso;
- (2)  $a^{**} \simeq a$ , indeed, we may define  $\cup_{a^*} : I \rightarrow a \otimes a^*$  and  $\cap_{a^*} : a^* \otimes a \rightarrow I$  as

$$\cup_{a^*} = \sigma_{a^*, a} \circ \cup_a, \quad \cap_{a^*} = \cap_a \circ \sigma_{a^*, a},$$

so that  $a$  is dual to  $a^*$ , and use (1);

- (3) if we fix  $a^*$  and  $\cup_a$  ( $\cap_a$ ), then  $\cap_a$  ( $\cup_a$ ) is uniquely determined;
- (4) any assignment  $a \mapsto a^*$  defines a functor  $C \rightarrow C^{op}$  (if  $f : a \rightarrow b$ , we can use  $\cup_a$  and  $\cap_b$  to "bend the wires" to obtain a map  $b^* \rightarrow a^*$ , this is obviously functorial);
- (5)  $(a \otimes b)^* \simeq a^* \otimes b^*$ , we can clearly put (using symmetry)

$$\cup_{a \otimes b} = \cup_a \otimes \cup_b, \quad \cap_{a \otimes b} = \cap_a \otimes \cap_b$$

- (5)  $C$  is closed, with  $[b, c] = b^* \otimes c$ : the iso  $i : C(a \otimes b, c) \simeq C(a, b^* \otimes c)$  can be obtain (roughly) as

$$i(g) = g \circ \cup_b, \quad i^{-1}(h) = \cap_b \circ h,$$

the maps are inverses by the snake identities, naturality is obvious, since  $i$  does nothing on  $a$  or  $c$ . The unit and counit of the adjunction are given as

$$\eta_a^b = a \otimes \cup_b : a \rightarrow b^* \otimes a \otimes b, \quad \epsilon_a^b = \cap_b \otimes a : b^* \otimes a \otimes b \rightarrow a.$$

- (6) Can we state a theorem like:  $C$  is compact if and only if for each  $b \in C$  there is some  $b^* \in C$  such that  $b^* \otimes -$  is the right adjoint of  $- \otimes b$  and ...? What should be the additional conditions?

## 2 Kleisli categories and monoidal monads

A **monad** on  $C$  is a triple  $(P, \eta, \mu)$ , where:

- $P : C \rightarrow C$  is an endofunctor;
- $\eta : Id_C \rightarrow P$ ,  $\mu : P^2 \rightarrow P$  are natural transformations satisfying some triangles and squares.

The **Kleisli category**  $C_P$  has the same objects as  $C$ , with morphisms:

$$C_p(a, b) = C(a, P(b)),$$

and for  $f \in C_p(a, b)$ ,  $g \in C_p(b, c)$ , the composition is defined as

$$g \circ f := \mu_c \circ P(g) \circ f.$$

A monad is **monoidal** [?] if there are maps

$$\kappa_{a,b} : Pa \otimes Pb \rightarrow P(a \otimes b), \quad a, b \in C,$$

natural in  $a, b$  and such that

- $(P, \eta, \kappa)$  is a **monoidal functor**, that is, some diagrams involving  $P$ ,  $\alpha$ ,  $\lambda$ ,  $\rho$ ,  $\kappa$  and  $\eta$  commute;
- additional diagrams containing  $\mu$  commutes;
- **symmetric**: additionally a diagram with  $\sigma$  commutes.

**Proposition 1.** [?, Prop. 1.2.2] *There is a bijective correspondence between:*

- families of morphisms  $\{\kappa_{a,b}\}$  such that  $(P, \eta, \mu, \kappa)$  is a (symmetric) monoidal monad;*
- (symmetric) monoidal structures on  $C_P$  such that the left adjoint functor  $F_P : C \rightarrow C_P$  is strict monoidal.*