On the structure of higher order quantum maps

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1 Preliminaries

1.1 The category FinVect

Let FinVect be the category of finite dimensional real vector spaces with linear maps. We will denote the usual tensor product by $\otimes$, then $(\mathrm{FinVect}, \otimes, I = \mathbb{R})$ is a symmetric monoidal category, with the associators, unitors and symmetries given by the obvious isomorphisms

$$\alpha_{U,V,W} : (U \otimes V) \otimes W \simeq U \otimes (V \otimes W),$$
$$\lambda_V : I \otimes V \simeq V, \qquad \rho_V : V \otimes I \simeq V,$$
$$\sigma_{U,V} : U \otimes V \simeq V \otimes U.$$

Let $(-)^* : V \mapsto V^*$ be the usual vector space dual, with duality denoted by $\langle \cdot, \cdot \rangle : V^* \times V \to \mathbb{R}$. We will use the canonical identification $V^{**} = V$ and $(V_1 \otimes V_2)^* = V_1^* \otimes V_2^*$. With this duality, FinVect is compact closed. This means that for each object V, there are maps $\eta_V : I \to V^* \otimes V$ (the "cup") and $\epsilon_V : V \otimes V^* \to I$ (the "cap") such that the following snake identities hold:

$$(\epsilon_V \otimes V) \circ (V \otimes \eta_V) = V, \qquad (V^* \otimes \epsilon_V) \circ (\eta_V \otimes V^*) = V^*, \tag{1}$$

here we denote the identity map on the object V by V. Indeed, η_V can be identified with an element $\eta_V(1) \in V^* \otimes V$ and $\epsilon_V \in (V \otimes V^*)^* = V^* \otimes V$ is again an element of the same space. Choose a basis $\{e_i\}$ of V, let $\{e_i^*\}$ be the dual basis of V^*, that is, $\langle e_i^*, e_j \rangle = \delta_{i,j}$. Let us then define

$$\eta_V(1) = \epsilon_V := \sum_i e_i^* \otimes e_i.$$

It is easy to see that this definition does not depend on the choice of the basis, indeed, ϵ_V is the linear functional on $V \otimes V^*$ defined by

$$\langle \epsilon_V, x \otimes x^* \rangle = \langle x^*, x \rangle, \qquad x \in V, \ x^* \in V^*.$$

It is also easily checked that the snake identities (1) hold.

For two objects V and W in FinVect, we will denote the set of all morphisms (i.e. linear maps) $V \to W$ $L(V, W)$ by $\mathrm{FinVect}(V, W)$. Then $\mathrm{FinVect}(V, W)$ is itself a real linear space and we have the well-known identification $\mathrm{FinVect}(V, W) \simeq V^* \otimes W$. This can be given as follows: for each $f \in FV(V, W)$, we have $C_f := (V^* \otimes f)(\epsilon_V) = \sum_i e_i^* \otimes f(e_i) \in V^* \otimes W$. Conversely, since $\{e_i^*\}$ is a basis of V^*, any element $w \in V^* \otimes W$ can be uniquely written as $w = \sum_i e_i^* \otimes w_i$ for $w_i \in W$,

and since $\{e_i\}$ is a basis of V, the assignment $f(e_i) := w_i$ determines a unique map $f : V \to W$. The relations between $f \in \mathrm{FinVect}(V, W)$ and $C_f \in V^* \otimes W$ can be also written as

$$\langle\, C_f, x \otimes y^* \,\rangle = \langle\, \epsilon_V, x \otimes f^*(y^*) \,\rangle = \langle\, f^*(y^*), x \,\rangle = \langle\, y^*, f(x) \,\rangle, \qquad x \in V, \; y^* \in W^*,$$

here $f^* : W^* \to V^*$ is the adjoint of f. Note that by compactness, the internal hom in FinVect satisfies $[V, W] \simeq V^* \otimes W$, so that in the case of FinVect, the object $[V, W]$ can be identified with the space of linear maps $\mathrm{FinVect}(V, W)$.

Example 1. Let $V = \mathbb{R}^N$. In this case, we fix the canonical basis $\{|i\,\rangle, \; i = 1, \ldots, N\}$. We will identify $(\mathbb{R}^N)^* = \mathbb{R}^N$, with duality $\langle\, x, y \,\rangle = \sum_i x_i y_i$, in particular, we identify $I = I^*$. We then have $\epsilon_V = \sum_i |i\,\rangle \otimes |i\,\rangle$ and if $f : \mathbb{R}^N \to \mathbb{R}^M$ is given by the matrix A in the two canonical bases, then $C_f = \sum_i |i\,\rangle \otimes A|i\,\rangle$ is the vectorization of A.

Example 2. Let $V = M_n^h$ be the space of $n \times n$ complex hermitian matrices. We again identify $(M_n^h)^* = M_n^h$, with duality $\langle\, A, B \,\rangle = \mathrm{Tr}\, A^T B$, where A^T is the usual transpose of the matrix A. Let us choose the basis in M_n^h, given as

$$\left\{ |j\,\rangle\langle k| + |k\,\rangle\langle j|, \; j \le k, \; i\Big(|j\,\rangle\langle k| - |k\,\rangle\langle j| \Big), \; j < k \right\}.$$

Then one can check that

$$\left\{ \frac{1}{2}\Big(|j\,\rangle\langle k| + |k\,\rangle\langle j| \Big), \; j \le k, \; \frac{i}{2}\Big(|k\,\rangle\langle j| - |j\,\rangle\langle k| \Big), \; j < k \right\}$$

is the dual basis and we have

$$\epsilon_V = \sum_{j,k} |j\,\rangle\langle k| \otimes |j\,\rangle\langle k|.$$

For any $f : M_n^h \to M_m^h$,

$$C_f = \sum_{j,k} |j\,\rangle\langle k| \otimes f(|j\,\rangle\langle k|)$$

is the Choi matrix of f.

1.2 Boolean functions

We will work with certain boolean functions, which are defined as functions from binary strings $\{0, 1\}^*$ to $\{0, 1\}$. We now list some basic notations used below.

For $s \in \{0, 1\}$, we denote $\bar{s} := 1 - s$. For binary strings of fixed length n, that is, elements of $\{0, 1\}^n$, we will denote by 0_n or just 0 the string $00 \ldots 0$ and

$$e^i = \delta_{i,1} \ldots \delta_{i,n}.$$

For $m, n \in \mathbb{N}$, the concatenation of strings $s \in \{0, 1\}^m$ and $t \in \{0, 1\}^n$ will be denoted by st, that is,

$$st = s_1 \ldots s_m t_1 \ldots t_n \in \{0, 1\}^{m+n}.$$

For a string $x \in \{0, 1\}^n$ and any set of indices $\{i_1, \ldots, i_k\} \subseteq [n]$, we will denote by $x^{i_1 \ldots i_k}$ the string in $\{0, 1\}^{n-k}$ obtained from x by removing $x_{i_1}, \ldots, x_{i_k}$. For any permutation $\sigma \in S_n$, we will denote by the same symbol the obvious action on $\{0, 1\}^n$, that is

$$\sigma(s_1 \ldots s_n) = s_{\sigma^{-1}(1)} \ldots s_{\sigma^{-1}(n)}.$$

Let us introduce the subset

$$\mathcal{F}_n := \{f : \{0,1\}^n \to \{0,1\}, \ f(0) = 1\}.$$

With the poitwise ordering, $\mathcal{F}_n$ is a (finite) distributive lattice, with top element the constant 1 function and the bottom element $p_n := \chi_0$, the characteristic function of the zero string. We may also define complementation in $\mathcal{F}_n$ as

$$f^* := 1 - f + p_n.$$

It can be easily checked that with these structures $\mathcal{F}_n$ is a boolean algebra.

We now introduce some more operations in $\mathcal{F}_n$. For $f \in \mathcal{F}_n$ and any permutation $\sigma \in S_n$, we see that $f \circ \sigma \in \mathcal{F}_n$. For $f \in \mathcal{F}_m$ and $g \in \mathcal{F}_n$, we define the function $f \otimes g \in \mathcal{F}_{m+n}$ as

$$(f \otimes g)(st) = f(s)g(t).$$

As it is, this tensor product is not symmetric, but there is a permutation $\sigma \in S_{m+n}$ such that $(g \otimes f) = (f \otimes g) \circ \sigma$ for any $f \in \mathcal{F}_m$ and $g \in \mathcal{F}_n$.

Lemma 1. *For $f \in \mathcal{F}_m$, $g \in \mathcal{F}_n$, we have*

$$f \otimes g \leq (f^* \otimes g^*)^*.$$

Equality holds if and only if f and g are either both the top or both the bottom elements in $\mathcal{F}_m$ resp. $\mathcal{F}_n$.

Proof. The inequality is easily checked, since $(f \otimes g)(st)$ can be 1 only if $f(s) = g(t) = 1$. If both s and t are the zero strings, then $st = 0_{m+n}$ and both sides are equal to 1. Otherwise, the condition $f(s) = g(t) = 1$ implies that $(f^* \otimes g^*)(st) = 0$, which implies that the right hand side must be 1. If f and g are both constant 1, then $(1 \otimes 1)^* = 1^* = p_{n+m} = 1^* \otimes 1^*$, in the case when both f and g are the minimal elements equality follows by duality. Finally, asume the equality holds and that $f \neq 1$, so that there is some s such that $f(s) = 0$. But then $s \neq 0$ and for any t,

$$0 = (f \otimes g)(st) = (f^* \otimes g^*)^*(st) = 1 - f^*(s)g^*(t) + p_{m+n}(st) = 1 - g^*(t),$$

which implies that $g(t) = 0$ for all $t \neq 0$, that is, $g = p_n$. By the same argument, $f = p_m$ if $g \neq 1$, which implies that either $f = 1$ and $g = 1$, or $f = p_m$ and $g = p_n$. $\qquad\square$

We now show an important example.

Example 3. Let $S \subseteq [n]$ be any subset. We will denote

$$p_S(s) := \prod_{i \in S} \bar{s}_i.$$

It is clear that $p_S \in \mathcal{F}_n$, $p_\emptyset = 1$, $p_{[n]} = \chi_0 = p_n$. The following properies are also easy to see for $S, T \subseteq [n]$:

   (i) if $S \subseteq T$, then $p_T \leq p_S$,

   (ii) $p_S \wedge p_T = p_S p_T = p_{S \cup T}$,

(iii) $p_S \vee p_T = p_S + p_T - p_{S \cup T}$.

(iv) let $S \subseteq [m]$ and $T \subseteq [n]$, then
$$p_S \otimes p_T = p_{S \cup (m+T)}.$$

We will use the above functions to introduce a convenient parametrization to $\mathcal{F}_n$. For this, we first include $\mathcal{F}_n$ into a larger set
$$\mathcal{F}_n \subseteq \{f : \{0,1\}^n \to \mathbb{R}\} =: \mathcal{V},$$

which is a 2^n-dimensional real vector space. It becomes a Hilbert space with respect to the inner product
$$\langle f, g \rangle = \sum_{s \in \{0,1\}^n} f(s) g(s).$$

Lemma 2. *The set $\{p_S, \ S \subseteq [n]\}$ is a basis of $\mathcal{V}$. Any $f \in \mathcal{V}$ can be uniquely written as*
$$f = \sum_{S \subseteq [n]} \hat{f}_S p_S,$$

where the coefficients $\hat{f}_S \in \mathbb{R}$ are obtained as
$$\hat{f}_S = \sum_{\substack{x \in \{0,1\}^n \\ x_i = 1, \forall i \in S^c}} (-1)^{\sum_{i \in S} x_i} f(x).$$

Proof. For $T \subseteq [n]$, let us define the function $p_T^{\perp}$ as
$$p_T^{\perp}(x) := (-1)^{\sum_{i \in T} x_i} \prod_{i \in T^c} x_i.$$

We prove that for $S, T \subseteq [n]$,
$$\langle p_S, p_T^{\perp} \rangle = \delta_{S,T},$$

which shows that $\{p_S, \ S \subseteq [n]\}$ is a basis and $\{p_T^{\perp}, \ T \subseteq [n]\}$ is the dual basis. We compute
$$\langle p_S, p_T^{\perp} \rangle = \sum_x p_S(x) p_T^{\perp}(x) = \sum_x (-1)^{\sum_{i \in T} x_i} \prod_{i \in S} \bar{x}_i \prod_{j \in T^c} x_j.$$

This expression can be nonzero only if $S \cap T^c = \emptyset$, that is, $S \subseteq T$. In this case, the last sum is equal to
$$\sum_{\substack{x \in \{0,1\}^n \\ x_i = 0, \forall i \in S \\ x_i = 1, \forall i \in T^c}} (-1)^{\sum_{j \in T \setminus S} x_j} = \begin{cases} 0 & \text{if } S \subsetneq T \\ 1 & \text{if } S = T \end{cases}$$

It is now clear that the coefficients
$$\hat{f}_S = \langle f, p_S^{\perp} \rangle$$

have the given form.

$\square$

Remark 1. This can be also obtained using Möbius inversion formula, see [Stanley, Sec. 3.7].

It may be useful to visualise the lattice $\mathcal{L}_n = \{S \subseteq [n]\}$ as a hypercube, and the coefficients of f as labels for its vertices. The fact that the function f has values in $\{0,1\}$ means that for a string $x \in \{0,1\}^n$ such that $x_j = 1$ if and only if $j \in T$, we must have

$$f(x) = \sum_{\substack{S \subseteq [n] \\ S \cap T = \emptyset}} \hat{f}_S \in \{0,1\},$$

that is, the sum of labels $\hat{f}_S$ over any face of the hypercube $\mathcal{L}_n$ containing the vertex $\emptyset$ must be 0 or 1. In particular, $\hat{f}_\emptyset = f(11\ldots 1) \in \{0,1\}$, which restricts the values of $\hat{f}_{\{i\}} \in \{0,1,-1\}$, etc. The fact that $f \in \mathcal{F}_n$ means that in addition,

$$f(0) = \sum_{S \subseteq [n]} \hat{f}_S = 1.$$

2   The category of affine subspaces

2.1   The category Af

We now introduce the category Af, whose objects are of the form $X = (V_X, A_X)$, where V_X is an object in FinVect and $A_X \subseteq V_X$ a proper affine subspace, which means that $0 \notin A_X \neq \emptyset$. Morphisms $X \xrightarrow{f} Y$ in Af are linear maps $f : V_X \to V_Y$ such that $f(A_X) \subseteq A_Y$. For any object X, we put

$$L_X = \mathrm{Lin}(A_X) := \{a_1 - a_2, \ a_1, a_2 \in A_X\} = \{a - a_X, \ a \in A_X\}, \qquad S_X := \mathrm{Span}(A_X).$$

Here a_X is any element in A_X and L_X does not depend on this choice. Then L_X and S_X are linear subspaces such that $d_X := \dim(L_X) = \dim(S_X) - 1$. We will also denote $D_X = \dim(V_X)$. For any element $a_X \in A_X$, the affine subspace is determined as

$$A_X = a_X + L_X.$$

Let us now define the duality of affine subspaces as follows. Let V be an object in FinVect and let $C \subseteq V$ be any subset. Let

$$\tilde{C} := \{v^* \in V^*, \ \langle v^*, c \rangle = 1\}.$$

The following lemma collects some properties that are easily proven.

Lemma 3.   *(i) $\tilde{C}$ is an affine subspace.*

  (ii) $0 \in \tilde{C}$ if and only if $C = \emptyset$ and $\tilde{C} = \emptyset$ if and only if $0 \in \mathrm{Aff}(C)$.

  (iii) Let $0 \notin \mathrm{Aff}(C)$, then $\mathrm{Aff}(C) = \tilde{\tilde{C}}$ and we have

$$\mathrm{Lin}(C) = \mathrm{Lin}(\tilde{\tilde{C}}) = \tilde{C}^\perp = \mathrm{Span}(\tilde{C})^\perp, \qquad \mathrm{Lin}(\tilde{C}) = C^\perp = \mathrm{Span}(C)^\perp$$
$$\mathrm{Span}(C) = C^{\perp\perp} = \mathrm{Lin}(\tilde{C})^\perp, \qquad\qquad \mathrm{Span}(\tilde{C}) = \mathrm{Lin}(C)^\perp.$$

For any $\tilde{a}_X \in \tilde{A}_X$, the subspace A_X is determined as

$$A_X = S_X \cap \{\tilde{a}_X\}^{\sim}.$$

The relation between the subspaces L_X and S_X is given as

$$S_X = L_X \oplus \mathbb{R}a_X, \qquad L_X = S_X \cap \{\tilde{a}_X\}^{\perp}.$$

By Lemma 3 above, $\tilde{A}_X$ is a proper affine subspace in V_X^*, so that $X^* := (V_X^*, \tilde{A}_X)$ is an object in Af. We have $X^{**} = X$ and the corresponding subspaces are related as

$$L_{X^*} = S_X^{\perp}, \qquad S_{X^*} = L_X^{\perp}. \tag{2}$$

Note also that for $X \xrightarrow{f} Y$, the adjoint map satisfies $f^*(\tilde{A}_Y) \subseteq \tilde{A}_X$, so that $Y^* \xrightarrow{f^*} X^*$ and the duality $(-)^*$ is a full and faithful functor $\mathrm{Af}^{op} \to \mathrm{Af}$.

We will next introduce a monoidal structure $\otimes$ as follows. For two objects X and Y, we put $V_{X \otimes Y} = V_X \otimes V_Y$ and construct the affine subspace $A_{X \otimes Y}$ as the affine span of

$$A_X \otimes A_Y = \{a \otimes b, \ a \in A_X, \ b \in A_Y\}.$$

Fix any $\tilde{a}_X \in \tilde{A}_X$ and $\tilde{a}_Y \in \tilde{A}_Y$. Since $A_X \otimes A_Y \subseteq \{\tilde{a}_X \otimes \tilde{a}_Y\}^{\sim}$, the affine span of $A_X \otimes A_Y$ is a proper affine subspace and we have by Lemma 3

$$A_{X \otimes Y} := \mathrm{Aff}(A_X \otimes A_Y) = \{A_X \otimes A_Y\}^{\approx}.$$

Lemma 4. *For any $a_X \in A_X$, $a_Y \in A_Y$, we have*

$$L_{X \otimes Y} = \mathrm{Lin}(A_X \otimes A_Y) = \mathrm{Span}(\{x \otimes y - a_X \otimes a_Y, \ x \in A_X, \ y \in A_Y\}) \tag{3}$$
$$= (a_X \otimes L_Y) + (L_X \otimes a_Y) + (L_X \otimes L_Y) \tag{4}$$

(here $+$ denotes the direct sum of subspaces). We also have

$$S_{X \otimes Y} = S_X \otimes S_Y.$$

Proof. The equality (3) follows from Lemma 3. For any $x \in A_X$, $y \in A_Y$ we have

$$x \otimes y - a_X \otimes a_Y = a_X \otimes (y - a_Y) + (x - a_X) \otimes a_Y + (x - a_X) \otimes (y - a_Y),$$

so that $L_{X \otimes Y} = \mathrm{Lin}(A_X \otimes A_Y)$ is contained in the subspace on the RHS of (4). Let d be the dimension of this subspace, then clearly

$$d_{X \otimes Y} \le d \le d_X + d_Y + d_X d_Y.$$

On the other hand, any element of S_X has the form tx for some $t \in \mathbb{R}$ and $x \in A_X$, so that it is easily seen that $S_X \otimes S_Y = S_{X \otimes Y}$. Hence

$$d_{X \otimes Y} = \dim(L_{X \otimes Y}) = \dim(S_{X \otimes Y}) - 1 = \dim(S_X)\dim(S_Y) - 1 = (d_X + 1)(d_Y + 1) - 1$$
$$= d_X + d_Y + d_X d_Y.$$

This completes the proof.

$\square$

Lemma 5. *Let $I = (\mathbb{R}, \{1\})$. Then $(\mathrm{Af}, \otimes, I)$ is a symmetric monoidal category.*

Proof. Note that this structure is inherited from the symmetric monoidal structure in FinVect. To show that $\otimes$ is a functor, we have to check that for $X_1 \xrightarrow{f} Y_1$ and $X_2 \xrightarrow{g} Y_2$ in Af, we have $X_1 \otimes Y_1 \xrightarrow{f \otimes g} X_2 \otimes Y_2$ which amounts to showing that

$$(f \otimes g)(A_{X_1 \otimes Y_1}) \subseteq A_{X_2 \otimes Y_2}.$$

Let $x \in A_{X_1}$, $y \in A_{Y_1}$, then $f(x) \otimes g(y) \in A_{X_2} \otimes A_{Y_2} \subseteq A_{X_2 \otimes Y_2}$. Since $A_{X_1 \otimes Y_1}$ is the affine subspace generated by $A(X_1) \otimes A(Y_1)$, the above inclusion follows by linearity of $f \otimes g$.

It only remains to prove that the associators, unitors and symmetries from FinVect are morphisms in Af. We will prove this for the associators $\alpha_{X,Y,Z} : V_X \otimes (V_Y \otimes V_Z) \to (V_X \otimes V_Y) \otimes V_Z$, the other proofs are similar. We need to check that $\alpha_{X,Y,Z}(A_{X \otimes (Y \otimes Z)}) \subseteq A_{(X \otimes Y) \otimes Z}$. It is easily checked that $A_{X \otimes (Y \otimes Z)}$ is the affine span of elements of the form $x \otimes (y \otimes z)$, $x \in A_X$, $y \in A_Y$ and $z \in A_Z$, and we have

$$\alpha_{X,Y,Z}(x \otimes (y \otimes z)) = (x \otimes y) \otimes z \in A_{(X \otimes Y) \otimes Z}$$

for all such elements. The desired inclusion follows by linearity.

$\square$

Theorem 1. *$(\mathrm{Af}, \otimes, I)$ is a *-autonomous category, with duality $(-)^*$, such that $I^* = I$.*

Proof. By Lemma 5, we have that $(\mathrm{Af}, \otimes I)$ is a symmetric monoidal category. We have also seen that the duality $(-)^*$ is a full and faithful contravariant functor. We only need to check the natural isomorphisms

$$\mathrm{Af}(X \otimes Y, Z^*) \simeq \mathrm{Af}(X, (Y \otimes Z)^*).$$

Since FinVect is compact, we have the natural isomorphisms

$$\mathrm{FinVect}(V_X \otimes V_Y, V_Z^*) \simeq \mathrm{FinVect}(V_X, V_Y^* \otimes V_Z^*),$$

determined by the equalities

$$\langle\, f(x \otimes y), z \,\rangle = \langle\, h(x), y \otimes z \,\rangle, \qquad x \in V_X, \ y \in V_Y, \ z \in V_Z,$$

for $f \in \mathrm{FinVect}(V_X \otimes V_Y, V_Z^*)$ and $h \in \mathrm{FinVect}(V_Z, V_Y^* \otimes V_Z^*)$. Since $A_{X \otimes Y}$ is an affine span of $A_X \otimes A_Y$, we see that f is in $\mathrm{Af}(X \otimes Y, Z^*)$ if and only if $f(x \otimes y) \in \tilde{A}_Z$ for all $x \in A_X$, $y \in A_Y$, that is,

$$1 = \langle\, f(x \otimes y), z \,\rangle = \langle\, h(x), y \otimes z \,\rangle \qquad \forall x \in A_X, \ \forall y \in A_Y, \ \forall z \in A_Z.$$

But this is equivalent to

$$h(x) \in (A_Y \otimes A_Z)^{\sim} = \tilde{A}_{Y \otimes Z}, \qquad \forall x \in A_X,$$

which means that $h \in \mathrm{Af}(X, (Y \otimes Z)^*)$.

$\square$

A *-autonomous category is compact closed if it satisfies $(X \otimes Y)^* = X^* \otimes Y^*$. In general, $X \odot Y = (X^* \otimes Y^*)^*$ defines a dual symmetric monoidal structure that is different from $\otimes$. We next show that Af is not compact.

Proposition 1. *For objects in* Af, *we have* $(X \otimes Y)^* = X^* \otimes Y^*$ *exactly in one of the following situations:*

(i) $X \simeq I$ *or* $Y \simeq I$,

(ii) $d_X = d_Y = 0$,

(iii) $d_{X^*} = d_{Y^*} = 0$.

Proof. It is easily seen by definition that $A_{X^*} \otimes A_{Y^*} = \tilde{A}_X \otimes \tilde{A}_Y \subseteq \tilde{A}_{X \otimes Y} = A_{(X \otimes Y)^*}$. Hence the equality holds if and only if $d_{X^* \otimes Y^*} = d_{(X \otimes Y)^*}$. From Lemma 4, we see that

$$d_{X^* \otimes Y^*} = d_{X^*} + d_{Y^*} + d_{X^*} d_{Y^*}.$$

On the other hand, we have using (2) that $L_{(X \otimes Y)^*} = S_{X \otimes Y}^\perp = (S_X \otimes S_Y)^\perp$, so that

$$d_{(X \otimes Y)^*} = D_X D_Y - \dim(S_X)\dim(S_Y) = D_X D_Y - (d_X + 1)(d_Y + 1).$$

Taking into account that by (2) we have $d_{X^*} = D_X - d_X - 1$, similarly for d_{Y^*}, we obtain

$$d_{(X \otimes Y)^*} - d_{X^* \otimes Y^*} = d_X d_{Y^*} + d_{X^*} d_Y.$$

This is equal to 0 iff $d_X d_{Y^*} = d_Y d_{X^*} = 0$, which amounts to the conditions in the lemma.

$\square$

In a *-autonomous category, the internal hom can be identified as $[X, Y] = (X \otimes Y^*)^*$. The underlying vector space is $V_{[X,Y]} = (V_X \otimes V_Y^*)^* = V_X^* \otimes V_Y$ and we have seen in Section 1.1 that we may identify this space with $\mathrm{FinVect}(V_X, V_Y)$, by $f \leftrightarrow C_f$. This property is extended to Af, in the following sense.

Proposition 2. *For any objects X, Y in* Af, *the map $f \mapsto C_f$ is a bijection of* $\mathrm{Af}(X, Y)$ *onto* $A_{[X,Y]}$.

Proof. Let $f \in \mathrm{FinVect}(V_X, V_Y)$. Since by definition $A_{[X,Y]} = \tilde{A}_{X \otimes Y^*}$ and $A_{X \otimes Y^*}$ is an affine span of $A_X \otimes A_Y^*$, we see that $C_f \in A_{[X,Y]}$ if and only if for all $x \in A_X$ and $y^* \in \tilde{A}_Y$, we have

$$1 = \langle C_f, x \otimes y^* \rangle = \langle y^*, f(x) \rangle.$$

This latter statement is clearly equivalent to $f(A_X) \subseteq A_Y$, so that $f \in \mathrm{Af}(X, Y)$. $\square$

In the next result, we restrict the objects to spaces of hermitian matrices, as in Example 2 and morphisms to completely positive maps. We show that this restriction amounts to taking an intersection of $A_{[X,Y]}$ with the cone of positive semidefinite matrices. This, and subsequent examples, shows that for characterization of sets of quantum objects such as states, channels, combs and transformations between them, it is enough to work with the category Af.

An object X of Af will be called quantum if $V_X = M_n^h$ for some n and A_X is an affine subspace such that both A_X and $\tilde{A}_X$ have nonempty intersection with the interior of the positive cone $int(M_n^+)$ (recall that we identify $(M_n^h)^* = M_n^h$).

Proposition 3. *Let X, Y be quantum objects in* Af. *Then*

(i) X^* *and* $X \otimes Y$ *are quantum objects as well*

(ii) *Let $V_X = M_n^h$, $V_Y = M_m^h$. Then for any $f \in \mathrm{FinVect}(M_n^h, M_m^h)$, we have $C_f \in A_{[X,Y]} \cap M_{mn}^+$ if and only if f is completely positive and*

$$f(A_X \cap M_n^+) \subseteq A_Y \cap M_m^+.$$

Proof. The statement (i) is easily seen from $A_X \otimes A_Y \subseteq A_{X \otimes Y}$ and $\tilde{A}_X \otimes \tilde{A}_Y \subseteq \tilde{A}_{X \otimes Y}$, together with the fact that $int(M_n^+) \otimes int(M_m^+) \subseteq int(M_{mn}^+)$. To show (ii), let $C_f \in A_{[X,Y]} \cap M_{mn}^+$. By the properties of the Choi isomorphism f is completely positive and by Proposition 2, $f(A_X) \subseteq A_Y$, this proves one implication. For the converse, note that we only need to prove that under the given assumptions, $f(A_X) \subseteq A_Y$, for which it is enough to show that $A_X \subseteq \mathrm{Aff}(A_X \cap M_n^+)$. To see this, pick some $a_X \in A_X \cap int(M_n^+)$. Any element in A_X can be written in the form $a_X + v$ for some $v \in L_X$. Since $a_X \in int(M_n^+)$, there is some $s > 0$ such that $a_\pm := a_X \pm sv \in M_n^+$, and since $\pm sv \in L_X$, we see that $a_\pm \in A_X \cap M_n^+$. It is now easily checked that

$$a_X + v = \frac{1+s}{2s} a_+ + \frac{s-1}{2s} a_- \in \mathrm{Aff}(A_X \cap M_n^+).$$

$\square$

We can define classical obejcts in Af in a similar way, replacing M_n^h by $\mathbb{R}^N$ and the positive cone by $\mathbb{R}_+^N$. A similar statement holds in this case, with complete positivity replaced by positivity. We can similarly treat classical-to-quantum and quantum-to-classical maps as morphisms between these types of objects, satisfying appropriate positivity assumptions.

Example 4. States, channels, combs, nonsignaling, etb, dual, process matrices

Example 5. POVMs, instruments, multimeters.

2.2 First order and higher order objects

We say that an object X in Af is first order if $d_X = D_X - 1$, equivalently, $S_X = V_X$. Another equivalent condition is $d_{X^*} = 0$, which means that A_X is determined by a single element $\tilde{a}_X \in V_X^*$ as

$$A_X = \{\tilde{a}_X\}^\sim, \qquad \tilde{A}_X = \{\tilde{a}_X\}.$$

In the case of first order quantum objects we additionally require that $\tilde{a}_X \in int(M_n^+)$, similarly for classical first order objects. Note that first order objects, resp. their duals, are exactly those satisfying condition (iii), resp. condition (ii), in Proposition 1, in particular, $(X \otimes Y)^* = X^* \otimes Y^*$ for first order objects X and Y.

For a first order object $X = (V_X, \{\tilde{a}_X\}^\sim)$, let us pick an element $a_X \in A_X$, then we have a direct sum decomposition

$$V_X = L_{X,0} \oplus L_{X,1},$$

where $L_{X,0} := \mathbb{R}\{a_X\}$, $L_{X,1} := \{\tilde{a}_X\}^\perp = L_X$. We also define the *conjugate object* $\tilde{X} = (V_X^*, \{a_X\}^\sim)$, note that we always have $\tilde{a}_X \in A_{\tilde{X}}$ and with the choice $a_{\tilde{X}} = \tilde{a}_X$, we have $\tilde{\tilde{X}} = X$ and

$$L_{\tilde{X},u} = L_{X,\bar{u}}^\perp, \qquad u \in \{0,1\}. \tag{5}$$

These definitions depend on the choice of a_X, but we will assume below that this choice is fixed and that we choose $a_{\tilde{X}} = \tilde{a}_X$. For quantum objects we will assume that $a_X \in int(M_n^+)$.

Example 6. Example: quantum states, multiple of identity...

Higher order objects are those obtained from a finite set $\{X_1, \ldots, X_n\}$ of first order objects by taking tensor products and duals, and applying any permuations of the spaces. The above is indeed a set, so that all the objects are different (though they may be isomorphic) and the ordering is not essential. We will also assume that the tensor unit is not contained in this set. Of course, any first order object is also higher order with $n = 1$. Note that we cannot say that such an object is automatically "of order n", as the following lemma shows.

Lemma 6. *Let X, Y be first order, then $X \otimes Y$ is first order as well.*

Proof. We have
$$S_{X \otimes Y} = S_X \otimes S_Y = V_X \otimes V_Y = V_{X \otimes Y}.$$
$\square$

Example 7. (states quantum first order, channels, supermaps - quantum higher order)

Example 8. replacement $X^* \otimes Y$, quantum

2.3 Description of higher order objects

We start by noticing that there are certain objects in Af that can be constructed from a set of first order objects and functions in $\mathcal{F}_n$.

Let $X_1, \ldots, X_n$ be first order objects in Af. Let $a_{X_i} \in A_{X_i}$ be fixed and let $\tilde{X}_i$ be the conjugate first order objects. Let us denote $V_i = V_{X_i}$ and

$$L_{i,u} := L_{X_i, u}, \qquad \tilde{L}_{i,u} := L_{\tilde{X}_i, u} \qquad u \in \{0, 1\}, \ i \in [n].$$

For $s \in \{0, 1\}^n$, we define

$$L_s := L_{1, s_1} \otimes \cdots \otimes L_{n, s_n}, \qquad \tilde{L}_s := \tilde{L}_{1, s_1} \otimes \cdots \otimes \tilde{L}_{n, s_n},$$

then we have the direct sum decompositions

$$V := V_1 \otimes \cdots \otimes V_n = \bigoplus_{s \in \{0,1\}^n} L_s, \qquad V^* = V_1^* \otimes \cdots \otimes V_n^* = \bigoplus_{s \in \{0,1\}^n} \tilde{L}_s.$$

Lemma 7. *For any $s \in \{0, 1\}^n$, we have*

$$L_s^\perp = \bigoplus_{t \in \{0,1\}^n} \bar{\chi}_s(t) \tilde{L}_t, \qquad \tilde{L}_s^\perp = \bigoplus_{t \in \{0,1\}^n} \bar{\chi}_s(t) L_t.$$

Here $\chi_s : \{0, 1\}^n \to \{0, 1\}$ is the characteristic function of s, $\bar{\chi}_s = 1 - \chi_s$.

Proof. Using (5) and the direct sum decomposition of V_i^*, we get

$$(L_{1, s_1} \otimes \cdots \otimes L_{n, s_n})^\perp = \bigvee_j \left(V_1^* \otimes \cdots \otimes V_{j-1}^* \otimes \tilde{L}_{j, \bar{s}_j} \otimes V_{j+1}^* \otimes \cdots \otimes V_n^* \right)$$

$$= \bigvee_j \left(\bigoplus_{\substack{t \in \{0,1\}^n \\ t_j \neq s_j}} \tilde{L}_{1, t_1} \otimes \cdots \otimes \tilde{L}_{n, t_n} \right)$$

$$= \bigoplus_{\substack{t \in \{0,1\}^n \\ t \neq s}} \left(\tilde{L}_{1, t_1} \otimes \cdots \otimes \tilde{L}_{n, t_n} \right).$$

The proof of the other equality is the same.

$\square$

Lemma 8. *Put $a := a_1 \otimes \cdots \otimes a_n$, $\tilde{a} := \tilde{a}_1 \otimes \cdots \otimes \tilde{a}_n$. For $f \in \mathcal{F}_n$ define*

$$S_f := \bigoplus_{s \in \{0,1\}^n} f(s) L_s, \qquad A_f = A_f(X_1, \ldots, X_n) := S_f \cap \{\tilde{a}\}^\sim.$$

Then A_f is a proper affine subspace in V containing a. Moreover,

$$L_{A_f} = \bigoplus_{s \in \{0,1\}^n \setminus \{0\}} f(s) L_s, \qquad S_{A_f} = S_f$$

and the dual affine subspace satisfies

$$\tilde{A}_f(X_1, \ldots, X_n) = A_{f^*}(\tilde{X}_1, \ldots, \tilde{X}_n) = \bigoplus_{s \in \{0,1\}^n} f^*(s) \tilde{L}_s \cap \{a\}^\sim.$$

Proof. It is clear from definition that A_f is an affine subspace. Since $f(0) = 1$, the space S_f always contains the subspace $L_0 = L_{1,0} \otimes \cdots \otimes l_{n,0} = \mathbb{R}\{a\}$ and it is clear that $L_s \subseteq \{\tilde{a}\}^\perp$ for any $s \neq 0$. It follows that $a \in A_f$, so that $A_f \neq \emptyset$, and since $A_f \subseteq \{\tilde{a}\}^\sim$, we see that A_f is proper and $\tilde{a} \in \tilde{A}_f$. The expressions for L_{A_f} and S_{A_f} are immediate from the definition and $L_{A_f} = S_{A_f} \cap \{\tilde{a}\}^\sim$. To obtain the dual affine subspace, we compute using Lemma 7 and the fact that the subspaces form an independent decomposition,

$$S_{\tilde{A}_f} = L_{A_f}^\perp = \left(\bigoplus_{s \in \{0,1\}^n \setminus \{0\}} f(s) L_s \right)^\perp = \bigwedge_{\substack{s \in \{0,1\}^n \\ s \neq 0, f(s) = 1}} L_s^\perp = \bigwedge_{\substack{s \in \{0,1\}^n \\ s \neq 0, f(s) = 1}} \left(\bigoplus_{t \in \{0,1\}^n} \bar{\chi}_s(t) \tilde{L}_t \right)$$

$$= \bigoplus_{t \in \{0,1\}^n} \left(\bigwedge_{\substack{s \in \{0,1\}^n \\ s \neq 0, f(s) = 1}} \bar{\chi}_s(t) \tilde{L}_t \right) = \bigoplus_{t \in \{0,1\}^n} f^*(t) \tilde{L}_t.$$

To see the last equality, note that

$$\bigwedge_{\substack{s \in \{0,1\}^n \\ s \neq 0, f(s) = 1}} \bar{\chi}_s(t) = \begin{cases} 1 & \text{if } t = 0 \\ 1 - f(t) & \text{if } t \neq 0 \end{cases} = f^*(t).$$

$\square$

Since L_s, $s \in \{0, 1\}$ is an independent decomposition, the map $f \mapsto S_f$, and hence also $f \mapsto A_f$, is injective. This map has the following further properties, which are easily checked:

(i) For the bottom and top elements in $\mathcal{F}_n$ we have

$$A_{p_n} = \{a\}, \qquad A_{1_n} = \{\tilde{a}\}^\sim,$$

(ii) We have $f \leq g$ if and only if $A_f \subseteq A_g$,

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(iii) $A_{f \wedge g} = A_f \cap A_g$,

(iv) $A_{f \vee g} = A_f \vee A_g := \mathrm{Aff}(A_f \cup A_g)$.

It follows that the set $\{A_f, \ f \in \mathcal{F}_n\}$ is a distributive lattice, with respect to the lattice operations $\cap$ and $\vee$.

Since all the affince subspaces $A_f \subseteq V$ are proper, there are objects $X_f := (V, A_f)$ in Af. The above relations can be rephrased as follows:

(i) $X_{1_n} = (V, \{\tilde{a}\}^\sim)$ is a first order object, $X_{p_n} = (V^*, \{a\}^\sim)^\sim$ is a dual first order object.

(ii) We have $f \le g$ if and only if id_V is a morphism $X_f \to X_g$ in Af,

(iii) Let $f, g \le h$. The following is a pullback diagram:

$$
\begin{array}{ccc}
X_{f \wedge g} & \xrightarrow{\ id_V\ } & X_f \\
{\scriptstyle id_V}\downarrow & & \downarrow{\scriptstyle id_V} \\
X_g & \xrightarrow[\ id_V\ ]{} & X_h
\end{array}
$$

(iv) Let $h \le f, g$. The following is a pushout diagram:

$$
\begin{array}{ccc}
X_h & \xrightarrow{\ id_V\ } & X_f \\
{\scriptstyle id_V}\downarrow & & \downarrow{\scriptstyle id_V} \\
X_g & \xrightarrow[\ id_V\ ]{} & X_{f \vee g}
\end{array}
$$

In particular, it follows that $\{X_f, \ f \in \mathcal{F}_n\}$, is a distributive lattice, with pullbacks and pushouts as lattice operations. Furthermore, using the conjugate objects, we may construct

$$\tilde{X}_f := (V^*, A_f(\tilde{X}_1, \ldots, \tilde{X}_n))$$

and we see from Lemma 8 that

$$X_f^* = \tilde{X}_{f^*}, \qquad f \in \mathcal{F}_n.$$

We next observe that the higher order objects are of the form X_f, for some choice of the first order objects $X_1, \ldots, X_n$ and a function f that belongs to a special subclass of $\mathcal{F}_n$. So assume that Y is a higher order object constructed from a set of distinct first order objects $Y_1, \ldots, Y_n$, $Y_i = (V_{Y_i}, \{\tilde{a}_{Y_i}\}^\sim)$, we will write $Y \sim \{Y_1, \ldots, Y_n\}$ in this case. Let us fix some element $a_{Y_i} \in A_{Y_i}$ and construct the objects $\tilde{Y}_i$.

By compactness of FinVect, we may assume (relabeling the objects if necessary) that the vector space of Y has the form

$$V_Y = V := V_1 \otimes \cdots \otimes V_n,$$

where V_i is either V_{Y_i} or $V_{Y_i}^*$, according to whether Y_i was subjected to taking duals an even or odd number of times. The indices such that the first case is true will be called the outputs and the subset of outputs in $[n]$ will be denoted by O, or O_Y, when we need to specify the object. The set $I = I_Y := [n] \setminus O_Y$ is the set of inputs. The reason for this terminology will become clear later.

Proposition 4. *For $i \in [n]$, let $X_i = Y_i$ if $i \in O_Y$ and $X_i = \tilde{Y}_i$ for $i \in I_Y$. There is a unique function $f \in \mathcal{F}_n$ such that*

$$Y = X_f = (V, A_f(X_1, \ldots, X_n)).$$

Proof. Since the map $f \mapsto X_f$ is injective, uniqueness is clear. To show existence of this function, we will proceed by induction on n. For $n = 1$, the assertion is easily seen to be true, since in this case, we we have either $Y = Y_1$ or $Y = Y_1^*$. In the first case, $O = [1]$, $X_1 = Y_1$ and

$$S_Y = V_Y = V_1 = L_{1,0} \oplus L_{1,1} = 1(0)L_{1,0} \oplus 1(1)L_{1,1},$$

so in this case f is the constant 1. If $Y = Y_1^*$, we have $O = \emptyset$, $X_1 = \tilde{Y}_1$, and then

$$S_Y = \mathbb{R}\{\tilde{a}_{Y_1}\} = L_{1,0} = 1^*(0)L_{1,0} \oplus 1^*(1)L_{1,1},$$

so that $f = 1^*$.

Assume now that the assertion is true for all $m < n$. By construction, Y is either the tensor product $Y = Z_1 \otimes Z_2$, with

$$Z_1 \sim \{Y_1, \ldots, Y_m\}, \qquad Z_2 \sim \{Y_{m+1}, \ldots, Y_n\},$$

or Y is the dual of such a product. Let us assume the first case. It is clear that $O_{Z_1} \cup O_{Z_2} = O_Y$, and similarly for I, so that the corresponding objects $X_1, \ldots, X_m$ and $X_{m+1}, \ldots, X_n$ remain the same. By the induction assumption, there are functions $f_1 \in \mathcal{F}_m$ and $f_2 \in \mathcal{F}_{n-m}$ such that

$$S_Y = S_{Z_1} \otimes S_{Z_2} = \bigoplus_{\substack{s \in \{0,1\}^m \\ t \in \{0,1\}^{n-m}}} f_1(s)f_2(t)L_{1,s_1} \otimes \cdots \otimes L_{m,s_m} \otimes L_{m+1,t_1} \otimes \cdots \otimes L_{n,t_{n-m}}.$$

This implies the assertion, with $f = f_1 \otimes f_2$.

To finish the proof, it is now enough to observe that if the assertion holds for Y then it also holds for Y^*. So assume that $Y = X_f = (V, A_f(X_1, \ldots, X_n))$ for some $f \in \mathcal{F}_n$, then $Y^* = X_f^* = \tilde{X}_{f^*} = (V^*, A_{f^*}(\tilde{X}_1, \ldots, \tilde{X}_n))$. In is now enough to notice that $\tilde{X}_i = \tilde{\tilde{Y}}_i = Y_i$ if $i \in I_Y$ and $\tilde{X}_i = \tilde{Y}_i$ if $i \in O_Y$. Since by definition $O_{Y^*} = I_Y$, this proves the statement.

$\square$

Let us stress that in general, the objects X_f depend on the choice of the elements a_{X_i}. From the above proof, it is clear that the description in Proposition 4 does not depend on the choice of the elements $a_{Y_i} \in A_{Y_i}$.

Example 9. Quantum higher order objects from state spaces, [**?**].

2.4 Type funtions and higher order objects

Let $\mathcal{T}_n \subseteq \mathcal{F}_n$ be defined as the subset generated from the constant function 1 on $\{0, 1\}$ by taking duals and tensor products. For example, we have

$$\mathcal{T}_1 = \mathcal{F}_1 = \{1, 1^*\}, \quad \mathcal{T}_2 = \{1 \otimes 1, (1 \otimes 1)^*, 1 \otimes 1^*, 1^* \otimes 1, (1^* \otimes 1)^*, (1 \otimes 1^*)^*\},$$

etc. Elements of $\mathcal{T}_n$ will be called *type functions*. Similarly as for the higher order objects, the indexes in $[n]$ such that the corresponding component was subjected to taking the dual an even number of times will be called the outputs (of f) and denoted by $O = O_f$, indexes in $I = I_f := [n] \setminus O_f$ will be called inputs. From the proof of Proposition 4, it is easily seen that a higher order object is of the form $Y = X_f$ for a function $f \in \mathcal{T}_n$ with the same outputs (and of course also inputs) as Y. We next show that the converse is true.

Proposition 5. *Let $\{X_1, \ldots, X_n\}$ be first order objects and let $f \in \mathcal{T}_n$, $O = O_f$, $I = I_f$. Then $Y = X_f$ is a higher order objects with $O_Y = O$ and $Y \sim \{Y_1, \ldots, Y_n\}$, where $Y_i = X_i$ for $i \in O$ and $Y_i = (V_i^*, \{a_i\}^{\sim})$ for $i \in I$.*

Proof. As before, we will proceed by induction on n. For $n = 1$, we only have the possibilities $f = 1$ or $f = 1^*$. In the first case, $O = [1]$ and we get

$$S_f = 1L_{1,0} \oplus 1L_{1,1} = V_1,$$

so that $X_f = (V_1, \{\tilde{a}_1\}^{\sim}) = X_1$. In the second case, $O = \emptyset$ and

$$S_f = 1L_{1,0} = \mathbb{R}\{a_1\},$$

so that $X_f = (V_1, \{a_1\}) = (V_1^*, \{a_1\}^{\sim})^* = Y_1^*$. Assume next that the statement is true for all $m < n$ and assume that $f = f_1 \otimes f_2$ for some $f_1 \in \mathcal{F}_m$, $f_2 \in \mathcal{F}_{n-m}$, then it is easily seen that $Y = Z_1 \otimes Z_2$ for $Z_1 = X_{f_1}$ and $Z_2 = X_{f_2}$, constructed from $\{x_1, \ldots, X_m\}$ resp. $\{x_{m+1}, \ldots, X_n\}$. By the induction assumption, Z_1 and Z_2 are higher order objects, with $O_{Z_i} = O_{f_i}$, it follows that Y is a higher order object with $O_Y = O_{Z_1} \cup O_{Z_2} = O_{f_1} \cup O_{f_2} = O_f$.

Finally, assume that the statement is true for $f \in \mathcal{F}_n$, we will show that it holds for f^*. Using Lemma 8 and its proof, we can see that X_{f^*} is the dual of the object $\tilde{X}_f$, constructed by replacing $L_{i,u}$ by $\tilde{L}_{i,u}$, for each i and $u \in \{0, 1\}$, which is the same as replacing each X_i by the object $\tilde{X}_i = (V_i^*, \{a_i\}^{\sim})$ (and then choosing $\tilde{a}_i$ as the special fixed element of $A_{\tilde{X}_i}$). Hmmm.

$$S_f^* =$$

$\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad$ $\square$

The above theorem shows that any element $Y \sim \{X_1, \ldots, X_n\}$ is (up to a possible permutation of the involved first order objects) determined by a unique boolean function $f = f_Y : \{0, 1\}^n \to \{0, 1\}$. . Functions obtained in this way will be called type functions and studied in more detail in Section **??**. It is rather obvious that not all boolean functions are type functions, for example, for $Y \sim \{X\}$ we only have $Y = X$ and $Y = X^*$, which corresponds to constant 1 and the negation, respectively.

In general, one can see from the above proof that we necessarily have

$$f \in \mathcal{F}_n = \{g : \{0, 1\}^n \to \{0, 1\}, \ g(0) = 1\}$$

for any type function f. Note that the set $\mathcal{F}_n$ is a boolean algebra, with smallest element 1^* and largest element 1, see Section **??**. Pick any element $f \in \mathcal{F}_n$ and any permutation $\pi \in S_n$ and let $\sigma_\pi : X_1 \otimes \cdots \otimes X_n \to X_{\sigma^{-1}(1)} \otimes \ldots X_{\sigma^{-1}(n)}$ be the corresponding symmetry. Keeping the notations in (**??**), let

$$S_{\pi,f} := \bigoplus_{s \in \{0,1\}^n} f(s)\sigma_\pi(L_s), \qquad A_{\pi,f} := S_{\pi,f} \cap \{\sigma_\pi(\tilde{a}_Y)\}^{\sim}, \qquad V_\pi := \sigma_\pi(\otimes_j V_j).$$

Proposition 6. *Let $f, f_1, f_2 \in \mathcal{F}_n$, $\pi, \pi_1, \pi_2 \in S_n$. Then*

(i) $X_{\pi,f} := (V_\pi, A_{\pi,f})$ is an object in Af.

(ii) For the smallest and the largest element, we have

$$X_{\pi,1^*} = (V_\pi, \{\sigma_\pi(a_Y)\}), \qquad X_{\pi,1} = (V_\pi, \{\sigma_\pi(\tilde{a}_Y)\}^{\sim}).$$

(iii) If $f_1 \leq f_2$, then $\sigma_{\pi_2} \circ \sigma_{\pi_1}^{-1}$ is a morphism $X_{\pi_1,f_1} \to X_{\pi_2,f_2}$ in Af.

(iv) The diagram

$$
\begin{array}{ccc}
X_{\pi,f_1 \wedge f_2} & \xrightarrow{\sigma_{\pi_1} \circ \sigma_\pi^{-1}} & X_{\pi_1,f_1} \\
{\scriptstyle \sigma_{\pi_2} \circ \sigma_\pi^{-1}} \downarrow & & \downarrow {\scriptstyle \sigma_\pi^{-1}} \\
X_{\pi_2,f_2} & \xrightarrow{\sigma_\pi^{-1}} & X_{id,1}
\end{array}
$$

is a pullback.

(v) The diagram

$$
\begin{array}{ccc}
X_{id,1^*} & \xrightarrow{\sigma_{\pi_1}} & X_{\pi_1,f_1} \\
{\scriptstyle \sigma_{\pi_2}} \downarrow & & \downarrow {\scriptstyle \sigma_\pi \circ \sigma_{\pi_1}^{-1}} \\
X_{\pi_2,f_2} & \xrightarrow{\sigma_\pi \circ \sigma_{\pi_1}^{-1}} & X_{\pi,f_1 \vee f_2}
\end{array}
$$

is a pushout.

Then $A_{\pi,f}$ is a proper affinne subspace in $V_\pi := \sigma_\pi(\otimes_j V_j)$, this folows from the fact that $f(0) = 1$, so that S contains the subspace $\mathbb{R}a$. Then $Y_f := (\otimes_j V_{i_j}, A_f)$ defines an object in Af such that $a \in A_f$ and $\tilde{a} \in \tilde{A}_f$. Such objects might not belong to AfH in general. With the pointwise ordering, $\mathcal{F}_n$ is a distributive lattice, with the smallest element χ_0 and largest element 1. It is easy to see that for $f, g \in \mathcal{F}_n$ and some corresponding objects Y_f, Y_g, we have $f \leq g$ if and only if there is some permutation $\sigma \in S_n$ such that $Y_f \xrightarrow{\sigma} Y_g$. In particular, since $\chi_0 \leq f \leq 1$ for all $f \in \mathcal{F}_n$, there is some permutation σ such that

$$
Y_{\min} \xrightarrow{\sigma} Y_f \xrightarrow{\sigma^{-1}} Y_{\max},
$$

where

$$
Y_{\min} := (V_1 \otimes \cdots \otimes V_n, \{a_1 \otimes \cdots \otimes a_n\}), \qquad Y_{\max} := (V_1 \otimes \cdots \otimes V_n, \{\tilde{a}_1 \otimes \cdots \otimes \tilde{a}_n\}^\sim).
$$

If Y_g is an object such that

$$
Y_{\min} \xrightarrow{\rho} Y_g \xrightarrow{\rho^{-1}} Y_{\max},
$$

for a permutation ρ, then we may define an object corresponding to $f \wedge g$ as the pullback of the two arrows $f \xrightarrow{\sigma^{-1}} Y_{\max}$ and $g \xrightarrow{\rho^{-1}} Y_{\max}$, similarly, $Y_{f \vee g}$ can be found as a pushout.

———

It is also shown that we can introduce a tensor product $\otimes$ and duality $(-)^*$ such that Af with these structures becomes a *-autonomous category. We will show that we may use this category to describe classical and quantum higher order maps.

To this end, we introduce a subcategory in Af, consisting of objects $X = (V_X, A_X)$, where the vector space V_X is restricted to be either $\mathbb{R}^n$ or M_n^h (see Example **??**). Accordingly, let also V_X^+ be either $\mathbb{R}_+^n$ (the cone of elements in $\mathbb{R}^n$ with nonnegative coordinates) or M_n^+ (the cone of positive semi-definite matrices). We then also require that both A_X and $\tilde{A}_X$ contain some interior

elements in V_X^+ (note that in this case we identify $V_X = V_X^*$). The morphisms are restricted so that we additionaly require that $(V_X, A_X) \xrightarrow{f} (V_Y, A_Y)$ are completely positive.

and morphisms $(\mathbb{R}^n, A) \xrightarrow{f} (\mathbb{R}^m, B)$ in Af such that we also have $f(\mathbb{R}_+^n) \subseteq \mathbb{R}_+^m$. The category Quant consists of objects of the form $X = (M_n, A)$ and morphisms $(M_n, A) \xrightarrow{f} (M_m, B)$ in Af such that $f(M_n^+) \subseteq M_m^+$.

For any object X, we also put

$$L_X := \text{Lin}(A_X) \qquad S_X := \text{Span}(A_X), \qquad d_X := \dim(L_X), \qquad D_X := \dim(V_X).$$

Note that X is uniquely determined also by the triple (V_X, L_X, a_X) with an element $a_X \in A_X$, or by $(V_X, S_X, \tilde{a}_X)$ with an element $\tilde{a}_X \in \tilde{A}_X$.

We will consider the following special kind of morphisms in Af. A morphism $X \xrightarrow{f} Y$ is a monomorphism if $f \circ h = f \circ g$ implies $h = g$ for any morphisms g, h, and an epimorphism if $h \circ f = g \circ f$ implies $h = g$. A morphism that is both mono and epi is called a bimorphism.

Lemma 9. *A morphism $X \xrightarrow{f} Y$ is a monomorphism if and only if it is injective as a map $f : V_X \to V_Y$. Similarly, f is an epimorphism if and only if it is surjective.*

Consequently, f is a bimorphism if and only if it is an isomorphism of V_X and V_Y. Note that a bimorphism is not necessarily an isomorphism in Af, which would mean that the inverse map satisfies $f^{-1}(A_Y) \subseteq A_X$.

Proof. Let f be a monomorphism in Af and let $K = Ker(f)$. Let $Z = (V_X \times K, A_X \times \{0\})$, then Z is an object in Af. Let $g, h : V_Z \to V_X$ be defined as $g(x, y) = x$, $h(x, y) = x + y$, then $g, h : Z \to X$ are morphisms in Af and we have

$$f \circ g(x, y) = f(x) = f(x) + f(y) = f \circ h(x, y), \qquad \forall (x, y) \in V_Z.$$

Hence $h = g$, so that we must have $K = \{0\}$ and f is injective. The converse is clear.

Similarly, let f be an epimorphism and let $R = f(V_X) \subseteq V_Y$. Let $Z = (V_Y \times V_Y|_R, A_Y \times \{[0]\})$ and let $g, h : V_Y \to V_Z$ be given by $g(y) = (y, [0])$, $h(y) = (y, q(y))$, where $q : V_Y \to V_Y|_R$ is the quotient map. Since $A_Y \subseteq R$, we have $q(A_Y) = \{[0]\}$, so that both g, h are morphisms in Af. Moreover,

$$g \circ f(x) = (f(x), [0]) = (f(x), q(f(x))) = h \circ f,$$

so that $g = h$, but this implies that $R = V_Y$ and f is surjective. The converse is clear. $\qquad\square$

Let X, Y, Z be objects in Af such that there are bimorphisms

$$Z \xrightarrow{f} X, \qquad Z \xrightarrow{g} Y.$$

Note that in particular $\psi := f \circ g^{-1}$ is an isomorphism of V_Y onto V_X.

Let us define $X \sqcup_{f,g} Y := (V_X, A_{X \sqcup_{f,g} Y})$, with

$$A_{X \sqcup_{f,g} Y} = \{sa + (1 - s)\psi(b), \ a \in A_X, \ b \in A_Y, \ s \in \mathbb{R}\}.$$

Note first that this is a proper object in Af if and only if

$$\forall b \in A_Y, \quad t\psi(b) \in A_X \implies t = 1. \tag{6}$$

Indeed, we only have to check that $0 \notin A_{X \sqcup_{f,g} Y}$ which is easily seen to be equivalent to (6).

Assume (6), then $X \sqcup_{f,g} Y$ together with the morphisms given by the linear maps $id : V_X \to V_X$ and $\psi : V_Y \to V_X$, is the **pushout** of the above diagram. Indeed, these are clearly bimorphisms $X \to X \sqcup_{f,g} Y$ and $Y \to X \sqcup_{f,g} Y$ in Af, and we have

$$id \circ f = f = \psi \circ g.$$

Also, if W is an object in Af and $X \xrightarrow{i} W$ and $Y \xrightarrow{j} W$ are such that $i \circ f = j \circ g$, then $i = i \circ id$, $j = i \circ \psi$, so the map i defines a morphism $X \sqcup_{f_0,g_0} Y \to W$, obviously unique, with the required properties. We have

$$L_{X \sqcup_{f,g} Y} = L_X \vee \psi(L_Y), \qquad S_{X \sqcup_{f,g} Y} = S_X \vee \psi(S_Y).$$

Let us also note that if (6) is not satisfied, there is some $z \in V_Z$ such that for some $t \neq 1$,

$$tf(z) \in A_X, \qquad g(z) \in A_Y.$$

If there are some $X \xrightarrow{i} W$ and $Y \xrightarrow{j} W$ as above, then $ti \circ f(z) \in A_W$, but also $i \circ f(z) = j \circ g(z) \in A_W$, so that W is not a proper object, in this case the pushout is the terminal object 0, with the unique arrows $X \xrightarrow{!} 0$, $Y \xrightarrow{!} 0$.

Similarly, let

$$X \xrightarrow{f} Z, \qquad Y \xrightarrow{g} Z$$

be bimorphisms and let $\psi = f^{-1} \circ g$. If

$$\phi(A_Y) \cap A_X \neq \emptyset, \tag{7}$$

the **pullback** of f, g is $X \sqcap_{f,g} Y = (V_X, A_X \cap \phi(A_Y))$, with the bimorphisms given by id_X and ϕ^{-1}. In this case

$$L_{X \sqcap_{f,g} Y} = L_X \cap \phi(L_Y), \qquad S_{X \sqcap_{f,g} Y} = S_X \cap \phi(S_Y).$$

Without condition (7), the above is not a proper object and in this case the pullback is the initial object $\emptyset$.

.1 Affine subspaces

A subset $A \subseteq V$ of a finite dimensional vector space V is an affine subspace if $\sum_i \alpha_i a_i \in A$ whenever all $a_i \in A$ and $\sum_i \alpha_i = 1$. We say that A is proper if $0 \neq A$ and $A \neq \emptyset$. We will always mean that an affine subspace is proper (if not explicitly stated otherwise).

.1.1 Description

An affine subspace can be determined in two ways:

(i) Let $L \subseteq V$ be a linear subspace and $a_0 \neq L$. Then

$$A = a_0 + L$$

is a proper affine subspace. Note that $a_0 \in A$ and $A \cap L = \emptyset$. Conversely, any proper affine subspace A can be given in this way, with a_0 an arbitrary element in A and

$$L = \mathrm{Lin}(A) := \{a_1 - a_2, \ a_1, a_2 \in A\} = \{a - a_0, \ a \in A\}.$$

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(ii) Let $S \subseteq V$ be a linear subspace and $a_0^* \in V^* \setminus S^\perp$. Then

$$A = \{a \in S, \langle a_0^*, a \rangle = 1\}$$

is a proper affine subspace. Conversely, any proper affine subspace A is given in this way, with $S = \mathrm{Span}(A)$ and a_0^* an arbitrary element in the dual

$$\tilde{A} = \{a^* \in V^*, \ \langle a^*, a \rangle = 1, \ \forall a \in A\}.$$

For an affine subspace A, the relation of $L = \mathrm{Lin}(A)$ and $S = \mathrm{Span}(A)$ is as follows:

$$S = L + \mathbb{R}a, \qquad L = S \cap \{\tilde{a}\}^\perp,$$

here $a \in A$ and $\tilde{a} \in \tilde{A}$ are arbitrary elements.

.1.2 Duality

For an affine subspace A, $\tilde{A}$ is an affine subspace as well. If A is proper, then $\tilde{A}$ is proper and we have $\tilde{\tilde{A}} = A$. More generally, if $\emptyset \neq C \subseteq A$ is any subset of a proper affine subspace A, then $\tilde{C}$ is a proper affine subspace and $\tilde{\tilde{C}}$ is the affine hull of C, that is,

$$\tilde{\tilde{C}} = \mathrm{Aff}(C) := \{\sum_i \alpha_i c_i, \ c_i \in C, \ \sum_i \alpha_i = 1\}.$$

In this case, we may write $\tilde{\tilde{C}}$ as

$$\tilde{\tilde{C}} = c_0 + \mathrm{Lin}(C) = c_0 + \mathrm{Span}(\{c_1 - c_2, \ c_1, c_2 \in C\})$$

with an arbitrary element $c_0 \in C$, or as

$$\tilde{\tilde{C}} = \{c \in \mathrm{Span}(C), \ \langle a_0^*, c \rangle = 1\}$$

for an arbitrary element $a_0^* \in \tilde{A}$.

Lemma 10. *Let A be a proper affine subspace and let $C \subseteq A$ be any subset. Then*

$$\mathrm{Lin}(C) = \mathrm{Lin}(\tilde{\tilde{C}}) = \tilde{C}^\perp = \mathrm{Span}(\tilde{C})^\perp, \qquad \mathrm{Lin}(\tilde{C}) = C^\perp = \mathrm{Span}(C)^\perp$$
$$\mathrm{Span}(C) = C^{\perp\perp} = \mathrm{Lin}(\tilde{C})^\perp, \qquad \mathrm{Span}(\tilde{C}) = \mathrm{Lin}(C)^\perp.$$

.1.3 The lattice of affine subspaces

Let $\mathcal{A}(V)$ be the set of all affine subspaces in a finite dimensional vector space V. Then $\mathcal{A}(V)$ can be ordered by inclusion and it is a complete lattice, with

$$\wedge \mathcal{A} = \cap \mathcal{A}, \qquad \vee \mathcal{A} = \{\sum_i \alpha_i a_i, \ a_i \in A_i \in \mathcal{A}, \sum_i \alpha_i = 1\}$$

for any subset $\mathcal{A} \subseteq \mathcal{A}(V)$. Let us choose any nonzero elements $a \in V$, $\tilde{a} \in V^*$ and put

$$\mathcal{A}_{a,\tilde{a}}(V) = \{A \in \mathcal{A}(V), \ a \in A, \ \tilde{a} \in \tilde{A}\}.$$

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Note that any subspace in $\mathcal{A}_{a,\tilde{a}}$ is proper and it is a complete sublattice in $\mathcal{A}(V)$. Moreover, we have
$$\text{Lin}(\wedge\mathcal{A}) = \wedge\{\text{Lin}(A),\ A \in \mathcal{A}\}, \qquad \text{Lin}(\vee\mathcal{A}) = \vee\{\text{Lin}(A),\ A \in \mathcal{A}\}$$
and similarly for Span.

We say that $A, B \in \mathcal{A}_{a,\tilde{a}}(V)$ are independent if $A \cap B = \{a\}$, equivalently, $\text{Lin}(A) \cap \text{Lin}(B) = \{0\}$, that is, $\text{Lin}(A)$ and $\text{Lin}(B)$ are independent linear subspaces. A family $\{A_1, \ldots, A_n\} \subseteq \mathcal{A}_{a,\tilde{a}}(V)$ is independent if A_i and $\vee_{j\in I}A_j$ are independent for any $i \in \{1, \ldots, n\}$ and $i \neq I \subseteq \{1, \ldots, n\}$. Equivalently, $\{\text{Lin}(A_1), \ldots, \text{Lin}(A_n)\}$ is an independent family of subspaces in V.

Lemma 11. *Let $\{A_1, \ldots, A_n\} \subseteq \mathcal{A}_{a,\tilde{a}}(V)$ be an independent family. Then the sublattice generated by $\{A_1, \ldots, A_n\}$ is distributive.*

Proof. Clear from a similar property of linear subspaces.

$\square$

.1.4 Limits and colimits

Limits and colimits should be obtained from those in FinVect, we have to spectify the other structures and check whether the corresponding arrows are in Af.

First, note that $\{0\}$ is both initial and terminal in FinVect. In Af, it is easily seen that $\emptyset$ is initial and 0 is terminal in Af.

Let X, Y be two objects in Af. Assume first that both are proper. We define their product as
$$X \times Y := (V_X \times V_Y, A_X \times A_Y, (a_X, a_Y), \frac{1}{2}(\tilde{a}_X, \tilde{a}_Y)),$$
where
$$A_X \times A_Y := \{(x, y) \in V_X \times V_Y,\ x \in A_X, y \in A_Y\}$$
is the direct product of A_X and A_Y. It is easily verified that this is indeed an affine subspace and the usual projections $\pi_X : V_X \times V_Y \to V_X$ and $\pi_Y : V_X \times V_Y \to V_Y$ are in Af. Moreover, for $f : Z \to X$ and $g : Z \to Y$, the map $f \times g(z) = (f(z), g(z))$ is also clearly a morphism $Z \to X \times Y$ in Af. We have
$$L_{X\times Y} = L_X \times L_Y, \qquad S_{X\times Y} = (L_X \times L_Y) \vee \mathbb{R}(a_X, a_Y) = (S_X \times S_Y) \wedge \{(\tilde{a}_X, -\tilde{a}_Y)\}^{\perp}$$
for an arbitrary choice $a_X \in A_X$, $a_Y \in A_Y$ and $\tilde{a}_X \in \tilde{A}_X$, $\tilde{a}_Y \in \tilde{A}_Y$.

Next, we put $X \times \emptyset = \emptyset$, with the unique morphisms $\pi_X : \emptyset \to X$ and $\pi_\emptyset : \emptyset \to \emptyset$. If $Y \xrightarrow{f} X$ and $Y \xrightarrow{g} \emptyset$, then it is clear that $Y = \emptyset$, this shows that this is indeed the product. Further, put $X \times 0 = X$, with $\pi_X = id_X$ and $\pi_0 : X \xrightarrow{!} 0$. It is also readily verified that this is the product.

The coproduct for proper objects X, Y is defined as
$$X \oplus Y = (V_X \times V_Y, A_X \oplus A_Y)$$
where
$$A_X \oplus A_Y := \{(tx, (1-t)y),\ x \in A_X, y \in A_Y,\ t \in \mathbb{R}\}$$
is the direct sum. To check that this is an affine subspace, let $x_i \in A_X$, $y_i \in A_Y$, $s_i \in \mathbb{R}$ and let $\sum_i \alpha_i = 1$, then
$$\sum_i \alpha_i(s_i x_i, (1-s_i)y_i) = (\sum_i s_i\alpha_i x_i, \sum_i (1-s_i)\alpha_i y_i) = (sx, (1-s)y) \in A_X \oplus A_Y,$$

where $s = \sum_i s_i \alpha_i$, $x = s^{-1} \sum_i s_i \alpha_i x_i$ if $s \neq 0$ and is arbitrary in A_X otherwise, similarly $y = (1 - s)^{-1} \sum_i (1 - s_i) \alpha_i y_i$ if $s \neq 1$ and is arbitrary otherwise. The usual embeddings $p_X : V_X \to V_X \times V_Y$ and $p_Y : V_Y \to V_X \times V_Y$ are easily seen to be morphsims in Af.

Let $f : X \to Z$, $g : Y \to Z$ be any morphisms in Af and consider the map $V_X \times V_Y \to V_Z$ given as $f \oplus g(u, v) = f(u) + g(v)$. We need to show that it preserves the affine subspaces. So let $x \in A_X$, $y \in A_Y$, then since $f(x), g(y) \in A_Z$, we have for any $s \in \mathbb{R}$,

$$f \oplus g(sx, (1 - s)y) = sf(x) + (1 - s)g(y) \in A_Z.$$

We also have

$$L_{X \oplus Y} = (L_X \times L_Y) \vee \mathbb{R}\{(a_X, -a_Y)\}, \qquad S_{X \oplus Y} = S_X \times S_Y$$

for some $a_X \in A_X$, $a_Y \in A_Y$.

Similarly as in the case of products, it is verified that $X \oplus \emptyset = X$ and $X \oplus 0 = 0$. (All the statements for coproducts can be obtained from duality defined below).

One can also discuss equalizers and coequalizers, here we only note that these may be trivial even for proper objects. We will consider pullbacks and pushouts for some special morphisms that will be needed below. (We may also add two special objects: the initial object $\emptyset := (\{0\}, \emptyset)$ and the terminal object $0 := (\{0\}, \{0\})$, here the affine subspaces are obviously not proper.)

.1.5 Pullbacks and pushouts

.1.6 Monoidal structure

Let X, Y be objects in Af. Let us define

$$A_{X \otimes Y} := \{x \otimes y, x \in A_X, y \in A_Y\}^{\approx}.$$

In other words, $A_{X \otimes Y}$ is the smallest affine subspace in $V_X \otimes V_Y$ containing $A_X \otimes A_Y$.

.1.7 Duality

We define $X^* := (V_X^*, \tilde{A}_X)$. Note that we have

$$L_{X^*} = S_X^{\perp}, \qquad S_{X^*} = L_X^{\perp}, \qquad d_{X^*} = D_X - d_X - 1.$$

It is easily seen that $(-)^*$ defines a full and faithful functor $\mathrm{Af}^{op} \to \mathrm{Af}$, moreover, $X^{**} = X$ (we will use the canonical identification of any V in FinVect with its second dual).

.1.8 The dual tensor product

Let us define the dual tensor product by $\odot$, that is

$$X \odot Y = (X^* \otimes Y^*)^*.$$

We then have

$$L_{X \odot Y} = S_{X^* \otimes Y^*}^{\perp} = (S_{X^*} \otimes S_{Y^*})^{\perp} = (L_X^{\perp} \otimes L_Y^{\perp})^{\perp}$$
$$S_{X \odot Y} = L_{X^* \otimes Y^*}^{\perp} = (\tilde{a}_X \otimes S_Y^{\perp})^{\perp} \wedge (S_X^{\perp} \otimes \tilde{a}_Y)^{\perp} \wedge (S_X^{\perp} \otimes S_Y^{\perp})^{\perp}$$

In particular,

$$d_{X \odot Y} = D_X d_Y + d_X D_Y - d_X d_Y.$$

Lemma 12. *Let X, Y be nontrivial. Then $X \otimes Y = X \odot Y$ exactly in one of the following situations:*

1. *$X \simeq I$ or $Y \simeq I$,*

2. *$d_X = d_Y = 0$,*

3. *$D_X = d_X + 1$ and $D_Y = d_Y + 1$ (Objects with this property will be called first order).*

Proof. It is easy to see that (when identifying $X = X^{**}$), we have $A_X \otimes A_Y \subseteq \tilde{A}_{X^* \otimes Y^*}$, hence $A_{X \otimes Y} \subseteq A_{X \odot Y}$. We see from the above computatons that

$$d_{X \odot Y} - d_{X \otimes Y} = d_X(D_Y - 1) + (D_X - 1)d_Y - 2d_X d_Y \geq 0,$$

with equality if and only if the conditions of the lemma hold.

$\square$

.1.9 The no signaling product

Lemma 13. *The space $A_{X \otimes Y}$ is precisely the affine subspace of elements $w \in A_{X \odot Y}$, such that $\langle w, \cdot \otimes \tilde{a}_Y \rangle$ and $\langle w, \tilde{a}_X \otimes \cdot \rangle$ do not depend on the choice of $\tilde{a}_Y \in \tilde{A}_Y$ and $\tilde{a}_X \in \tilde{A}_X$.*

Proof. Any element $w \in A_{X \otimes Y}$ has the form $w = \sum_i \alpha_i x_i \otimes y_i$, for $x_i \in A_X$, $y_i \in A_Y$ and $\sum_i \alpha_i = 1$. It follows that for any $\tilde{a}_X \in \tilde{A}_X$ and $\tilde{a}_Y \in \tilde{A}_Y$,

$$\langle w, \cdot \otimes \tilde{a}_Y \rangle = \sum_i \alpha_i x_i, \qquad \langle w, \tilde{a}_X \otimes \cdot \rangle = \sum_i \alpha_i y_i.$$

Conversely, assume that $w \in A_{X \odot Y}$ has this property, then for any $\tilde{x} \in L_{X^*}$ and $\tilde{y} \in L_{Y^*}$, we have

$$\langle w, \cdot \otimes \tilde{y} \rangle = 0, \qquad \langle w, \tilde{x} \otimes \cdot \rangle = 0.$$

It follows that

$$w \in (V_X^* \otimes L_{Y^*})^{\perp} \cap (L_{X^*} \otimes V_Y^*)^{\perp} = (V_X \otimes S_Y) \cap (S_X \otimes V_Y) = S_X \otimes S_Y$$

Since $w \in A_{X \odot Y}$, we have $\langle w, \tilde{a}_X \otimes \tilde{a}_Y \rangle = 1$ for any choice of $\tilde{a}_X \in \tilde{A}_X$, $\tilde{a}_Y \in \tilde{A}_Y$. Since $\tilde{a}_X \otimes \tilde{a}_Y \in \tilde{A}_{X \otimes Y}$ and $S_{X \otimes Y} = S_X \otimes S_Y$, this implies $w \in A_{X \otimes Y}$.

$\square$

We will now define one-sided variants of this property. Namely, let

$$A_{X \prec Y} := \{w \in A_{X \odot Y}, \ \langle w, \cdot \otimes \tilde{a}_Y \rangle \text{ does not depend on } \tilde{a}_Y \in \tilde{A}_Y\}$$
$$A_{X \succ Y} := \{w \in A_{X \odot Y}, \ \langle w, \tilde{a}_X \otimes \cdot \rangle \text{ does not depend on } \tilde{a}_X \in \tilde{A}_X\}.$$

We then put $X \prec Y = (V_X \otimes V_Y, A_{X \prec Y})$ and $X \succ Y = (V_X \otimes V_Y, A_{X \succ Y})$.

Lemma 14. *For any choice of $a_Y \in A_Y$, we have*

$$L_{X \prec Y} = V_X \otimes L_Y + L_X \otimes a_Y, \qquad S_{X \prec Y} = V_X \otimes L_Y + S_X \otimes a_Y.$$

Similarly, for any $a_X \in A_X$, we have

$$L_{X \succ Y} = L_X \otimes V_Y + a_X \otimes L_Y, \qquad S_{X \succ Y} = L_X \otimes V_Y + a_X \otimes S_Y.$$

Proof. By the definition, we have

$$L_{X \prec Y} = L_{X \odot Y} \cap (V_X^* \otimes L_{Y^*})^\perp = (S_{X^*} \otimes S_{Y^*})^\perp \cap (V_X \otimes S_Y) = V_X \otimes L_Y + L_X \otimes a_Y,$$

for any element $a_Y \in A_Y$. The proof for $\succ$ is similar.

□

For the rest of this section, we fix some $a_X \in A_X$ and $\tilde{a}_X \in \tilde{A}_X$. We will use the notation $X_{\min} := (V_X, \{a_X\})$ and $X_{\max} = (V_X, \{\tilde{a}_X\}^{\sim})$.

We also introduce a decomposition of V_X into an independent family of subspaces L_X^0, L_X^1, L_X^2 as $L_X^0 := \mathbb{R}a_X$, $L_X^1 := L_X$ and L_X^2 is any complement of L_X in the subspace $\{\tilde{a}_X\}^\perp$. We see that the L or S spaces of any of the objects discussed in this paragraph is a union of some of the subspaces $L_X^i \otimes L_Y^j$, $i, j = 0, 1, 2$. We may therefore represent the subspaces in question by 3×3 matrices such that the i,j-th element is 1 if the subspace contains $L_X^i \otimes L_Y^j$ and 0 otherwise. For example, we have

$$L_{X \otimes Y} \equiv \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 0 \end{bmatrix}, \ L_{X \odot Y} \equiv \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 0 \end{bmatrix}, \ L_{X \prec Y} \equiv \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \end{bmatrix}, \ L_{X \succ Y} \equiv \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 0 & 0 \end{bmatrix}. \quad (8)$$

The S-spaces can be obtained from these by putting the top left element to 1. Notice also that all these objects belong to the set of objects $Z = (V_X \otimes V_Y, A_Z)$ such that

$$X_{\min} \otimes Y_{\min} \hookrightarrow Z \hookrightarrow X_{\max} \otimes Y_{\max} (= X_{\max} \odot Y_{\max})$$

and this set forms a lattice under inclusion ordering with $Z_0 \wedge Z_1 = Z_0 \sqcap Z_1$, $Z_0 \vee Z_1 = Z_0 \sqcup Z_1$, where $\sqcap$ and $\sqcup$ are the pullback and pushout of the inclusions $X_{\min} \otimes Y_{\min} \hookrightarrow Z_i \hookrightarrow X_{\max} \otimes Y_{\max}$. Furthermore, since $\{L_X^i \otimes L_Y^j\}$ is an independent decomposition of $V_X \otimes V_Y$, all the objects in (8) are contained in a distributive sublattice of objects such that $A_Z = a_X \otimes a_Y + L$, where L is a subspace represented by a matrix M_Z with the top left element equal to 0. For such elements Z_1 and Z_2, the representing matrices $M_{Z_1 \sqcap Z_2}$ resp. $M_{Z_1 \sqcup Z_2}$ are given by poinwise minimum resp. maximum of M_{Z_1} and M_{Z_2}.

Some further useful elements of this sublattice are represented as

$$L_{X_{\max} \otimes Y} \equiv \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \end{bmatrix}, \ L_{X \otimes Y_{\max}} \equiv \begin{bmatrix} 0 & 1 & 1 \\ 1 & 1 & 1 \\ 0 & 0 & 0 \end{bmatrix}, \ L_{X_{\min} \odot Y} \equiv \begin{bmatrix} 0 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \end{bmatrix}, \ L_{X \odot Y_{\min}} \equiv \begin{bmatrix} 0 & 0 & 0 \\ 1 & 1 & 1 \\ 0 & 0 & 0 \end{bmatrix}.$$

From these consideration, the following is immediate.

Lemma 15. *We have*

$$\begin{aligned} X \prec Y &= (X \otimes Y) \sqcup (X_{\min} \odot Y) = (X \odot Y) \sqcap (X_{\max} \otimes Y) \\ X \succ Y &= (X \otimes Y) \sqcup (X \odot Y_{\min}) = (X \odot Y) \sqcap (X \otimes Y_{\max}) \\ X \otimes Y &= (X \prec Y) \sqcap (X \succ Y) \\ X \odot Y &= (X \prec Y) \sqcup (X \succ Y). \end{aligned}$$

We have the inclusions

$$X_a \to X \to X^{\tilde{a}}.$$

We also fix $b \in A_Y$, $\tilde{b} \in \tilde{A}_Y$. By inclusions, we have the following diagrams

$$X_a \otimes Y_b \to X \otimes Y, \qquad X_a \otimes Y_b \to X_a \odot Y$$

and

$$X \odot Y \to X^{\tilde{a}} \odot Y^{\tilde{b}} = X^{\tilde{a}} \otimes Y^{\tilde{b}}, \qquad X \otimes Y^{\tilde{b}} \to X^{\tilde{a}} \otimes Y^{\tilde{b}}.$$

Lemma 16. $X \prec Y = X \otimes Y \sqcup X \odot Y_b = X \odot Y \sqcap X \otimes Y^{\tilde{b}}.$

We see that the identity map $id_{V_X \otimes V_Y}$ defines bimorphisms

$$X \otimes Y \to X \prec Y \to X \odot Y, \qquad X \otimes Y \to X \succ Y \to X \odot Y.$$

Lemma 17. *The pushout and pullback of the above diagram are*

$$X \otimes Y = (X \prec Y) \sqcap (X \succ Y), \qquad X \odot Y = (X \prec Y) \sqcup (X \succ Y).$$

Proof. We have by Lemma 13 that

$$A_{X \otimes Y} = A_{X \prec Y} \cap A_{X \succ Y} = A_{(X \prec Y) \sqcap (X \succ Y)},$$

we clearly have the last equality since the intersection $A_{X \prec Y} \cap A_{X \succ Y}$ is nonempty. For the second part,

$\square$

To each object $X = (V_X, A_X, a_X, \tilde{a}_X)$ we may define two object

$$X_{\min} := (V_X, \{a_X\}, a_X, \tilde{a}_X), \qquad X_{\max} := (V_X, \{\tilde{a}_X\}^{\sim}, a_X, \tilde{a}_X).$$

It is easily seen that $X_{\min} = (X^*_{\max})^*$ and $X_{\max} = (X^*_{\min})^*$, moreover, $X_{\max}$ and $(X_{\min})^*$ are first order objects. We have the inclusions

$$X_{\min} \xrightarrow{id} X \xrightarrow{id} X_{\max}.$$

We also have the inclusions

$$X \otimes Y \to X \odot Y \to X_{\max} \odot Y_{\max} = X_{\max} \otimes Y_{\max}$$

and

$$X \otimes Y \to X_{\max} \otimes Y \to X_{\max} \otimes Y_{\max}, \quad X \otimes Y \to X \otimes Y_{\max} \to X_{\max} \otimes Y_{\max}.$$

We can therefore define pullbacks and pushouts, which then becomes

$$(X_{\max} \otimes Y) \sqcap (X \otimes Y_{\max}) = X \otimes Y, \qquad (X_{\max} \otimes Y) \sqcup (X \otimes Y_{\max}) = X_{\max} \otimes Y_{\max}.$$

Hence we may decompose $X \odot Y$ into two parts

$$X \prec Y := (X \odot Y) \sqcap (X_{\max} \otimes Y), \qquad X \succ Y := (X \odot Y) \sqcap (X \otimes Y_{\max}).$$

Note that these forms do not depend on the choice of the elements $a_X, a_Y...$!

Lemma 18. *We have*

$$X \otimes Y = (X \prec Y) \sqcap (X \succ Y), \qquad X \odot Y = (X \prec Y) \sqcup (X \succ Y).$$

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Proof. We have

$$(X \prec Y) \sqcap (X \succ Y) = ((X \odot Y) \sqcap (X_{\max} \otimes Y)) \sqcap ((X \odot Y) \sqcap (X \otimes Y_{\max}))$$
$$= (X \odot Y) \sqcap ((X_{\max} \otimes Y) \sqcap (X \otimes Y_{\max})) = (X \odot Y) \sqcap (X \otimes Y) = X \otimes Y.$$

The second equality follows easily from Lemma 14 □

We next show that $- \prec -$ and $- \succ -$ define a functor $\mathrm{Af} \times \mathrm{Af} \to \mathrm{Af}$. Let $X_1 \xrightarrow{f} Y_1$ and $X_2 \xrightarrow{g} Y_2$, we will show that $X_1 \prec Y_1 \xrightarrow{f \otimes g} X_2 \prec Y_2$. For this, we need to prove that $(f \otimes g)(A_{X_1 \prec X_2}) \subseteq A_{X_2 \prec Y_2}$ and $(f \otimes g)(A_{X_1 \succ X_2}) \subseteq A_{X_2 \succ Y_2}$. This is clear from Lemma 14.

Lemma 19. $(\mathrm{Af}, \prec, I)$ *is a monoidal category.*

Proof. It is easily checked from Lemma 14 that $\alpha_{X,Y,X}(L_{(X \prec Y) \prec Z}) = L_{X \prec (Y \prec Z)}$ and clearly also $\alpha_{X,Y,Z}(a_X \otimes a_Y \otimes a_Z) = a_X \otimes a_Y \otimes a_Z$, so that α is the associator. Since $I \otimes X = I \odot X$ and $X \odot I = X \otimes I$, we have $I \prec X = I \otimes X$ and $X \prec I = X \otimes I$, so λ and ρ are the unitors. But note that $\sigma_{X,Y}(A_{X \prec Y}) = A_{Y \succ X}$, so this structure is not symmetric. □

We have $(X \prec Y)^* = X^* \prec Y^*$. Indeed, by duality,

$$(X \prec Y)^* = ((X \odot Y) \sqcap (X_{\max} \odot Y))^* = (X \odot Y)^* \sqcup (X_{\max} \otimes Y)^*$$
$$= (X^* \otimes Y^*) \sqcup (X^*_{\max} \odot Y^*)$$

.1.10 Internal hom

The internal hom has the form

$$[X, Y] = (X \otimes Y^*)^* = X^* \odot Y. \tag{9}$$

We then have

$$L_{[X,Y]} = (S_X \otimes L_Y^\perp)^\perp = (V_X^* \otimes L_Y) \vee (L_{X^*} \otimes V_Y),$$
$$S_{[X,Y]} = (\tilde{a}_X \otimes S_Y^\perp)^\perp \wedge (L_X \otimes \tilde{a}_Y)^\perp \wedge (L_X \otimes S_Y^\perp)^\perp = (V_X^* \otimes L_Y) \vee (L_{X^*} \otimes V_Y) \vee (\tilde{a}_X \otimes a_Y).$$

and

$$d_{[X,Y]} = D_X D_Y - (d_X + 1)(D_Y - d_Y).$$

As we have seen in FinVect, the space $V_{[X,Y]} = V_X^* \otimes V_Y$ is identified with the space of all linear maps $V_X \to V_Y$, by (**??**). We will show that $A_{[X,Y]}$ corresponds to the affine subspace of maps mapping A_X into A_Y, that is, morphisms in Af. Indeed, we see from (**??**) that f is in Af if and only if

$$\langle f(x), y^* \rangle = \langle w, x \otimes y^* \rangle = 1, \qquad x \in A_X, \ y^* \in \tilde{A}_Y,$$

which is equivalent to $w \in (A_X \otimes \tilde{A}_Y)^\sim = \tilde{A}_{X \otimes Y^*}$.

.1.11 The no signaling product

For two objects X, Y we define

$$X \prec Y := (V_X \otimes V_Y, A_{X \prec Y}, a_X \otimes a_Y, \tilde{a}_X \otimes \tilde{a}_Y),$$

where $A_{X \prec Y}$ is determined by

$$S_{X \prec Y} = V_X \otimes S_Y \cap S_{X \odot Y} = (V_X \otimes S_Y) \cap (L_{X \odot Y} \vee \{a_X \otimes a_Y\}) = (V_X \otimes S_Y) \cap ((V_X \otimes L_Y) \vee (L_X \otimes V_Y) \vee \{a_{X}$$

Lemma 20.

We may similarly define $X \succ Y$. Setting $f \prec g = f \otimes g$ for $X_1 \xrightarrow{f} X_2$, $Y_1 \xrightarrow{g} Y_2$, we see that $(- \prec -)$ is functorial. Indeed, to show that $X_1 \prec Y_1 \xrightarrow{f \otimes g} X_2 \prec Y_2$, we need to show that $f \otimes g(A_{X_1 \prec Y_1}) \subseteq A_{X_2 \prec Y_2}$. Assume $w \in A_{X_1 \prec Y_1}$, that is, $w \in V_{X_1} \otimes S_{Y_1}$ and $\langle w, \tilde{a}_{X_1} \otimes \tilde{a}_{Y_1} \rangle = 1$.

.1.12 Dualizable (nuclear) objects

An object in Af is nuclear if the natural map $X^* \otimes X \to [X, X]$ is an isomorphism (santocanale). That is, the inclusion $X^* \otimes X \subseteq X^* \odot X$ that comes from the embedding

$$\tilde{A}_X \otimes A_X \subseteq (A_X \otimes \tilde{A}_X)\tilde{}$$

becomes an equality. As we have seen in Lemma 12, for proper objects we have $X^* \otimes X = X^* \odot X$ if and only if

$$d_X + 1 = D_X = D_{X^*} = d_{X^*} + 1 = D_X - d_X.$$

It follows that $d_X = 0$ and $D_X = 1$, so that $X \simeq I$. Hence the tensor unit is the unique dualizable (or nuclear) object in Af.

.1.13 No signaling

We say that $X \xrightarrow{f} Y$ is no signaling if

$$y^* \circ f = \tilde{a}_Y \circ f, \qquad \forall y \in Y^* = [Y, I],$$

in other words

$$y \circ f = 0, \qquad \forall y^* \in L_{Y^*} = S_Y^\perp.$$

Taking $w \in A_{[X,Y]}$ be the corresponding elements, this means that

$$\langle w, x \otimes y^* \rangle = 0, \qquad \forall x \in V_X, \ y^* \in S_Y^\perp,$$

in other words

$$w \in (V_X \otimes S_Y^\perp)^\perp = V_X^* \otimes S_Y,$$

so that

$$w \in A_{[X,Y]} \cap (V_X^* \otimes S_Y).$$

Since $a_{[X,Y]} = \tilde{a}_X \otimes a_Y \in V_X^* \otimes S_Y$, we have that

$$a_{[X,Y]} - w \in L_{[X,Y]} \cap V_X^* \otimes S_Y = (S_X \otimes L_Y^\perp)^\perp \cap V_X^* \otimes S_Y = (V_X^* \otimes L_Y) + (L_{X^*} \otimes a_Y).$$

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We can also define no signaling in the oposite way, that is,

$$f(x) = f(a_X), \qquad \forall x \in A_X.$$

This is of course the same as

$$f(x) = 0, \qquad \forall x \in L_X,$$

or

$$\langle w, x \otimes y^* \rangle = 0, \qquad \forall x \in L_X, \ y^* \in V_Y^*,$$

that is,

$$w \in (L_X \otimes V_Y^*)^\perp = L_X^\perp \otimes V_Y = S_{X^*} \otimes V_Y.$$

It follows that

$$\tilde{a}_X \otimes a_Y - w \in L_{[X,Y]} \cap S_{X^*} \otimes V_Y = L_{X^*} \otimes V_Y + \tilde{a}_X \otimes L_Y.$$

.2   Once more on the monoidal structures

.2.1   Tensor product

We have

$$L_{X \otimes Y} = (a_X \otimes L_Y) + (L_X \otimes a_Y) + (L_X \otimes L_Y) = (S_X \otimes L_Y) + (L_X \otimes a_Y) = (L_X \otimes S_Y) + (a_X \otimes L_Y)$$

A closed symmetric monoidal structure. We have

$$L_{X_{\max} \otimes Y} = V_X \otimes L_Y + \{\tilde{a}_X\}^\perp \otimes a_Y, \qquad L_{X_{\min} \otimes Y} = a_X \otimes L_Y.$$

Lemma 21. *We have*

$$X \otimes Y = (X_{\max} \otimes Y) \sqcap (X \otimes Y_{\max}).$$

Proof. This is easy, since

$$S_{X_{\max} \otimes Y} \cap S_{X \otimes Y_{\max}} = (V_X \otimes S_Y) \cap (S_X \otimes V_Y) = S_X \otimes S_Y.$$

$\square$

.2.2   Dual product

By definition, $X \odot Y = (X^* \otimes Y^*)^*$. We have

$$L_{X \odot Y} = (L_X^\perp \otimes L_Y^\perp)^\perp = (V_X \otimes L_Y) \vee (L_X \otimes V_Y).$$

We have

$$L_{X_{\max} \odot Y} = (V_X \otimes L_Y) \vee (\{\tilde{a}_X\}^\perp \otimes V_Y), \qquad L_{X_{\min} \odot Y} = V_X \otimes L_Y.$$

Lemma 22. *We have*

$$
\begin{aligned}
X_{\max} \otimes Y_{\max} &= (X_{\max} \odot Y) \sqcup (X \odot Y_{\max}) \\
X \odot Y &= (X_{\min} \odot Y) \sqcup (X \odot Y_{\min}) \\
X \otimes Y &= (X_{\min} \otimes Y) \sqcup (X \otimes Y_{\min}) \sqcup (X_{\min} \odot Y \sqcap X \odot Y_{\min})
\end{aligned}
$$

Proof. The first is easy, the seconf follows from Lemma 21 by duality, the third is also easy.

$\square$

.3 The no signalling product

Let us define $X \prec Y := (X \odot Y) \sqcap (X_{\max} \otimes Y)$. We have

$$L_{X \prec Y} = V_X \otimes L_Y + L_X \otimes a_Y = L_{X_{\min} \odot Y} + L_{Y \otimes Y_{\min}}.$$

So that

$$X \prec Y := (X \odot Y) \sqcap (X_{\max} \otimes Y) = (X_{\min} \odot Y) + (X \otimes Y_{\min})$$

Lemma 23. *We have*

$$(X \otimes Y) \sqcup (X_{\min} \odot Y) = (X \odot Y) \sqcap (X_{\max} \otimes Y)$$
$$= (X_{\min} \odot Y) + (X \otimes Y_{\min}) = (X_{\max} \otimes Y) \sqcap (X \odot Y_{\max}).$$

Let us denote the above object by $X \prec Y$. Then $A_{X \prec Y}$ is the set of elements in $V_X \otimes V_Y$ such that $\langle w, \cdot \otimes y^* \rangle$ is a fixed element in A_X, independently of $y^* \in \tilde{A}_Y$.

Blbe uvedenie, definicia!

Proof. We see that $A_{X \prec Y} \subseteq A_{X \odot Y}$, moreover,

$$A_{X \prec Y} = \{w \in A_{X \odot Y}, \ \langle w, id_X \otimes y^* \rangle = 0, \ \forall y^* \in L_{Y^*}\}.$$

In other words, since clearly $a_X \otimes a_Y \in A_{X \prec Y}$,

$$L_{X \prec Y} = \{w - a_X \otimes a_Y, \ w \in A_{X \prec Y}\} = L_{X \odot Y} \cap (V_X^* \otimes L_Y^*)^\perp = (L_X^\perp \otimes S_{Y^*})^\perp \cap (V_X^* \otimes L_{Y^*})^\perp$$
$$= ((L_X^\perp \otimes S_{Y^*}) \vee (V_X^* \otimes L_{Y^*}))^\perp = ((V_X^* \otimes L_{Y^*}) + (S_X^* \otimes \tilde{a}_Y))^\perp = S_{(X \prec Y)^*}^\perp$$

But also

$$L_{X \odot Y} \cap (V_X^* \otimes L_Y^*)^\perp = ((V_X \otimes L_Y) \vee L_X \otimes V_Y) \cap (V_X \otimes S_Y)$$
$$= ((V_X \otimes L_Y) + (L_X \otimes a_Y) + (L_X \otimes S_Y')) \cap ((V_X \otimes L_Y) + (L_X \otimes a_Y) + (L_X' \otimes a_Y))$$
$$= ((V_X \otimes L_Y) + (L_X \otimes a_Y).$$

First note that we have

$$L_{X_{\min} \odot Y} = V_X \otimes L_Y, \qquad L_{X \otimes Y_{\min}} = L_X \otimes a_Y$$

and therefore

$$L_{X \otimes Y_{\min}} \cap L_{X_{\min} \odot Y} = \{0\}.$$

Further,

$$L_{X \otimes Y} = (S_X \otimes L_Y) + (L_X \otimes a_Y) = (S_X \otimes L_Y) + L_{X \otimes Y_{\min}}$$

and

$$L_{X \odot Y} = (V_X \otimes L_Y) \vee (L_X \otimes V_Y), \qquad L_{X_{\max} \otimes Y} = (V_X \otimes L_Y) + (\{\tilde{a}_X\}^\perp \otimes a_Y)$$

We have

$$S_{(X \otimes Y) \sqcup (X_{\min} \odot Y)} = S_{X \otimes Y} \vee S_{X_{\min} \odot Y} = S_{X \otimes Y} \vee L_{X_{\min} \odot Y}$$

$\square$

.4 The category AfH

The category AfH will be constructed as a subcategory in Af.

.4.1 First order objects

It is easily seen that the following are equivalent:

1. $D_X = d_X + 1$;

2. $S_X = V_X$;

3. $L_X = \{\tilde{a}_X\}^\perp$;

4. $S_{X^*} = \mathbb{R}\tilde{a}_X$;

5. $L_{X^*} = \{0\}$.

We say that an object X is first order if any of these conditions is fulfilled. We have seen that for proper objects, $X \otimes Y = X \odot Y$ if and only if both X and Y are first order. We also have

Lemma 24. *X is first order if and only if $X \otimes X = X \odot X$.*

ʿ

Lemma 25. *Let X, Y be first order, then $X \otimes Y$ is first order.*

Proof. We have
$$S_{X \otimes Y} = S_X \otimes S_Y = V_X \otimes V_Y = V_{X \otimes Y}.$$
$\square$

.4.2 Channels

A channel is an object $[X, Y]$ where X and Y are first order. As we have seen,
$$X^* \otimes Y \subseteq X^* \odot Y = [X, Y].$$

If X is first order, $\tilde{A}_X = \{\tilde{a}_X\}$ and the elements of $A_{X^* \otimes Y} = \tilde{a}_X \otimes A_Y$ are identified with channels of the form
$$f(x) = \langle \tilde{a}_X, x \rangle y, \qquad x \in V_X,$$
for some $y \in A_Y$. Such maps will be called replacement channels.

Lemma 26. *Let X, Y be first order and let $w \in V_X^* \otimes V_Y$. Then $w \in A_{[X,Y]}$ if and only if*
$$\circ_Y : w_{X^*Y} \otimes \tilde{a}_Y \mapsto \tilde{a}_X.$$

Proof. Let $f : V_X \to V_Y$ be the map corresponding to w, then
$$\circ_Y(w \otimes \tilde{a}_Y) = (V_X^* \otimes e_{V_Y})(w \otimes \tilde{a}_Y) = \tilde{a}_Y \circ f,$$

where $\tilde{a}_Y \in V_Y^*$ is seen as a map $V_y \to \mathbb{R}$. So $\tilde{a}_Y \circ f : V_X \to \mathbb{R}$ is an element in V_X^*. We know that $w \in A_{[X,Y]}$ iff $f(A_X) \subseteq A_Y$, which is equivalent to $\tilde{a}_Y \circ f(x) = 1$ for all $x \in A_X$, so that $\tilde{a}_Y \circ f \in \tilde{A}_X = \{\tilde{a}_X\}$, since X is first order. $\square$

Lemma 27. *Let Y be first order and $w \in V_X^* \otimes V_Y$. Then $w \in A_{[X,Y]}$ if and only if*

$$\circ_Y(w_{X^*Y} \otimes \tilde{a}_Y) \in \tilde{A}_X.$$

Moreover,

$$\tilde{A}_{[X,Y]} = A_X \otimes \{\tilde{a}_Y\}.$$

Proof. Since Y is first order, we have $A_{Y^*} = \tilde{A}_Y = \{\tilde{a}_Y\}$ and by (9)

$$\tilde{A}_{[X,Y]} = A_{X \otimes Y^*} = A_X \otimes \{\tilde{a}_Y\}.$$

As in the above proof, let $f : V_X \to V_Y$ be the map corresponding to w. Then $\tilde{a}_Y \circ f \in V_X^*$ and $w \in A_{[X,Y]}$ iff $f(A_X) \subseteq A_Y$. This means that

$$\tilde{a}_Y \circ f(x) = 1, \qquad \forall x \in A_X,$$

which means that $\tilde{a}_Y \circ f \in \tilde{A}_X$.

$\square$

.4.3 AfH

The category AfH is the full subcategory in Af created from first order objects by taking tensor products and duals. We will add more later. We will use the notation V_{XY^*} for $V_X \otimes V_Y^*$, etc.

Any object X in AfH is created from first order objects $X_1, \ldots, X_k$, so that $V_X = \tilde{V}_{X_1} \otimes \cdots \otimes \tilde{V}_{X_k}$, where $\tilde{V}_{X_i}$ is either V_{X_i} or $V_{X_i}^*$, $i = 1, \ldots, k$. We will next show that any object is a set of channels that contains all replacement channels.

Proposition 7. *Let X be an object in* AfH. *Then there are first order objects Y_I and Y_O and inclusions f, g such that*

$$Y_I^* \otimes Y_O \xrightarrow{f} X \xrightarrow{g} [Y_I, Y_O]. \tag{10}$$

Proof. Let X be first order, then since I is first order,

$$I^* \otimes X = I \otimes X \xrightarrow{\lambda_X} X \xrightarrow{\lambda_X^{-1}} I \otimes X = I \odot X = [I, X].$$

Clearly, $f = \lambda_X$ and $g = \lambda_X^{-1}$ are inclusions. Now assume that Z satisfies (10) and let $X = Z^*$. Taking duals and composing with symmetries, we get

$$Y_O^* \otimes Y_I \xrightarrow{\sigma_{V_{Y,O}^*, V_{Y_I}}} Y_I \otimes Y_0^* = [Y_I \otimes Y_O]^* \xrightarrow{g^*} X \xrightarrow{f^*} (Y_I^* \otimes Y_O)^* \xrightarrow{\sigma_{V_{Y_I}, V_{X_O}^*}} (Y_O \otimes Y_I)^* = [Y_O, Y_I].$$

Since the compositions of f^* and g^* with symmetries are inclusions, we see that X satisfies (10).

Next, let X_1 and X_2 satisfy (10) with some first order objects Y_I^i, Y_O^i and inclusions f^i, g^i, $i = 1, 2$, and let $X = X_1 \otimes X_2$. We then have, using the appropriate symmetries

$$Y_I^1 Y_I^2 \otimes (Y_O^1 Y_O^2)^* \xrightarrow{\sigma_{Y_I^2, Y_O^1}} Y_I^1 \otimes (Y_O^1)^* \otimes Y_I^2 \otimes (Y_O^1)^* \xrightarrow{f^1 \otimes f^2} X \xrightarrow{g^1 \otimes g^2} [Y_I^1, Y_O^1] \otimes [Y_I^2, Y_O^2] \xrightarrow{\sigma_{Y_O^1, Y_I^2}} [Y_I^1 Y_I^2, I_O^1 Y_O^2].$$

Perhaps the last arrow needs some checking, so lt us do it properly. We need to show that for $w \in A_{[Y_I^1, Y_O^1] \otimes [Y_I^2, Y_O^2]}$, we have $\sigma_{Y_O^1, Y_O^2}(w) \in A_{[Y_I^1 Y_I^2, I_O^1 Y_O^2]}$, but this is clear using Lemma 26.

$\square$

The pair (Y_I, Y_O) for an object X will be called the setting of X. For objects of the same setting we may take pullbacks and pushouts of the corresponding inclusions.

Pullbacks are intersections, pushouts the affine mixture.

Channels into (from) products and coproducts

We define AfH as the full subcategory of Af containing all first order objects and closed under (finite products,) duals and tensor products .