

CSE2121: Data Structure

2019 4c

Given the following input (4322, 1334, 1471, 9679, 1989, 6171, 6173, 4199) and the hash function $x \bmod 10$, which of the following statements are true? Explain.

- i. 9679, 1989, 4199 hashes to the same value
- ii. 1471, 6171 has to the same value
- iii. All elements hash to the same value
- iv. Each element hashes to a different value
- v. None of the above.

Now, let's check each statement:

i. 9679, 1989, 4199 hash to the same value

Yes, they all hash to 9. So, this statement is true.

ii. 1471, 6171 has to the same value

Yes, they both hash to 1. So, this statement is also true.

iii. All elements hash to the same value

No, this is not true. The elements hash to different values.

iv. Each element hashes to a different value

No, this is also not true. Some elements hash to the same value.

Therefore, the correct answer is option 'C': i and ii only.

The given input is: 4322, 1334, 1471, 9679, 1989, 6171, 6173, 4199

The hash function is: $x \bmod 10$

So, the possible hash values are: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9

Let's apply the hash function to each input:

- $4322 \bmod 10 = 2$
- $1334 \bmod 10 = 4$
- $1471 \bmod 10 = 1$
- $9679 \bmod 10 = 9$
- $1989 \bmod 10 = 9$
- $6171 \bmod 10 = 1$
- $6173 \bmod 10 = 3$
- $4199 \bmod 10 = 9$

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3(c) Simulate the postfix expression evaluation algorithm using 12, 6, /, 6, 2, +, *, 12, 4, /, - by showing Stack's contents as each element is scanned (2.75)

2016

element	Stack	Operation
12	12	
6	12, 6	
/	2	[12/6]
6	2, 6	
2	2, 6, 2	
+	2, 8	[6+2]
*	16	[2*8]
12	16, 12	
4	16, 12, 4	
/	16, 3	[12/4]
-	13	[16-3]

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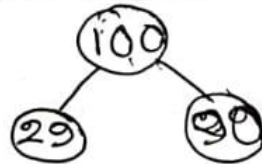
(c) Simulate the maxheap algorithm for the following values: 100, 29, 90, 148, 12, 34, 90, 89, 120, 99 and 12. 3.75

2019

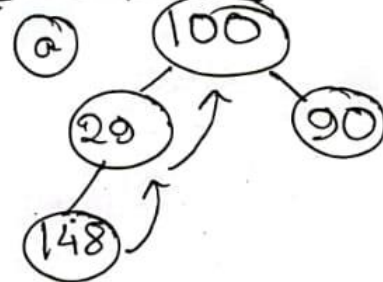
6) © Maxheap algorithm-

Given values:- 100, 29, 90, 148, 12, 34, 90, 89, 120, 99 & 12

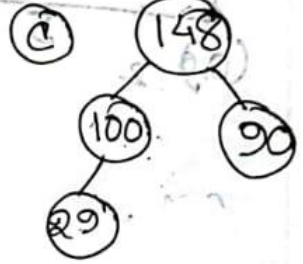
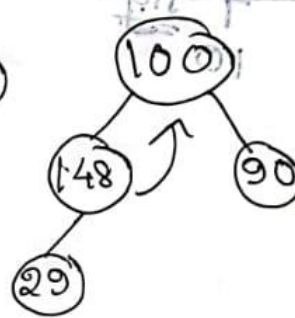
Step-1:



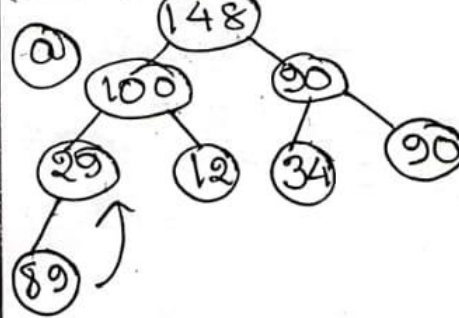
Step-2:



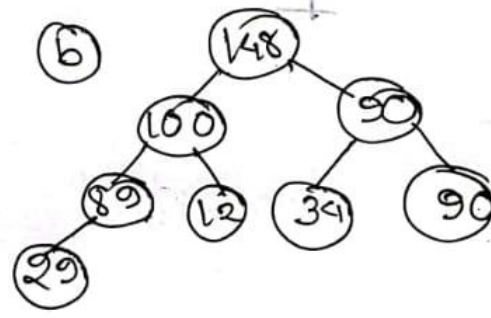
(b)



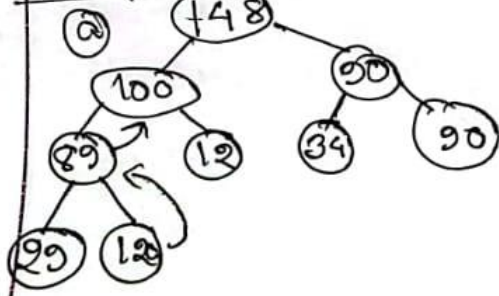
Step-3:



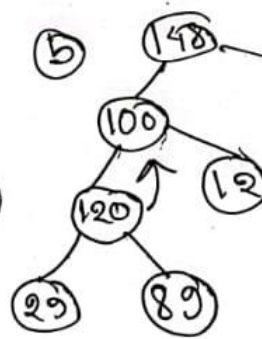
(b)



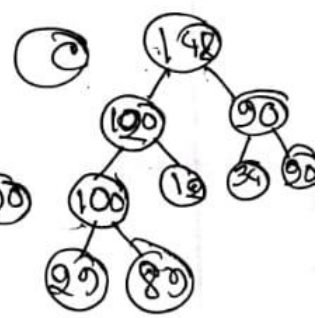
Step-4:-



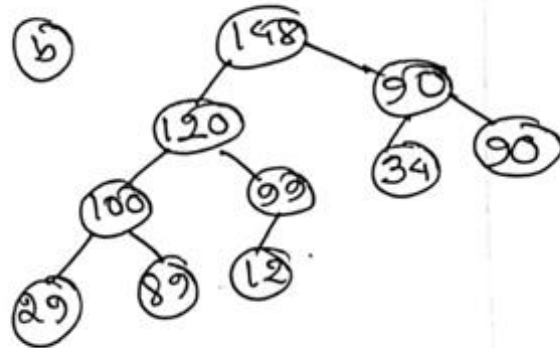
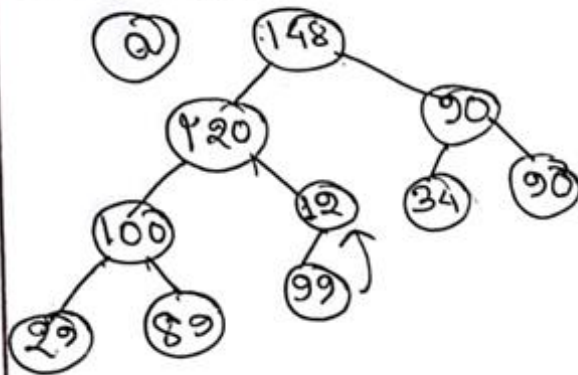
(b)



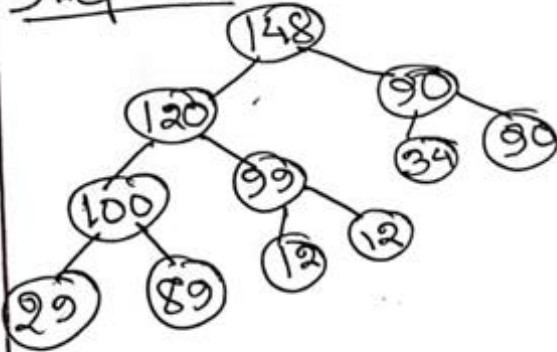
(c)



Step -5:-



Step -6:-



2019 7-c

Briefly discuss the advantages of binary search trees. Use the Warshall's algorithm to find the shortest path matrix of the weighted matrix given below. ...4.75

$$W = \begin{pmatrix} 2 & 5 & 0 & 0 \\ 2 & 0 & 1 & 4 \\ 5 & 3 & 3 & 3 \\ 4 & 2 & 3 & 0 \end{pmatrix}$$

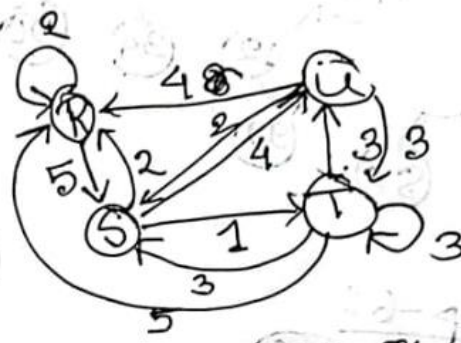
Warshall's algorithm to find shortest path matrix of weighted matrix -

$$W = \begin{matrix} & \begin{matrix} v_1 & v_2 & v_3 & v_4 \end{matrix} \\ \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{matrix} & \begin{pmatrix} 2 & 5 & \infty & \infty \\ 2 & 0 & 14 & \\ 5 & 3 & 3 & 3 \\ 4 & 2 & 3 & 0 \end{pmatrix} \end{matrix}$$

$$R=v_1, S=v_2, T=v_3, U=v_4$$

$$Q_k[i,j] = \min(Q_{k-1}[i,j], Q_{k-1}[i,k] + Q_{k-1}[k,j])$$

$$Q_0 = \begin{pmatrix} 2 & 5 & \infty & \infty \\ 2 & \infty & 14 & \\ 5 & 3 & 3 & 3 \\ 4 & 2 & 3 & \infty \end{pmatrix}$$



RR	RS	RT	RU
SR	SS	ST	SU
TR	TS	TT	TU
UR	US	UT	UU

$$Q_1 = \begin{pmatrix} 2 & 5 & \infty & \infty \\ 2 & \infty & 14 & \\ 5 & 3 & 3 & 3 \\ 4 & 2 & 3 & \infty \end{pmatrix}$$

Q_1 এর জন্য 1st row & 1st column as it is নিবে।

Min $(D^0[2,2], D^0[2,1] + D^0[1,2])$ এখন

$$\infty, 2 + 5 = 7$$

$$\infty > 7 \text{ So take } 7$$

1 দিলে গেলে & direct গেলে যোগ করে ছোট সেটা নিবে।

$$D^0[2,3], D^0[2,1] + D^0[1,3]$$

$$1, 2 + \infty$$

$$1, \infty$$

$$1 < \infty$$

$$Q_2 = \begin{pmatrix} 2 & 5 & 6 & 11 \\ 2 & 7 & 1 & 4 \\ 5 & 3 & 3 & 3 \\ 4 & 2 & 3 & 6 \end{pmatrix}$$

$$Q_3 = \begin{pmatrix} 2 & 5 & 6 & 6 \\ 2 & 7 & 1 & 4 \\ 5 & 3 & 3 & 3 \\ 4 & 2 & 3 & 6 \end{pmatrix}$$

$$Q_4 = \begin{pmatrix} 2 & 5 & 6 & 6 \\ 2 & 7 & 1 & 4 \\ 5 & 3 & 3 & 3 \\ 4 & 2 & 3 & 6 \end{pmatrix}$$

2019 8c

How many ways a graph can be traversed? Consider the adjacency list of the Graph G in the following table. Find the nodes that are reachable from node C using Depth First Search.

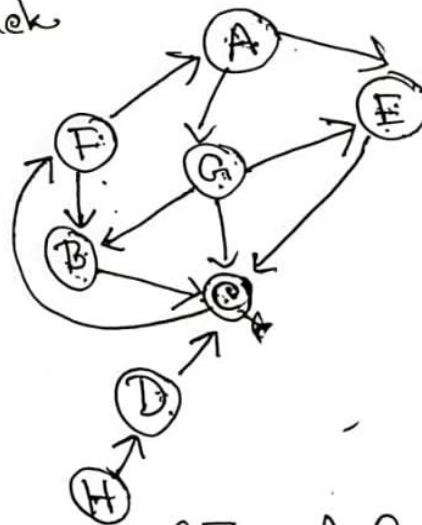
Node	Adjacency	Node	Adjacency
A	G, E	E	C
B	C	F	A, B
C	F	G	B, C, E
D	C	H	D

8 ©:- Reachable from node C using DFS. (3)

DFS traversal: Stack



consider C as root



Result: C, F, B, A, G, E

∴ Reachable from C are - C, F, B, A, G & E

Considering C & to find out all nodes reachable from C (including C itself) use stack.

(a) Initially push C onto stack.

Stack: C.

(b) Pop & Print C, Put F onto stack (neighbour)

Print: C Stack: F

(c) Pop & print F, Push A & B onto stack.

Print: F, Stack: A B.

(d) Pop & print B. There are no nodes ~~from~~ in neighborhood which has not been on stack.

Print: B Stack: - A, -

(e) Pop & print A. Push ~~A & E~~ neighbors ~~E & G~~ onto stack

Print: A Stack: G E

(f) Pop E & Print ~~E~~ it. No ^{new} neighbors.

Print: E

(g) Pop & print G. No new neighbours.

Print: G.

∴ Reachable from C →
C, F, B, A, E, G