

# CS5800: Algorithms — Virgil Pavlu

## Homework 8

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Collaborators:

### Instructions:

- Make sure to put your name on the first page. If you are using the  $\text{\LaTeX}$  template we provided, then you can make sure it appears by filling in the `yourname` command.
- Please review the grading policy outlined in the course information page.
- You must also write down with whom you worked on the assignment. If this changes from problem to problem, then you should write down this information separately with each problem.
- Problem numbers (like Exercise 3.1-1) are corresponding to CLRS 3<sup>rd</sup> edition. While the 2<sup>nd</sup> edition has similar problems with similar numbers, the actual exercises and their solutions are different, so make sure you are using the 3<sup>rd</sup> edition.

1. Exercise 16.3-3

Amortized Analysis of Binomial Heaps

**Solution:** Suppose  $f(n)$  represents the runtime of the EXTRACT-MIN operation for a heap containing  $n$  items. If we define the potential function  $\phi = \sum_{k=1}^n f(k)$ , the potential decreases by  $f(n)$  whenever an EXTRACT-MIN operation is executed. Furthermore, it is permissible to increase the potential by  $f(n)$  during each INSERT operation, since  $f(n) = O(\log n)$ , leading to an amortized cost of  $O(\log n) + O(\log n) = O(\log n)$  for the INSERT operation.

Let  $\phi = \sum_{k=1}^n f(k)$ , where  $n$  is the number of items in the heap after the  $i$ -th operation. Initially,  $\phi_0 = 0$  and  $\phi_i \geq 0$  for all  $i$ . For an INSERT operation, we have  $c'_i = c_i + \Delta\phi_i = \log n + f(n) = O(\log n)$ . For an EXTRACT-MIN operation, the amortized cost is  $c'_i = c_i + \Delta\phi_i = f(n) - f(n) = 0 = O(1)$ .