

1) Algorithm:

```

bool containsSimpleCycle (Matrix graph[noOfVertices]
                           [noOfVertices])
    for i = 0 to i != noOfVertices
        flag[noOfVertices] = { 0 }
        for j = 0 to j != noOfVertices
            if (j == i OR graph[i][j] == 0)
                continue
            for k = 0 to k != noOfVertices
                if (k == j OR k == i OR
                    graph[j][k] == 0)
                    continue
                if (flag[k] == 1)
                    return true
                else
                    flag[k] = 1
            return false
    
```

Time Complexity: →

- ↳ If two vertices share more than one neighbour, then the graph contains a square
- ↳ for any vertices v_1 and v_2 , we will compare the row of v_1 with row of v_2 in $O(V)$ time
- ↳ We will iterate all through V_1 and V_2 which will take $O(V^2)$ time

↳ Thus, the overall time complexity of the algorithm is $O(V^3)$ as $O(V^2)$ for the iteration and $O(V)$ for the comparison.

Explanation: →

↳ We will use 3 loops and the flag array to check if the ^{vertices} graph contains more than one common neighbour.

↳ We will iterate till index is not equal to number of vertices. We will ~~mark the index~~ update the value of element in flag array as 1 to check if the ~~row~~ has common neighbour.

↳ first, we will create adjacency matrix and update the value of element as 1 if there is an edge from vertex 1 to vertex. The size of matrix will be no. of vertices * no. of vertices.

CLASSMATE
Date _____
Page _____

↳ Flag array is used to keep track of the neighbour vertices.
~~If the~~

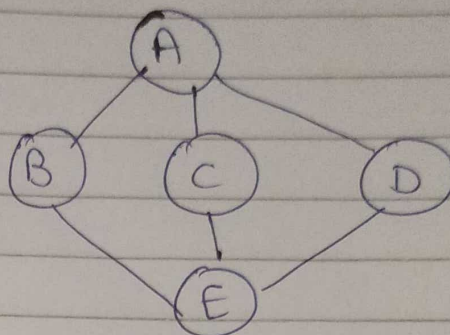
↳ We will compare the row of vertex 1 and vertex 2; if there is more than one common neighbour then we will return true else false.

↳ If ~~there~~ 1 is already present in the flag array and the value of another row is also 1, then ~~we check~~ that means both the vertices share the common neighbour and hence a square cycle is found.

↳ We will iterate first for loop till it does not become equal to no of vertices. We will reset the flag array in every iteration of first for loop. We will iterate second for loop till it does not become equal to the no of vertices. If index of first for loop & second for loop is equal, we will continue the loop. We will use third loop for comparison. The index of third for loop is used to update the value of flag array.

Example

for Graph: \rightarrow



Adjacency Matrix: \rightarrow

	A	B	C	D	E
A	0	1	1	1	0
B	1	0	0	0	1
C	1	0	0	0	1
D	1	0	0	0	1
E	0	1	1	1	0

By using Algorithm: \rightarrow

Two Vertices, A & E share more than neighbour and B, C, D also share more than neighbour. So in flag array (Refer the above algorithm), the flag[k] will be set to 1 for vertex A and for vertex E, we will get the flag of k^{th} index as 1, the result will be returned as true.

\therefore The Algorithm mentioned is correct.