

# COMP 330/543: The Relational Algebra

Sinan Kockara

Luis Guzman

Chris Jermaine

Rice University

# Relational Calculus vs. Algebra

In Relational Calculus

- You say what you want
- And not how to compute it

But obviously...

- This needs to be compiled into an actual computational plan
- And in relational DBs, the plan is expressed in relational algebra

RA is the “abstract machine” of relational databases

# What Is An Algebra?

Many Definitions!

- Simplest: it is a set (domain) with a number of operations
- The domain is closed under those operations

In RA...

- The domain is the set of all valid relations
- The set of operations includes  $\pi$ ,  $\sigma$ ,  $\times, \bowtie$ ,  $\cup$ ,  $\cap$ ,  $-$

# Projection

Projection removes attributes

$\pi_A(R)$ ...

- $A$  is a set of attributes of relation  $R$
- This simply removes all atts not in  $A$  from  $R$
- Note: cardinality of output can differ from  $R$
- Output is a relation

COURSE

CRN	NAME	DOW	STARTTIME	ENDTIME
12809	COMP 430	MWF	14:00:00	14:50:00
16954	COMP 543	MWF	10:00:00	10:50:00
10396	COMP 140	TTH	10:50:00	12:05:00
13970	COMP 436	TTH	14:30:00	15:45:00

$\pi_{NAME,DOW}(COURSE)$

Answers the query

“what are the names and week days of the courses ?”

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NAME	DOW
COMP 430	MWF
COMP 543	MWF
COMP 140	TTH
COMP 436	TTH

# Selection

Selection removes tuples

$\sigma_B(R)$ ...

- $B$  is a boolean pred that can be applied to a single tuple from  $R$
- This simply removes all tuples not accepted by  $B$
- Again: output is a relation

COURSE	CRN	NAME	DOW	STARTTIME	ENDTIME
	12809	COMP 430	MWF	14:00:00	14:50:00
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$\sigma_{NAME="COMP\ 543"}(COURSE)$

Answers the query “which courses have name COMP 543?”

CRN	NAME	DOW	STARTTIME	ENDTIME
16954	COMP 543	MWF	10:00:00	10:50:00

# Question

What is the query for:

“which courses meet for less than an hour at a time?”

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$\sigma_{ENDTIME-STARTTIME < 1:00:00}(COURSE)$

CRN	NAME	DOW	STARTTIME	ENDTIME
12809	COMP 430	MWF	14:00:00	14:50:00
16954	COMP 543	MWF	10:00:00	10:50:00



# Selection/Projection Example

LIKES (DRINKER, BEER)

FREQUENTS (DRINKER, BAR)

SERVES (BAR, BEER)

Query: Who likes the beer 'PBR'?

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- $\pi_{\text{DRINKER}}(\sigma_{\text{BEER}=\text{'PBR'}}(\text{LIKES}))$
- What is equivalent RC?

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Query: Who likes the beer 'PBR'?

➤  $\pi_{\text{DRINKER}}(\sigma_{\text{BEER}='PBR'}(\text{LIKES}))$

➤ What is equivalent RC?

$\{ l.\text{DRINKER} \mid l \in \text{LIKES} \wedge l.\text{BEER} = 'PBR' \}$

# Rename

➤  $\rho_{A/B}(R)$

Renames attribute B to A in relation R

Output is a relation

➤  $\rho_S(A_1, A_2, \dots, A_n)(R)$

Renames relation R to S and renames all attributes as specified

Output is a relation

# Assignment

➤  $X \leftarrow (\text{RA statement})$

- Assigns the output of an RA statement, which is a relation, to a temporary relation
  - For convenience

Example – assign the courses that meet on Monday-Wednesday-Friday to a temporary relation “MWF”

$$MWF \leftarrow \sigma_{\text{dow}=\text{"MWF"}}(COURSE)$$

# Join: Cartesian Product

Join combines tuples

Simplest join is Cartesian product (aka: cross product)

$R \times S \dots$

- Returns  $r \bullet s$  for all  $r \in R, s \in S$
- What is the output cardinality?

## Join: Theta Join

LIKES (DRINKER, BEER)

FREQUENTS (DRINKER, BAR)

SERVES (BAR, BEER)

Often you want  $\sigma_B(R \times S)$

Shorthand for this is  $R \bowtie_B S$

Query: Who likes a beer that 'Luis' likes?

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Query: Who likes a beer that 'Luis' likes?

- $\text{TEMP}(d_1, b_1, d_2, b_2) \leftarrow \text{LIKES} \bowtie_{\text{BEER}=\text{BEER}} (\sigma_{\text{DRINKER}=\text{'Luis'}}(\text{LIKES}))$
- $\pi_{d_1}(\text{TEMP})$



# Example Walkthrough

$$\text{LIKES} \quad \times \quad \sigma_{\text{DRINKER}='Luis'}(\text{LIKES}) \quad =$$

Drinker	Beer
Luis	PBR
Chris	Corona
Chris	<u>Budlight</u>
Luis	<u>Modelo</u>
Chris	PBR

x

Drinker	Beer
Luis	PBR
Luis	<u>Modelo</u>

=

Drinker	Beer	Drinker	Beer
Luis	PBR	Luis	PBR
Luis	PBR	Luis	<u>Modelo</u>
Chris	Corona	Luis	PBR
Chris	Corona	Luis	<u>Modelo</u>
Chris	<u>Budlight</u>	Luis	PBR
Chris	<u>Budlight</u>	Luis	<u>Modelo</u>
Luis	<u>Modelo</u>	Luis	PBR
Luis	<u>Modelo</u>	Luis	<u>Modelo</u>
Chris	PBR	Luis	PBR
Chris	PBR	Luis	<u>Modelo</u>

# Example Walkthrough

$\text{TEMP}(d_1, b_1, d_2, b_2) \leftarrow \text{LIKES} \times \sigma_{\text{DRINKER}='Luis'}(\text{LIKES})$

$\sigma_{b_1=b_2'}(\text{TEMP})$

TEMP

d1	b1	d2	b2
Luis	PBR	Luis	PBR
Luis	PBR	Luis	<u>Modelo</u>
Chris	Corona	Luis	PBR
Chris	Corona	Luis	<u>Modelo</u>
Chris	<u>Budlight</u>	Luis	PBR
Chris	<u>Budlight</u>	Luis	<u>Modelo</u>
Luis	<u>Modelo</u>	Luis	PBR
Luis	<u>Modelo</u>	Luis	<u>Modelo</u>
Chris	PBR	Luis	PBR
Chris	PBR	Luis	<u>Modelo</u>

d1	b1	d2	b2
Luis	PBR	Luis	PBR
Luis	<u>Modelo</u>	Luis	<u>Modelo</u>
Chris	PBR	Luis	PBR

$\pi_{d1}(\sigma_{b_1=b_2'}(\text{TEMP}))$

d1
Luis
Chris

$\sigma_{D1 \neq 'Luis'}(\pi_{d1}(\text{TEMP}))$

d1
Chris

# Join: Natural Join

LIKES (DRINKER, BEER)

FREQUENTS (DRINKER, BAR)

SERVES (BAR, BEER)

Often you want to join two relations

- Using an equality check on all attrs having the same name
- Then project away redundant attributes

Shorthand for this  $R * S$

Query: Who goes to a bar serving a beer that they like?

# Natural Join Example

## STUDENT

StudentID	CourseID	Instructor
S1	C101	Dr. Smith
S2	C102	Dr. Johnson
S3	C103	Dr. Brown

## COURSE

CourseID	Instructor	CourseName
C101	Dr. Smith	Math
C102	Dr. Johnson	Physics

## STUDENT \* COURSE

StudentID	CourseID	Instructor	CourseName
S1	C101	Dr. Smith	Math
S2	C102	Dr. Johnson	Physics

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- $LIKES * SERVES * FREQUENTS$

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Shorthand for this  $R * S$

Query: Who goes to a bar serving a beer that they like?

- $LIKES * SERVES * FREQUENTS$
- $\pi_{\text{DRINKER}}(LIKES * SERVES * FREQUENTS)$

# Natural join example

FREQUENTS (DRINKER, CAFE)

SERVES (CAFE, COFFEE)

LIKES (DRINKER, COFFEE)

Query – who goes to a cafe serving cold brew ?

RC

RA

# Natural join example

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Query – who goes to a cafe serving cold brew ?

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$$\{ f.DRINKER / f \in FREQ \wedge \\ \exists(s) (s \in SERVES \wedge \\ s.CAFE = f.CAFE \wedge \\ s.COFFEE = 'Cold Brew') \}$$

RA



# Natural join example

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Query – who goes to a cafe serving cold brew ?

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RA

$$\pi_{DRINKER}(\sigma_{COFFEE="Cold Brew"}(FREQUENTS * SERVES))$$

# Set-Based Operations

LIKES (DRINKER, BEER)

FREQUENTS (DRINKER, BAR)

SERVES (BAR, BEER)

Can use standard set ops as well:  $\cup$ ,  $\cap$ ,  $-$

- To use, types and numbers of input attributes must match
- By convention, attribute names come from LHS
- $R \cup S$ : all tuples in  $R$  or in  $S$
- $R \cap S$ : all tuples in  $R$  and in  $S$
- $R - S$ : all tuples in  $R$  and not in  $S$

Query: Who does not like the beer 'PBR'?

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Query: Who does not like the beer 'PBR'?

- $\text{PBRGOOD} \leftarrow \pi_{\text{DRINKER}}(\sigma_{\text{BEER}='PBR'}(\text{LIKES}))$
- $(\pi_{\text{DRINKER}}(\text{LIKES})) - \text{PBRGOOD}$

# Complicated Set-Based Example

LIKES (DRINKER, BEER)

FREQUENTS (DRINKER, BAR)

SERVES (BAR, BEER)

Who only goes to bars where they can get a beer they like?

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LIKES (DRINKER, BEER)

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Who only goes to bars where they can get a beer they like?

- Use ‘all people’ – ‘those who go to a bar where they can’t get a beer they like’
- $ALLPEPES \leftarrow \pi_{DRINKER}(LIKES)$
- How about ‘those who go to a bar where they can’t get a beer they like’?
- Use FREQUENTS – ‘DRINKER, BAR combos where the person can get a beer they like’
- $GOODBEER \leftarrow \pi_{DRINKER, BAR}(LIKES * SERVES)$

Then the answer is

- $ALLPEPES - \pi_{DRINKER}(FREQUENTS - GOODBEER)$

# Questions?