Effective Caching of Shortest Paths for Location-Based Services

Jeppe Rishede Thomsen¹, Man Lung Yiu¹, Christian S. Jensen²

¹Department of Computing Hong Kong Polytechnic University

²Department of Computer Science Aarhus University

May 17, 2012

Motivation

- Many Shortest Path Queries
 - 40 % of Google Maps usage is mobile
 - 200+ million active mobile users as of may 2011
 - Users drive 12 billion miles a year with Google Maps Navigation
 - 20 % of Google searches are for local information

http://techcrunch.com/2011/03/11/marissa-mayer-40 of-google-maps-usage-is-mobile-and-there-are-150-million-mobile-users/

http://techcrunch.com/2011/05/25/google-maps-formobile-stats/

http://blog.hubspot.com/blog/tabid/6307/bid/10829/5 Google-Local-Stats-Every-Marketer-Should-Know-Data aspx



Why is caching necessary

• Shortest Path calculation is expensive

Target	Travel	Response
	time (s)	time (ms)
The Smithsonian	372	101.24
Washington, DC	419	110.94
War Memorials	41831	168.44
Capitol Building	75805	278.44
Statue of Liberty	88947	362.8
Mount Rushmore	99456	364.68
Golden Gate Bridge	108353	342.8
	The Smithsonian Washington, DC War Memorials Capitol Building Statue of Liberty Mount Rushmore	The Smithsonian 372 Washington, DC 419 War Memorials 41831 Capitol Building 75805 Statue of Liberty 88947 Mount Rushmore 99456

Google Directions API

Setting

Web search scenario:

[Maarkatos et al., Computer Communications 2001]

- Can have a cache at either a proxy or server site
- Saves response- or computation- time
- Cache stores web search results
- Existing cache algorithms include:
 - -Least Recently Used
 - —Highest Query Frequency

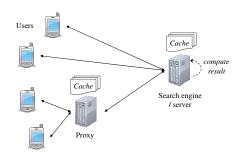
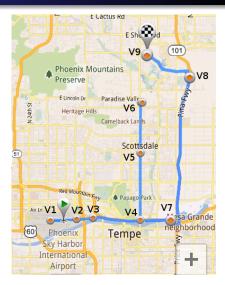


Figure: Web Search

How does a Shortest Path Cache Work?

	Cache Content				
Path	Shortest Path				
1	V1,V2,V3,V4,V7,V8,V9				
2	V1,V2,V3,V4,V5,V6				
3	V5,V4,V7				

Queries							
Query	Result	Path					
$Q_{V1,V9}$	HIT	1					
$Q_{V2,V5}$	HIT	2					
$Q_{V5,V9}$							



How do we optimize the Cache Performance?

Benefit is expected cost saved:

- On server: computation time
- On proxy: communication time

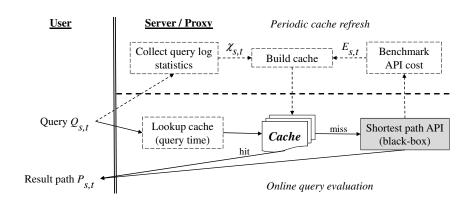
We need to answer:

- Which queries $Q_{s,t}$ can be answered by the path $P_{a,b}$?
- For query $Q_{s,t}$, what are the cost savings?

Given a:

- Cache Size Budget
- Query Log

Then build a cache Ψ with max benefit $\gamma(\Psi)$



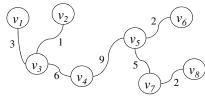
- Systematic model for benefit
- Techniques for query log statistics extraction

- Efficient caching structure
- Shortest Path benchmarking techniques

Statistics Extraction

Timestamp	Query
T_1	$Q_{3,6}$
T_2	$Q_{1,6}$
T_3	$Q_{2,7}$
T_4	$Q_{1,4}$
T_5	$Q_{4,8}$
T_6	$Q_{2,5}$
T_7	$Q_{3,6}$
T_8	$Q_{3,6}$

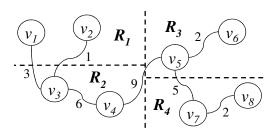
$\chi_{s,t}$	v_1	v_2	v_3	v_4	v_5	v_6	v_7	v_8
v_1	/	0	0	1	0	1	0	0
v_2	0	/	0	0	1	0	1	0
v_3	0	0	/	0	0	3	0	0
v_4	1	0	0	/	0	0	0	1
v_5	0	1	0	0	/	0	0	0
v_6	1	0	3	0	0	/	0	0
v_7	0	1	0	0	0	0	/	0
v_8	0	0	0	1	0	0	0	/



Grouping

$\chi_{s,t}$	v_1	v_2	v_3	v_4	v_5	v_6	v_7	v_8
v_1	/	0	0	1	0	1	0	0
v_2	0	/	0	0	1	0	1	0
v_3	0	0	/	0	0	3	0	0
v_4	1	0	0	/	0	0	0	1
v_5	0	1	0	0	/	0	0	0
v_6	1	0	3	0	0	/	0	0
v_7	0	1	0	0	0	0	/	0
v_8	0	0	0	1	0	0	0	/

χ_{R_i,R_j}	R_1	R_2	R_3	R_4
R_1	0	1	2	1
R_2	1	0	3	1
R_3	2	3	0	0
R_4	1	1	0	0



SP Call Cost Estimation

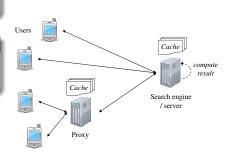
Proxy Scenario

$$E_{s,t}(Proxy) = 1$$

Server Scenario

Intuition: Longer query results incur higher cost.

- Cost only estimated in server scenario
- Estimation methods developed for Server scenario



Incremental Benefit per size

Benefit formula

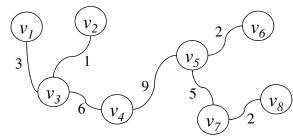
$$\Delta \overline{\gamma}(P_{a,b}, \Psi) = \sum\limits_{P_{s,t} \in \mathfrak{U}(P_{a,b}) - \mathfrak{U}(\Psi)} \frac{\chi_{s,t} * E_{s,t}}{|P_{a,b}|}$$

- $P_{a,b}$: a shortest path
- \bullet Ψ : the cache
- $\mathfrak{U}(P_{a,b})$: all sub-paths of $P_{a,b}$.
- $\chi_{s,t}$ frequency of query s to t.
- ullet $E_{s,t}$: cost of calculating $P_{s,t}$

Cache: SP Result Ranking

Greedy algorithm

diccay algo	
Timestamp	Query
T_1	$Q_{3,6}$
T_2	$Q_{1,6}$
T_3	$Q_{2,7}$
T_4	$Q_{1,4}$
T_5	$Q_{4,8}$
T_6	$Q_{2,5}$
T_7	$Q_{3,6}$
T_8	$Q_{3,6}$

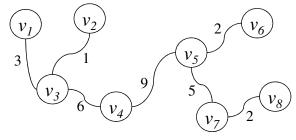


Round			Cach	те Ψ				
	$P_{1,4}$	$P_{1,4} \mid P_{1,6} \mid P_{2,5} \mid P_{2,7} \mid P_{3,6} \mid P_{4,8}$						After
1	1/3	$ \overline{\mathbf{5/5}} $	1/4	2/5	3/4	1/4	empty	$P_{1,6}$
2							$P_{1,6}$	$P_{1,6},?$

Cache: SP Result Ranking - Incremental benefit

Incremental benefit calculation

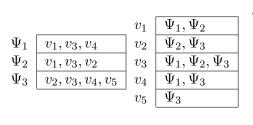
	 	~ ~
Tin	nestamp	Query
	T_1	$Q_{3,6}$
	T_2	$Q_{1,6}$
	T_3	$Q_{2,7}$
	T_4	$Q_{1,4}$
	T_5	$Q_{4,8}$
	T_6	$Q_{2,5}$
	T_7	$Q_{3,6}$
	T_8	$Q_{3,6}$
_	- 1	

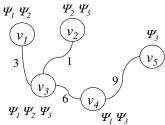


Round			Ca	che Ψ				
	$P_{1,4}$	$P_{1,6}$	$P_{2,5}$	$P_{2,7}$	$P_{3,6}$	$P_{4,8}$	Before	After
1	1/3	$ \overline{\mathbf{5/5}} $	1/4	2/5	3/4	1/4	empty	$P_{1,6}$
2	0	0	1/4	$ \overline{2/5} $	0	1/4	$P_{1,6}$	$P_{1,6}, P_{2,7}$

Efficient Data Structure for the Cache

- Lookup time grows with size cache
- Support return of full or partial cache items
- Compact storage of shortest paths





Paths

Inverted List

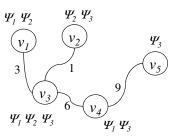
Visualization

Efficient Data Structure for the Cache

- Support efficient lookup and return of full or partial cache items
- Compact storage of shortest paths

v_1	v_3
v_2	v_3
v_3	v_1, v_2, v_4
v_4	v_3, v_5
v_5	v_4

	content	parent
v_1	Ψ_1,Ψ_2	NIL
v_2		
v_3	Ψ_3	v_1
v_4		
v_5		



Graph representation

Prefix compressed

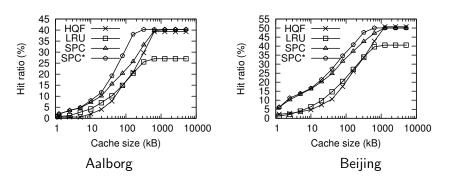
Visualization

Experimental Setup

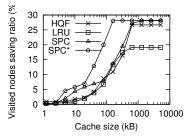
- Query logs divided equally into 2 sets:
 - historical
 - query workload
- Comparison with:
 - Least Recently Used (LRU)
 - Higest Query Frequency (HQF)

Dataset	Trajectories	Road network
Aalborg	Infati GPS data	From downloads.cloudmade.com
	4,401 trajectories	129k nodes, 137k edges
Beijing	Geo-Life GPS data	From downloads.cloudmade.com
	12,928 trajectories	76k nodes, 85k edges

Hit ratio

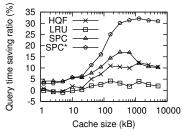


Performance Savings — Dijkstra



Visited nodes savings, Aalborg

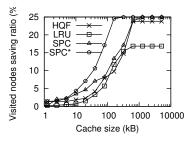
Visited nodes savings, Beijing



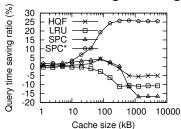
Query time savings, Aalborg

Query time savings, Beijing

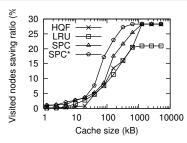
Performance Savings — A*



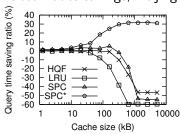
Visited nodes savings, Aalborg



Query time savings, Aalborg



Visited nodes savings, Beijing



Query time savings, Beijing

Conclusion

- Introduced benefit model, capturing the benefit of adding a shortest path result, relative to other results.
- Designed statistics extraction techniques
- Developed techniques to estimate cost of calculating a shortest path
- Designed efficient data structure for the cache storage
- Experimental results show:
 - High hit ratio
 - Small lookup overhead
 - Low query time

End of Presentation

Thank You For Listening