

# Graph Theory & Surface Reconstruction

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## Abstract

The purpose of this paper is to utilize a technique for recreating models of 3-D objects. This process is referred to as surface reconstruction, which replaces a set of sample points using a faceted geometric model. Concepts of graph theory are used, specifically duals of graphs,  $n$ -regular graphs, bridgeless graphs, and matchings.

## Introduction

The paper will refer to the triangulation of a 3-D surface given in the figure below.

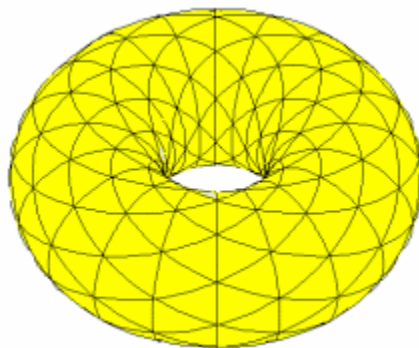


Figure 1: A torus.

The figure resembles a torus, which is created by revolving a circle along a line made by another circle. Divisions of polygonal faces that form the torus produces triangular faces throughout the torus.

## Dual Graph

By definition, the planar dual of any planar graph,  $G$ , can be constructed by placing a vertex in each region of  $G$  and then adding edges between regions that share a border. The planar dual is denoted as  $D(G)$ .

If a dual of the torus is to be constructed, the same procedure to create the dual of a planar graph will be used. A point will be placed in each of the triangular faces, then lines

between pairs of vertices would be placed, if and only if, the corresponding faces share a common side. The dual of the torus,  $G$ , is constructed below.

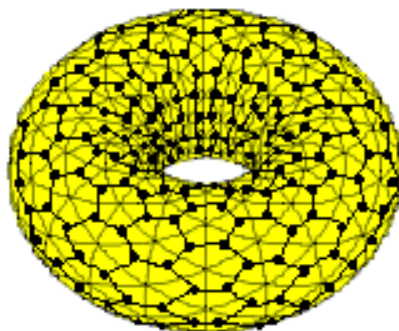


Figure 2: Dual of a torus,  $G$ .

As shown in the figure, the triangular faces of the original torus are hexagonal faces in the dual of the torus.

## n-regular Graph

A graph is considered  $n$ -regular if every vertex in the graph is incident to exactly  $n$  edges.

The dual of the torus is a 3-regular graph. According to the definition, that would mean that every vertex in the dual of the torus,  $G$ , is incident to exactly 3 edges. The graph of the dual is a 3-regular because each face in the original graph is a complete graph,  $K_3$ , in which each vertex is connected to all the other vertices. A  $K_3$  graph consists of 3 vertices that all have degree 2. The dual of a  $K_3$  is just one vertex inside the cycle and one vertex on the outside of the cycle, which represents the infinite region. Each edge in  $K_3$  has an edge in the dual graph crossing the edge to connect those two vertices, as shown below.

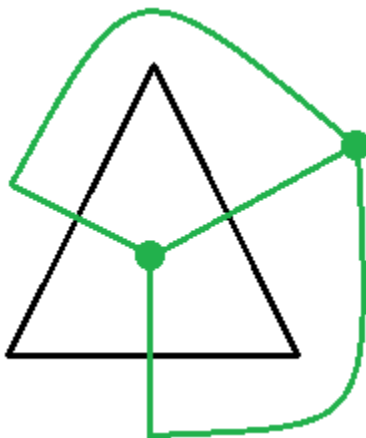


Figure 2: Dual of a complete graph,  $K_3$ .

By taking the dual of groups of  $K_3$ , a similar concept is used to connect the vertices for each region, excluding the infinite region. As a result of the three edges in each  $K_3$ , each vertex representing one region has three incident edges in the planar dual graph of each  $K_3$ . Thus, the dual of the torus is a 3-regular graph.

## Bridgeless Graph