Stack Applications

- Parentheses Matching
- Maze Router
- Infix to Postfix

Parentheses Matching

- (((a+b)*c+d-e)/(f+g)-(h+j)*(k-l))/(m-n)
 - 출력 값 (i, j) 는 i 번째 여는 괄호는 j 번째 닫는 괄호와 쌍이라는 것을 의미하고 있다, 그러므로 위의 수식에서는 다음과 같은 괄호 매칭 쌍이 출력된다
 - **(2,6)** (1,13) (15,19) (21,25) (27,31) (0,32) (34,38)
- (a+b))*((c+d)
 - **(0,4)**
 - right parenthesis at 5 has no matching left parenthesis
 - **(8,12)**
 - left parenthesis at 7 has no matching right parenthesis

Parentheses Matching

- 알고리즘
 - 수식을 왼쪽부터 오른쪽으로 scan 하면서 ...
 - □ 여는 괄호를 만나면 stack에 PUSH한다
 - □ 닫는 괄호를 만나면 stack으로부터 POP을 한다
 - 이때 stack이 empty이거나
 - Scan 완료 후에도 stack에 괄호 위치가 남아 있으면 괄호의 mismatching

Example

(((a+b)*c+d-e)/(f+g)-(h+j)*(k-l))/(m-n)

2 1 0

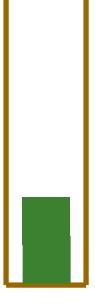
15 0 (2,6) (1,13)

21

(2,6) (1,13) (15,19)

27

(2,6) (1,13) (15,19) (21,25)



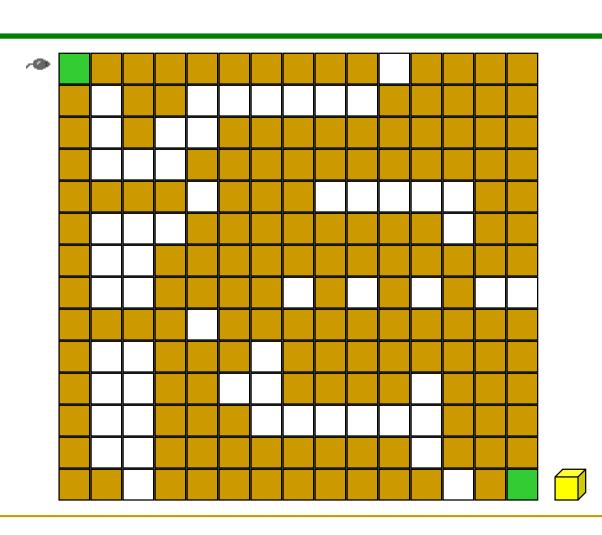
(2,6) (1,13) (15,19) (21,25)(27,31) (0,32)

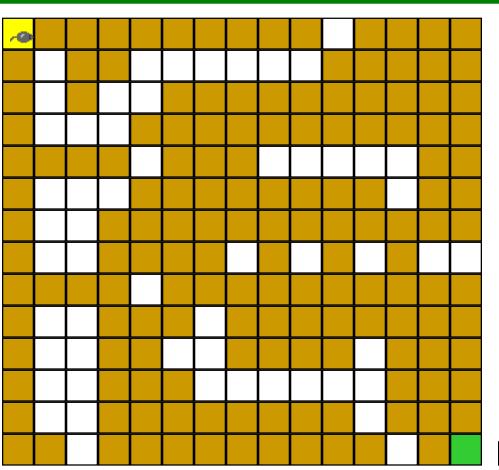
and so on

More Stack Applications

- Parentheses Matching
- Maze Router
- Infix to Postfix

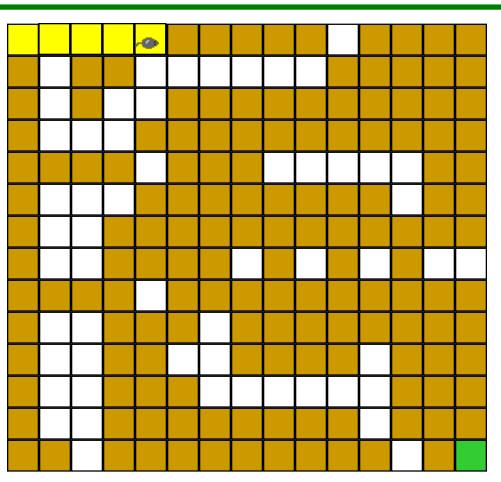
Rat In A Maze





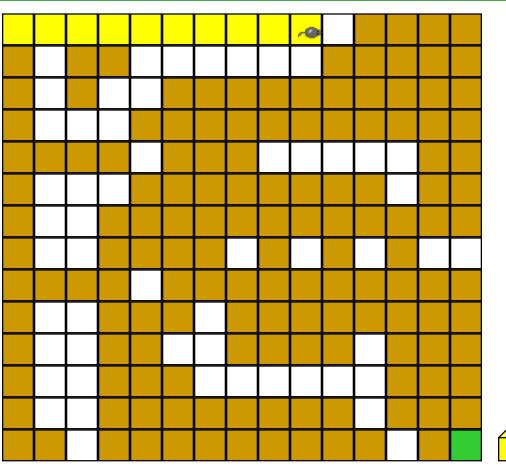


- Move order is: right, down, left, up
- Block positions to avoid revisit.



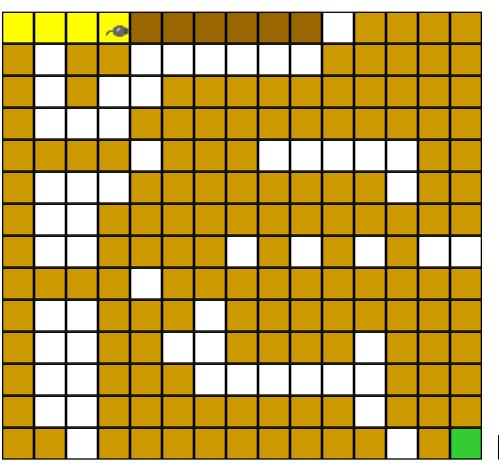


- Move order is: right, down, left, up
- Block positions to avoid revisit.



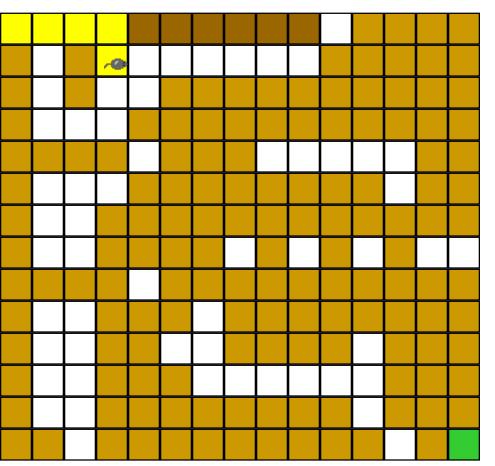


 Move backward until we reach a square from which a forward move is possible.



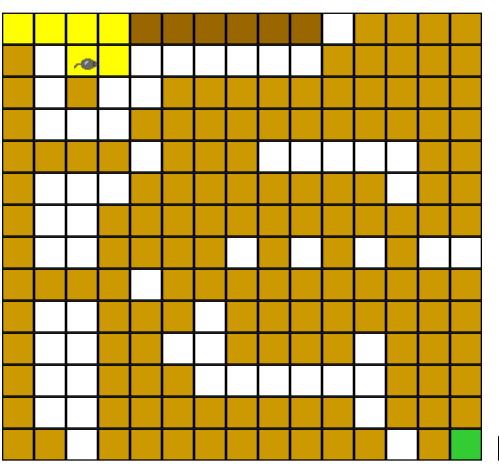


Move down.



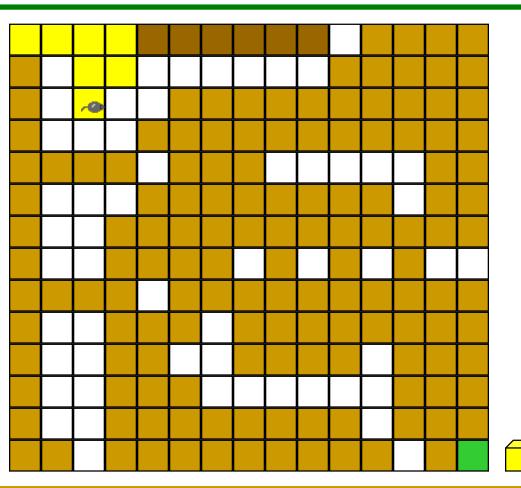


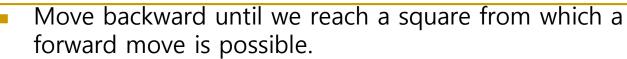
Move left.

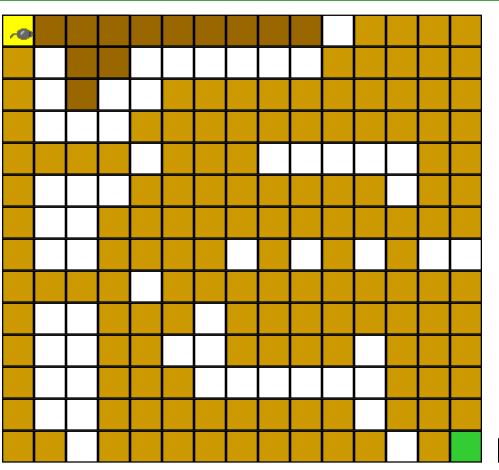




Move down.

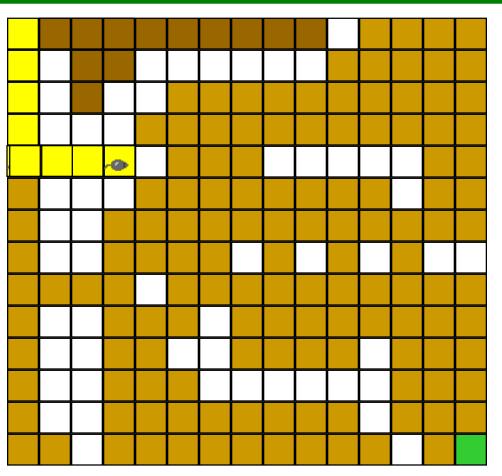






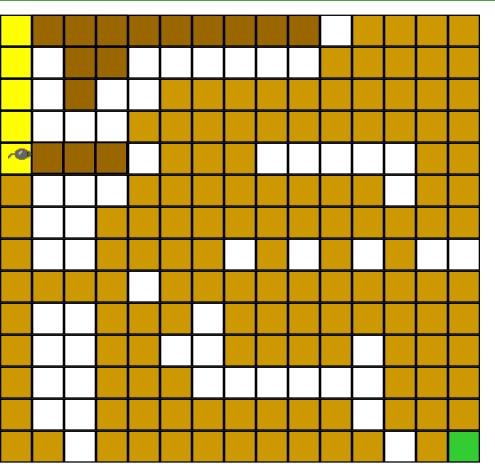


- Move backward until we reach a square from which a forward move is possible.
- Move downward



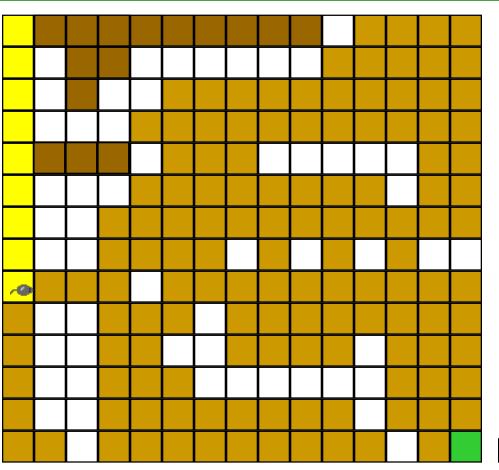


- Move right.
- Backtrack



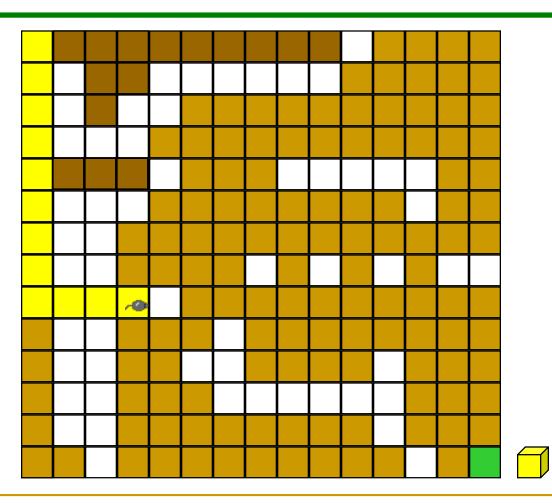


Move downward.

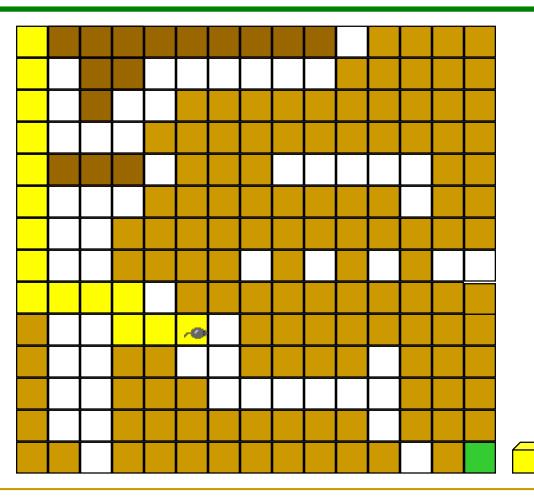




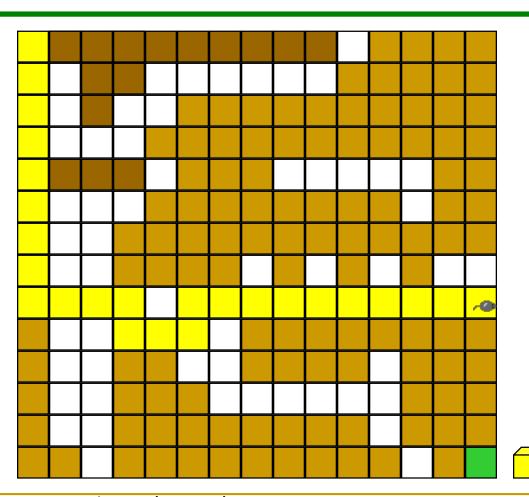
Move right.







Move one up and then right.



- Move down to exit and eat cheese.
- Path from maze entry to current position operates as a stack.

```
#include <stdio.h>
#include <stdlib.h>
#define MAX_SIZE 100
#define MAZE_SIZE 5
typedef struct Pos{
short x;
short y;
}Pos;
typedef struct Stack
Pos data[MAX_SIZE];
int top;
}Stack;
```

```
char maze[MAZE_SIZE][MAZE_SIZE]={ \ \'1','1','1','1','1'}, \ \ \ \'e','0','1','0','1'}, \ \ \ \ \'1','0','0','1'}, \ \ \ \ \ \'1','1','1','1','1'}, \ \ \ \ \ \'1','1','1','1','1'}};
```

/* 미로 찾기의 핵심은 방문한 곳을 표기하고 다음 방문할 곳을 탐색 한 후 스택에 가능한 곳 전부를 Push하고, 다시 Pop하면서 현재 경로로 변경하는 것을 반복하는 것이다. 이동이 가능한 곳은 길 또는 방문하지 않은 곳이다 */

```
void Init(Stack *p) {
           p->top=-1;
int Is_full(Stack *p) {
           return (p->top == MAX_SIZE-1);
int Is_empty(Stack *p) {
           return (p->top == -1);
void push(Stack *p,Pos data) {
  if(Is_full(p)) {
           printf("Stack Full !!\n"); return ;
  else {
           p->top++;
           p->data[p->top].x=data.x;
           p->data[p->top].y=data.y;
```

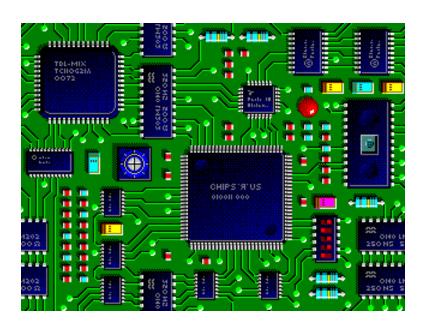
```
Pos pop(Stack *p) {
  if(ls_empty(p)) {
          printf("스택이 비어있습니다\n"); exit(1); }
  return p->data[(p->top)--];
void Push_Loc(Stack *s,int x,int y) {
  if(x < 0 || y < 0 || x > MAZE_SIZE || y > MAZE_SIZE) return;
  if(maze[x][y] != '1' && maze[x][y] != '.') {
          Pos tmp;
          tmp.x=x;
          tmp.y=y;
         Push(s,tmp);
```

```
int main() {
  Stack s;
  Pos here;
  int i,j,x,y;
  Init(&s);
// 시작점 탐색
  for(i=0;i<MAZE_SIZE;i++) {</pre>
          for(j=0;j<MAZE\_SIZE;j++) {
            if(maze[i][j]=='e') {
                    here.x=i;
                    here.y=j;
printf("시작 점 (%d,%d) ₩n",here.x,here.y);
```

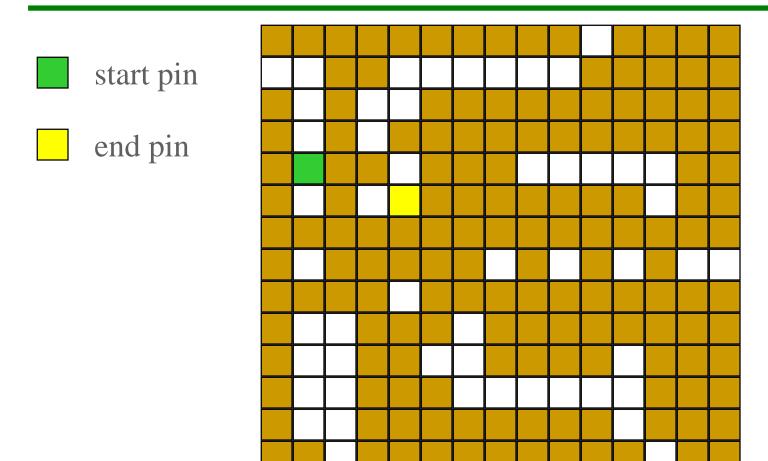
```
while(maze[here.x][here.y] != 'x') {
          x=here.x;
          y=here.y;
          maze[x][y]='.'; // 방문한 곳을 표시
          // 좌,우,위,아래중 이동 가능한 곳을 탐색
          Push_Loc(\&s,x+1,y);
          Push_Loc(&s,x-1,y);
          Push_Loc(&s,x,y+1);
          Push_Loc(&s,x,y-1);
          if(ls_empty(&s)) {
                     printf("실패₩n");
                     return 0;
          else {
                     here=Pop(&s); // 현재 좌표를 변경
                     printf("(%d,%d)₩n",here.x,here.y);
printf("도착 점 (%d,%d)₩n", here.x, here.y);
```

printf("탐색 성공₩n");

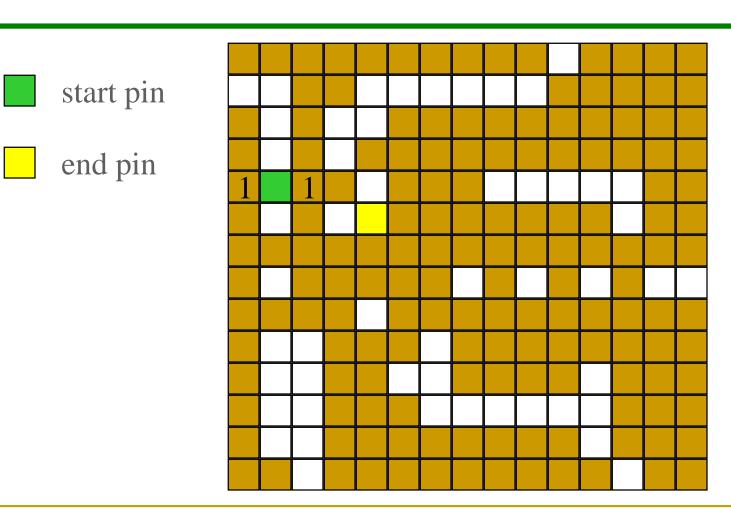
Wire Routing



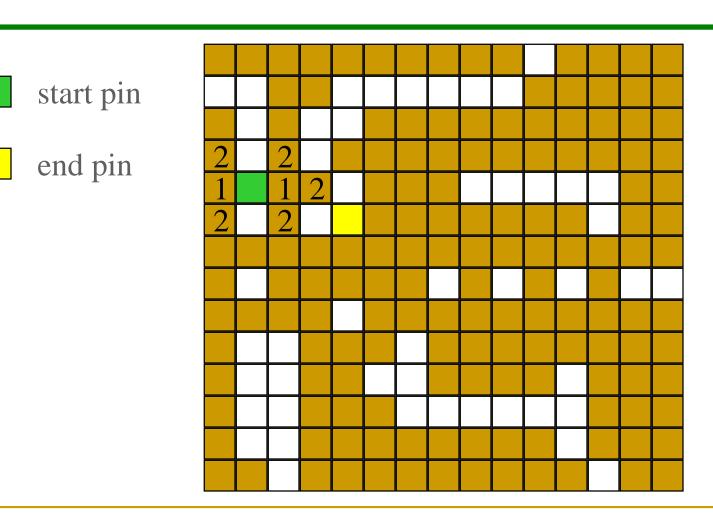
Lee's Wire Router



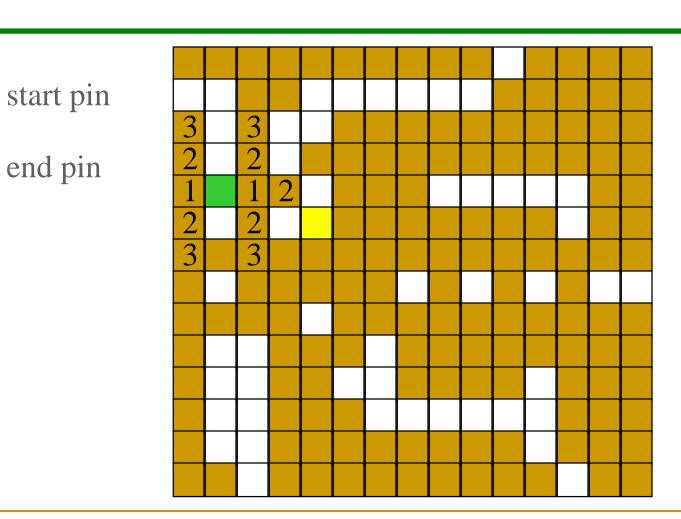
Label all reachable squares by 1 unit from start.



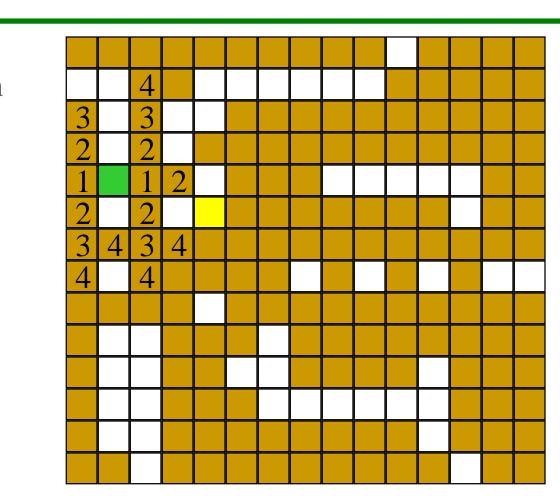
Label all reachable unlabeled squares by 2 units from start.



Label all reachable unlabeled squares by 3 units from start.



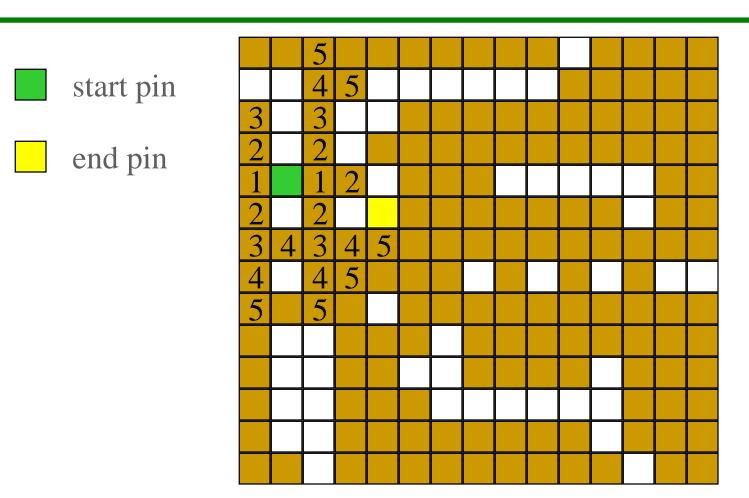
Label all reachable unlabeled squares by 4 units from start.



start pin

end pin

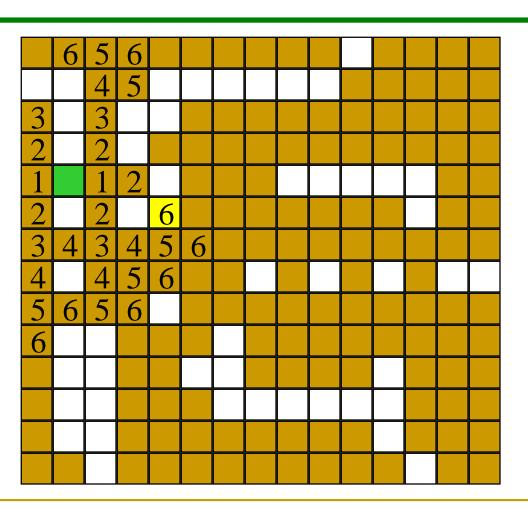
Label all reachable unlabeled squares by 5 units from start.



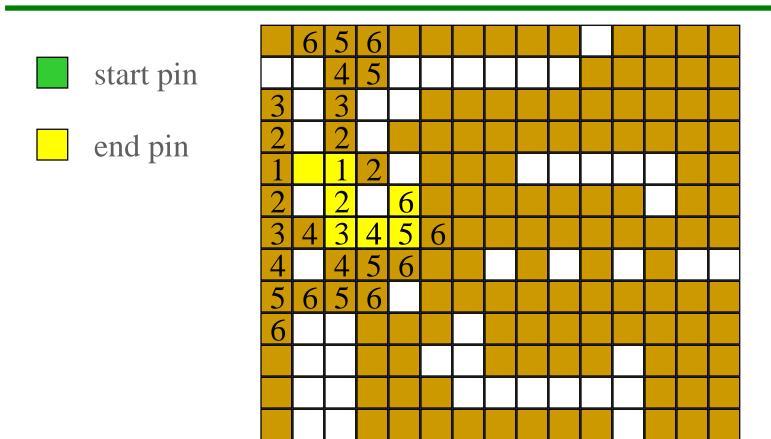
Label all reachable unlabeled squares by 6 units from start.



end pin



End pin reached. Traceback.



End pin reached. Traceback.

More Stack Applications

- Parentheses Matching
- Maze Router
- Infix to Postfix

Evaluation of Expressions

$$X = a / b - c + d * e - a * c$$

$$a = 4$$
, $b = c = 2$, $d = e = 3$

Interpretation 1:

$$((4/2)-2)+(3*3)-(4*2)=0+8+9=1$$

Interpretation 2:

$$(4/(2-2+3))*(3-4)*2=(4/3)*(-1)*2=-2.66666...$$

How to generate the machine instructions corresponding to a given expression?

precedence rule + associative rule

Mathematical Expression

- Infix notation
- reverse Polish notation (RPN) : postfix notation
 - **345*+**
- Polish notation : prefix notation
 - + 3 * 45
- Infix 연산을 RPN으로 변환하면 stack을 이용하여 연산을 매우 효율
 적으로 수행할 수 있다
 - Shunting-yard algorithm
 - Stack machine

user

compiler

Infix	Postfix
2+3*4	234*+
a*b+5	ab*5+
(1+2)*7	12+7*
a*b/c	ab*c/
(a/(b-c+d))*(e-a)*c	abc-d+/ea-*c*
a/b-c+d*e-a*c	ab/c-de*ac*-

Postfix: no parentheses, no precedence

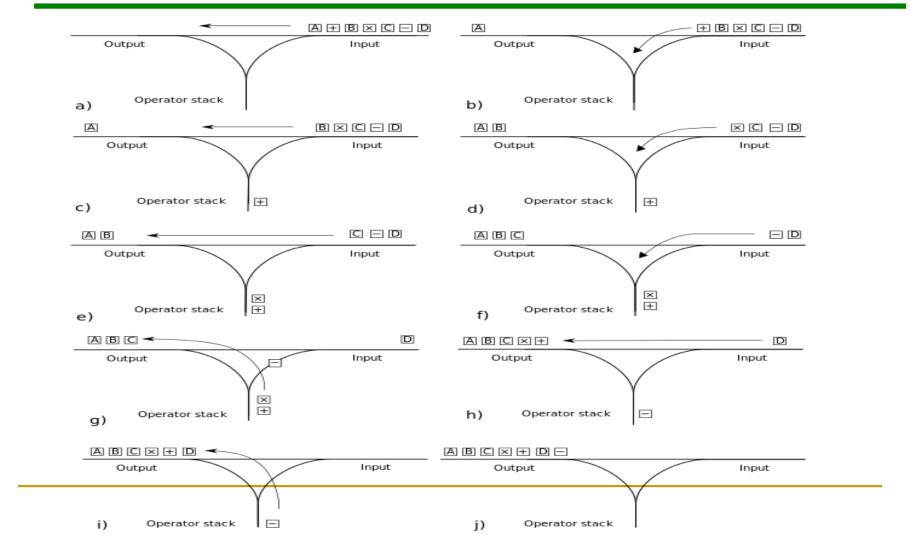
Token	Stack	Тор
	[0] [1] [2]	
6	6	0
2	6 2	1
/	6/2	0
3	6/2 3	1
_	6/2-3	0
4	6/2-3 4	1
2	6/2-3 4 2	2
*	6/2-3 4*2	1
+	6/2-3+4*2	0

г

_

Shunting-Yard Algorithm

Developed by E. Dijkstra





Infix to Postfix

```
Assumptions:
    operators: +, -, *, /, %
    operands: single digit integer
#define MAX_STACK_SIZE 100 /* maximum stack size */
#define MAX_EXPR_SIZE 100 /* max size of expression */
typedef enum
{1paran, rparen, plus, minus, times, divide, mod, eos, operand} precedence;
int stack[MAX_STACK_SIZE]; /* global stack */
char expr[MAX_EXPR_SIZE]; /* input string */
```

Evaluation of Postfix Expressions

```
int eval(void)
/* evaluate a postfix expression, expr, maintained as a global variable, 0 is the
the end of the expression. The stack and top of the stack are global variables.
get token is used to return the token type and the character symbol.
Operands are assumed to be single character digits */
 precedence token;
 char symbol; int op1, op2;
 int n = 0; /* counter for the expression string */
 int top = -1;
 token = get token(&symbol, &n);
 while (token != eos) {
   if (token == operand)
      push(&top, symbol-'0'); /* stack insert */
```

```
else { /* remove two operands, perform operation, and return result to the stack */
   op2 = pop(&top); /* stack delete */
   op1 = pop(\&top);
   switch(token) {
     case plus: push(&top, op1+op2); break;
      case minus: push(&top, op1-op2); break;
      case times: push(&top, op1*op2); break;
      case divide: push(&top, op1/op2); break;
     case mod: push(&top, op1%op2);
 token = get token (&symbol, &n);
return pop(&top); /* return result */
```

```
precedence get_token(char *symbol, int *n)
{
/* get the next token, symbol is the character representation, which is returned, the token is represented by its enumerated value, which is returned in the function name */
```

```
*symbol =expr[(*n)++];
switch (*symbol) {
 case '(': return lparen;
 case ')': return rparen;
 case '+': return plus;
 case '-': return minus;
 case '/': return divide;
 case '*': return times;
 case '%': return mod;
 case '\0': return eos;
 default: return operand;
          /* no error checking, default is operand */
```

Infix to Postfix Conversion

(Intuitive Algorithm)

(1) Fully parenthesized expression

(2) All operators replace their corresponding right parentheses.

(3) Delete all parentheses.

The orders of operands in infix and postfix are the same. a + b * c, * > +

Token	Stack			Тор	Output
	[0]	[1]	[2]		
а				-1	а
+	+			0	а
b	+			0	ab
*	+	*		1	ab
С	+	*		1	abc
eos				-1	abc*+

$$a *_{1} (b + c) *_{2} d$$

Token	Stack			Тор	Output	
	[0] [1]	[2]			
a				-1	а	
*1	*1			0	а	
(*1	(1	а	
b	*1	(1	ab	
+	*1	(+	2	ab	
С	*1	(+	2	abc	
)	*1			0	abc+	
*2	*2			0	abc+* ₁	
d	*2			0	abc+* ₁ d	
eos	*2			0	abc+* ₁ d* ₂	

Rules

- (1) Operators are taken out of the stack as long as their in-stack precedence is higher than or equal to the incoming precedence of the new operator.
- (2) (has low in-stack precedence, and high incoming precedence.

```
( ) + - * / % eos
isp 0 19 12 12 13 13 13 0
icp 20 19 12 12 13 13 13 0
```

```
precedence stack[MAX_STACK_SIZE];
/* isp and icp arrays -- index is value of precedence
lparen, rparen, plus, minus, times, divide, mod, eos */
static int isp[] = {0, 19, 12, 12, 13, 13, 13, 0};
static int icp[] = {20, 19, 12, 12, 13, 13, 13, 0};
```

isp: in-stack precedence

icp: incoming precedence

Infix to Postfix

```
void postfix(void)
/* output the postfix of the expression. The expression string, the stack,
and top are global */
 char symbol;
 precedence token;
 int n = 0;
 int top = 0; /* place eos on stack */
 stack[0] = eos;
 for (token = get _token(&symbol, &n); token != eos;
             token = get_token(&symbol, &n)) {
  if (token == operand)
    printf ("%c", symbol);
  else if (token == rparen ) {
```

Infix to Postfix (cont' d)

```
/*unstack tokens until left parenthesis */
  while (stack[top] != lparen)
    print_token(delete(&top));
  pop(&top); /*discard the left parenthesis */
 else {
  /* remove and print symbols whose isp is greater than or equal to the
     current token's icp */
  while(isp[stack[top]] >= icp[token])
    print_token(delete(&top));
  push(&top, token);
while ((token = pop(&top)) != eos)
  print token(token);
print("\n");
```