

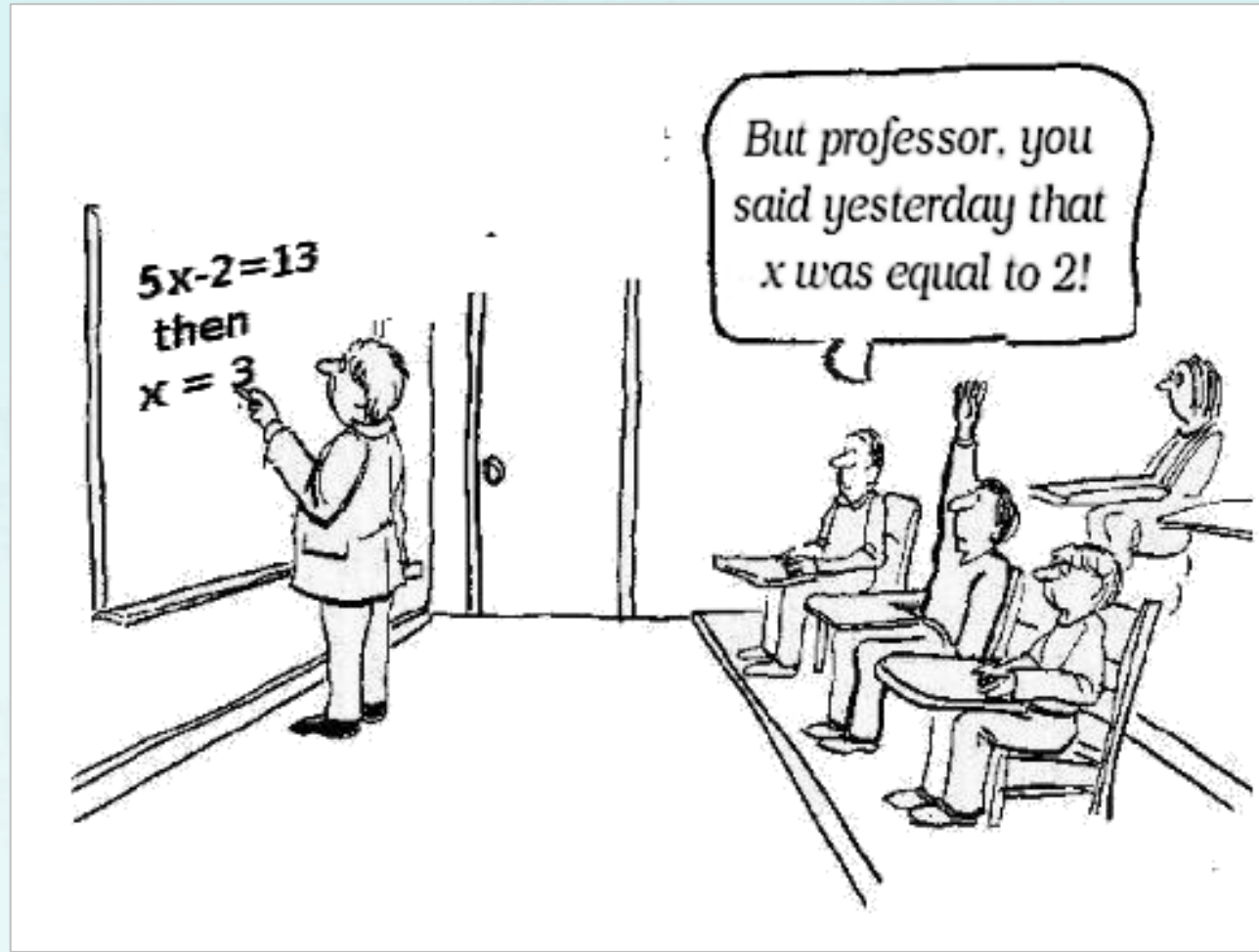
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Stack and Queue

Data Structures
C++ for C Coders

한동대학교 김영섭 교수
idebtor@gmail.com

stacks & queues using dynamic arrays

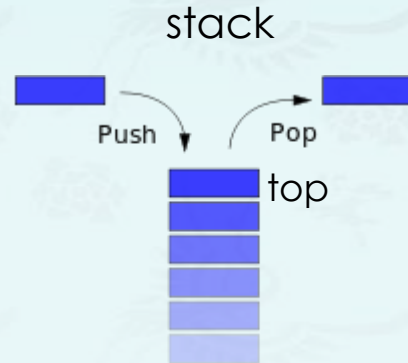


Source <http://eng.funiacs.com/funny-pictures/31079>

Stacks

Fundamental data types:

- **Value:** collections of objects
- **Operations:** insert, remove, iterate, test if empty, test if full
- Intent is clear when we insert.
- Which item do we remove?



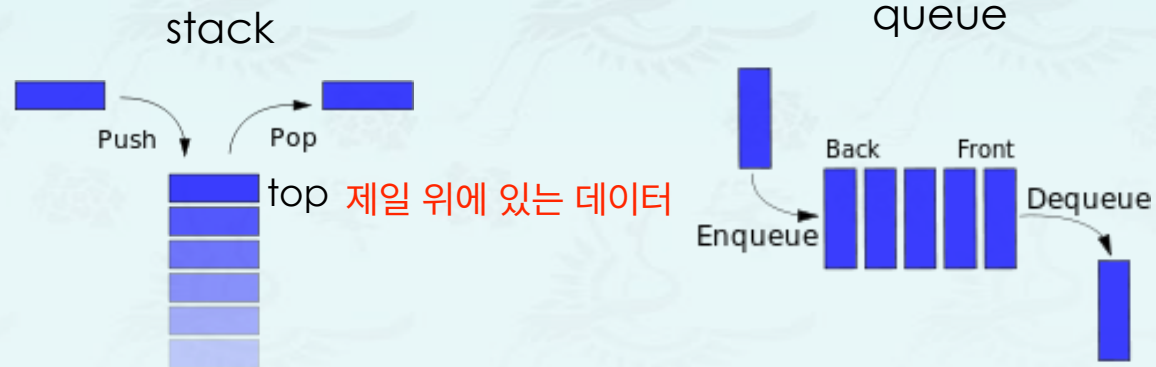
Stack: Examine the item most recently added. 마지막에 넣은 데이터가 가장 먼저 나온다. LIFO = "last in first out"



Stacks

Fundamental data types:

- **Value:** collections of objects
- **Operations:** insert, remove, iterate, test if empty, test if full
- Intent is clear when we insert.
- Which item do we remove?



Stack: Examine the item most recently added. LIFO = "last in first out"

Queue: Examine the item least recently added. FIFO = "first in first out"

Stacks

Abstract Data Type

ADT Stack is

objects: a finite ordered list with zero or more elements

functions:

Stack CreateStack(maxStackSize)

boolean IsFull(stack)

boolean IsEmpty(stack)

void Push(stack, item)

Element Pop(stack)

Why ADTs?

Stacks

Why ADTs again?

Separate interface and implementation.

Ex: stack, queue, bag, priority queue, symbol table, union-find,

Benefits.

- **Client** can't know details of implementation ⇒ client has many implementation from which to choose.
- **Implementation** can't know details of client needs ⇒ many clients can re-use the same implementation.
- **Design**: creates modular, reusable libraries.
- **Performance**: use optimized implementation where it matters.

Client: program using operations defined in interface. 테스트해보거나 실행해보는 파일

Implementation: actual code implementing operations. 실제 코드를 구현한 파일

Interface: description of data type, basic operations. 헤더파일 처럼 선언만 해두는것

Stacks

Example: Stack of strings data type (implemented in Java)

public class	StackOfStrings	스트링 타입을 쌓아서 스택으로 구현
	StackOfStrings()	<i>create an empty stack</i>
void	push(String item)	<i>insert a new string onto stack</i>
String	pop()	<i>remove and return the string most recently added</i>
boolean	isEmpty	<i>is the stack empty?</i>
boolean	isFull	<i>is the statck full?</i>
int	size()	<i>member of strings on the stack</i>

Warmup client: Reverse sequence of strings from standard input.

Stacks

Stack test client:

- Read string from standard input.
 - If string equals "-", pop string from stack and print.
 - Otherwise, push string onto stack.

```
public static void main (String[] args) {  
    StackOfStrings stack = new StackOfStrings();  
    while (!StdIn.isEmpty()) {  
        String s = StdIn.readString();  
        if (s.equals("-")  
            Stdout.print(stack.pop()));  
        else  
            stack.push(s);  
    }  
}
```

Exercise:

```
%more tobe.txt  
to be or not to - be - - that - - - is  
% java StackOfStrings < tobe.txt  
to be not that or be
```


Stacks using dynamic arrays

Array implementation of a stack:

- Use array `s[]` to store `N` items on stack.
 - `push()`: add new item at `s[N]`.
 - `pop()`: remove item from `s[N-1]`.
- 벡터에는 푸시 팝 메소드가 있기 때문에 따로 구현할 필요가 없음
Grace of C++!

<code>s[]</code>	to	be	or	not	to	be	null	null	null	null
	0	1	2	3	4	5	6	7	8	9

`N` capacity = 10

`N`은 빈자리의 가장 앞을 나타낸다.
capacity는 최대 길이.

Defect. Stack overflows when `N` exceeds capacity. *[stay tuned]*

Stacks using dynamic arrays

```
public class FixedSizeStackOfStrings {  
    private String[] s;  
    private int N = 0;  
  
    public FixedSizeStackOfStrings(int capacity) {  
        s = new String[capacity];  
    }  
    public boolean isEmpty() {  
        return N == 0;  
    }  
    public void push(String item) {  
        s[N++] = item;  
    }  
    public String pop() {  
        return s[--N];  
    }  
}
```

a shortcoming
(stay tuned)

use to index into array;
then increment N

decrement N:
then use to index into array

Stacks using dynamic arrays

Things to consider:

- **Overflow and underflow:**
 - Underflow: throw exception if pop from an empty stack or return null;
 - Overflow: use resizing array for array implementation. [stay tuned]
- **Null items:** Allow null items to be inserted or not. Clarify during the design.
- **Loitering:** Holding a reference to an object when it is no longer needed.

```
public String pop() {  
    return s[--N];  
}
```

loitering

```
public String pop() {  
    String item = s[--N];  
    s[N] = null;  
    return item;  
}
```

This version avoids "loitering": Garbage collector can reclaim memory only if no outstanding references. In C/C++ implementation. free the resources.

Stacks using dynamic arrays

Problem: Requiring client to provide **capacity** (size of stack) is inappropriate.

Question: How to grow and shrink array?

First try.

- **push():** increase size of array `s[]` by 1.
- **pop():** decrease size of array `s[]` by 1.

Too expensive.

- Need to copy all items to a new array.
- Inserting first N items takes time proportional to $1 + 2 + 3 + \dots + N \approx N^2/2$.


infeasible for large N

Challenge: Ensure that array resizing happens infrequently.

Stacks using dynamic arrays

Q. How to grow and shrink array?

"successive doubling"

↓ 알고리즘의 효율성을 위해서 두배씩 늘린다. $O(n)$ 알고리즘이 됨

A. If array is full, create a new array of **twice** the size, and copy items.

```
public class ResizingStackOfStrings {
    s = new String[1];
}

메모리를 늘릴때는 새로운 메모리를 새로운 사이즈로 할당해서 기존 데이터들을 복사한다.
public void push(String item) {
    if (N == s.length) resize(s.length * 2);
    s[N++] = item;
}

private void resize(int capacity) {
    String[] copy = new String[capacity];
    for (int i = 0; i < N; i++)
        copy[i] = s[i];
    s = copy;
}
```

Consequence: Inserting first N items takes time proportional to N , not N^2

Stacks using dynamic arrays

Q. Cost of inserting first N items by `resize(s.length + 10)`?

A. $T(N) = 1 + (10 + 20 + 30 + \dots + N)$

1 array access per push
k array accesses when memory is resized by increment of 10
(ignoring cost to create new array)
(assuming `realloc()` costs copying each item one by one)

When $N = 1$, Capacity = 1 \rightarrow 11	// (?) cost to copy the existing items into the new array
Cost: 1 + (0)	// (0) since no copy is needed
When $N = 2$, Capacity = 11	
Cost: 1 + (0)	// (0) items to copy into the new array
When $N = 3$, Capacity = 11	
Cost: 1 + (0)	// (0) since no copy is needed
When $N = 4$, Capacity = 11	
Cost: 1 + (0)	
....	
When $N = 11$, Capacity = 11 \rightarrow 21	
Cost: 1 + (10)	// (10) items to copy into the new array
When $N = 12$, Capacity = 21	
Cost: 1 + (0)	
....	
When $N = 21$, Capacity = 21 \rightarrow 31	
Cost: 1 + (20)	// (20) items to copy into the new array
When $N = 22$, Capacity = 31	
Cost: 1 + (0)	

Stacks using dynamic arrays

Q. Cost of inserting first N items by `resize(s.length + 10)`?

A. $T(N) = 1 + (10 + 20 + 30 + \dots + N) = ?$

How many terms? k terms, then $N = 10k$

$$T(N) = 1 + (10 + 20 + 30 + \dots + N)$$

Let $N = 10k$, then it becomes

$$\begin{aligned} T(N) &= 1 + (10 + 20 + 30 + \dots + 10k) \\ &= 1 + 10(1 + 2 + 3 + \dots + k) \end{aligned}$$

$$= 1 + 10 \frac{k(k+1)}{2}$$

$$= 1 + 10 \frac{\frac{N}{10}(\frac{N}{10} + 1)}{2}$$

$$\text{Therefore, } T(N) = 1 + \frac{N}{2} \left(\frac{N}{10} + 1 \right)$$

→ The time complexity of the algorithm is $O(n^2)$.

Stacks using dynamic arrays

Q. Cost of inserting first N items by `resize(s.length * 2)` ?

A. $T(N) = 1 + (1 + 2 + 4 + 8 + \dots + N)$

1 array access per push
k array accesses to double to size k
(ignoring cost to create new array)
(assuming `realloc()` costs copying each item one by one)

When $N = 1$, Capacity = 1

Cost: $1 + (0)$

When $N = 2$, Capacity = 1

Cost: $1 + (1)$

When $N = 3$, Capacity = 2

Cost: $1 + (2)$

When $N = 4$, Capacity = 4

Cost: $1 + (0)$

When $N = 5$, Capacity = 4

Cost: $1 + (4)$

When $N = 6$, Capacity = 8

Cost: $1 + (0)$

When $N = 7$, Capacity = 8

Cost: $1 + (0)$

When $N = 8$, Capacity = 8

Cost: $1 + (0)$

When $N = 9$, Capacity = 8

Cost: $1 + (8)$

// (?) cost to copy the existing items into the new array

// (1) items to copy

// (2) items to copy into the new array

// (0) since no copy is needed

// (4) items to copy into the new array

// (8) items to copy into the new array

Stacks using dynamic arrays

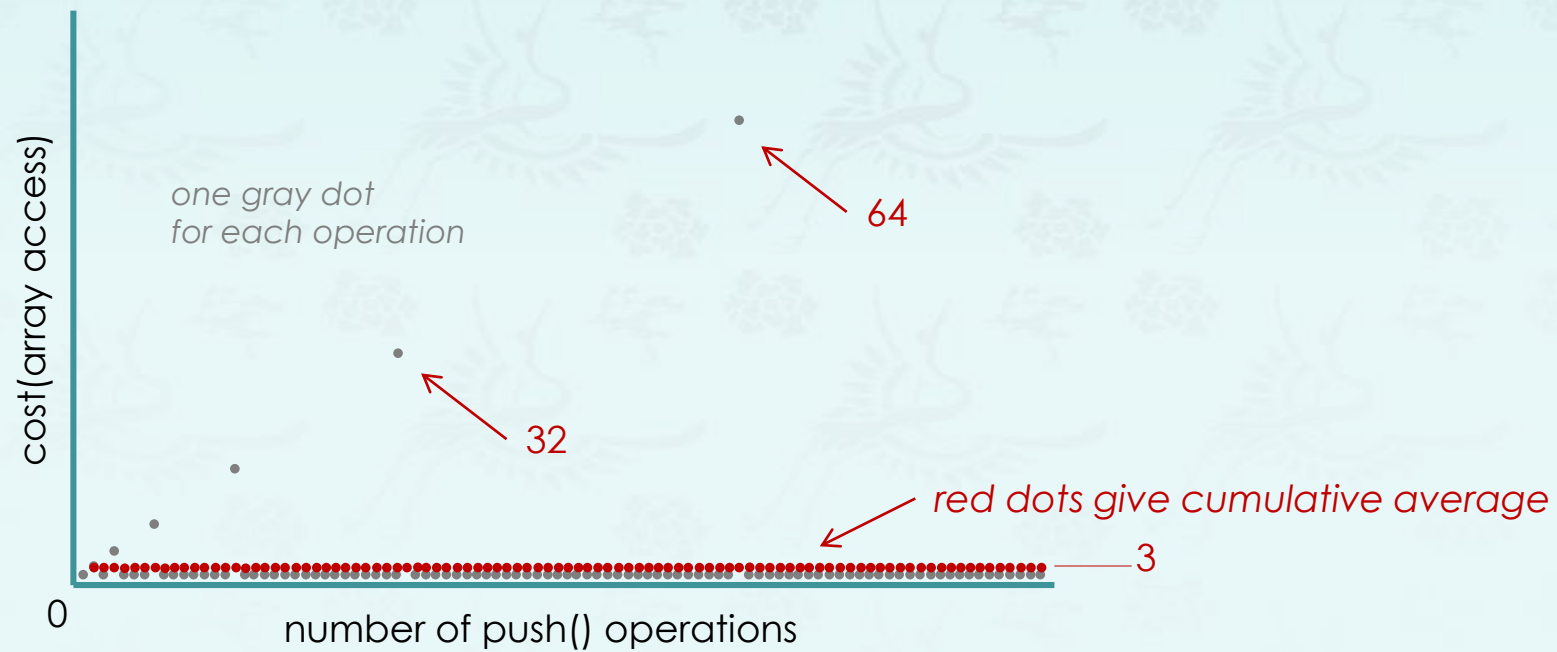
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Stacks using dynamic arrays

Q. Cost of inserting first N items by `resize(s.length * 2)` ?

A. $T(N) = 1 + (1 + 2 + 4 + 8 + \dots + N)$



Stacks using dynamic arrays

Q: How to shrink array?

First try.

- **push():** double size of array $s[]$ when array is full
- **pop():** halve size of array $s[]$ when array is one-half full.

Too expensive in worst case.

- Consider push-pop-push-pop- ... sequence when array is full
- Each operation takes time proportional to N .

$N=5$ to be or not to be null null

$N=4$ to be or not

$N=5$ to be or not to be null null

$N=4$ to be or not

Stacks using dynamic arrays

Q: How to shrink array?

Efficient solution

- **push()**: double size of array `s[]` when array is full
- **pop()**: **halve** size of array `s[]` when array is **one-quarter full**.

75%의 데이터가 사라지면 capacity를 반으로 줄인다.

```
public String pop() {  
    String item = s[--N];  
    s[N] = null;  
    if (N > 0 && N == s.length/4)  
        resize(s.length/2);  
    return item;  
}
```

❖ **Invariant**. Array is between 25% and 100% full.

Stacks using dynamic arrays

Amortized analysis: Average running time per operation over a worst-case sequence of operations.

Proposition: Starting from an empty stack, any sequence of N push and pop operations takes time proportional to N .

	best	worst	amortized
construct	$O(1)$	$O(1)$	$O(1)$
push	$O(1)$	$O(n)$	$O(1)$
pop	$O(1)$	$O(n)$	$O(1)$
size	$O(1)$	$O(1)$	$O(1)$

doubling and
halving operations

order of growth of running time
for resizing stack with N items

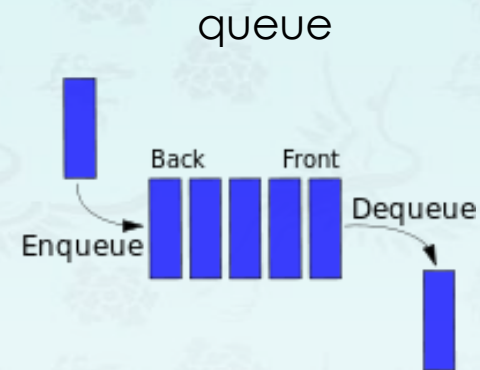
Data Structures

- ***stacks & queues***
using dynamic arrays
- *some applications*



Queues

Queue: An ordered list in which **enqueues** (insertion or add) at the **rear** and **dequeues** (deletion or remove) take place at different end or **front**. It is also known as a First-in-first-out(**FIFO**) list.



- ❖ Items can only be added at the **rear** of the queue and the only item that can be removed is the one at the **front** of the queue.

Queues

Queue: An ordered list in which **enqueues** (insertion or add) at the **rear** and **dequeues** (deletion or remove) take place at different end or **front**. It is also known as a First-in-first-out(FIFO) list.

ADT Queue is

objects: a finite ordered list with zero or more elements

functions:

```
Queue CreateQueue(maxQueueSize)
boolean IsFull(queue, maxQueueSize)
boolean IsEmpty(queue)
void Add(queue, item)           // Enqueue
Element Delete(queue)           // Dequeue
```


Queues

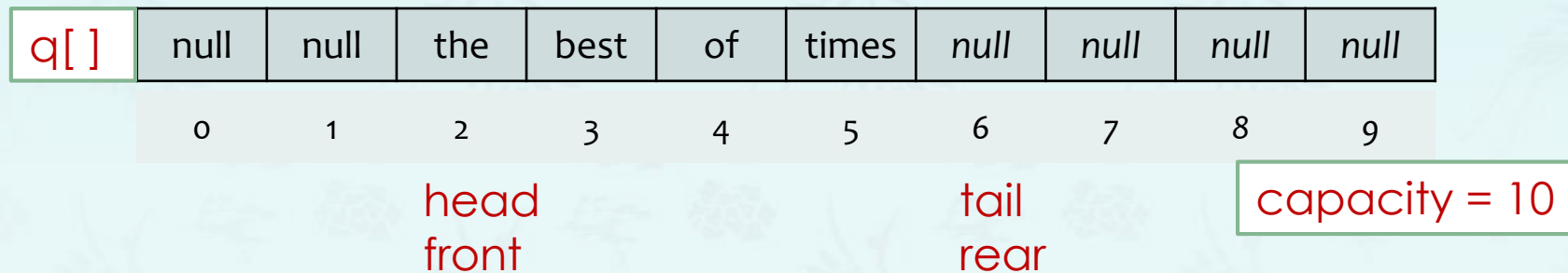
Example: Stack of strings data type (implemented in Java)

public class	QueueOfStrings	
	QueueOfStrings()	<i>create an empty queue</i>
void	enqueue(String item)	<i>insert a new string onto queue</i>
String	dequeue()	<i>remove and return the string least recently added</i>
boolean	isEmpty()	<i>is the queue empty?</i>
boolean	isFull()	<i>is the queue full?</i>
int	size()	<i>member of strings on the queue</i>

Queues

Array implementation of a queue:

- Use array `q[]` to store items in queue.
- `enqueue()`: add new item at `q[tail]`.
- `dequeue()`: remove item from `q[head]`.
- Update head and tail modulo the capacity.
- Add resizing array.

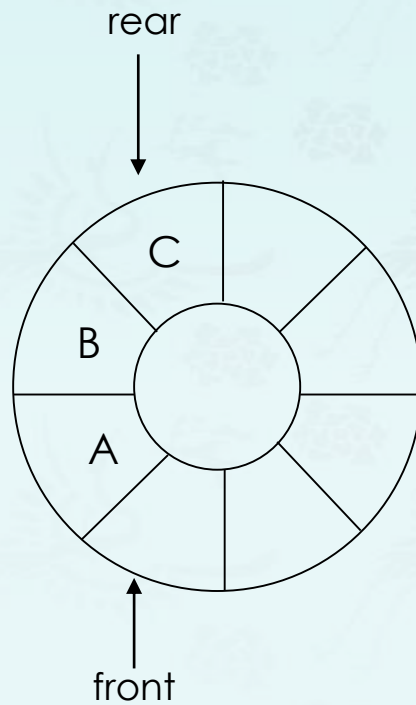


Q. How to resize?

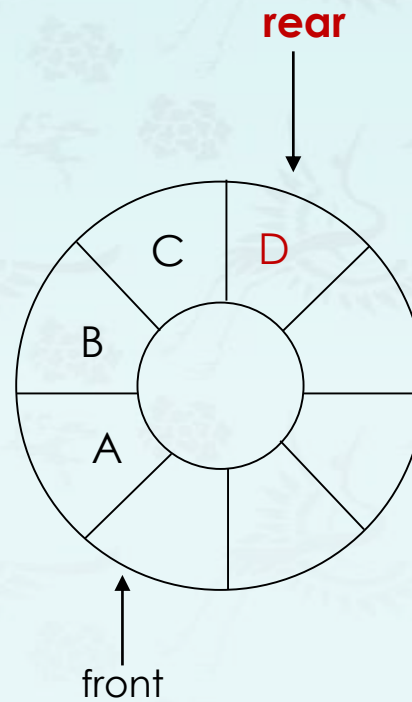
Queues

Circular queue:

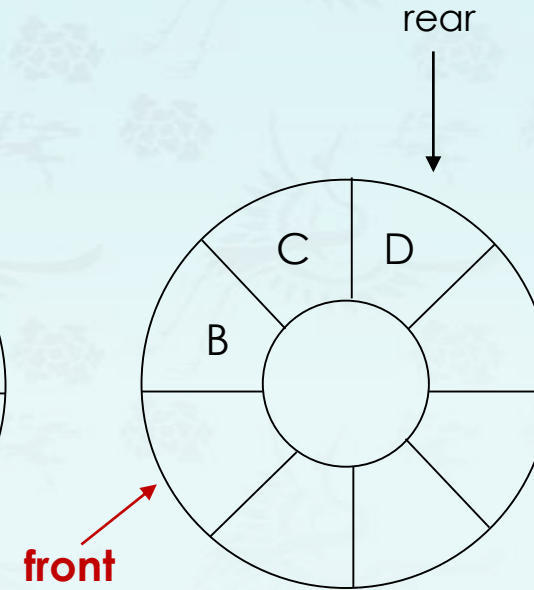
- To avoid shifting array and resizing array which is costly
- The array positions are handled in a circle rather than in a straight line.



(a) Initial



(b) Addition/Enqueue

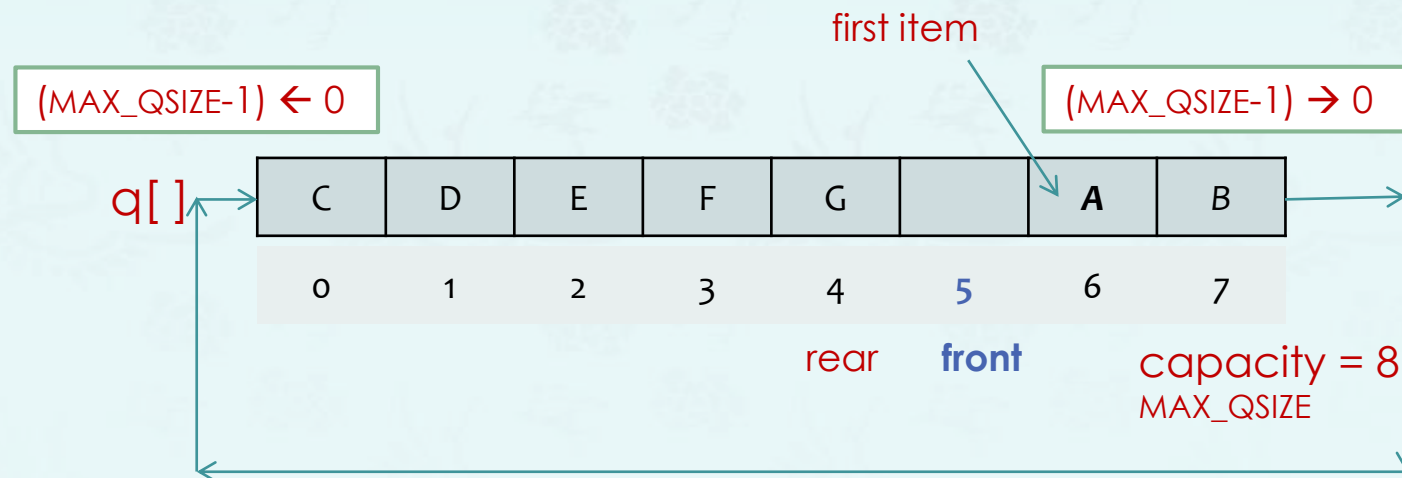


(c) Deletion/Dequeue

Queues

Circular queue:

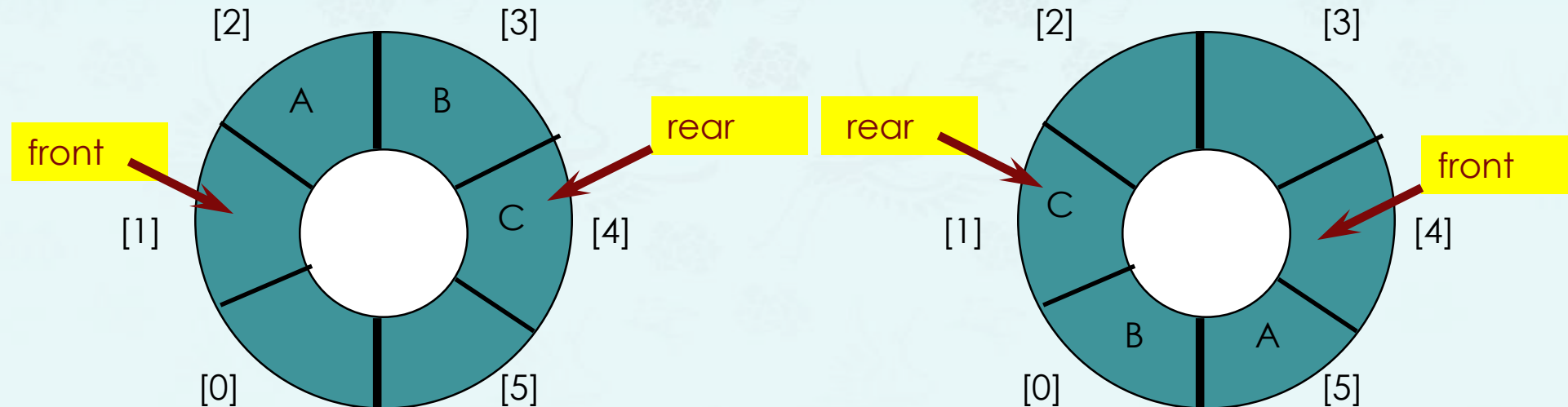
- To be circular, the position next to position $\text{MAX_QSIZE}-1$ is 0, the position that precedes 0 is $\text{MAX_QSIZE}-1$
- By convention, the **first item** is located at $(\text{front} + 1)$ position.
- If $\text{front} = \text{rear}$, then queue is **empty** or **full**?
To distinguish this case, double the queue size just before it becomes full.
- The initial value for **front** and **rear** is 0.



Queues

Circular queue:

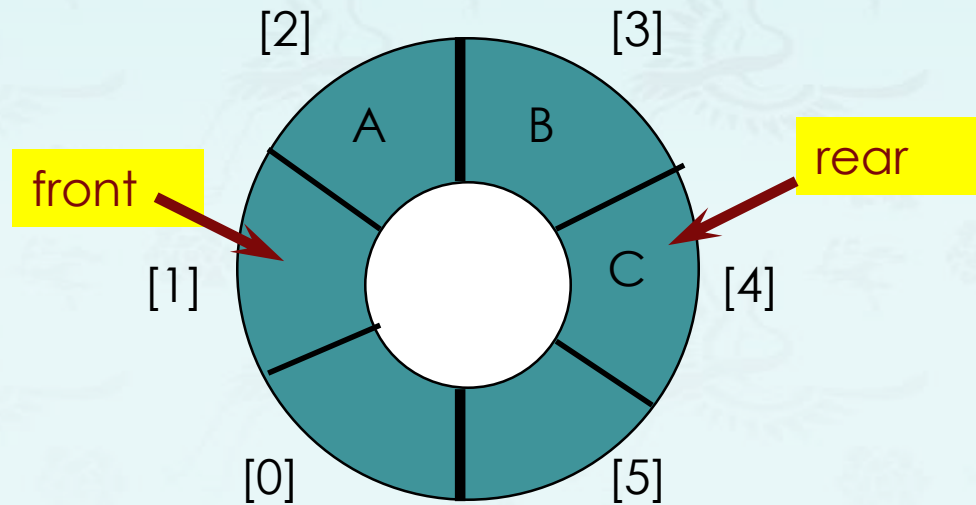
- Use integer variables **front** and **rear**.
 - **front** is one position counterclockwise from first element
 - **rear** gives position of last element



Queues

Circular queue:

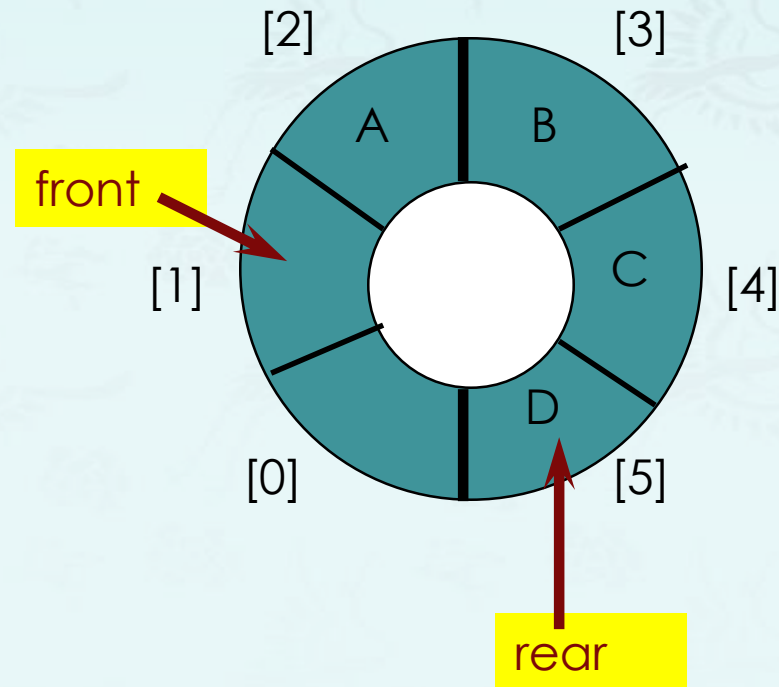
- **Add** an element
 - Move **rear** one clockwise.



Queues

Circular queue:

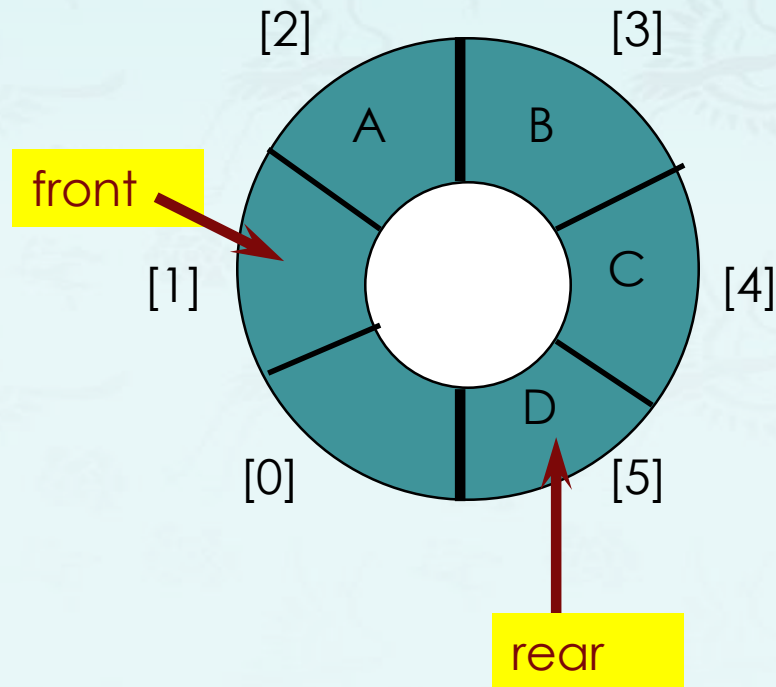
- **Add** an element
 - Move **rear** one clockwise.
 - Then put into **queue[rear]**



Queues

Circular queue:

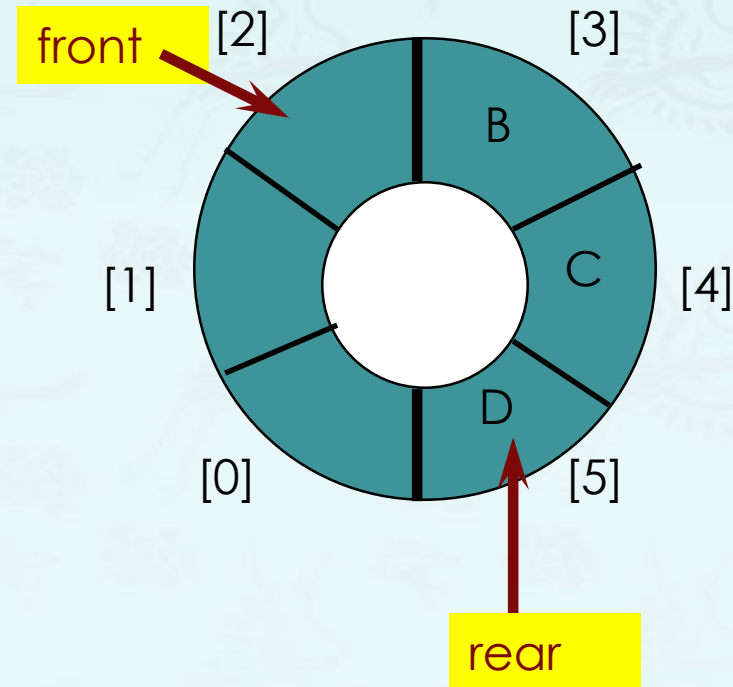
- **Delete** an element
 - Move **front** one clockwise.



Queues

Circular queue:

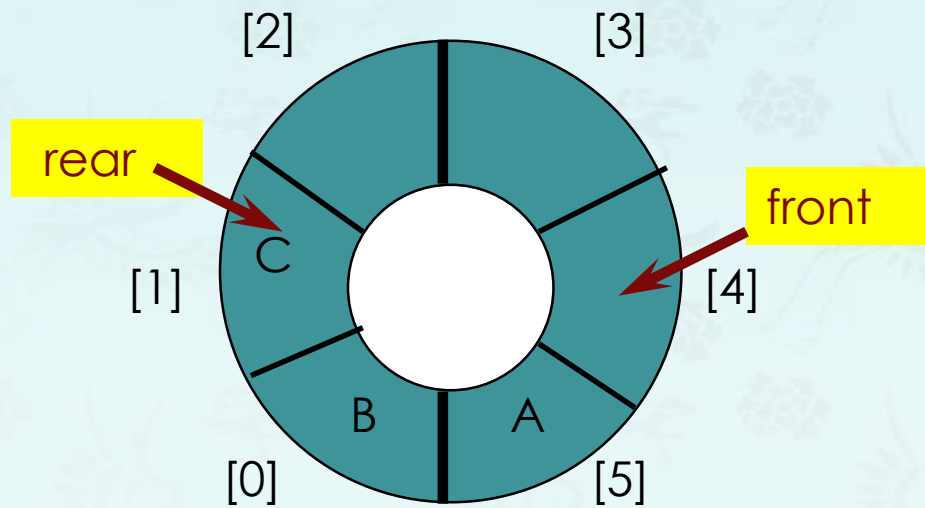
- **Delete** an element
 - Move **front** one clockwise.
 - Then extract from **queue[front]**.



Queues

Circular queue:

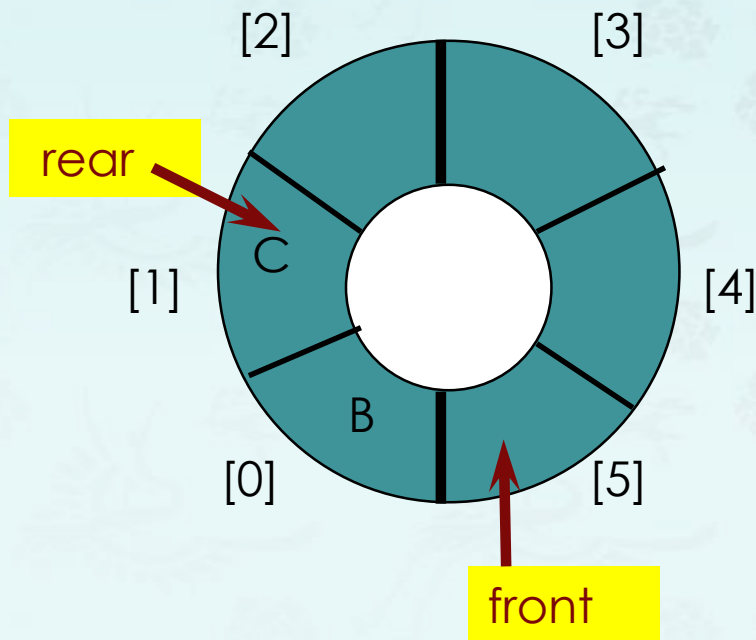
- Empty that queue



Queues

Circular queue:

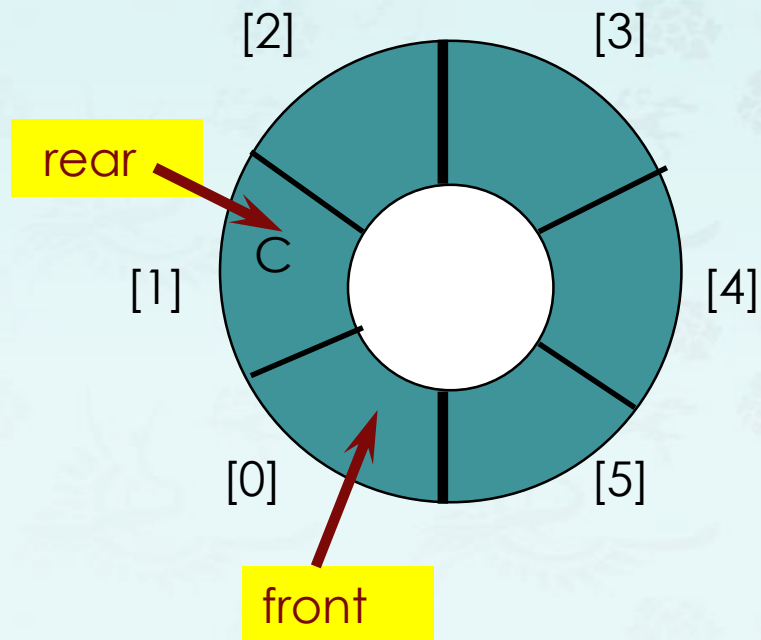
- Empty that queue



Queues

Circular queue:

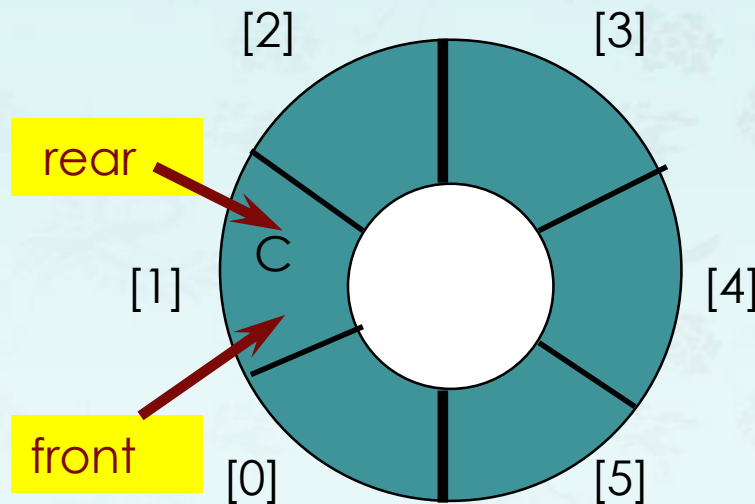
- Empty that queue



Queues

Circular queue:

- **Empty** that queue

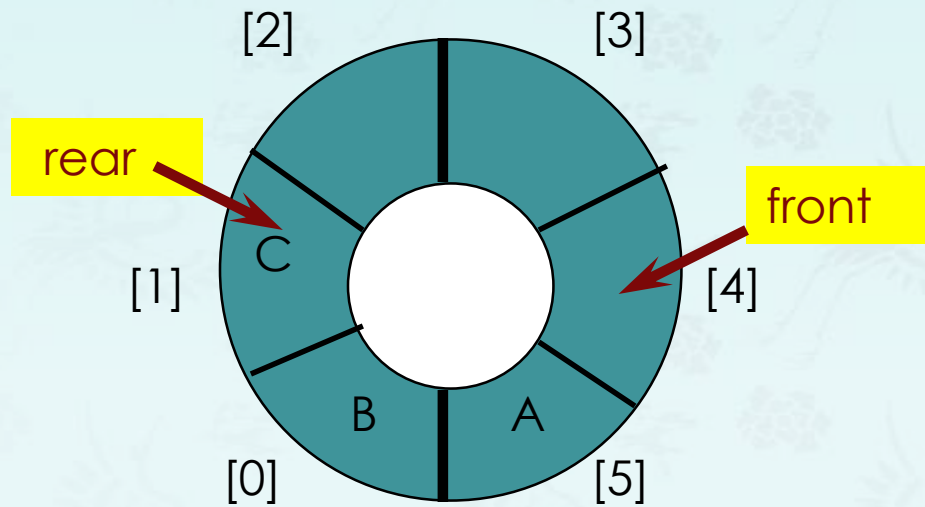


- When a series of removes causes the queue to become empty, **front = rear**.
- When a queue is constructed, it is empty.
- So initialize **front = rear = 0**.

Queues

Circular queue:

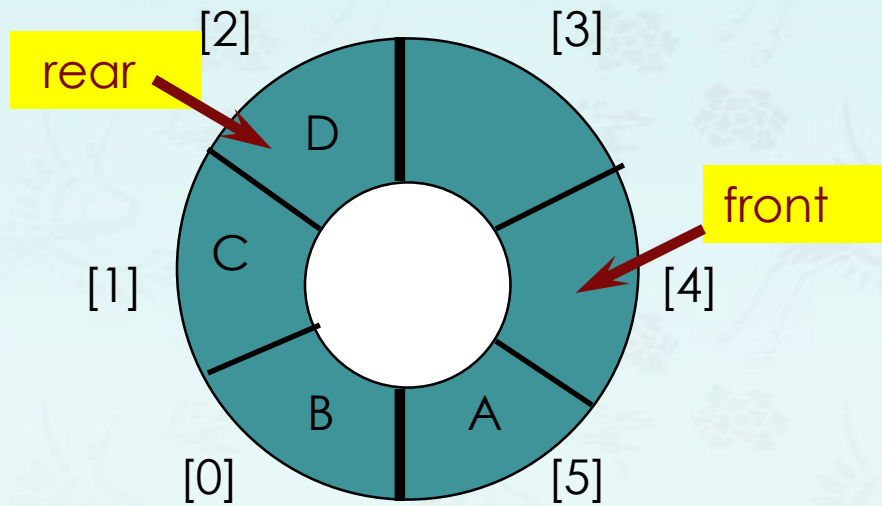
- Full that queue



Queues

Circular queue:

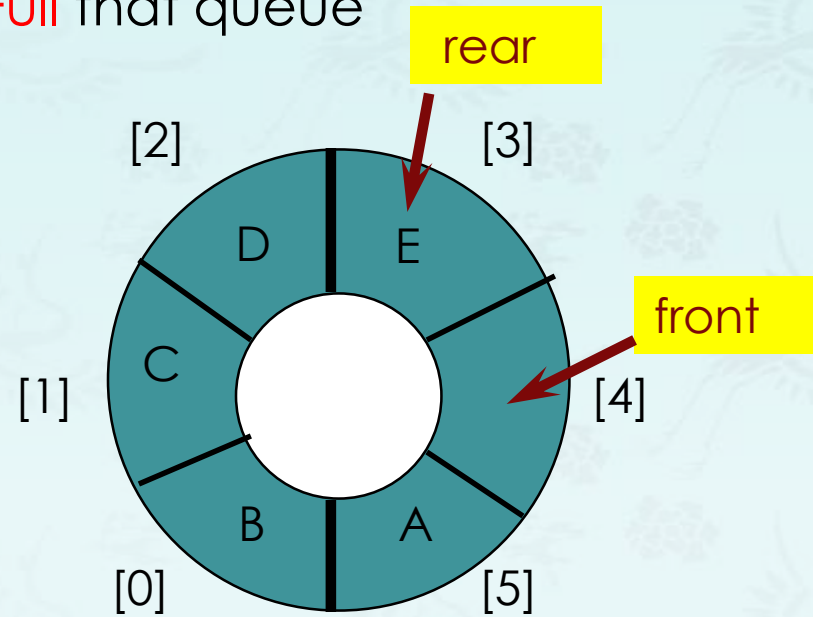
- Full that queue



Queues

Circular queue:

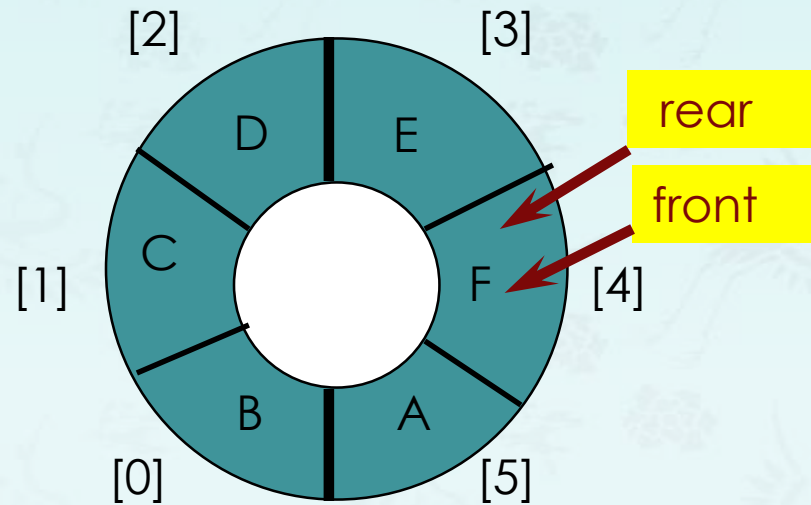
- Full that queue



Queues

Circular queue:

- **Full** that queue



- When a series of adds causes the queue to become full, **front = rear**.
- So we **cannot distinguish** between a full queue and an empty queue!

Queues

Circular queue:

- Challenge: **front = rear** when queue is empty and full.
- Solutions:
 - Don't let the queue get full.
 - When the addition of an element will cause the queue to be full, increase array size.
 - This is what the text does.
 - Define a boolean variable **lastOperationIsAddQ**.
 - Following each **AddQ** set this variable to true.
 - Following each **DeleteQ** set to false.
 - Queue is **empty** iff $(\text{front} == \text{rear}) \ \&\& \ !\text{lastOperationIsAddQ}$
 - Queue is **full** iff $(\text{front} == \text{rear}) \ \&\& \ \text{lastOperationIsAddQ}$

Queues

Circular queue:

- Challenge: **front = rear** when queue is empty and full.
- Solutions: (continued)
 - Define an integer variable **size**.
 - Following each **AddQ** do **size++**.
 - Following each **DeleteQ** do **size--**.
 - Queue is empty iff (**size == 0**)
 - Queue is full iff (**size == arrayLength**)
 - Performance is slightly better when first strategy is used.

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Stack and Queue

Data Structures
C++ for C Coders

한동대학교 김영섭 교수
idebtor@gmail.com

stacks & queues using dynamic arrays
some applications – infix and postfix