# Stack and Queue

Data Structures C++ for C Coders

한동대학교 김영섭 교수 idebtor@gmail.com

applications - infix to postfix



#### Queues

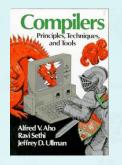
**Queue:** An ordered list in which **enqueues** (insertion or add) at the **rear** and **dequeues** (deletion or remove) take place at different end or **front**. It is also known as a Fist-in-first-out(FIFO) list.



Items can only be added at the rear of the queue and the only item that can be removed is the one at the front of the queue.

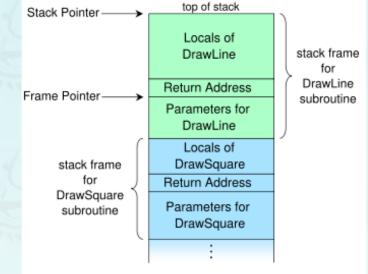


#### **Stack** and Queue Applications



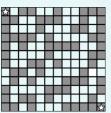






- Parsing in a compiler. (p.127)
- Undo in a word processor.
- Back button in a Web browser.
- PostScript language for printers.
- Backtracking as in a maze (p.121)
- Implementing function calls in a compiler. (p.108)

• ...





### Stack and Queue Applications

In a computer OS: Requests for services come in unpredictable order and timing, sometimes faster than they can be serviced.

- print a file
- need a file from the disk system
- send an email
- job scheduling

Goal: Convert an infix expression to a postfix expression using a stack.



Stack: ( Output:

Stack: ( Output: 1

Stack: (+ Output: 1

Stack: (+ Output: 12

Stack:

Output: 1 2 +

Stack: \*

Output: 1 2 +

Stack: \*

Output: 1 2 + 3

Stack:

Output: 1 2 + 3 \*

- Operands are output immediately
- Stack operators until right parens
- Unstack until left parens
   Delete left parens
- In general, higher precedence operator must be output before lower one.)

postfix

Goal: Convert an infix expression to a postfix expression using a stack.



Stack: ( Output:

Stack: ( Output: 1

Stack: (+ Output: 1

Stack: (+ Output: 12

Stack:

Output: 12+

Stack: \*

Output: 1 2 +

Stack: \*

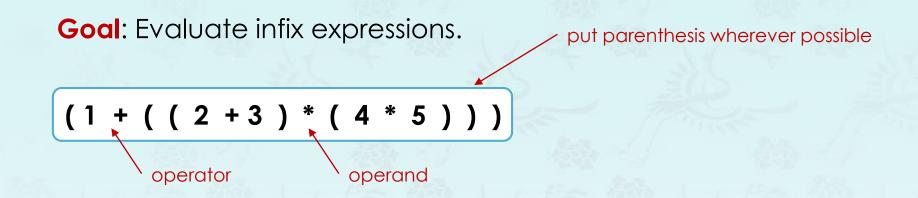
Output: 1 2 + 3

Stack:

Output: 1 2 + 3 \*

- Operands are output immediately
- Stack operators until right parens
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postfix 1 2 3 + 4 5 \* \* +



Two-stack algorithm. [E. W. Dijkstra]

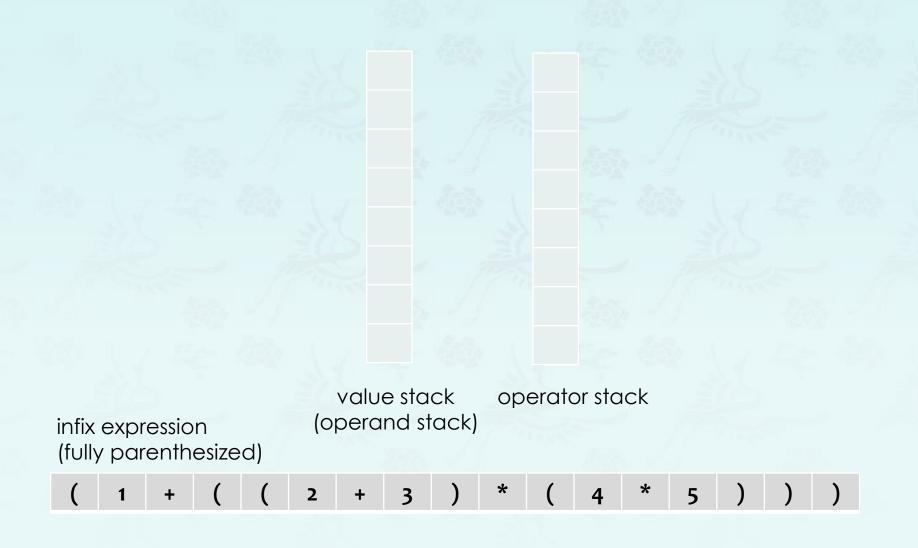
Goal: Evaluate infix expressions.

(1 + ( ( 2 + 3 ) \* ( 4 \* 5 ) ) )

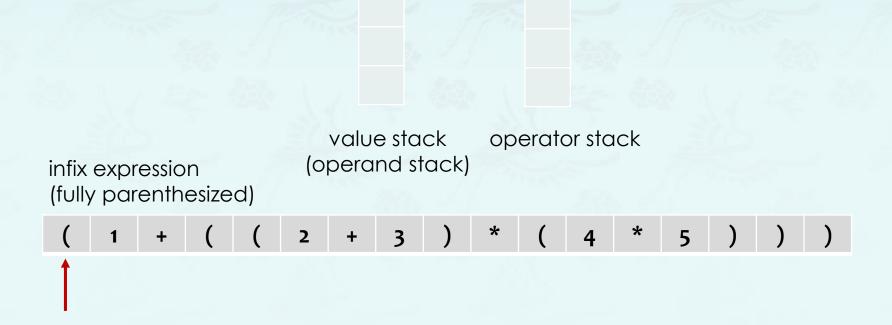
operator operand

Two-stack algorithm. [E. W. Dijkstra]

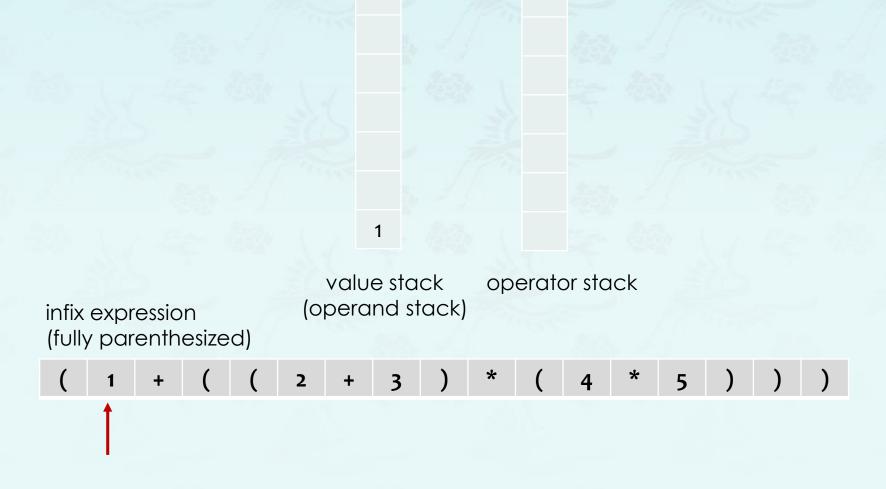
- Value: push onto the value stack.
- Operator: push onto the operator stack.
- Left parenthesis: ignore.
- Right parenthesis:
  - pop operator and two values;
  - push the result of applying that operator to those values onto the value stack.



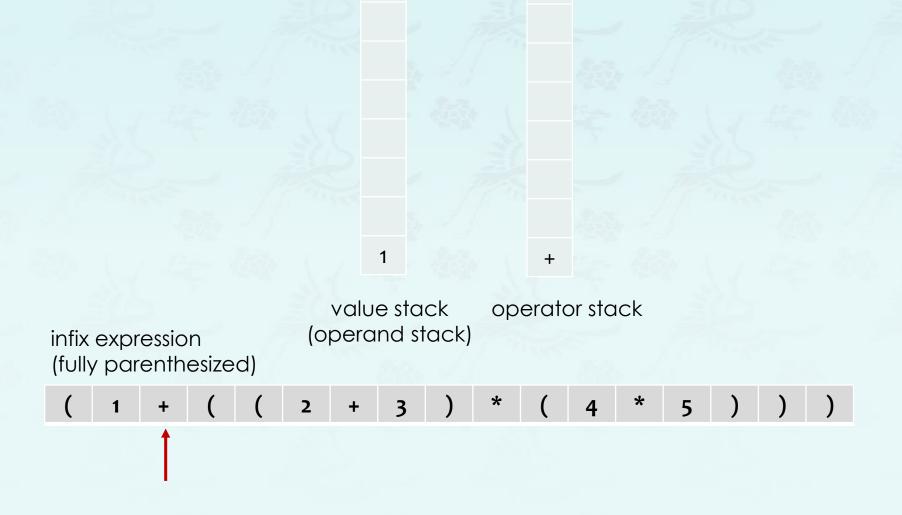
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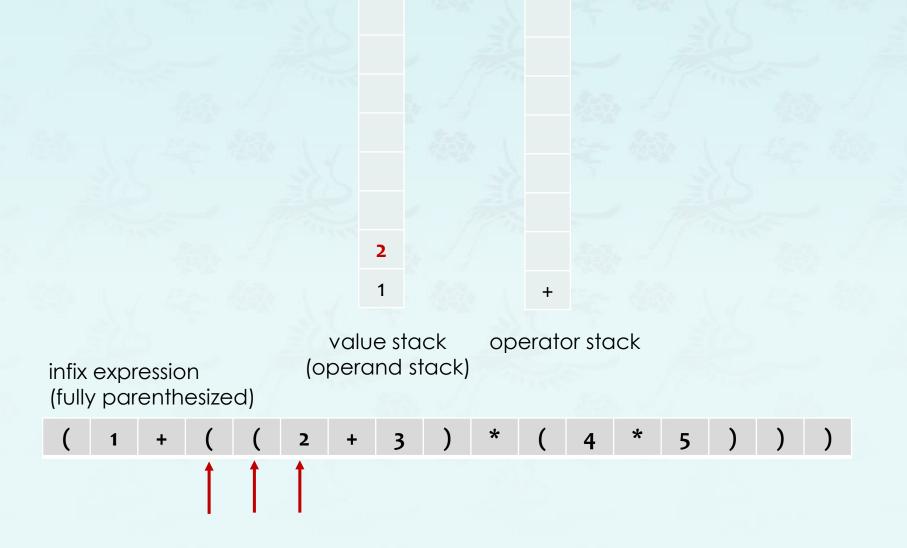
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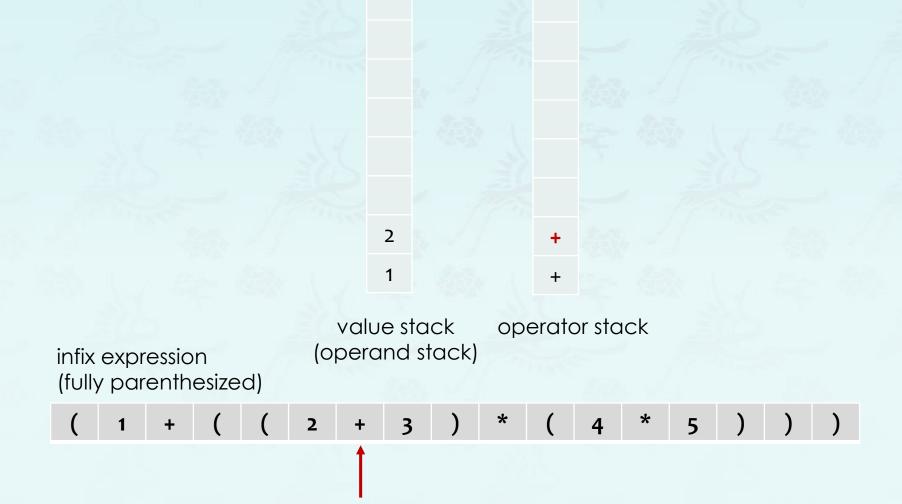
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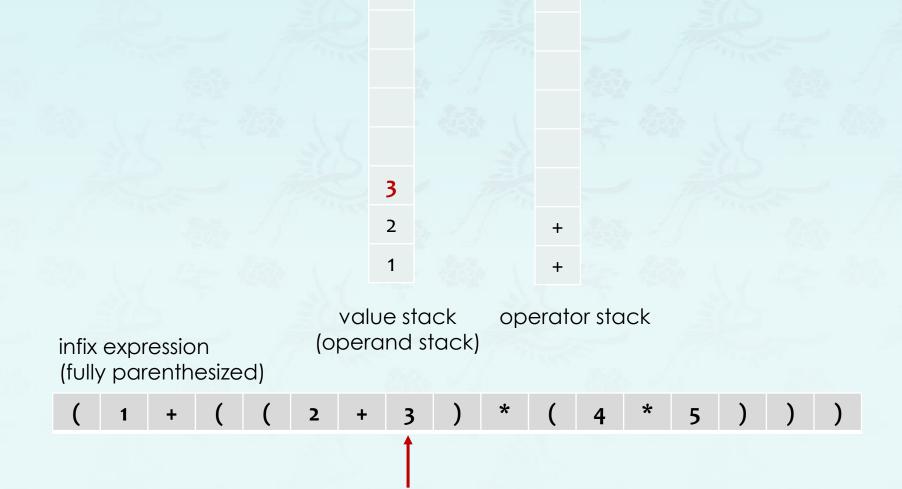
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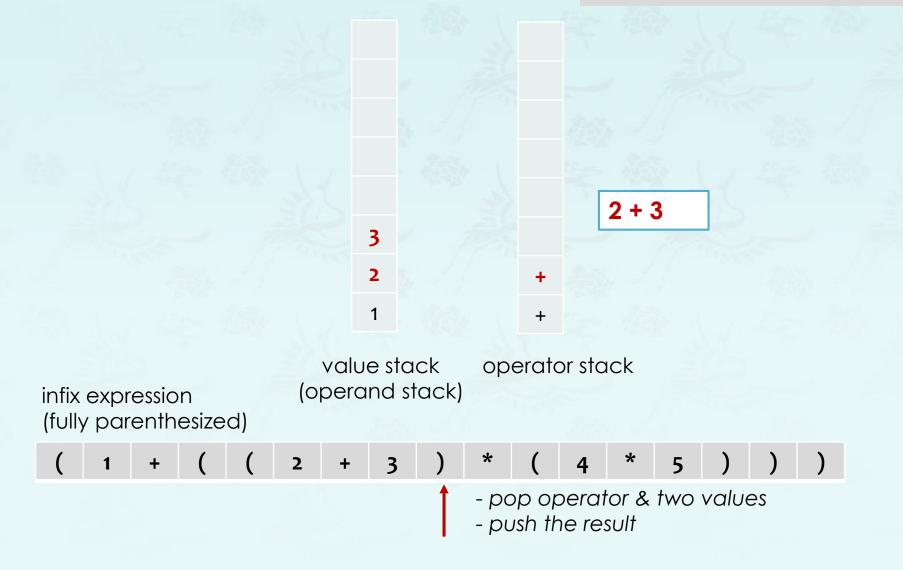
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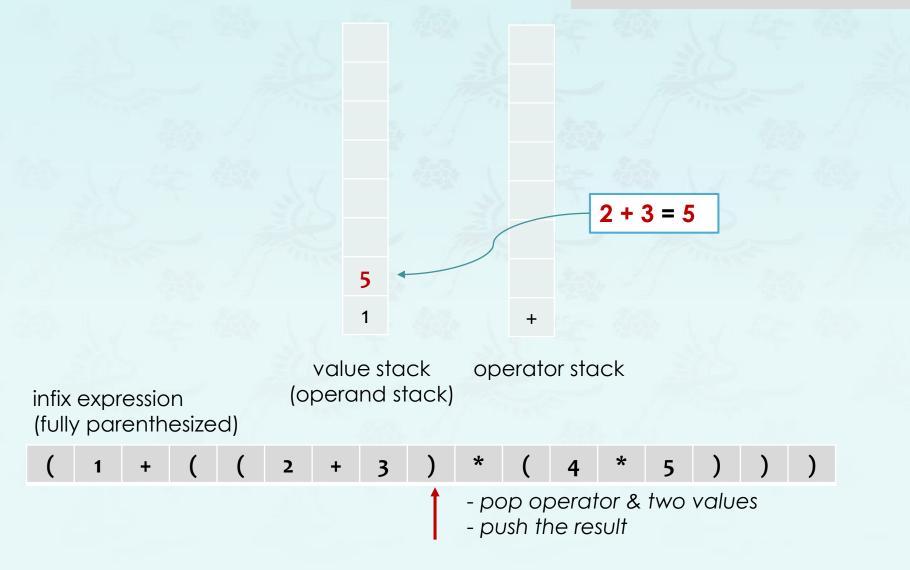
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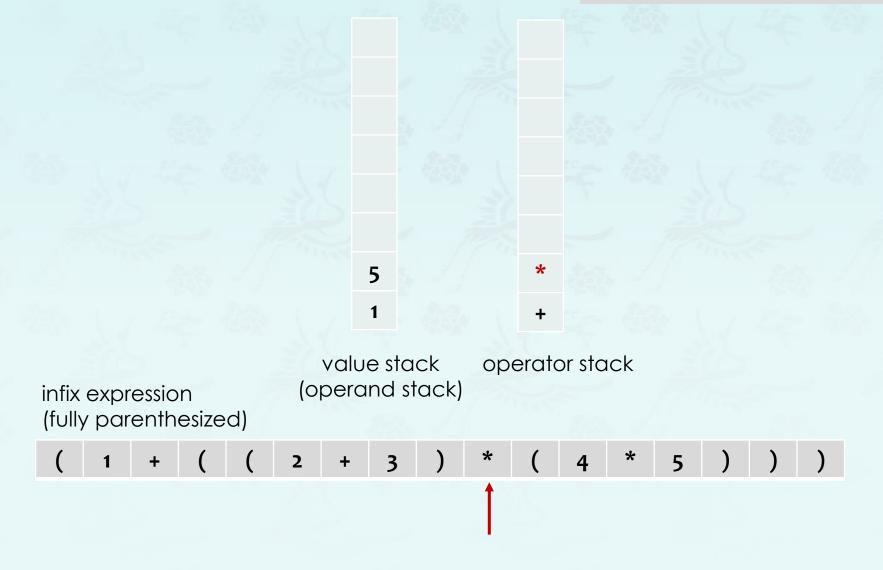
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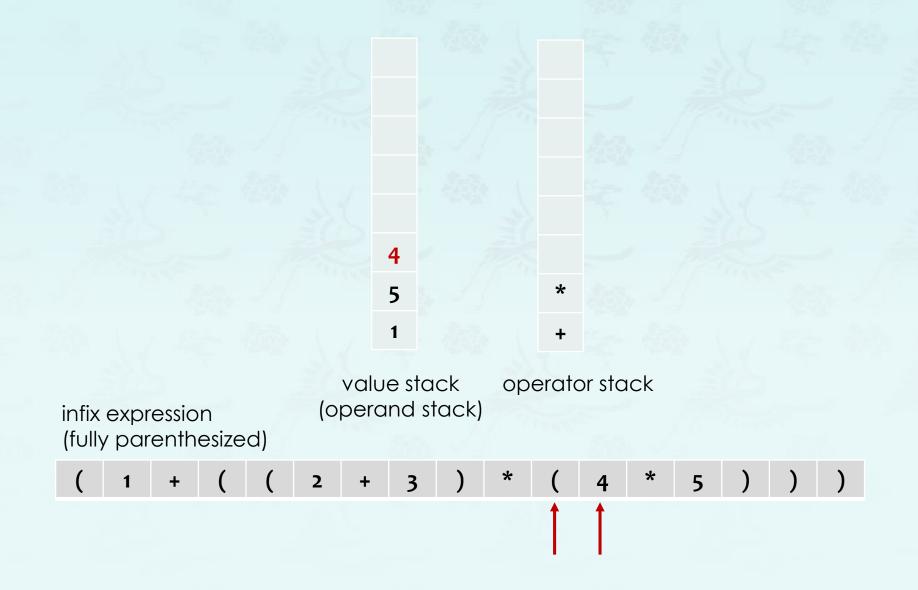


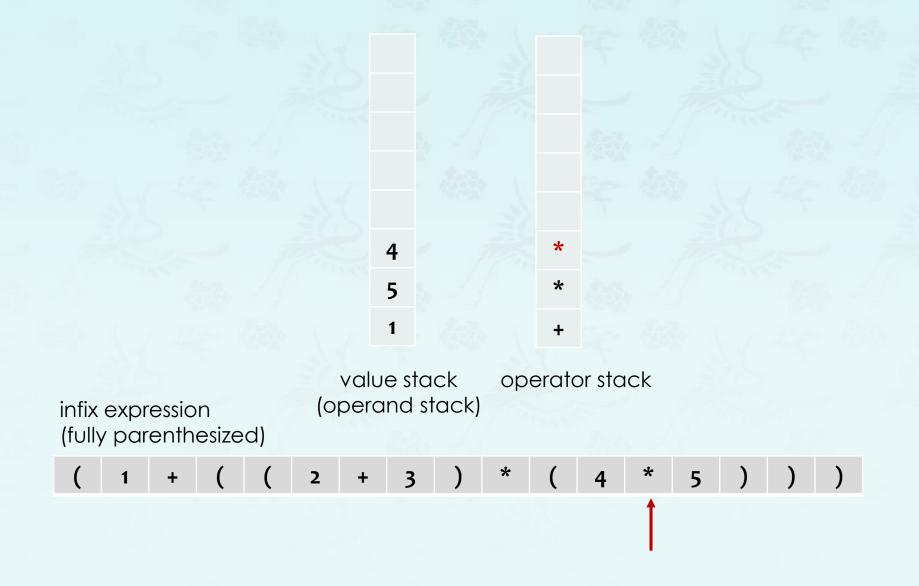
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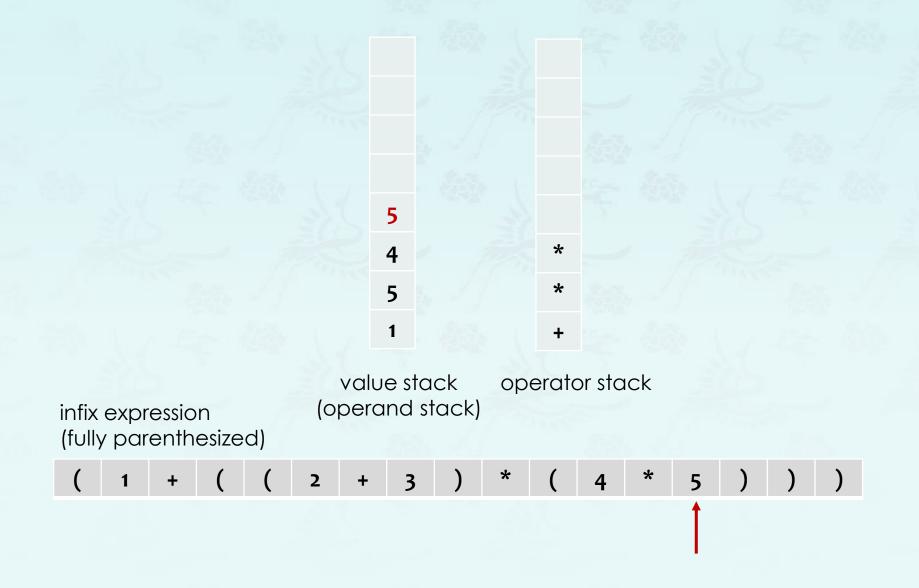


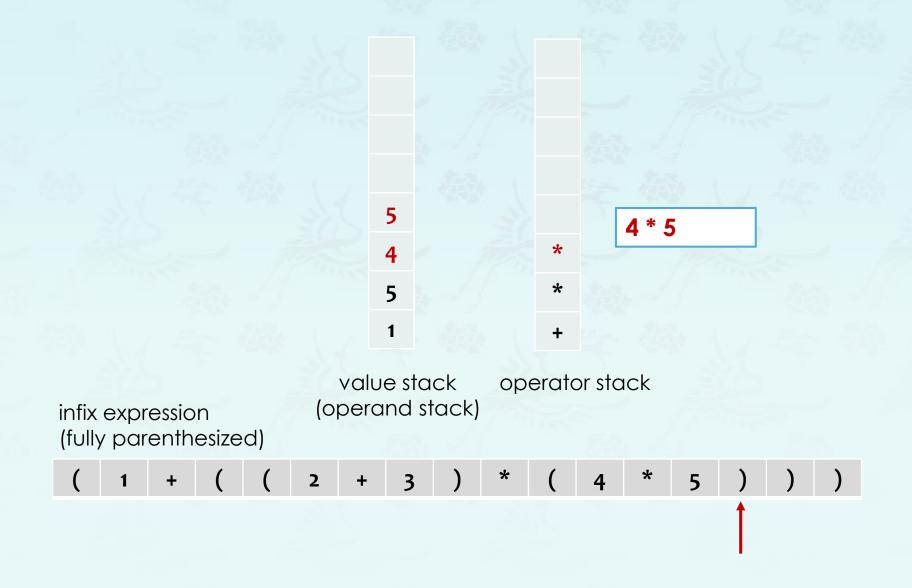
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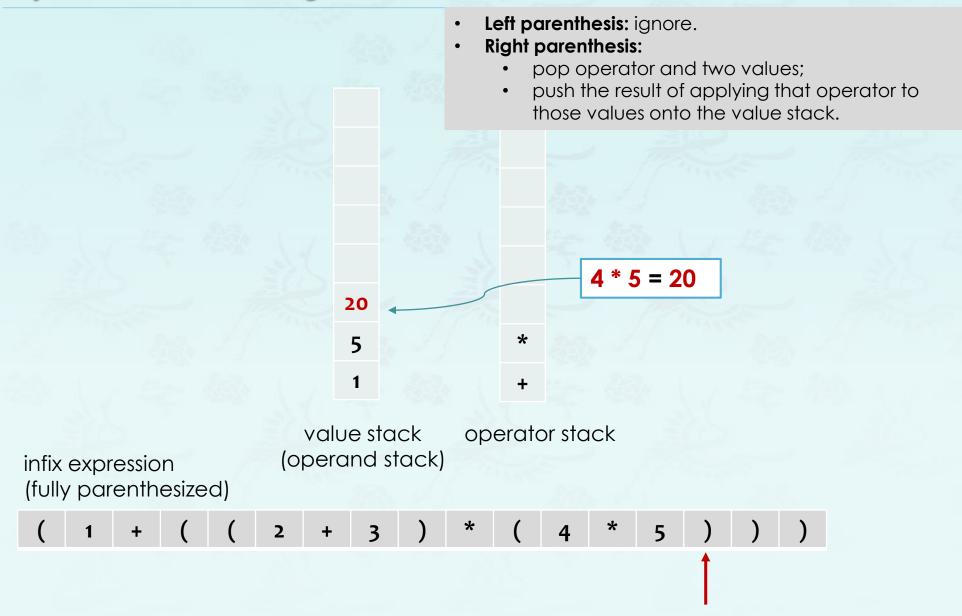


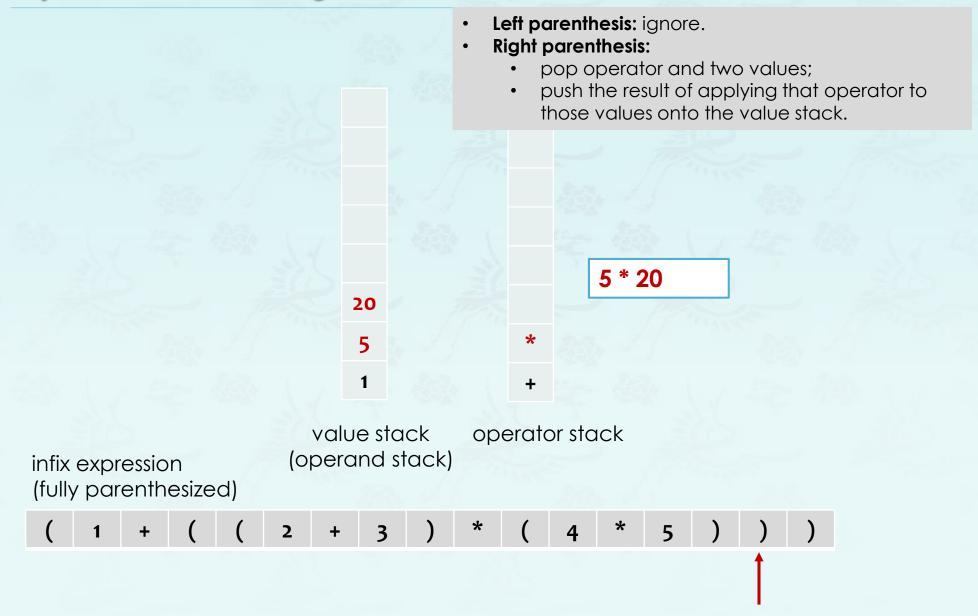


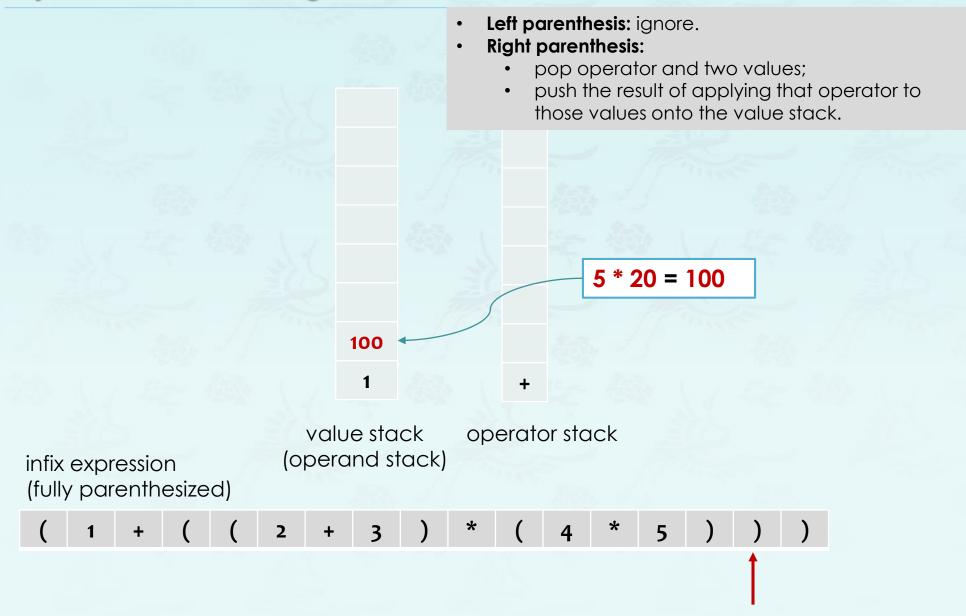


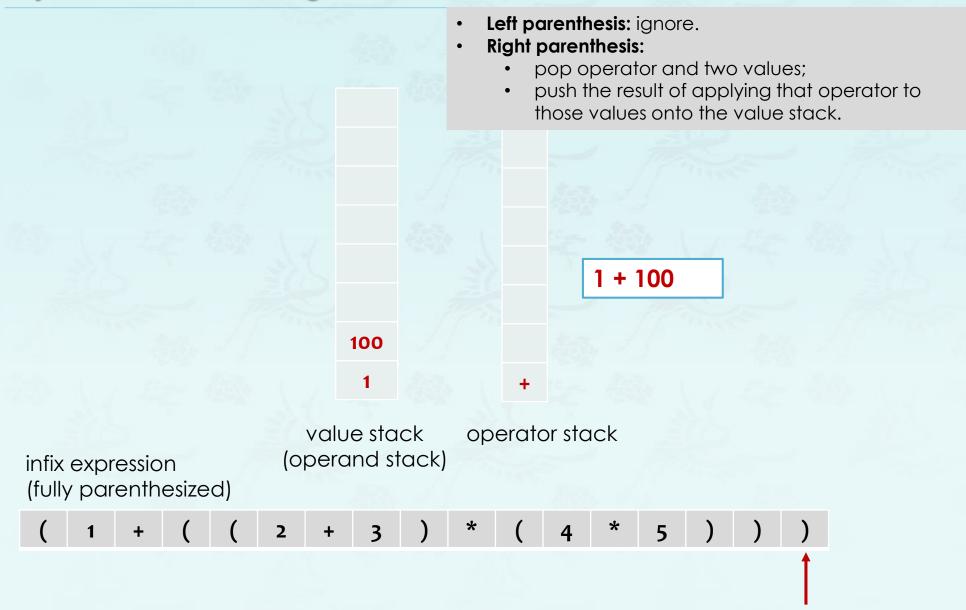


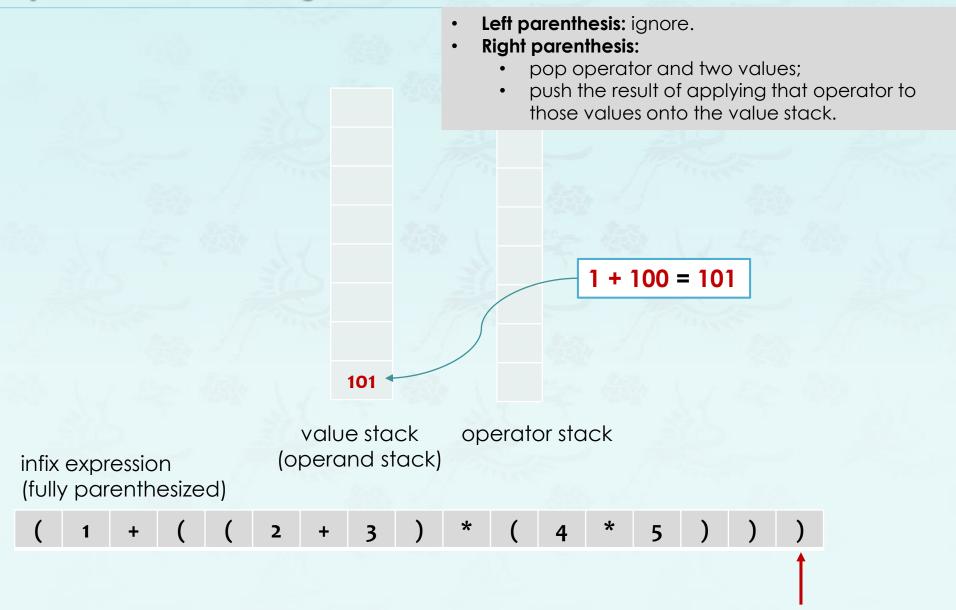


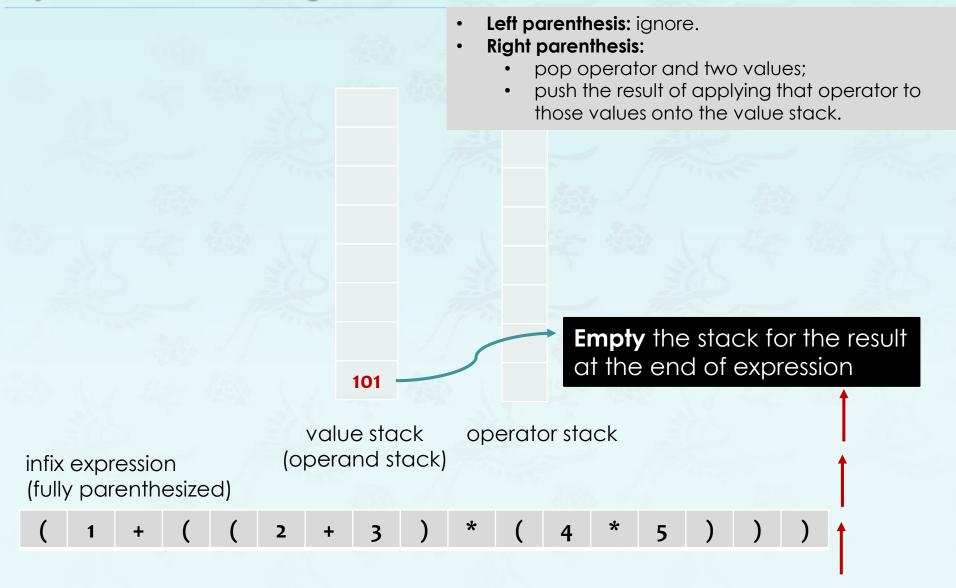












```
public class ArithmeticExpression {
  public static void main(String[] args) {
   Stack<Character> ops = new Stack<Character>();
   Stack<Double> vals = new Stack<Double>();
   String e = JOptionPane.showInputDialog(null,
     "Enter an expression", "Stack application", JOptionPane.QUESTION MESSAGE);
   if (e == null) return;  // Check "Cancel"
   for (int i = 0; i < e.length(); i++) {
     Character c = e.charAt(i);
     if (c.equals(' ') | c.equals('('));
     else if (c == '+') ops.push(c);
     else if (c == '*') ops.push(c);
     else if (c == ')') {
         Character op = ops.pop();
         if (op.equals('+')) vals.push(vals.pop() + vals.pop());
         else if(op.equals('*')) vals.push(vals.pop() * vals.pop());
     else {
         String s = "" + c;
         vals.push(Double.parseDouble(s));
   JOptionPane.showMessageDialog(null, e + " = " + vals.pop());
```

Q: How does it work?

A: When algorithm encounters an operator surrounded by two values within parentheses, it leaves the result on the value stack.

as if the original input were:

Repeating the argument:

Extensions: More ops, precedence order, associativity.

**Observation 1.** Dijkstra's two-stack algorithm computes the same value if the operator occurs **after** the two values.

Observation 2.

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Observation 2. All of the parentheses are redundant!

Bottom line: Postfix or "reverse Polish" notation. Applications: Postscript, calculators, JVM, ....

infix	postfix
2 + 3 * 4	2 3 4 * +
a * b + 5	a b * 5 +
(1 + 2) * 7	
a * b / c	
(a/(b-c+d))*(e-a)*c	
a / b - c + d * e - a * c	

Infix and postfix notation

infix	postfix
2 + 3 * 4	2 3 4 * +
a * b + 5	a b * 5 +
(1 + 2) * 7	1 2 + 7 *
a * b / c	a b * c /
( a / (b - c + d) ) * ( e - a ) * c	a b c - d + / e a - * c *
a / b - c + d * e - a * c	a b / c - d e * + a c * -

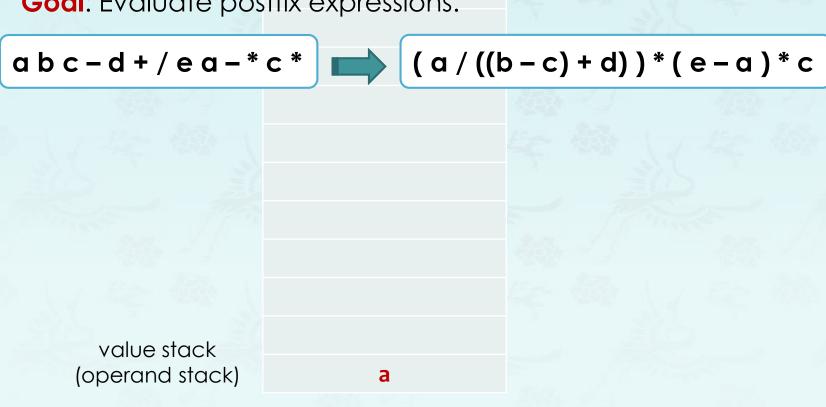
Infix and postfix notation

infix	postfix
2 + 3 * 4	2 3 4 * +
a * b + 5	a b * 5 +
	1 2 + 7 *
	a b * c /
	a b c - d + / e a - * c *
	a b / c - d e * + a c * -

Infix and postfix notation

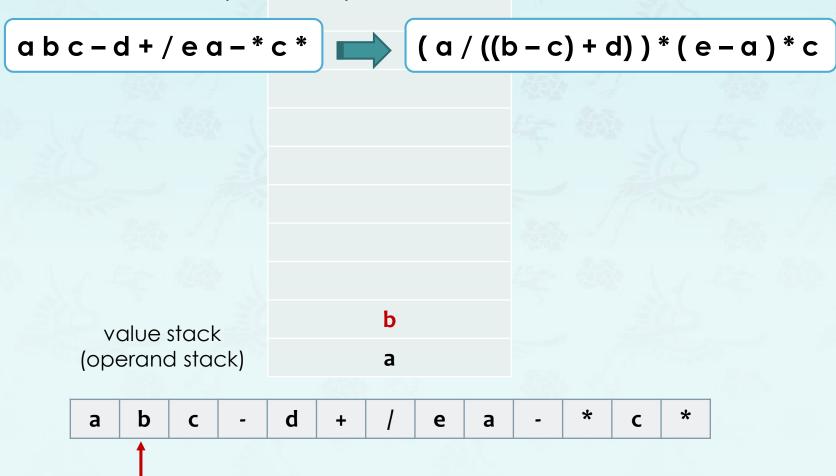


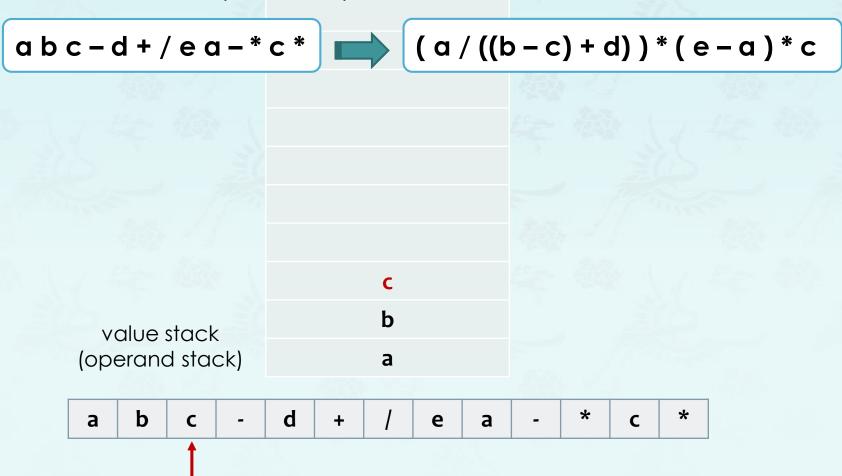
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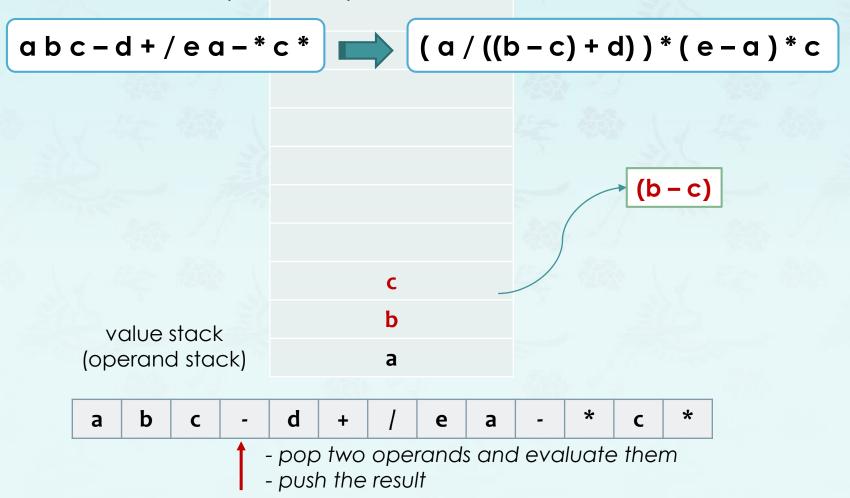


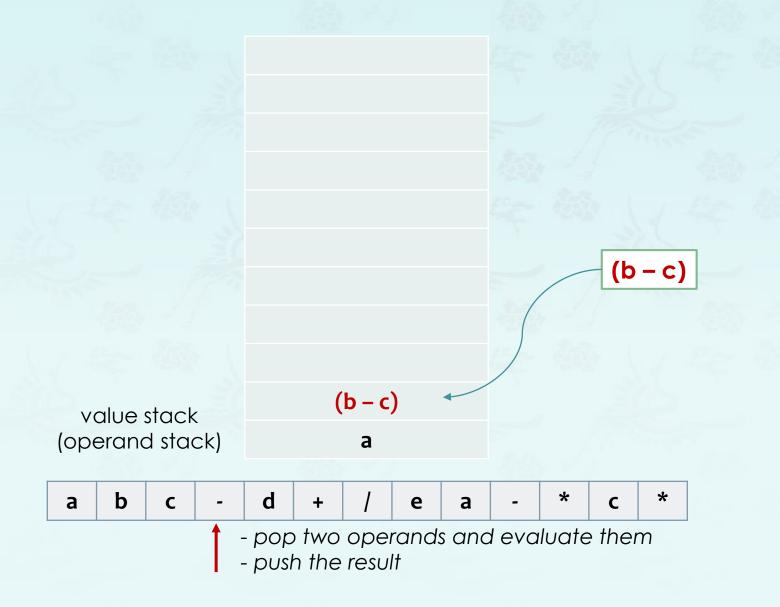
\* b d C e C a a

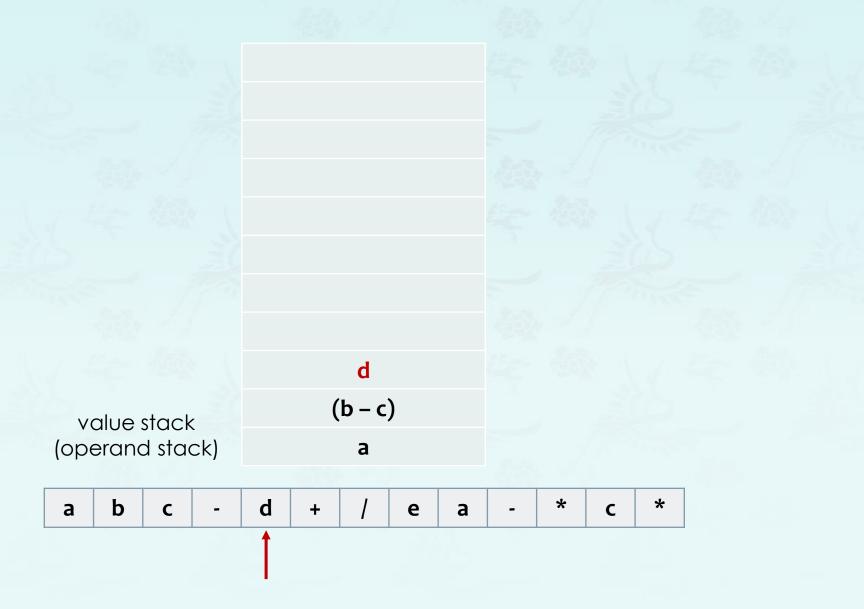
push the operands until an operator comes up.

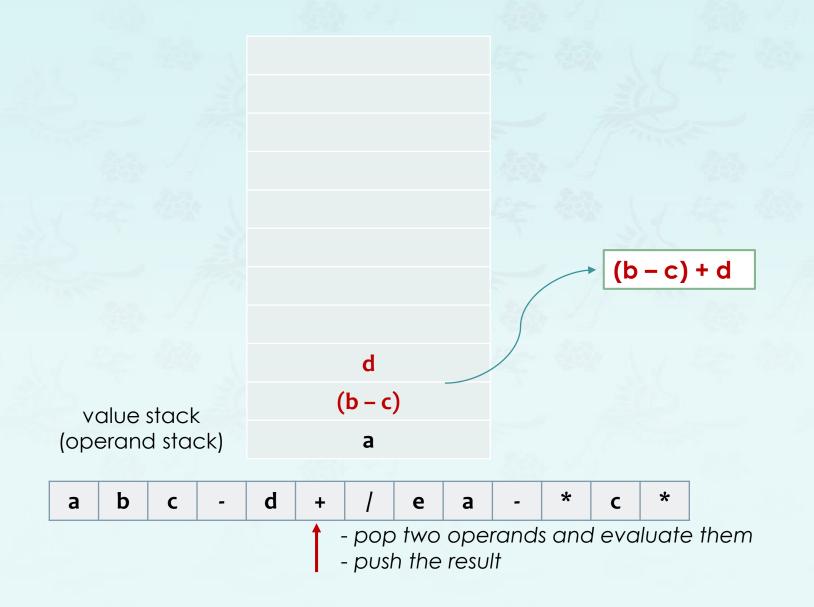


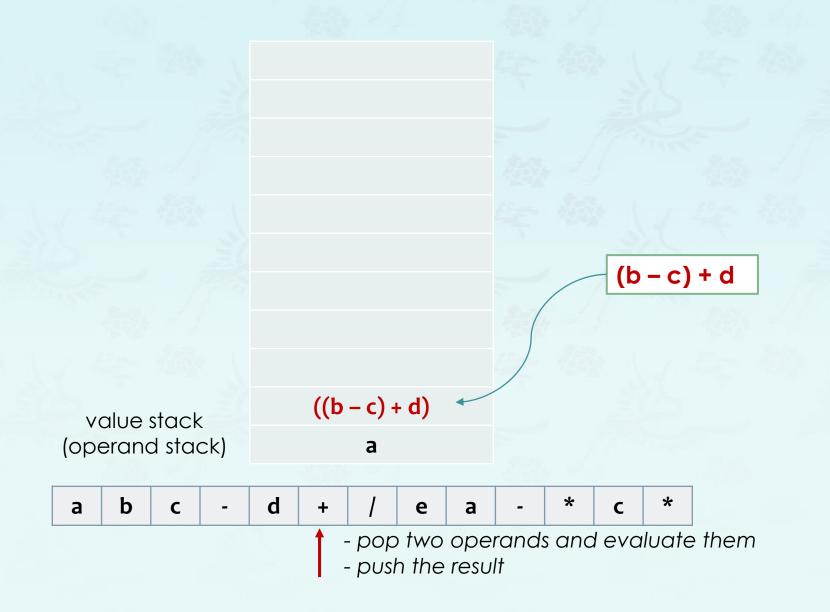


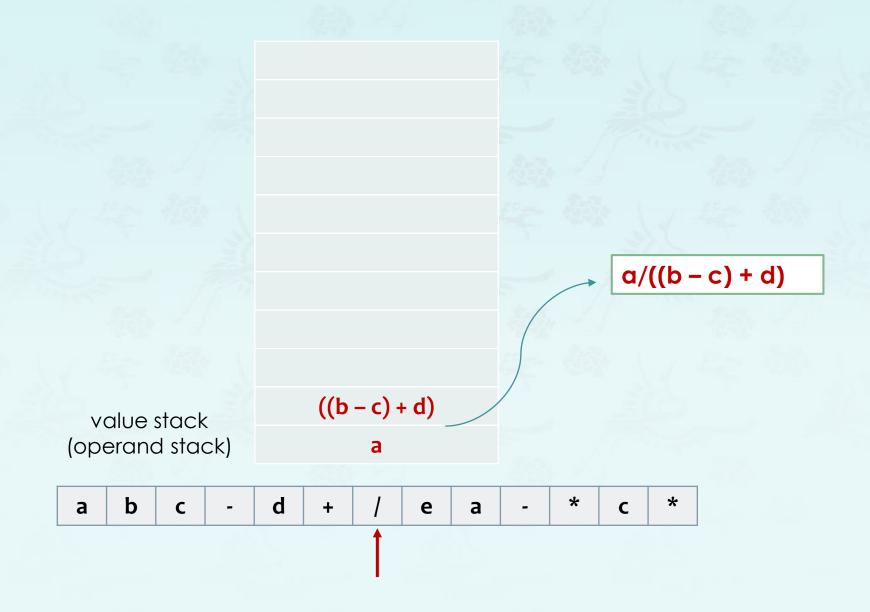


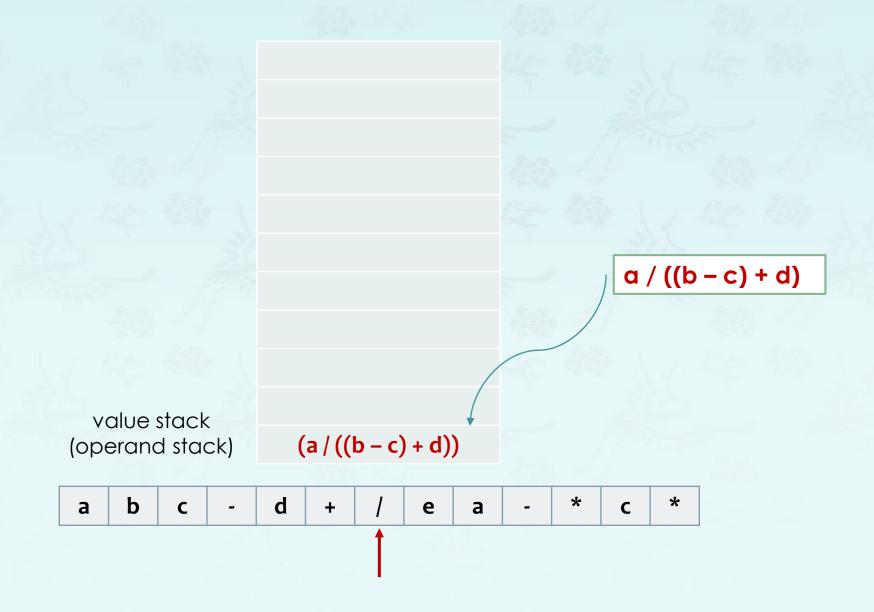


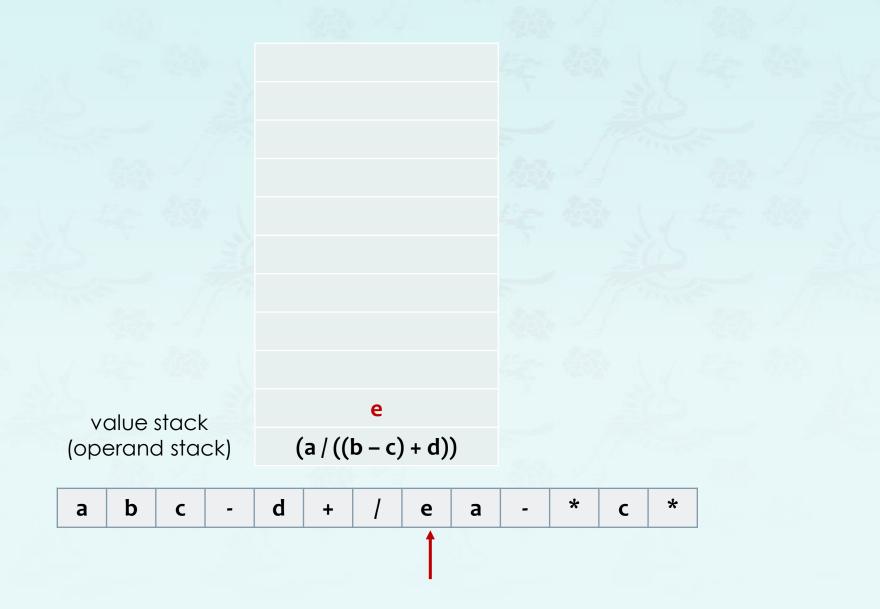


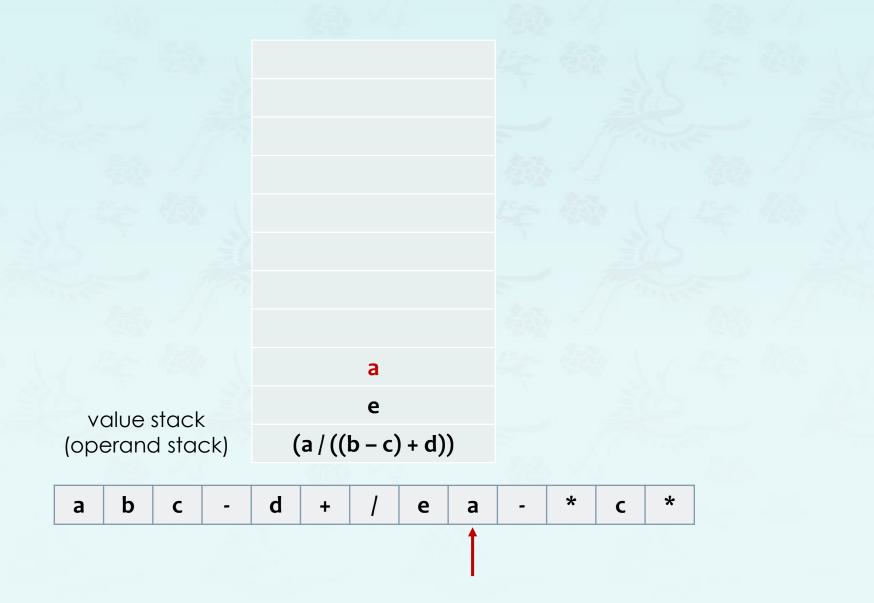


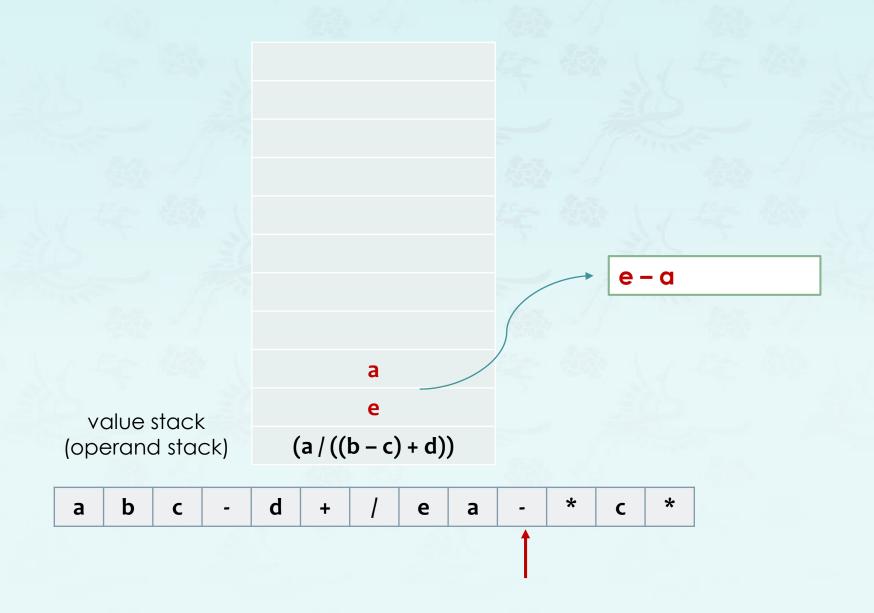


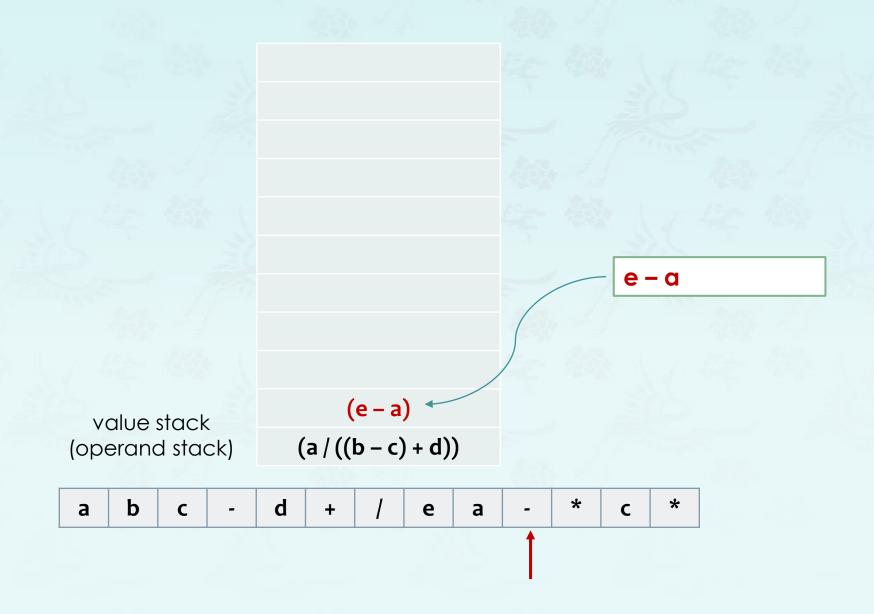


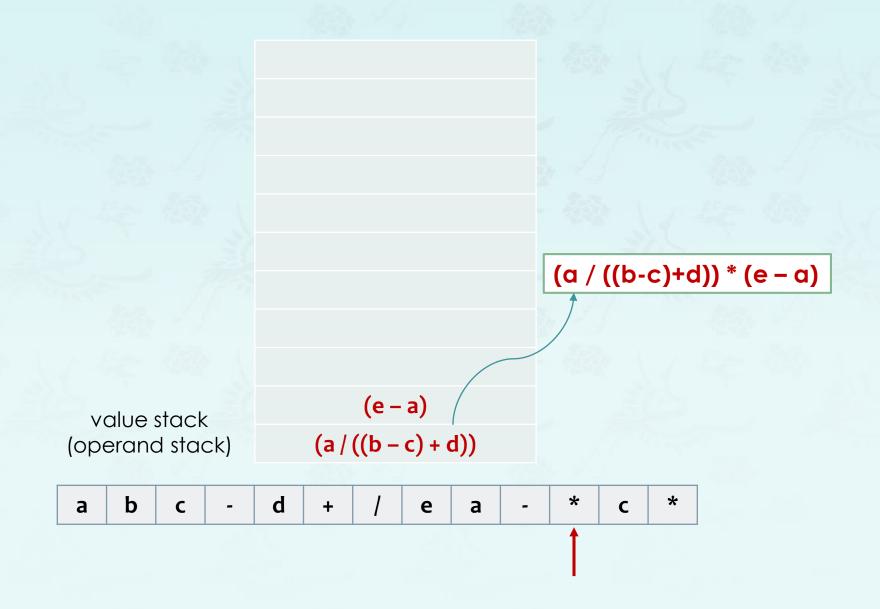


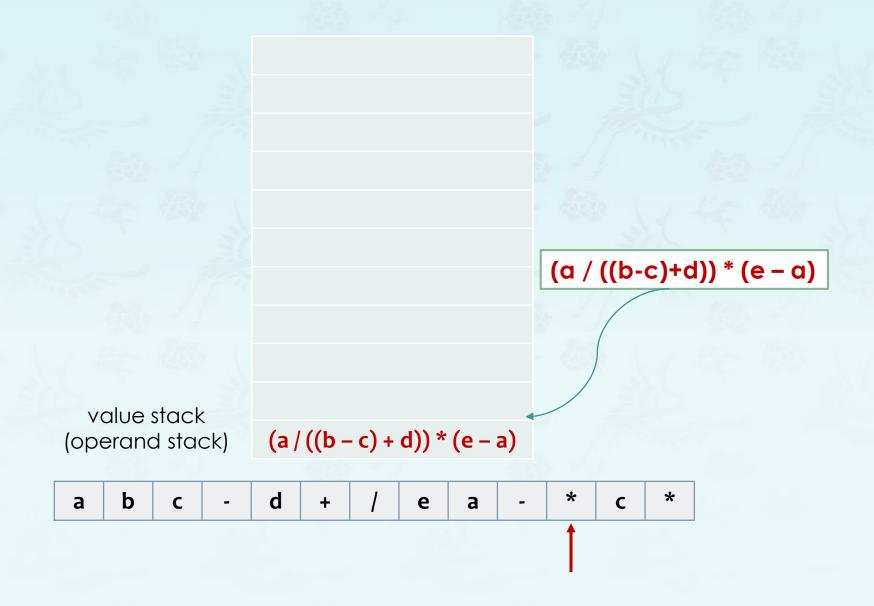


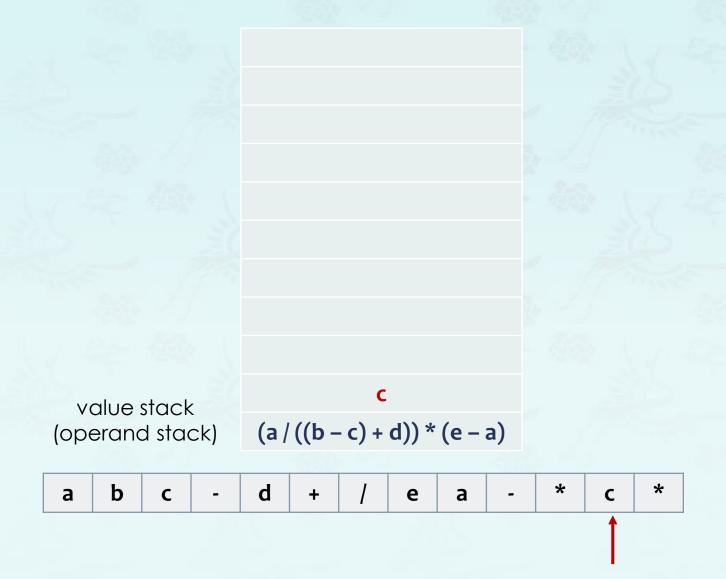


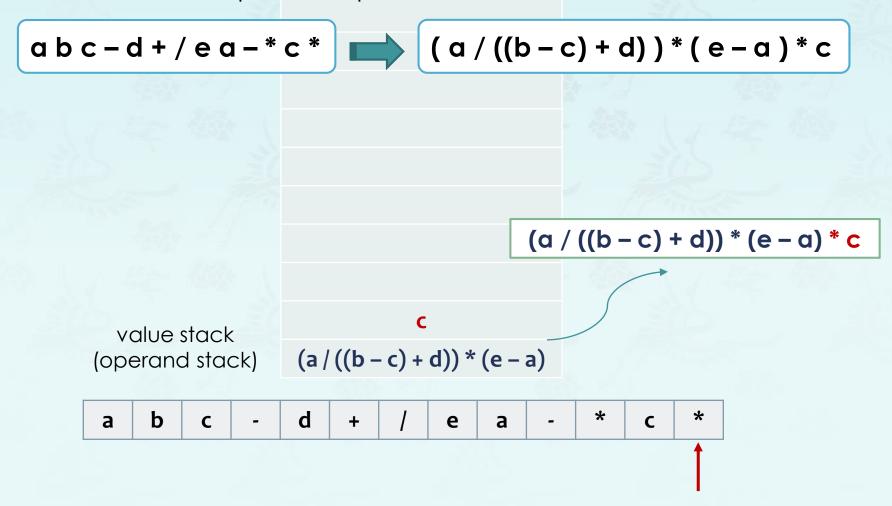












#### Arithmetic expression evaluation - Infix to Postfix Conversion

We use a stack.

- 1. When an **operand** is read, output it.
- 2. When an operator is read,
  - Pop until the top of the stack has an element of lower precedence.
  - Then push it.
- 3. When ) is found, pop until we find the matching (.
- ( has the lowest precedence when in the stack but has the highest precedence when in the input.
- 5. When we reach the end of input, pop until the stack is empty.

# Arithmetic expression evaluation - Infix to Postfix Conversion Example 1:

infix: 3+4\*5/6

in	stack(bottom to top)	postfix
3		3
+	+	
4		3 4
*	+ *	
5		3 4 5
/	+ /	3 4 5 *
6		
	3 4 5 *6/+	

# Arithmetic expression evaluation - Infix to Postfix Conversion Example 2:

infix: (1+3)\*(4-2)/(5+7)

in	stack (bottom to top)	postfix	in	stack	postfix
(	(		(	/ (	1 3 + 4 2 - *
1		1	5		1 3 + 4 2 - * 5
+	( +		+	/ ( +	
3		1 3	7		1 3 + 4 2 - * 5 7
)		13+	)		13+42-*57+
*	*				13+42-*57+/
(	* (				
4		1 3 + 4			
-	* ( -			- Operan	ds are output immediately
2		1 3 + 4 2		<ul> <li>Stack operators until right parens</li> <li>Unstack until left parens</li> <li>Delete left parens</li> </ul>	
)	*	1 3 + 4 2 -			
/	/	1 3 + 4 2 - *			ral, higher precedence opera output before lower one.)

### Arithmetic expression evaluation - Infix to Postfix Conversion

#### Example 3:

infix: 
$$a - (b + c * d) / e$$

in	stack(bottom to top)	postfix
а		
-		
(		
b		
+		
С		
*		
d		
)		
/		
е		

- Operands are output immediately
- Stack operators until right parens
- Unstack until left parensDelete left parens

In general, higher precedence operator must be output before lower one.)

# Arithmetic expression evaluation - Infix to Postfix Conversion Example 3:

infix: 
$$A * (B + C * D) + E$$

	in	stack(bottom to top)	postfix
1	Α		
2	*		
3	(		
4	В		
5	+		
6	С		
7	*		
8	D		
9	)		
10	+		
11	Е		
12			

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