

Marcel Solé Àvila, Razgar Ebrahimi, Georgios Tsaousoglou

# How inefficient can a market be?

# The goal

Characterize the **inefficiency** of the electricity market

$$\eta_{th} = \frac{W_{out}}{Q_{in}} < 1$$

$$\text{Price of Anarchy} = \frac{\text{Cost at equilibrium}}{\text{Minimum cost}} \geq 1$$

cost = burned fuel

**Price of Anarchy = 1**, we burn the **minimum** fuel to satisfy demand

**Price of Anarchy = 2**, we burn **twice as much** fuel to satisfy demand



# How is the power system operated?

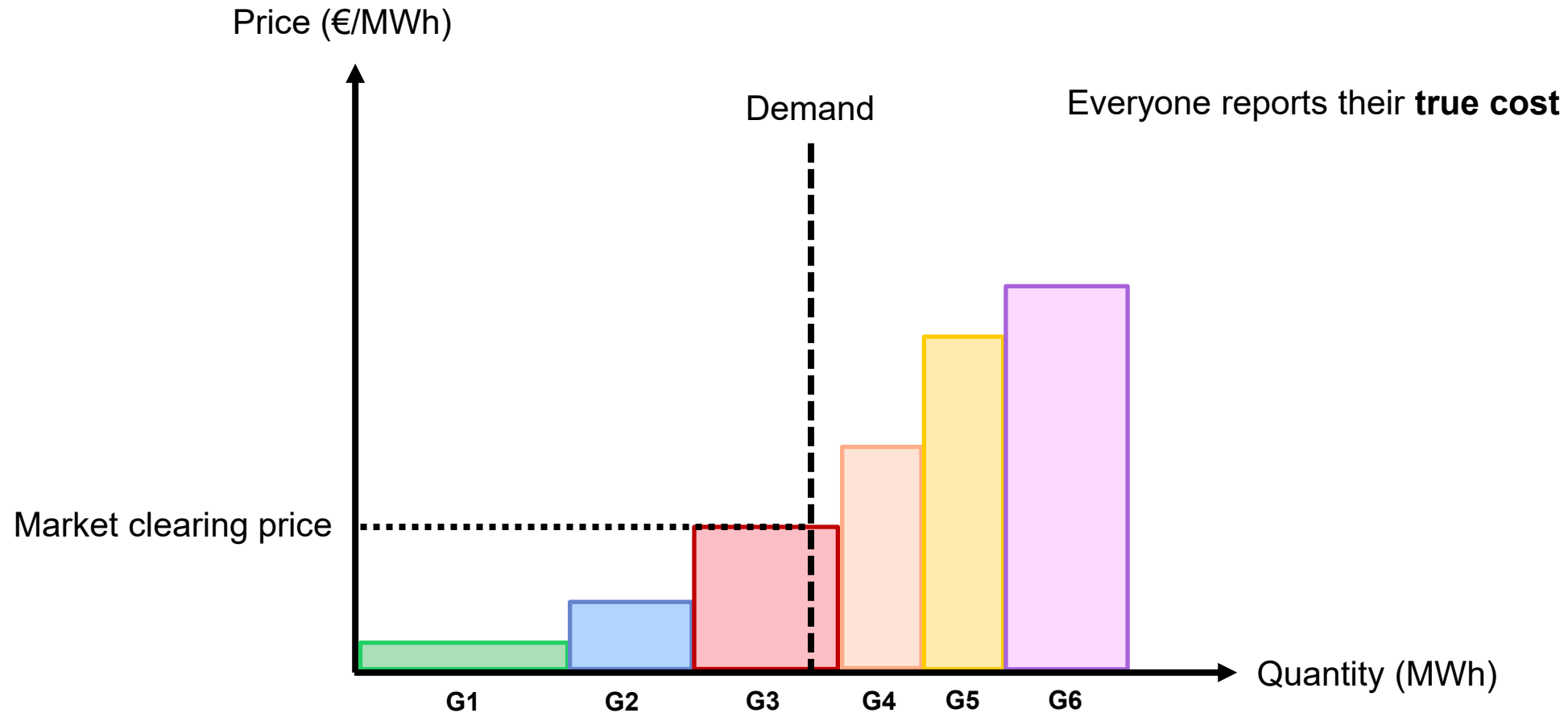


Image from: <https://www.britannica.com/>

# How is the power system operated?

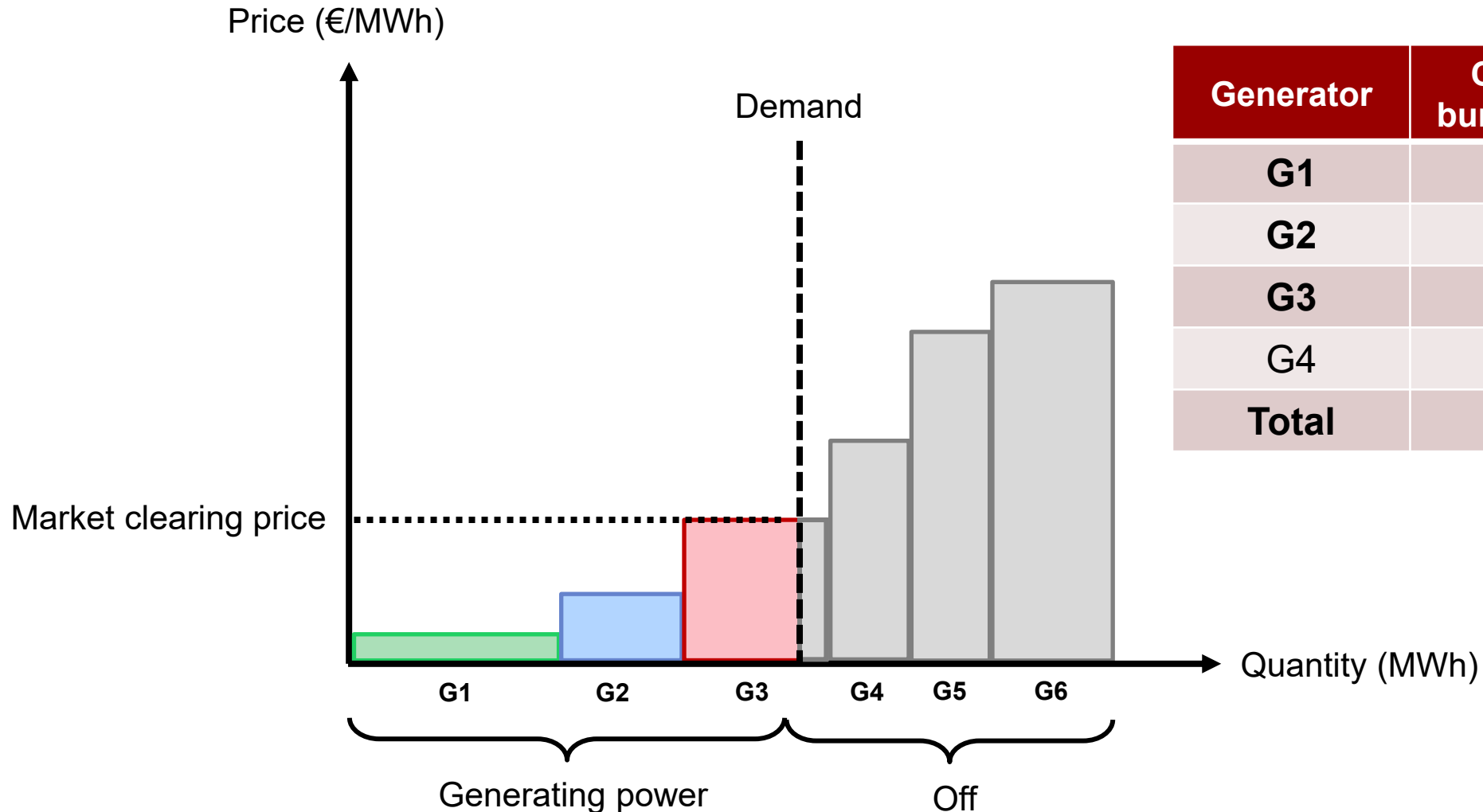


Image from: <https://www.britannica.com/>

# Strategic behaviour

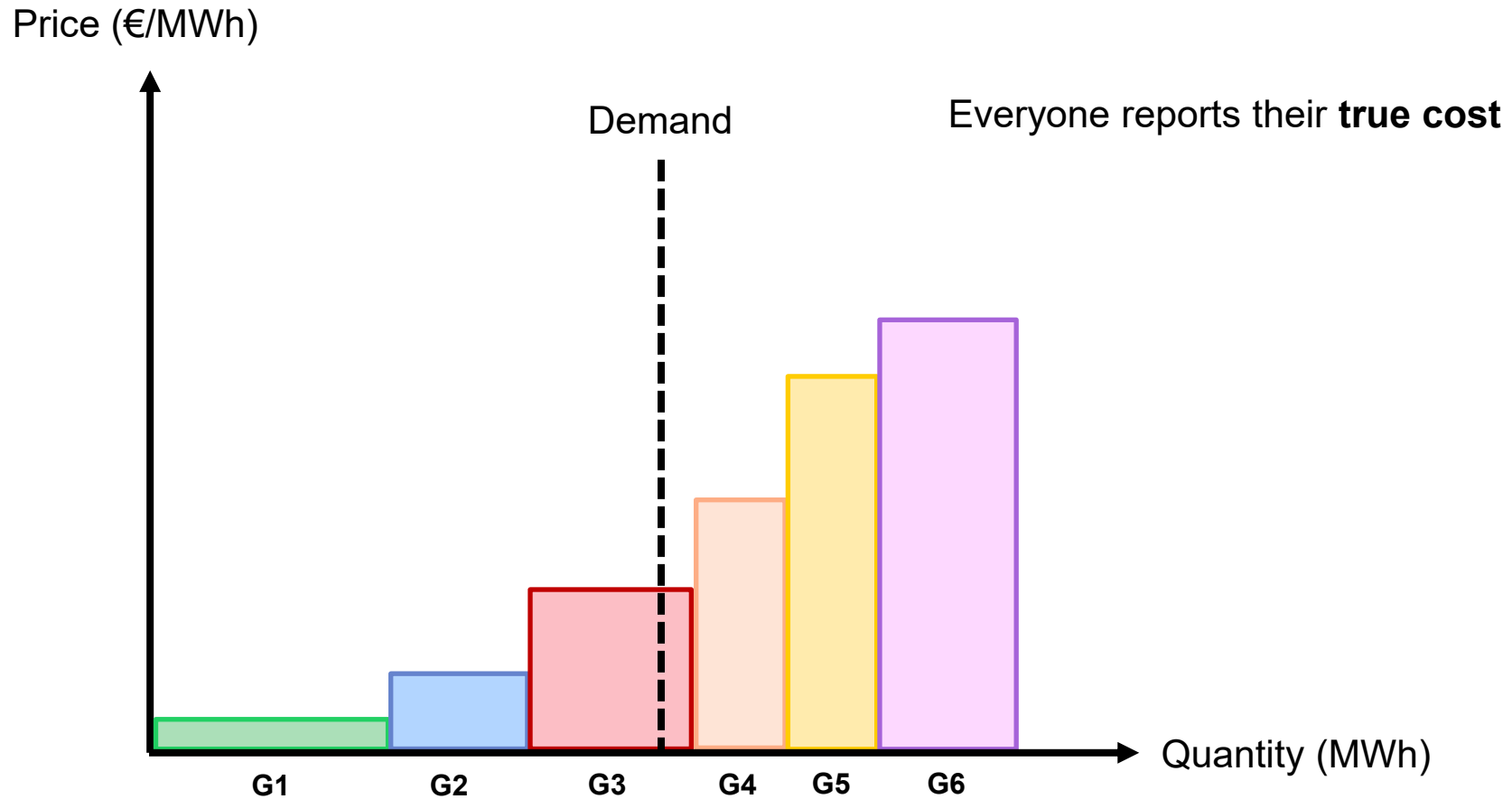
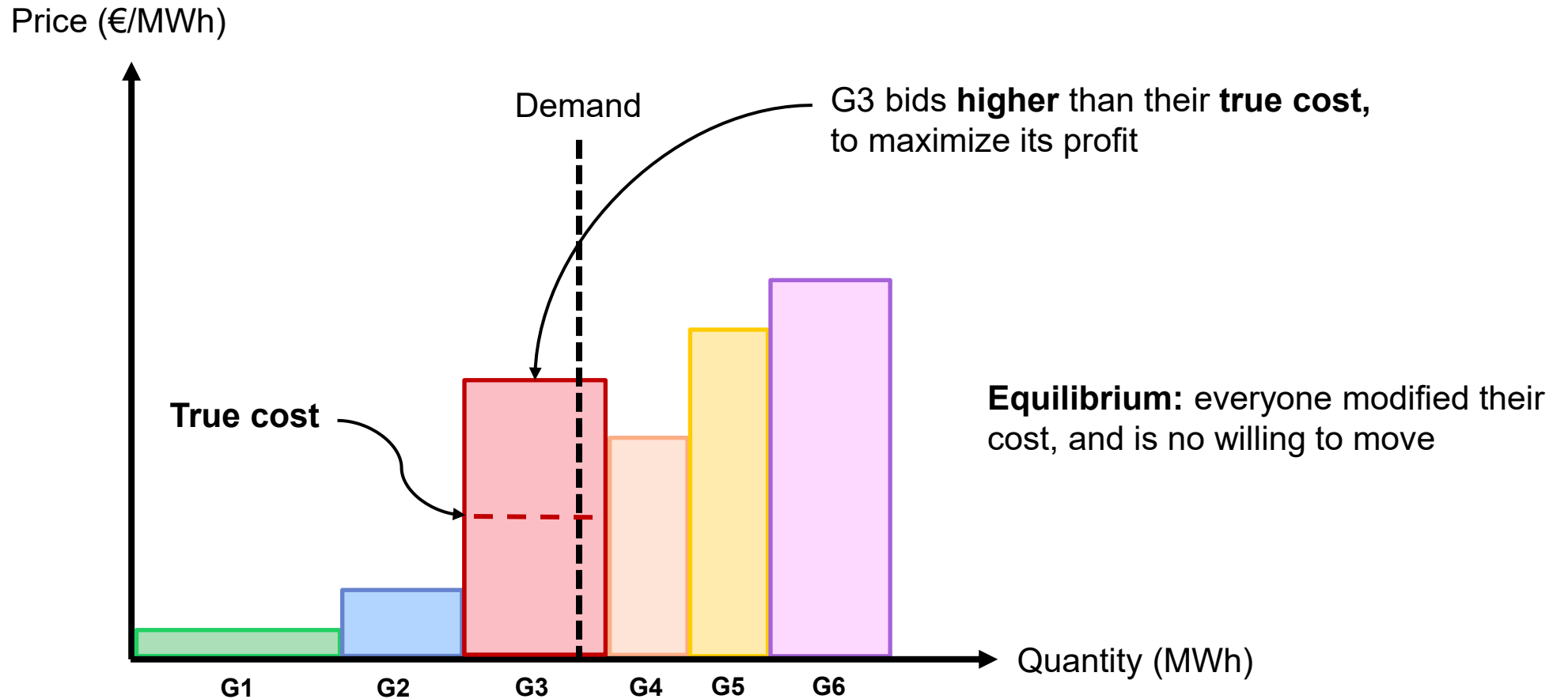


Image from: <https://www.britannica.com/>

# Strategic behaviour



# Strategic behaviour

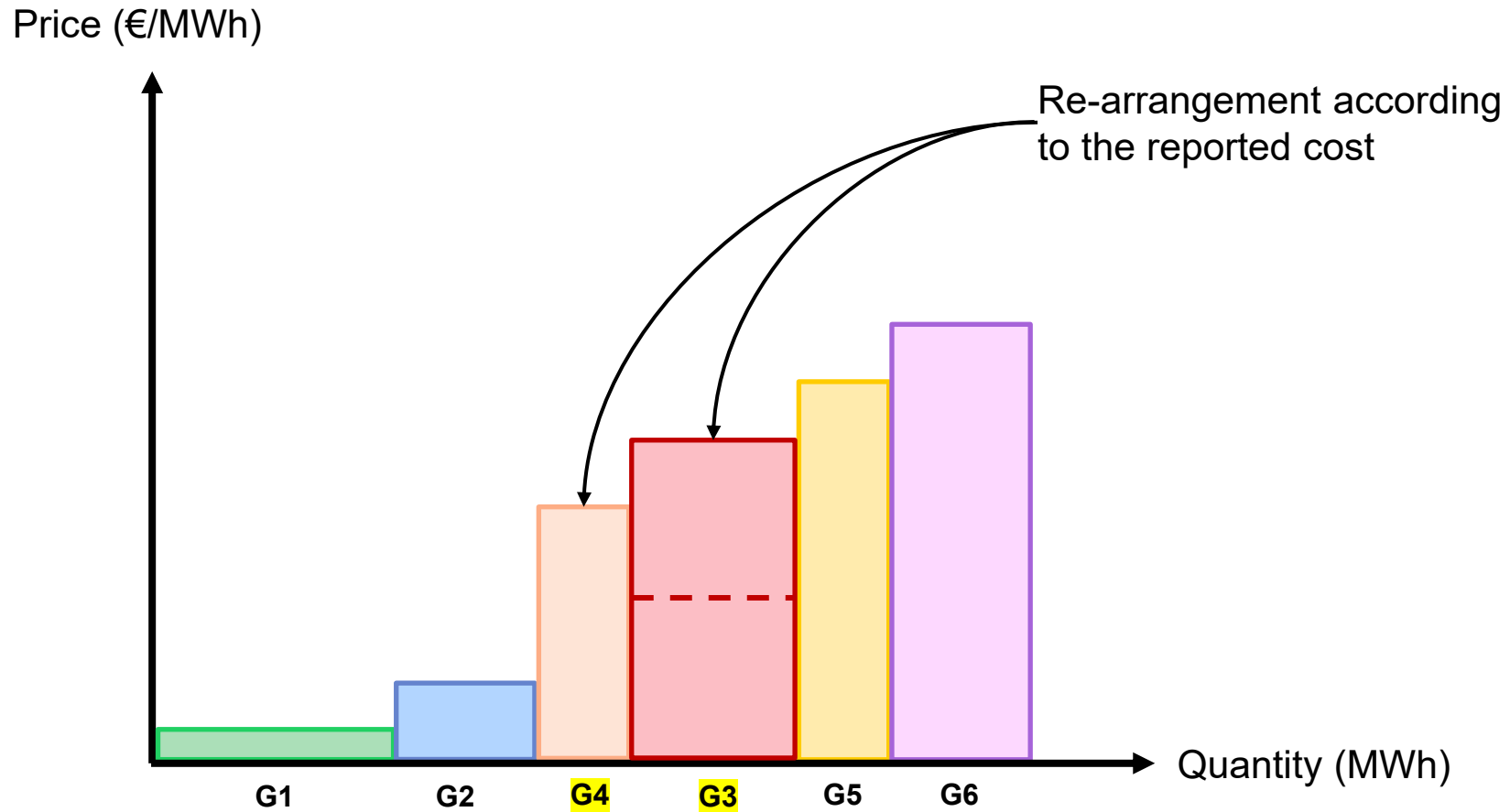
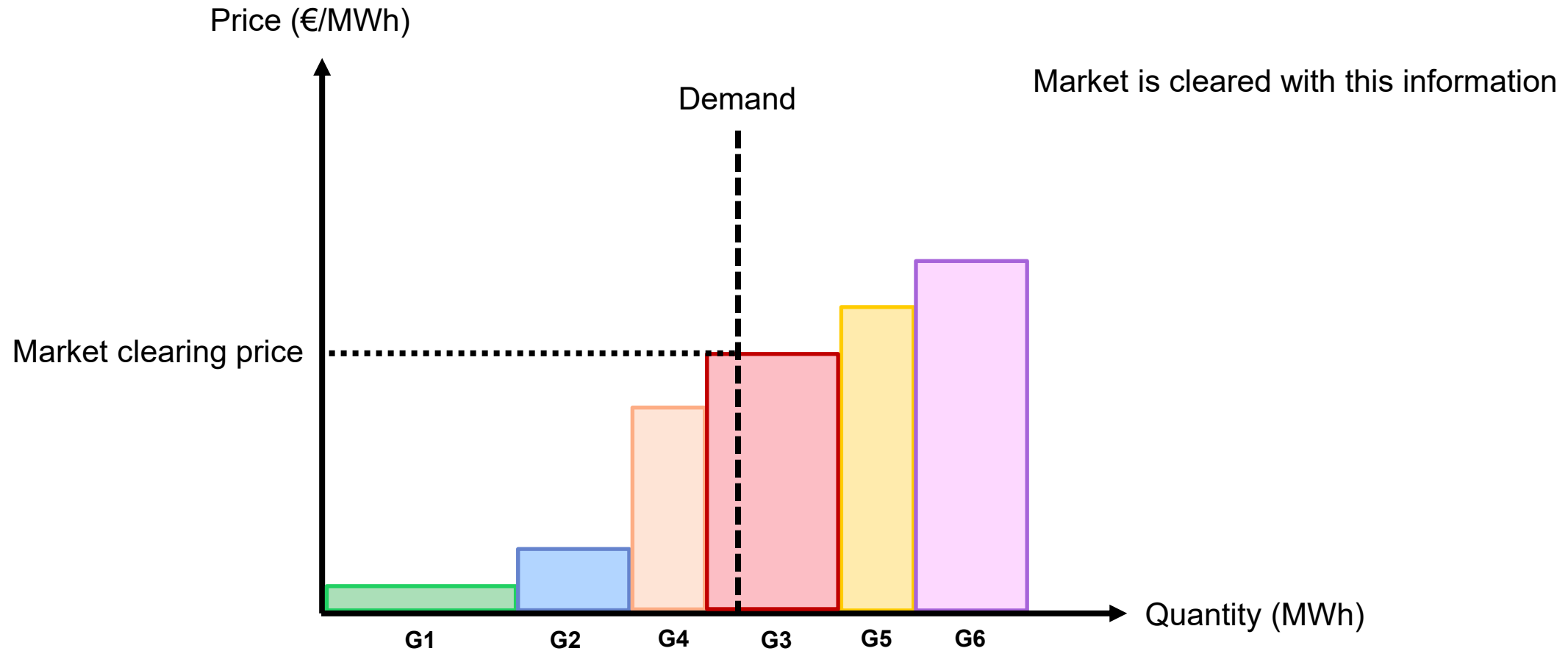


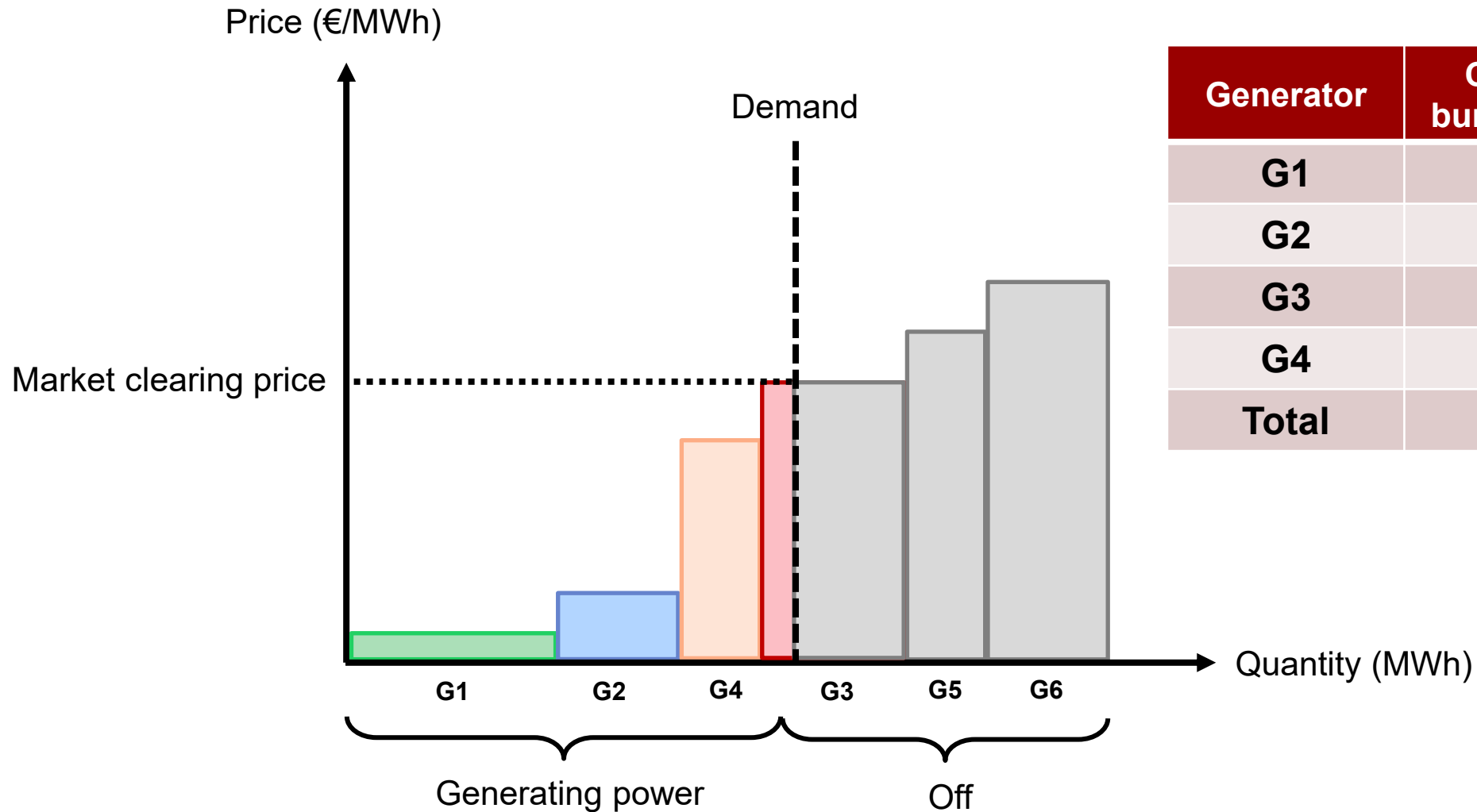
Image from: <https://www.britannica.com/>

# Strategic behaviour



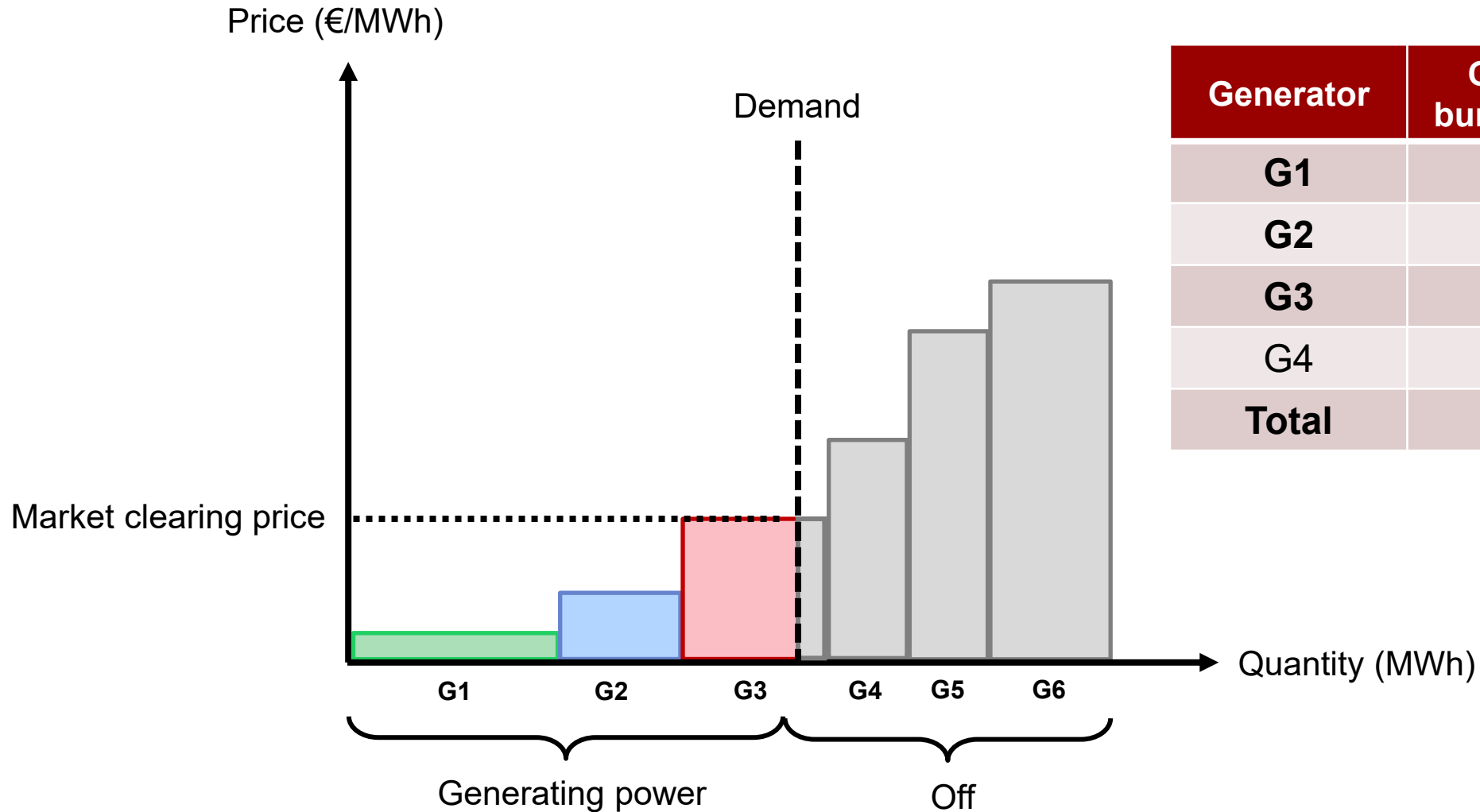


# Strategic behaviour



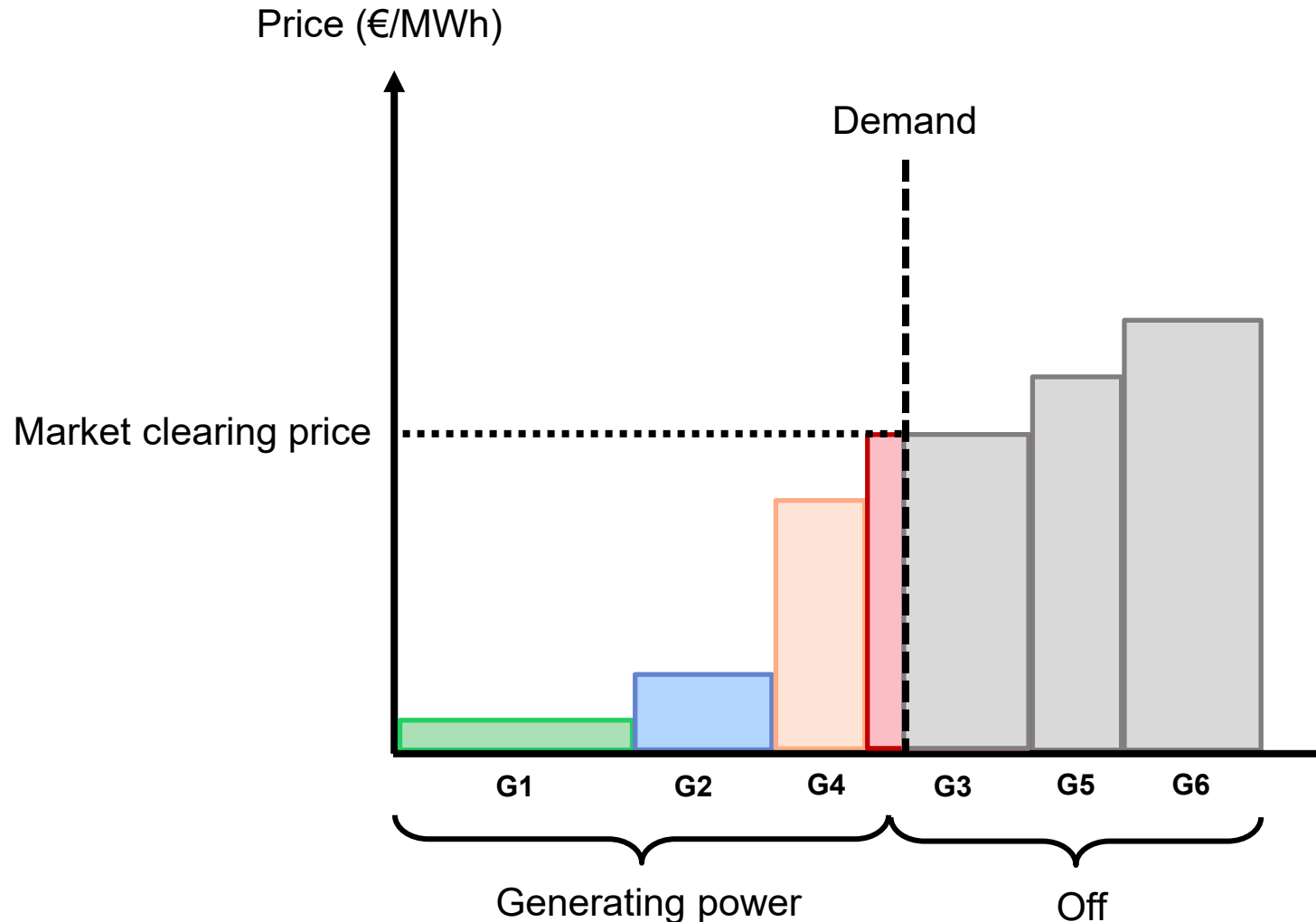
Generator	Cost of burned fuel	Generation
G1	10	100%
G2	15	100%
G3	6	20%
G4	35	100%
Total	66	-

# Original system



Generator	Cost of burned fuel	Generation
G1	10	100%
G2	15	100%
G3	24	80%
G4	0	0%
Total	50	-

# Strategic behaviour

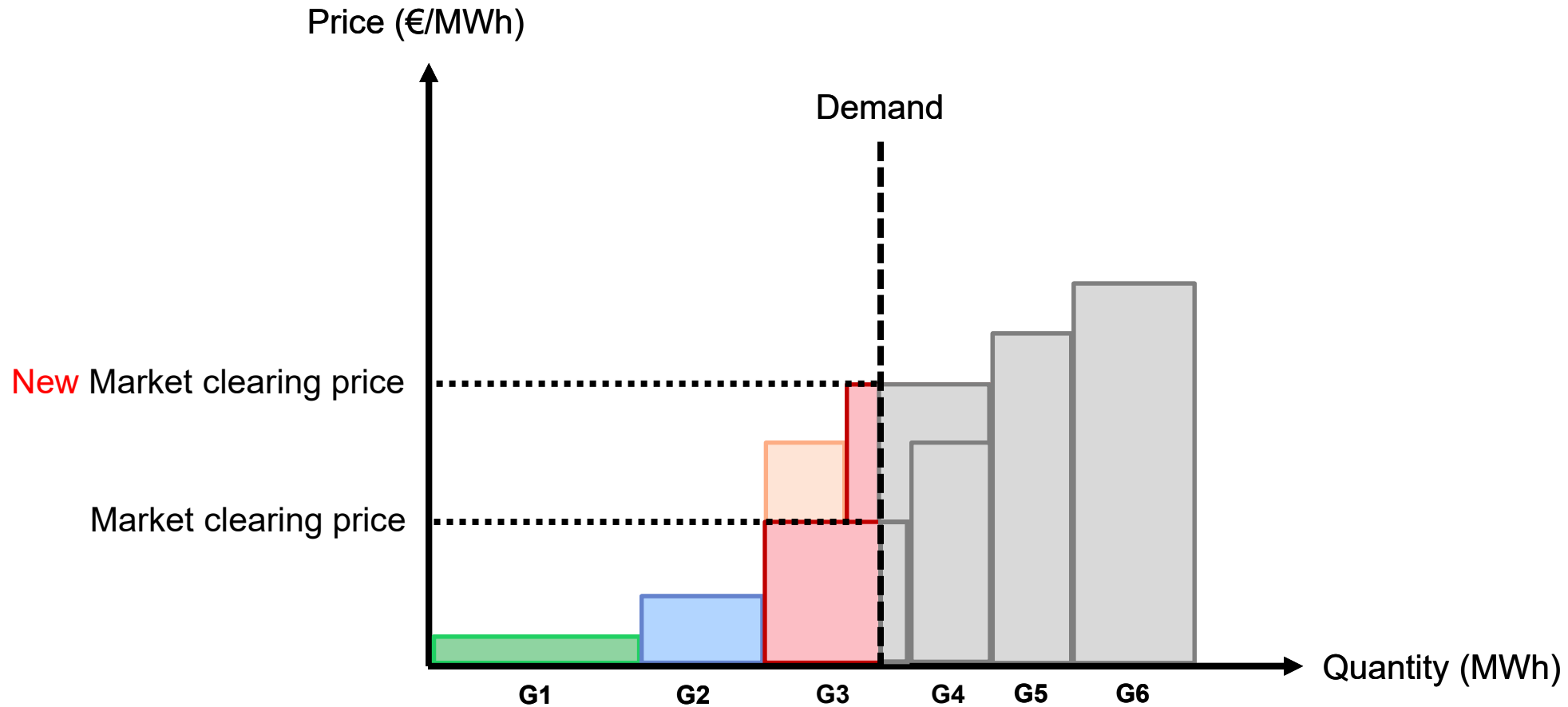


Generator	Cost of burned fuel	Generation
G1	10	100%
G2	15	100%
G3	6	20%
G4	35	100%
Total (old)	50	-
Total (new)	66	-

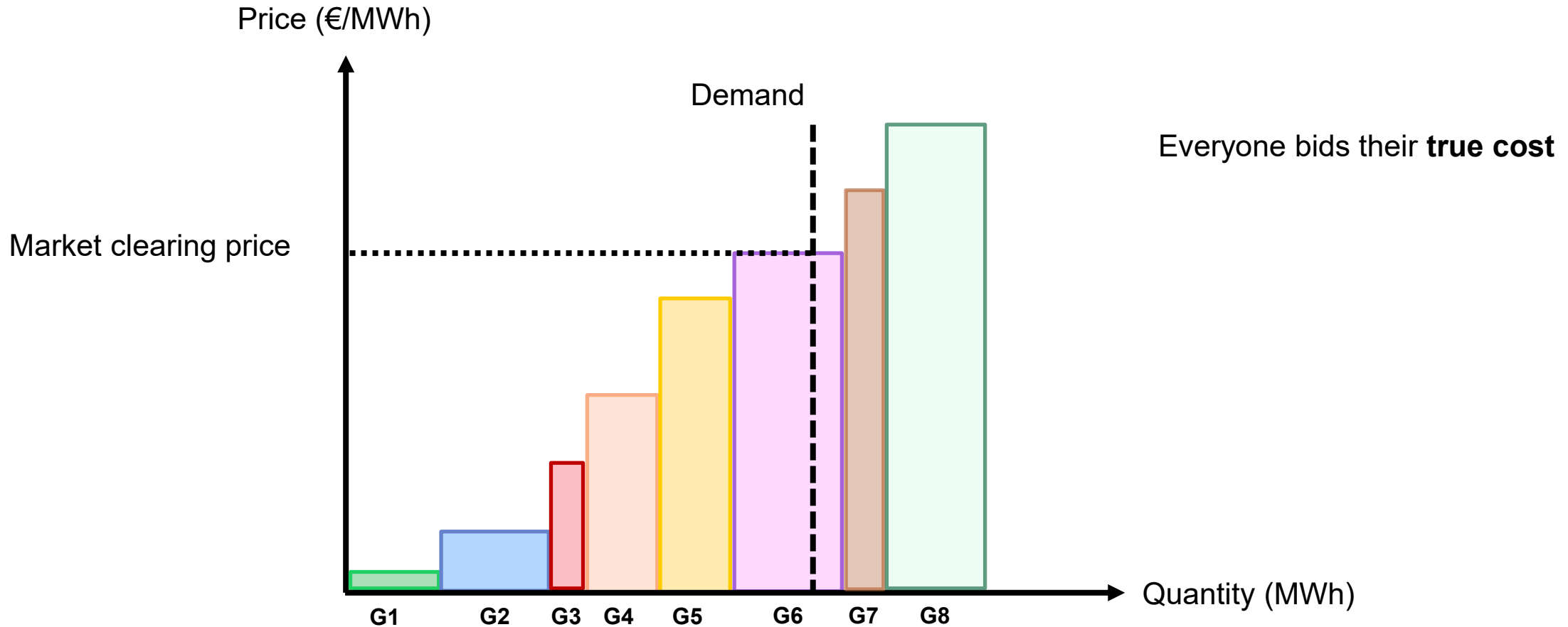
$$\text{Price of Anarchy} = \frac{66}{50} = 1.32$$

The cost is **32%** higher than needed

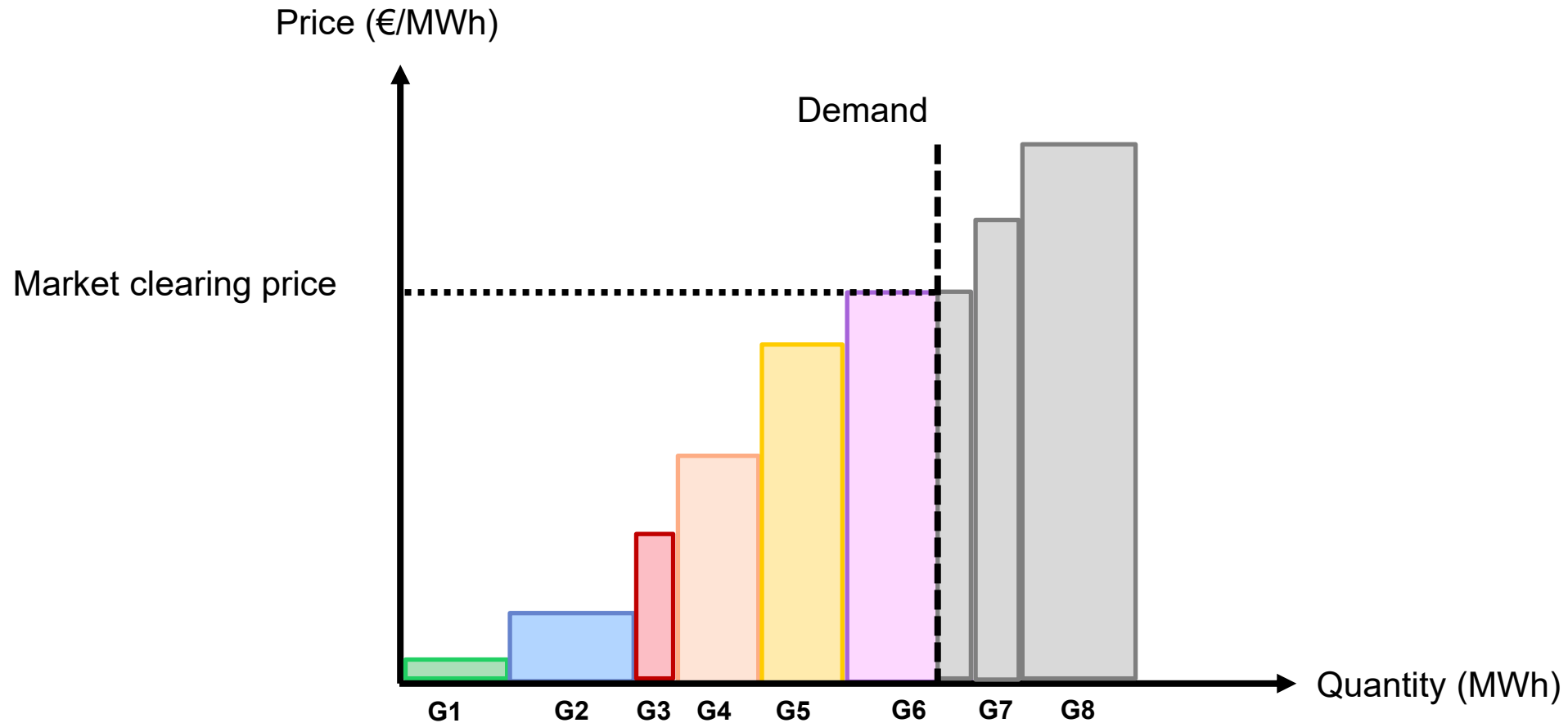
# Our first system



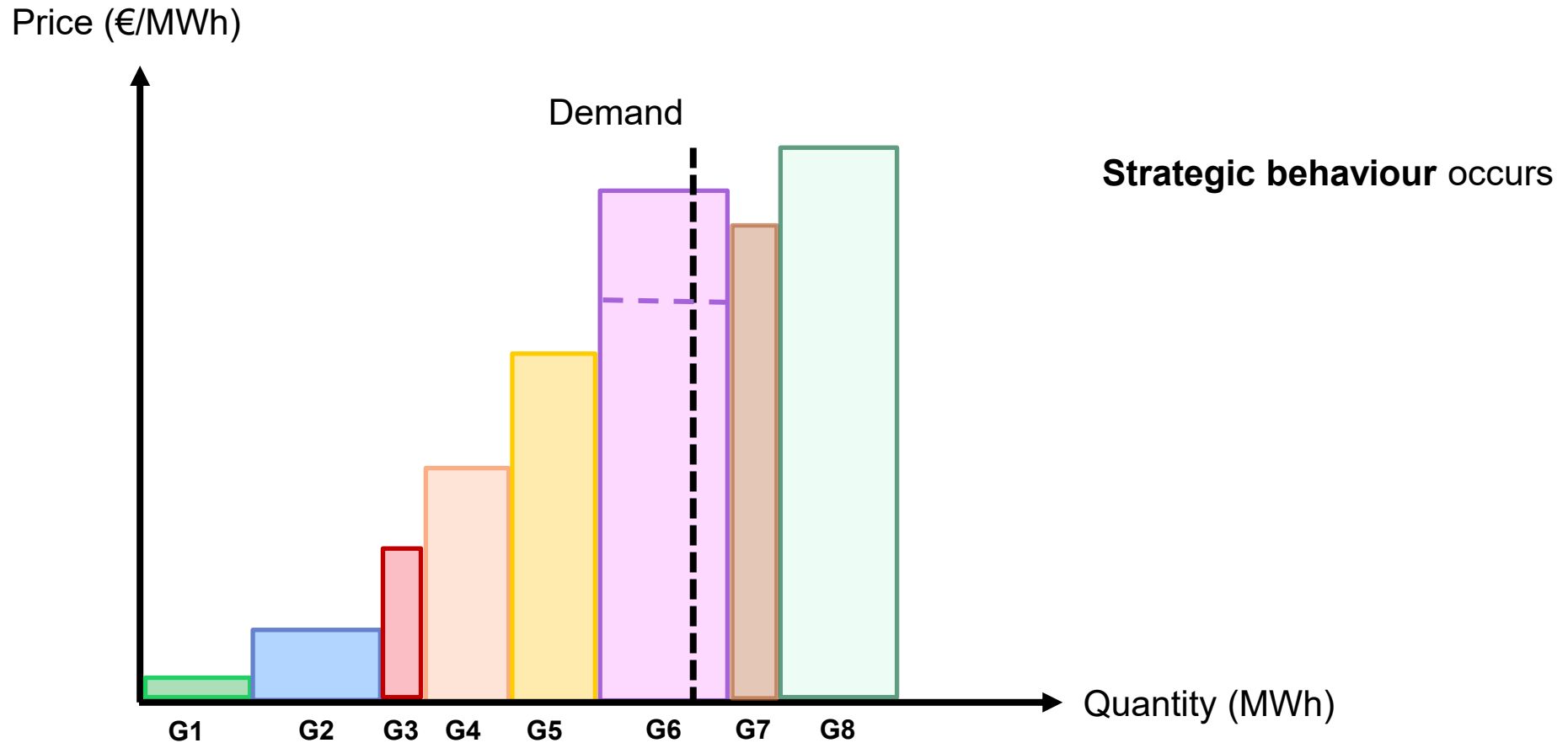
# A new system, different generators



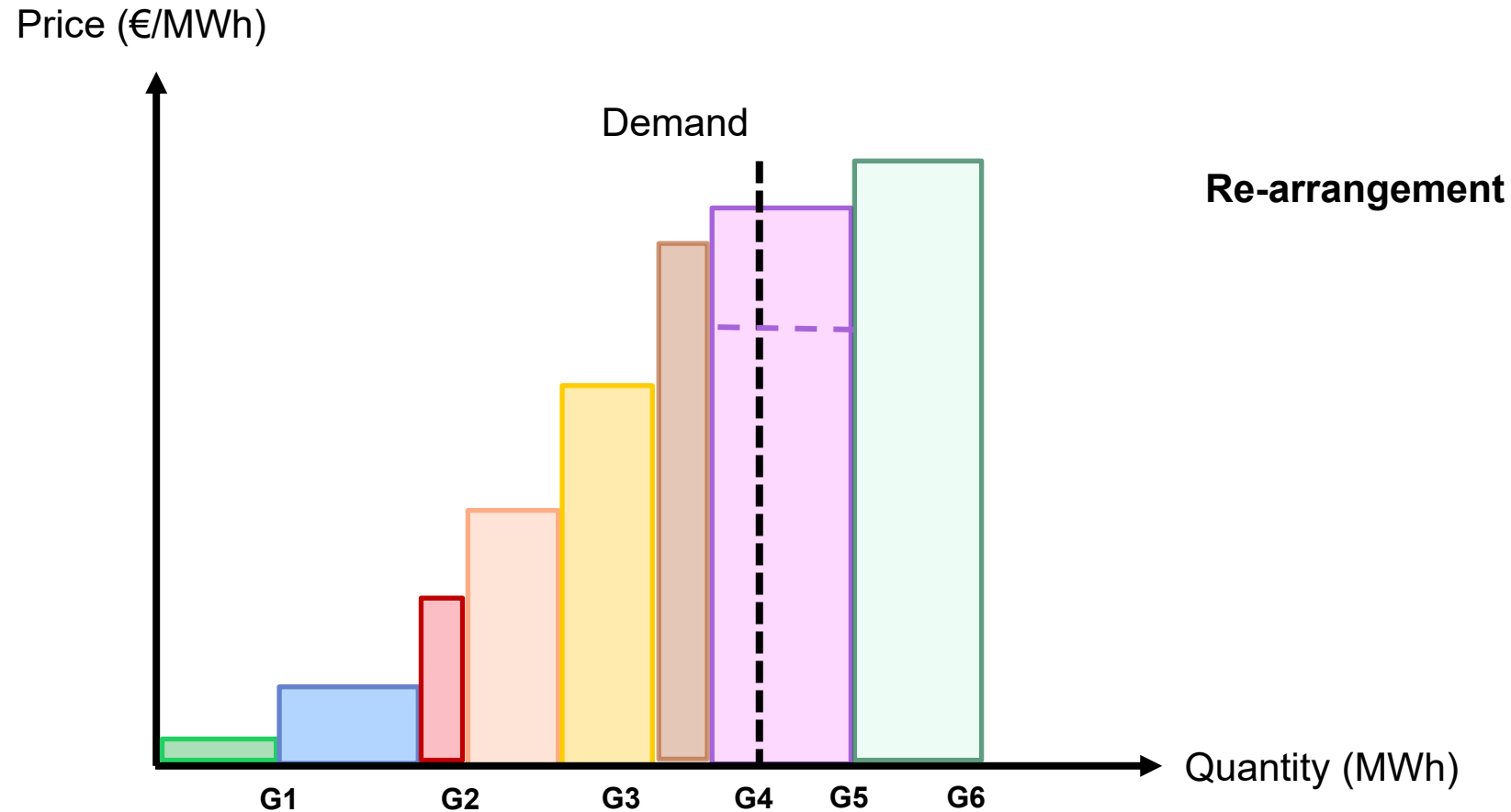
# A new system, different generators



# A new system, different generators

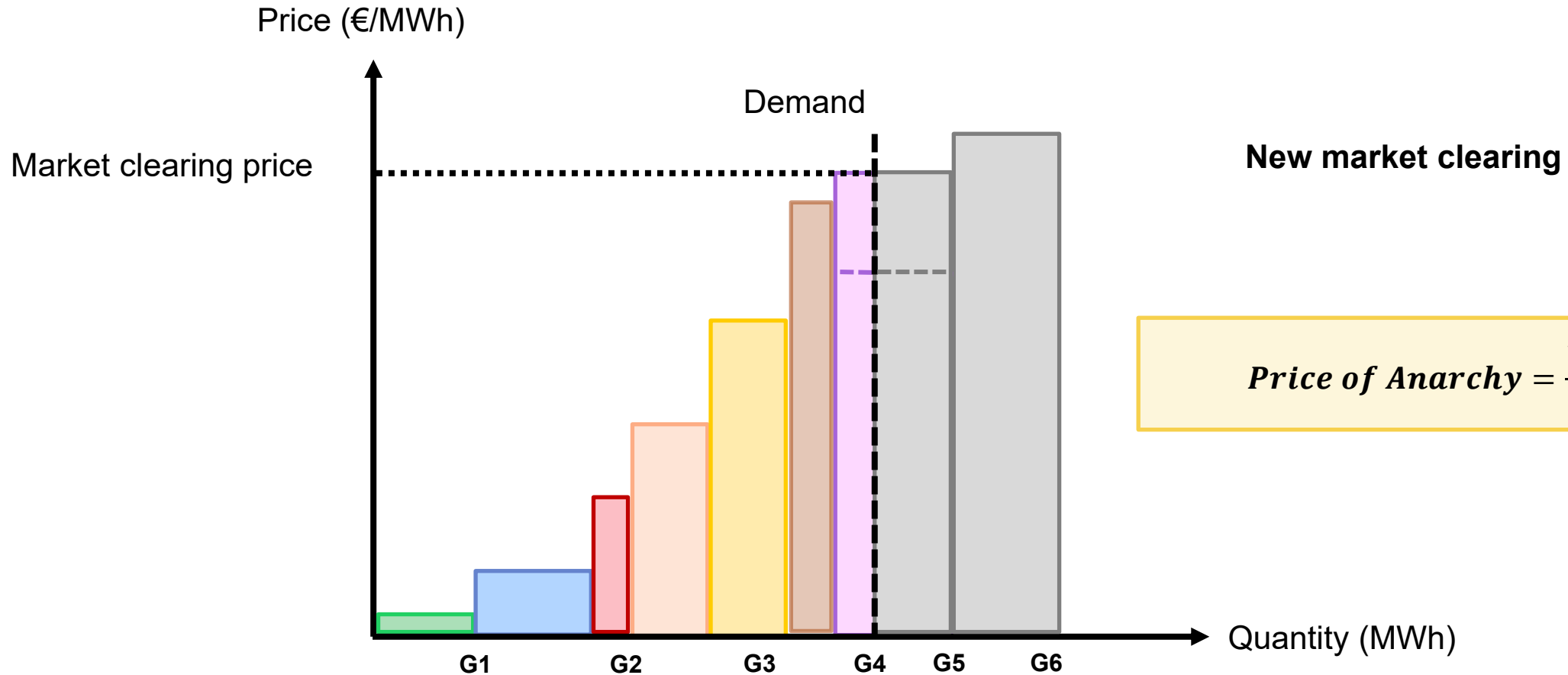


# A new system, different generators





# A new system, different generators



# The Robust PoA

**Robust PoA:** worst possible combination of generators that exists

$$\begin{aligned} & \max_{\text{cost functions}} \{ \textit{Price of Anarchy} \} \\ & \left\{ \begin{array}{l} \text{s.t. } \max_{\theta_n} \{ \sum_t [u_{n,t} \lambda_{n,t} - c_n(u_{n,t})] \} \\ \text{s.t. } u_{n,t} \in \operatorname{argmax} \{ \sum_{m,t} \hat{g}(u_{m,t}; \theta_{m,t}) \} \end{array} \right\} \forall n \end{aligned}$$

“Search in the space of possible generators to find the one that **maximizes the inefficiency**”

# Johari's upper bound

If enough assumptions are made, there exists a close form **upper bound** for the PoA.

$$\textit{Price of Anarchy} \leq 1 + \frac{1}{N-2}$$

$N$ : number of generators

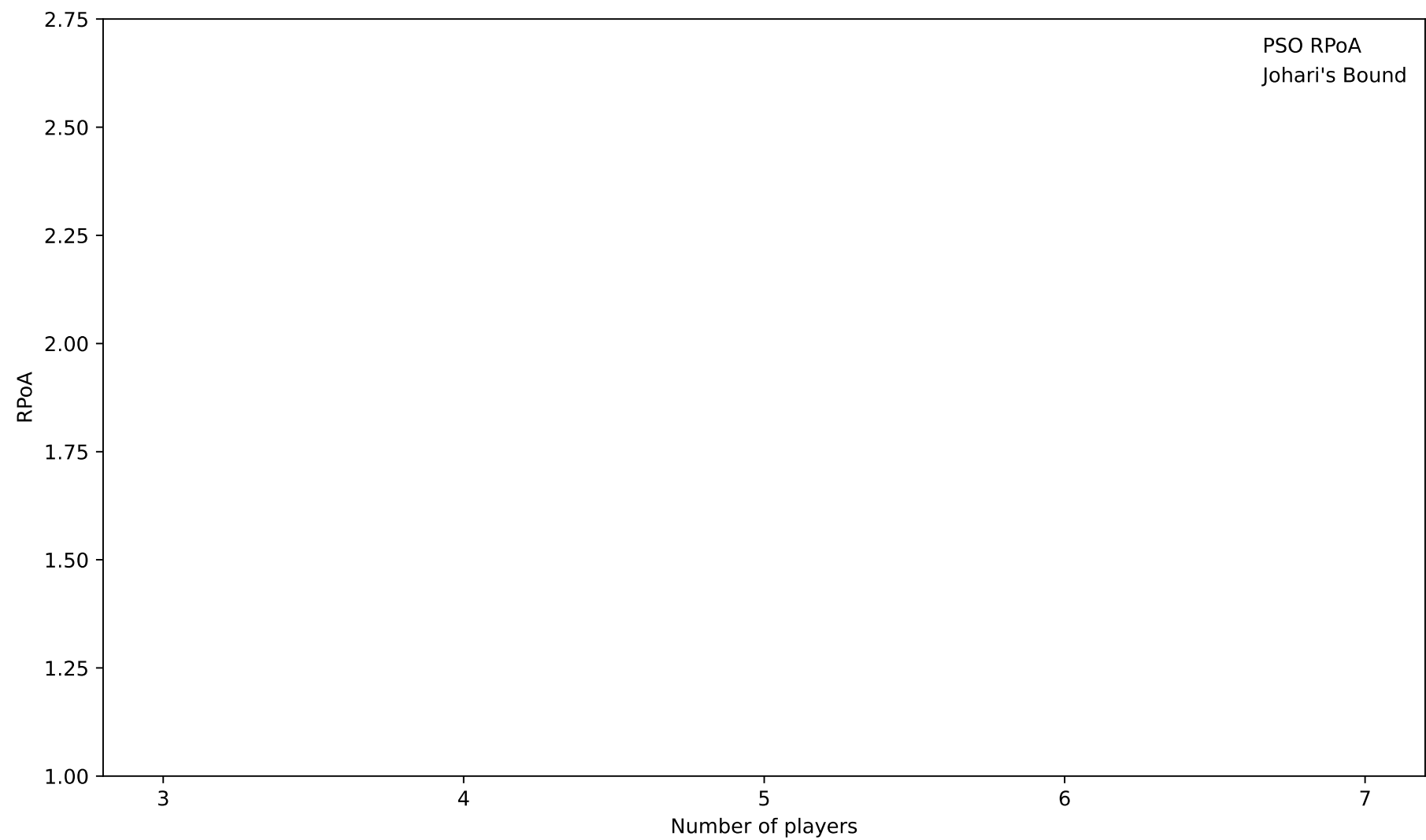
When:  $N \rightarrow \infty$ ,  $\textit{PoA} \rightarrow 1$  (perfect competition)

# Our contribution

We **drop** the assumptions to:

- Have a more realistic cost curve for generators
- Approximate better the market behavior

# Our results



**PSO: Particle Swarm Optimization**

# Exercises

# Exercise 1

1. Construct the market clearing problem (merit order curve) as a LP.
  1. Demand is inelastic.
  2. There needs to be at least 3 generators, bid = (price, quantity).
  3. Generators have a minimum and maximum power dispatch
  4. The cost of a generator depends linearly on the power dispatch.  
Assume:  $\text{Cost} = cx$ . Where  $C$  is the cost,  $c$  the per unit cost (EUR/MWh), and  $x$  the production (MWh).
  5. There are no start-up costs, no ramp constraints, no minimum up or down time. This keeps the model linear.
  6. Assume the grid to be a copper plate. I.e., there are no power lines.

## Exercise 2

1. Construct the dual of exercise 1.
2. Solve the dual.
3. Check that strong duality holds.



## Exercise 3

1. Construct the market clearing problem (merit order curve) as a LP.
  1. Demand bids into the market, bid = (price, quantity).
  2. At least 3 consumers.
3. There needs to be at least 3 generators, bid = (price, quantity).
4. Generators have a minimum and maximum power dispatch
5. The cost of a generator depends linearly on the power dispatch.  
Assume:  $\text{Cost} = cx$ . Where  $C$  is the cost,  $c$  the per unit cost (EUR/MWh), and  $x$  the production (MWh).
6. There are no start-up costs, no ramp constraints, no minimum up or down time. This keeps the model linear.
7. Assume the grid to be a copper plate. I.e., there are no power lines.

## Exercise 4

1. Take the model of exercise 3 and add a model for the power network. Use a linear approximation, e.g., DC-OPF.
2. Keep the amount of network nodes low (e.g., 3 nodes).
3. Populate the grid with enough generators and demands so the market can be cleared (i.e., willingness to pay  $\geq$  willingness to produce).
4. Report the local marginal price of each bus, i.e., the dual variable of the power balance of each bus.

## Primal (dummy LP with mixed constraints)

$$\begin{aligned}
 &\text{minimize } 2x_1 + x_2 \\
 &\text{subject to } x_1 + x_2 \geq 1 && \text{(inequality ')} \\
 &\quad -x_1 + 2x_2 \leq 3 && \text{(inequality ')} \\
 &\quad 2x_1 - x_2 = 0 && \text{(equality)} \\
 &\quad [2pt]x_1 \geq 0, \quad x_2 \text{ free.}
 \end{aligned}$$

We'll convert everything to the Lagrangian-friendly " $\leq 0$ " form and add a multiplier for **each** such constraint (including the variable bound for  $x_1$ ).

- $x_1 + x_2 \geq 1 \iff 1 - x_1 - x_2 \leq 0 \rightarrow$  multiplier  $y_1 \geq 0$ .
- $-x_1 + 2x_2 \leq 3 \rightarrow$  multiplier  $y_2 \geq 0$ .
- $2x_1 - x_2 = 0 \rightarrow$  multiplier  $y_3$  **free** (can be any real).
- $x_1 \geq 0 \iff -x_1 \leq 0 \rightarrow$  multiplier  $s_1 \geq 0$ .
- $x_2$  is free  $\rightarrow$  **no bound constraint**, hence no  $s_2$ .

## Lagrangian

$$\begin{aligned}\mathcal{L}(x_1, x_2, y_1, y_2, y_3, s_1) &= (2x_1 + x_2) \cdot y_1 (1 - x_1 - x_2) \cdot y_2 (-x_1 + 2x_2 - 3) \cdot y_3 (2x_1 - x_2) \cdot s_1 (-x_1) \\ &= \underbrace{(y_1 - 3y_2)}_{\text{constant}} \cdot \underbrace{(2 - y_1 - y_2 + 2y_3 - s_1)}_{\text{coeff of } x_1} x_1 \cdot \underbrace{(1 - y_1 + 2y_2 - y_3)}_{\text{coeff of } x_2} x_2.\end{aligned}$$

**Dual function:**  $g(y, s) = \inf_{x_1, x_2} \mathcal{L}$

Because we dualized all constraints, the infimum is over **free**  $x_1, x_2$ . For the infimum to be finite (not  $-\infty$ ), the coefficients of  $x_1$  and  $x_2$  must vanish (otherwise we could drive them to  $\pm\infty$ ):

$$(\text{stationarity w.r.t. } x_1): \quad 2 - y_1 - y_2 + 2y_3 - s_1 = 0, \quad s_1 \geq 0,$$

$$(\text{stationarity w.r.t. } x_2): \quad 1 - y_1 + 2y_2 - y_3 = 0 \quad (\text{no } s_2 \text{ since } x_2 \text{ is free}).$$

When these hold, the  $x$ -terms drop out and

$$g(y, s) = y_1 - 3y_2.$$

# Duality

## Dual problem

Maximize the dual function subject to the multiplier sign rules and the stationarity equations:

- Variables and signs:

$$y_1 \geq 0, \quad y_2 \geq 0, \quad y_3 \in \mathbb{R} \text{ (free)}, \quad s_1 \geq 0.$$

- Equalities/inequalities from stationarity:

$$1 - y_1 + 2y_2 - y_3 = 0 \quad \Longleftrightarrow \quad y_1 - 2y_2 + y_3 = 1,$$

$$2 - y_1 - y_2 + 2y_3 - s_1 = 0 \quad \Longleftrightarrow \quad s_1 = 2 - y_1 - y_2 + 2y_3 \geq 0$$

$$\Longleftrightarrow \quad y_1 + y_2 - 2y_3 \leq 2.$$

Putting it all together, the **dual** is:

$$\begin{aligned} &\text{maximize } y_1 - 3y_2 \\ &\text{subject to } y_1 - 2y_2 + y_3 = 1, \\ &\quad y_1 + y_2 - 2y_3 \leq 2, \\ &\quad y_1 \geq 0, \quad y_2 \geq 0, \quad y_3 \text{ free.} \end{aligned}$$

That's the full Lagrangian-based derivation for a primal with a mix of ' $\geq$ ', ' $\leq$ ', and ' $=$ ' constraints, plus a nonnegative variable and a free variable.

DTU

