

Divide y Conquista

Merge-Sort (A, p, r)

if p < r

q = L(p+r)/2

Merge-Sort (A, p, q)

Merge-Sort (A, q+1, r) ← **Dividir**

Merge (A, p, q, r) ← **Combinar**

Merge (A, p, q, r) //

Recibe un arreglo A [L...n], los índices p, q, r tales que A [p...q], A [q+1...r] están ordenados. Ordena A [p...r].

n₁ = q - p + 1

n₂ = r - q

for i = 1 to n₁

L[i] = A[p + i - 1]

for j = 1 to n₂

R[j] = A[q + j]

L[n₁ + 1] = ∞

R[n₂ + 1] = ∞

i = 1

j = 1

for k = p to r

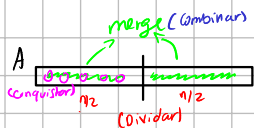
if L[i] ≤ R[j]

A[k] = L[i]

i = i + 1

else A[k] = R[j]

j = j + 1



C	1
C	1
C	n
C	n+1
C	n ₂ +1
C	n ₂
C	1
C	1
C	1
C	1
C	n ₁ + n ₂ + 1
C	n ₁ × n ₂
C	n ₁ × n ₂
C	n ₁ × n ₂
C	n ₁ × n ₂
C	n ₁ × n ₂

$$F(n) \geq F(2^k)$$

$$= 2^k (\lg(2^k) + 1)$$

$$= 2^k (k+1)$$

$$= \frac{2^{k+1}}{2} (k+1) > \frac{n \lg n}{2}$$

porque sabemos que $2^k \leq n < 2^{k+1}$

Concluimos que

$$\frac{1}{2} n \lg n < F(n) \leq 6n \lg n$$

Faltó demostrar que F(n) es creciente.

Recordar

$$F(n) = \begin{cases} 1 & n=1 \\ 2F(\lfloor n/2 \rfloor) + n & n>1 \end{cases}$$

Basta probar que $F(n) < F(n+1)$ para todo $n \geq 1$

Por ind en n. Si $n=1$, tenemos que $F(n) = F(1) = 1 < 4 = F(2) = F(n+1)$

Si $n > 1$

Caso 1: n es par. $F(n) < F(n) + 1 = 2F(n/2) + n + 1 = F(n+1)$

Caso 2: n es impar. $F(n) = 2F(\frac{n-1}{2}) + n$

$$\begin{aligned} & \text{Ch)} \\ & < 2F\left(\frac{n-1}{2} + 1\right) + n \\ & = 2F\left(\frac{n+1}{2}\right) + n \\ & < 2F\left(\frac{n+1}{2}\right) + n + 1 < 2F\left(\frac{n+1}{2}\right) + n + 1 = F(n+1) \end{aligned}$$

Sea $F: \mathbb{N} \rightarrow \mathbb{R}^+$ definido por:

$$F(n) = \begin{cases} 1 & n=1 \\ 2F(n-1) + 1 & \text{caso contrario} \end{cases}$$

¿Cuánto vale F(n)?

$$\begin{aligned} F(n) &= 2F(n-1) + 1 \\ &= 2(2F(n-2) + 1) + 1 \\ &= 2^2 F(n-2) + 2 + 1 \\ &= 2^2 (2F(n-3) + 1) + 2 + 1 \\ &= 2^3 F(n-3) + 2^2 + 2 + 1 \\ &= 2^3 (2F(n-4) + 1) + 2^2 + 2 + 1 \\ &= 2^4 F(n-4) + 2^3 + 2^2 + 2 + 1 \\ &\vdots \\ &= 2^i F(n-i) + \sum_{k=0}^{i-1} 2^k \\ &= 2^i F(n-i) + 2^i - 1 \\ &= 2^{n-1} F(1) + 2^{n-1} - 1 \\ &= 2^{n-1} + 2^{n-1} - 1 \\ &= 2^n - 1 \end{aligned}$$

Sea $F: \mathbb{N} \rightarrow \mathbb{R}^+$ definido por:

$$F(n) = \begin{cases} 1 & n=1 \\ F(n-1) + n & \text{caso contrario} \end{cases}$$

¿Cuánto vale F(n)?

$$\begin{aligned} F(n) &= F(n-1) + n \\ &= (F(n-2) + (n-1)) + n \\ &= F(n-2) + 2n - 1 \\ &= (F(n-3) + (n-2)) + n \\ &= F(n-3) + 3n - (1+2) \\ &= (F(n-4) + (n-3)) + n \\ &= F(n-4) + 4n - (1+2+3) \\ &\vdots \\ &= F(n-i) + in - (1+2+3+\dots+(i-1)) \end{aligned}$$

Analysis Tiempo de ejecución

$$T(n) = aT(n/b) + D(n) + C(n)$$

a: cantidad de sub problemas

b: división de sub problemas

Tm del MergeSort: $T(n) = 2T(n/2) + C_1 + C_2 n$

$$T(n) = \begin{cases} 2F(n/2) + n & n \geq 1 \\ 1 & n = 1 \end{cases}$$

Supongamos que $n = 2^k$ para algún $k \in \mathbb{N}$.

$$\begin{aligned} F(2^k) &= 2F(2^{k-1}) + 2^k \\ &= 2(2F(2^{k-2}) + 2^{k-1}) + 2^k \\ &= 2^2 F(2^{k-2}) + 2^k + 2^k \\ &= 2^2 (2F(2^{k-3}) + 2^{k-2}) + 2^k + 2^k + 2^k \\ &\vdots \\ &= 2^i F(2^{k-i}) + i \cdot 2^k \\ &\vdots \\ &= 2^k \cdot F(2^0) + k \cdot 2^k = 2^k + k \cdot 2^k = 2^k (k+1) = n(\lg n + 1) \end{aligned}$$

Cuando n es potencia de 2 $\Rightarrow F(n) = n(\lg n + 1)$

Ahora, suponga que $n \in \mathbb{N}$. Sea $x \in \mathbb{N}$ tal que $2^x \leq n < 2^{x+1}$

Como F es creciente (probaremos después), $F(2^x) \leq F(n) < F(2^{x+1})$

Luego, $F(n) < F(2^{x+1}) = 2^{x+1} (\lg(2^{x+1}) + 1)$

$$= 2^{x+1} \cdot (x+2)$$

$$\leq 2^{x+1} \cdot 3x$$

$$= 2 \cdot 2^x \cdot 3x$$

$$= 6x \cdot 2^x \leq 6n \lg n$$

$$= F(n-i) + in - (1+2+3+\dots+(i-1))$$

$$i=n-1$$

$$= F(1) + n(n-1) - (1+2+3+\dots+(n-2))$$

$$= 1 + n(n-1) - \frac{(n-2)(n-1)}{2}$$

$$= 1 + (n-1) \left(n - \frac{(n-2)}{2} \right)$$

$$= 1 + (n-1) \left(\frac{n+2}{2} \right)$$

$$= 1 + \frac{n^2+n-2}{2}$$

$$= \frac{n^2+n}{2} = \frac{n(n+1)}{2} \checkmark$$

$$T(1)=1$$

$$T(n) = 2T(\lfloor n/2 \rfloor) + n^2$$

Para n potencia de 2 $n=2^k$

$$= 2T(n/2) + n^2$$

$$= 2(2T(n/4) + (n/2)^2) + n^2$$

$$= 2^2 T(n/4) + \frac{n^2}{2} + n^2$$

$$= 2^2 (2T(n/8) + (n/4)^2) + \frac{n^2}{2} + n^2$$

$$= 2^3 T(n/8) + 2^2 (n/4)^2 + \frac{n^2}{2} + n^2$$

$$= 2^3 T(n/8) + \frac{n^2}{4} + \frac{n^2}{2} + n^2$$

\vdots

$$= 2^i T(n/2^i) + n^2 \left(\frac{1}{2^0} + \frac{1}{2^1} + \frac{1}{2^2} + \dots + \frac{1}{2^{i-1}} \right)$$

$$= 2^i T(n/2^i) + n^2 \left(1 + \left(\frac{1}{2}\right)^2 + \left(\frac{1}{2}\right)^3 + \dots + \left(\frac{1}{2}\right)^{i-1} \right)$$

$$= 2^i T(n/2^i) + n^2 \left(\frac{(1/2)^{i-1} - 1}{1/2 - 1} \right)$$

$$i=\lg n$$

$$= 2^{\lg n} T(n/2^{\lg n}) + n^2 \left(\frac{(1/2)^{\lg n} - 1}{1/2 - 1} \right)$$

$$= n T(1) + n^2 \left(\frac{1 - 1/2}{1/2} \right)$$

$$= n + 2n^2 = n(2n+1)$$

Demostrar por inducción

$$T(n) = \begin{cases} 1 & n=1 \\ 2T(\lfloor n/2 \rfloor) + n & \text{caso contrario.} \end{cases}$$

Probar $T(n) = \Omega(n \lg n)$

Prop $T(n) \geq \frac{1}{4} n \lg n$ para todo $n \geq 4$.

Caso base

$$T(4) = 2T(2) + 4 = 12 \geq \frac{1}{4} 4 \lg 4 = \frac{1}{4} n \lg n$$

Paso inductivo

$$T(n) = 2T(\lfloor n/2 \rfloor) + n$$

$$\geq 2 \left(\frac{1}{4} \lfloor \frac{n}{2} \rfloor \lg \lfloor \frac{n}{2} \rfloor \right) + n$$

$$\geq \frac{1}{2} \left(\frac{n}{2} - 1 \right) \lg \left(\frac{n}{2} - 1 \right) + n$$

$$\geq \frac{1}{2} \left(\frac{n}{2} - 1 \right) (\lg n - 2) + n$$

$$= \frac{1}{4} (n-2) (\lg n - 2) + n$$

$$= \frac{1}{4} n \lg n - \frac{1}{2} \lg n - \frac{n}{2} + 1 + n$$

$$= \frac{1}{4} n \lg n + \frac{1}{2} (n - \lg n) + 1 \geq \frac{1}{4} n \lg n$$

Dado,

$$T(n) = \begin{cases} 1 & n=1 \\ 2T(\lfloor n/2 \rfloor) + 1 & \text{caso contrario} \end{cases}$$

$$\lfloor n \rfloor \leq n \leq \lceil n \rceil$$

Probar por inducción que $T(n) = O(n)$.

Prop $T(n) \leq 2n-1$ para todo $n \geq 1$

Caso base

$$T(1) = 1 \leq 2 \cdot 1 - 1$$

Paso inductivo

$$T(n) = 2T(\lfloor n/2 \rfloor) + 1$$

$$\leq 2C \left\lfloor \frac{n}{2} \right\rfloor + 1$$

$$\leq 2C \left(\frac{n}{2} - 1 \right) + 1$$

$$\leq 2n - 2 + 1$$

$$= 2n - 1$$

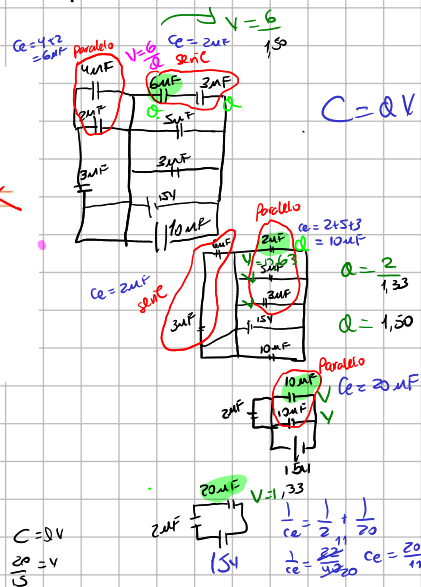
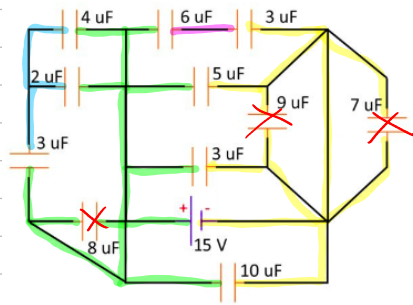
Corolario

$$T(n) \leq 2n \quad \forall n \geq 1$$

$$\text{Entonces } T(n) = O(n)$$

$$T(n) = n \left(\frac{3n}{2} - \frac{1}{2} \right)$$

$$T(n) = 1 + n^2 \left(\frac{3n-3}{2} \right)$$



$$C = QV$$
$$\frac{Q}{V} = V$$