Analisis y Diseño de Algoritmos

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October 5, 2022

Análisis probabilístico y Quicksort

Análisis probabilístico

```
HIRE-ASSISTANT (n)

1  best = 0  // candidate 0 is a least-qualified dummy candidate

2  for i = 1 to n

3  interview candidate i

4  if candidate i is better than candidate best

5  best = i

6  hire candidate i
```

Figure 1: Tomada del libro Cormen, Introduction to Algorithms

```
QUICKSORT(A, p, r)

1 if p < r

2 q = \text{PARTITION}(A, p, r)

3 QUICKSORT(A, p, q - 1)

4 QUICKSORT(A, q + 1, r)
```

Figure 2: Tomada del libro Cormen, Introduction to Algorithms

```
PARTITION(A, p, r)

1 x = A[r]

2 i = p - 1

3 for j = p to r - 1

4 if A[j] \le x

5 i = i + 1

6 exchange A[i] with A[j]

7 exchange A[i + 1] with A[r]

8 return i + 1
```

Figure 3: Tomada del libro Cormen, Introduction to Algorithms

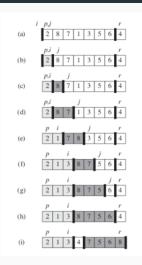


Figure 4: Tomada del libro Cormen, Introduction to Algorithms

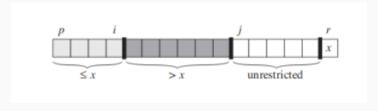


Figure 5: Tomada del libro Cormen, Introduction to Algorithms

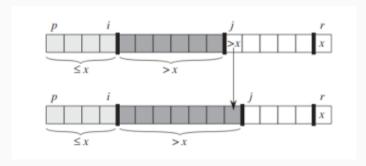


Figure 6: Tomada del libro Cormen, Introduction to Algorithms

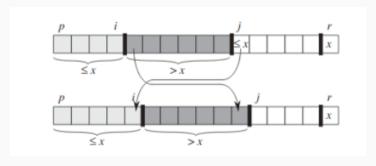


Figure 7: Tomada del libro Cormen, Introduction to Algorithms

T(n): cota superior para el tiempo de ejecución del Quicksort con n elementos.

$$T(n) \le \max_{0 < q < n-1} \{ T(q) + T(n-1-q) \} + O(n).$$

T(n): cota inferior para el tiempo de ejecución del Quicksort con n elementos.

$$T(n) \ge \min_{0 < q < n-1} \{ T(q) + T(n-1-q) \} + \Omega(n).$$

Peor caso.

$$T(n) = T(0) + T(n-1) + O(n).$$

En este caso, $T(n) = O(n^2)$

¿Por qué es el peor caso?

Recuerde que, si T(n) es una cota superior para cualquier n, entonces

$$T(n) \le \max_{0 \le q \le n-1} \{ T(q) + T(n-1-q) \} + O(n).$$

Podemos probar por inducción que $T(n) \le cn^2$ para alguna constante c > 0.

Gracias