

Analisis y Diseño de Algoritmos

Juan Gutiérrez

November 10, 2022

Ordenación en tiempo lineal

Algoritmos basados en comparaciones

- Mergesort
- Heapsort
- Quicksort

Todos son $\Omega(n \lg n)$. ¿Se podrá hacer mejor que eso?

Límites inferiores para algoritmos basados en comparaciones

Árbol de decisión

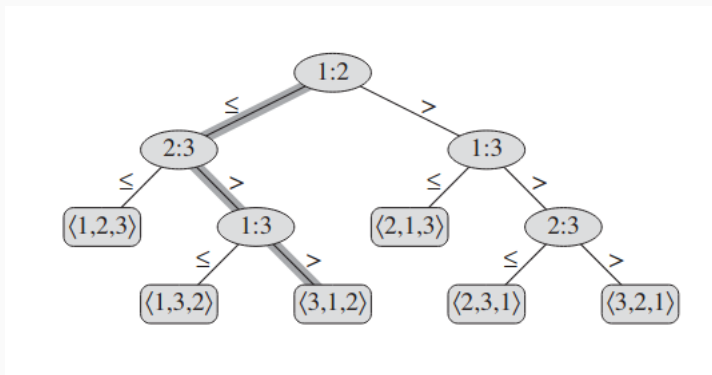


Figure 1: Tomada del libro Cormen, Introduction to Algorithms

Counting sort

```
COUNTING-SORT( $A, B, k$ )
1  let  $C[0..k]$  be a new array
2  for  $i = 0$  to  $k$ 
3       $C[i] = 0$ 
4  for  $j = 1$  to  $A.length$ 
5       $C[A[j]] = C[A[j]] + 1$ 
6  //  $C[i]$  now contains the number of elements equal to  $i$ .
7  for  $i = 1$  to  $k$ 
8       $C[i] = C[i] + C[i - 1]$ 
9  //  $C[i]$  now contains the number of elements less than or equal to  $i$ .
10 for  $j = A.length$  downto 1
11      $B[C[A[j]]] = A[j]$ 
12      $C[A[j]] = C[A[j]] - 1$ 
```

Figure 2: Tomada del libro Cormen, Introduction to Algorithms

Counting sort

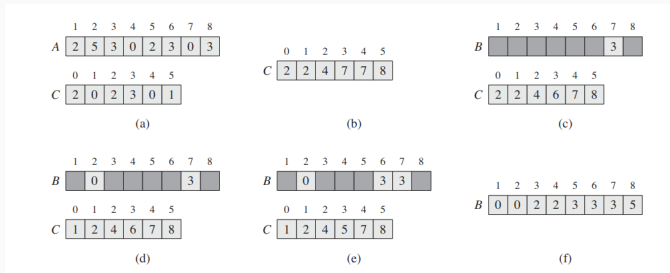


Figure 3: Tomada del libro Cormen, Introduction to Algorithms

RADIX-SORT(A, d)

1 **for** $i = 1$ **to** d

2 use a stable sort to sort array A on digit i

Figure 4: Tomada del libro Cormen, Introduction to Algorithms

Radix sort

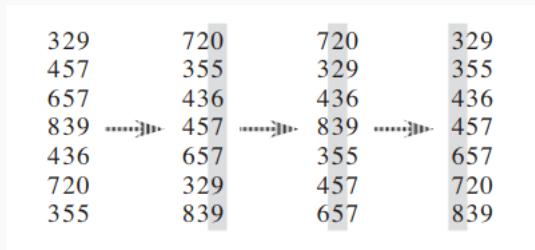


Figure 5: Tomada del libro Cormen, Introduction to Algorithms

Bucket sort

BUCKET-SORT(A)

```
1  let  $B[0..n-1]$  be a new array
2   $n = A.length$ 
3  for  $i = 0$  to  $n - 1$ 
4      make  $B[i]$  an empty list
5  for  $i = 1$  to  $n$ 
6      insert  $A[i]$  into list  $B[\lfloor nA[i] \rfloor]$ 
7  for  $i = 0$  to  $n - 1$ 
8      sort list  $B[i]$  with insertion sort
9  concatenate the lists  $B[0], B[1], \dots, B[n-1]$  together in order
```

Figure 6: Tomada del libro Cormen, Introduction to Algorithms

Bucket sort

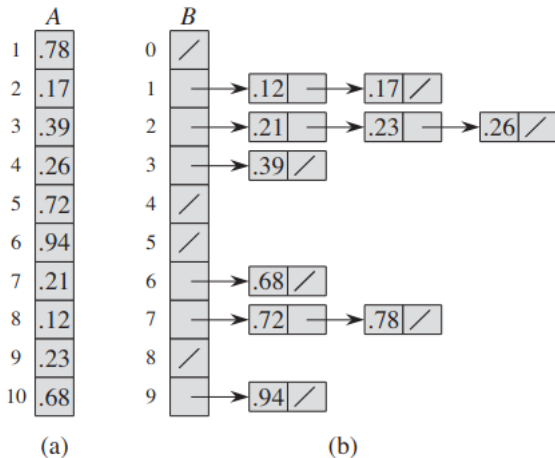


Figure 7: Tomada del libro Cormen, Introduction to Algorithms

Gracias