

Eliminar la recurrencia:

En cada caso suponer que $T(1)=1$.

③ $T(n) = 4T(n/2) + n$

Suponemos que n es potencia de 2. Es decir, $n=2^k$.

$$\begin{aligned} T(n) &= 4T(n/2) + n \\ &= 4(4T(n/4) + n/2) + n \\ &= 4^2 T(n/4) + 2n + n \\ &= 4^2 (4T(n/8) + n/4) + 2n + n \\ &= 4^3 T(n/8) + 4n + 2n + n \\ &= 4^3 (4T(n/16) + n/8) + 4n + 2n + n \\ &= 4^4 T(n/16) + 8n + 4n + 2n + n \\ &= 4^i T(n/2^i) + 2^i n + 2^{i-1} n + \dots + 2^2 n + 2^1 n \\ &= 2^i T(n/2^i) + n(2^i + 2^{i-1} + \dots + 2^2 + 2^1 + 2^0) \\ &= 2^i T(n/2^i) + n(2^i + 2^{i-1} + \dots + 2^0) \\ &= 2^i T(n/2^i) + n(2^{i+1} - 1) \\ &= 2^i T(n/2^i) + n(2^i - 1) \end{aligned}$$

$i = \lg n$

$$\begin{aligned} &= 2^{\lg n} T(n/2^{\lg n}) + n(2^{\lg n} - 1) \\ &= nT(1) + n(n-1) \\ &= \underline{n + n(n-1)} \quad \alpha. \end{aligned}$$

④ $T(n) = 2T(n/2) + n^3$

Suponemos que n es potencia de 2. Es decir, $n=2^k$.

$$\begin{aligned} T(n) &= 2T(n/2) + n^3 \\ &= 2(2T(n/4) + (n/2)^3) + n^3 \\ &= 2^2 T(n/4) + \frac{n^3}{2^2} + \frac{n^3}{2^2} \\ &= 2^2 (2T(n/8) + (n/8)^3) + \frac{n^3}{2^2} + \frac{n^3}{2^2} \\ &= 2^3 T(n/8) + \frac{n^3}{2^3} + \frac{n^3}{2^3} + \frac{n^3}{2^3} \\ &= 2^3 (2T(n/16) + (n/16)^3) + \frac{n^3}{2^3} + \frac{n^3}{2^3} + \frac{n^3}{2^3} \\ &= 2^4 T(n/16) + \frac{n^3}{2^4} + \frac{n^3}{2^4} + \frac{n^3}{2^4} + \frac{n^3}{2^4} \\ &= 2^i T(n/2^i) + n^3 \left(\frac{1}{2^i} + \frac{1}{2^i} + \dots + \frac{1}{2^i} \right) \\ &= 2^i T(n/2^i) + n^3 \left(\frac{1}{2^i} + \frac{1}{2^i} + \dots + \frac{1}{2^i} \right) \end{aligned}$$

$i = \lg n$

$$= 2^{\lg n} T(n/2^{\lg n}) + n^3 \left(\frac{1}{2^{\lg n}} + \frac{1}{2^{\lg n}} + \dots + \frac{1}{2^{\lg n}} \right)$$

$$= nT(1) + n^3 \left(\frac{1}{2^{\lg n}} + \frac{1}{2^{\lg n}} + \dots + \frac{1}{2^{\lg n}} \right)$$

$$= n + \frac{n^3}{2^{\lg n}} - \frac{n^3}{2^{\lg n+1}}$$

$$= n - \frac{4n}{3} + \frac{4n^3}{3}$$

$$= -\frac{n}{3} + \frac{4n^3}{3}$$

$$= \frac{n}{3} (4n^2 - 1)$$

$$= \frac{n}{3} (2n+1)(2n-1)$$

⑤ Suponga que $T(1)=1$,

$$T(n) = 2T(n/2) + n^2$$

$n=2^k$

$$\begin{aligned} T(n) &= 2T(n/2) + n^2 \\ &= 2(2T(n/4) + \frac{n^2}{2}) + n^2 \\ &= 2^2 T(n/4) + \frac{n^2}{2} + n^2 \\ &= 2^2 (2T(n/8) + \frac{n^2}{2^2}) + \frac{n^2}{2} + n^2 \\ &= 2^3 T(n/8) + \frac{n^2}{2^2} + \frac{n^2}{2} + n^2 \\ &\vdots \\ &= 2^i T(n/2^i) + \frac{n^2}{2^0} + \frac{n^2}{2^1} + \dots + \frac{n^2}{2^{i-1}} \end{aligned}$$

$i = \lg n$

$$\begin{aligned} &= nT(1) + \left(\frac{1}{2^0} + \frac{1}{2^1} + \dots + \frac{1}{2^{\lg n-1}} \right) n^2 \\ &= n + \sum_{k=0}^{\lg n-1} \left(\frac{1}{2^k} \right) n^2 \\ &= n + \left(\frac{1/2 - 1/2^{\lg n}}{1/2 - 1} \right) n^2 \\ &= n + \frac{2(n-1)}{n} n^2 \end{aligned}$$

$$\begin{aligned} T(n) &= n + 2(n-1)n \\ &= n + 2n^2 - 2n \\ &= 2n^2 - n \\ &= \underline{n(2n-1)} \quad \alpha. \end{aligned}$$

⑥

$$T(n) = 2T(n-1) + 3n - 2$$

$$\begin{aligned} &= 2(2T(n-2) + 3(n-1) - 2) + 3n - 2 \\ &= 2^2 T(n-2) + 2 \cdot 3(n-1) - 2^2 + 3n - 2 \\ &= 2^2 (2T(n-3) + 3(n-2) - 2) + 2 \cdot 3(n-1) - 2^2 + 3n - 2 \\ &= 2^3 T(n-3) + 3 \cdot 2^2 (n-2) - 2^3 + 3 \cdot 2^2 (n-1) - 2^2 + 3n - 2 \\ &\vdots \\ &= 2^i T(n-i) + 3 \left(2^{i-1} (n-(i-1)) + 2^{i-2} (n-(i-2)) + \dots + 2^1 (n-1) \right) - (2^i + 2^{i-1} + \dots + 2^1) \end{aligned}$$

$i = n-1$

$$= 2^{n-1} T(1) + 3 \left(2^{n-2} (2) + 2^{n-3} (3) + \dots + 2^{n-n} (n) \right) - (2^1 + 2^2 + \dots + 2^{n-1})$$

$$S = 2 \cdot 2^{n-2} + 3 \cdot 2^{n-3} + \dots + n \cdot 2^1$$

$$2S = 2 \cdot 2^{n-1} + 3 \cdot 2^{n-2} + \dots + n \cdot 2^0$$

$$S = 1 \cdot 2^{n-2} + 1 \cdot 2^{n-3} + \dots + 1 \cdot 2^1 + 2 \cdot 2^0 - n \cdot 2^0$$

$$= \frac{2^{n+1} - 2}{2-1} - 1 + 2^n - n$$

$$= 2^{n+1} - 2 + 2^n - n$$

$$= 2^n (1+2) - 2 - n$$

$$= 2^{n+1} - 2 - n$$

$$= 2^{n-1} + 3 \left(2^{n-2} + 2^{n-3} + \dots + 2^1 \right) - (2^{n-1} + 2^{n-2} + \dots + 2^1)$$

$$= 2^{n-1} + 3 \cdot 2^{n-2} - 3 \cdot 2^{n-2} - 2^n + 2$$

$$= 2^{n-1} + 2^{n-1} - 2^n - 3n - 4$$

$$= 2^{n-1} (1+2-2) - 3n - 4$$

$$= 2^{n-1} \cdot 2 - 3n - 4$$

$$= 2^n - 3n - 4$$

Exámenes Pasados

Ejercicio 3 (5 pts). Resuelva la recurrencia de manera explícita, usando los métodos vistos en clase. Debe dar la solución exacta cuando n es potencia de 3. Cuando no lo es, debe acotar adecuadamente.

$$T(n) = \begin{cases} 1 & n = 1 \\ 3T(\lfloor n/3 \rfloor) + n^2 & \text{caso contrario} \end{cases}$$

Puede suponer que es conocido que $T(n)$ es una función creciente.

$$n = 3^k$$

$$\begin{aligned} T(n) &= 3T(n/3) + n^2 \\ &= 3(3T(n/9) + \frac{n^2}{3}) + n^2 \\ &= 9T(n/9) + \frac{n^2}{3} + n^2 \\ &= 9(3T(n/27) + \frac{n^2}{9}) + \frac{n^2}{3} + n^2 \\ &= 27T(n/27) + \frac{n^2}{3} + \frac{n^2}{3} + n^2 \\ &\vdots \\ &= 3^i T(n/3^i) + \left(\frac{n^2}{3^{i-1}} + \frac{n^2}{3^{i-2}} + \dots + \frac{n^2}{3} + n^2 \right) \\ &= 3^i T(n/3^i) + \left(\frac{1}{3^{i-1}} + \frac{1}{3^{i-2}} + \dots + \frac{1}{3} + 1 \right) n^2 \\ &\stackrel{i = \log_3 n}{=} n T(1) + \left(\frac{1}{3^0} + \frac{1}{3^1} + \dots + \frac{1}{3^{i-2}} + \frac{1}{3^{i-1}} \right) n^2 \\ &= n \cdot 1 + \left(1 + \left(\frac{1}{3} \right)^1 + \dots + \left(\frac{1}{3} \right)^{i-2} + \left(\frac{1}{3} \right)^{i-1} \right) n^2 \\ &= n + \sum_{i=0}^{\log_3 n - 1} \left(\frac{1}{3} \right)^i n^2 \\ &= n + \frac{\left(\frac{1}{3} \right)^{\log_3 n - 1} - 1}{\frac{1}{3} - 1} n^2 \\ &= n + \frac{3(1-n)}{2n} n^2 \\ &= n + \left(\frac{3-3n}{2n} \right) n^2 \\ &= n + \left(\frac{3n-3}{2n} \right) n^2 \\ &= n + \frac{3n^2 - n}{2} \\ &= \frac{3n^2 - n}{2} \\ &= n \left(\frac{3n-1}{2} \right) \end{aligned}$$

$$\frac{1}{\frac{1}{3} - 1} = \frac{1}{\frac{1}{n} - 1} = \frac{1-n}{\frac{1}{3} - 1} = \frac{1-n}{-\frac{2}{3}} = \frac{3(1-n)}{-2n}$$

PRUEBAS POR INDUCCION

Ejercicio 3 (3 pts). Considere la siguiente recurrencia sobre los números naturales

$$T(n) = \begin{cases} 0 & n = 0 \\ 0 & n = 1 \\ T(n) = T(n-1) + T(n-2) + 1 & \text{caso contrario} \end{cases}$$

Pruebe por inducción que $T(n) = \Omega((3/2)^n)$. $0 \leq C(3/2)^n \leq T(n)$, C_1 positivo, $n \geq n_0 > 0$

Caso base, $n = 2$

$$0 \leq C \left(\frac{3}{2} \right)^2 = C \left(\frac{9}{4} \right) \leq T(2) = 1$$

$$\begin{aligned} \frac{n=2}{C \left(\frac{3}{2} \right)^2} &= \frac{C \cdot \frac{9}{4} \leq T(2) = 1}{\frac{9C}{4} \leq 1} \\ &\Rightarrow C \leq \frac{4}{9} \end{aligned}$$

Paso inductivo, $n > 2$

$$\begin{aligned} T(n) &= T(n-2) + T(n-1) + 1 \\ &\leq C \left(\frac{3}{2} \right)^{n-2} + C \left(\frac{3}{2} \right)^{n-1} + 1 \\ &= C \left(\frac{3}{2} \right)^{n-2} \left(1 + \frac{3}{2} \right) + 1 \\ &= C \left(\frac{3}{2} \right)^{n-2} \left(\frac{5}{2} \right) + 1 \\ &= C \left(\frac{10}{9} \right) \left(\frac{3}{2} \right)^n + 1 \\ &\leq C \left(\frac{10}{9} \right) \left(\frac{3}{2} \right)^n \\ &= \left(\frac{4}{9} \right) \left(\frac{10}{9} \right) \left(\frac{3}{2} \right)^n \\ &\leq \frac{4}{9} \left(\frac{3}{2} \right)^n \end{aligned}$$

Ejemplo 3.4. Sea $F: \mathbb{N} \rightarrow \mathbb{R}^+$ definido por

$$T(n) = \begin{cases} 1 & n = 1 \\ 2T(\lfloor \frac{n}{2} \rfloor) + n & \text{caso contrario} \end{cases}$$

Compruebe por inducción que $T(n) = O(n^2)$. Compruebe por inducción que $T(n) = \Omega(n \lg n)$.

Probaremos por inducción en n que $T(n) \leq 2n^2$ para $n \geq 1$

Si $n = 1$

$$T(1) = 1 \leq 2 \leq 2n^2$$

Si $n > 1$, tenemos que

$$\begin{aligned} T(n) &= 2T(n/2) + n \\ &\leq 2T(n/2) + n \\ &\stackrel{(n.i)}{\leq} 2 \left(\frac{n}{2} \right)^2 + n \\ &\leq 4 \left(\frac{n}{2} \right)^2 + n \\ &\leq n^2 + n^2 \\ &\leq 2n^2 \end{aligned}$$

$$\begin{aligned} 0 &\leq T(n) \leq Cn^2 \\ T(n/2) &\leq C \left(\frac{n}{2} \right)^2 \\ 2T(n/2) &\leq 2C \left(\frac{n}{2} \right)^2 \\ 2T(n/2) + n &\leq 2C \left(\frac{n}{2} \right)^2 + n \end{aligned}$$