





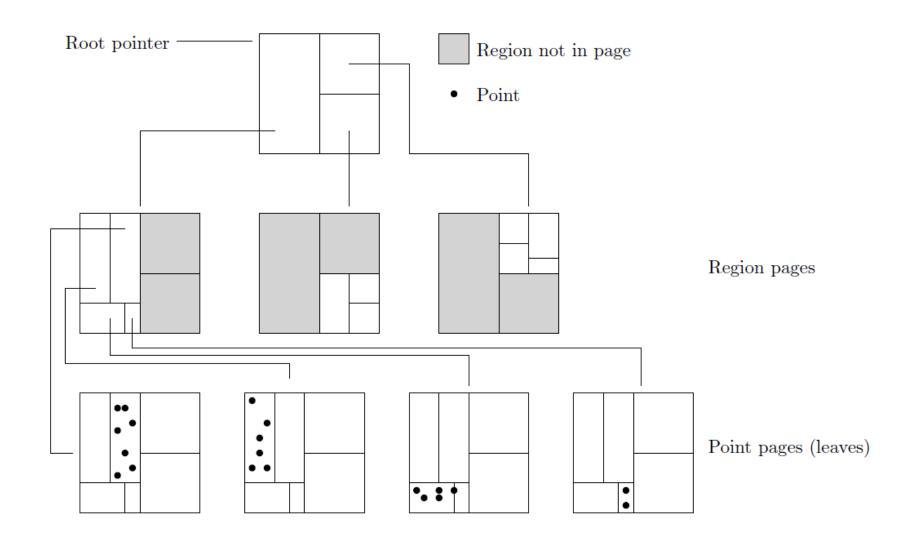
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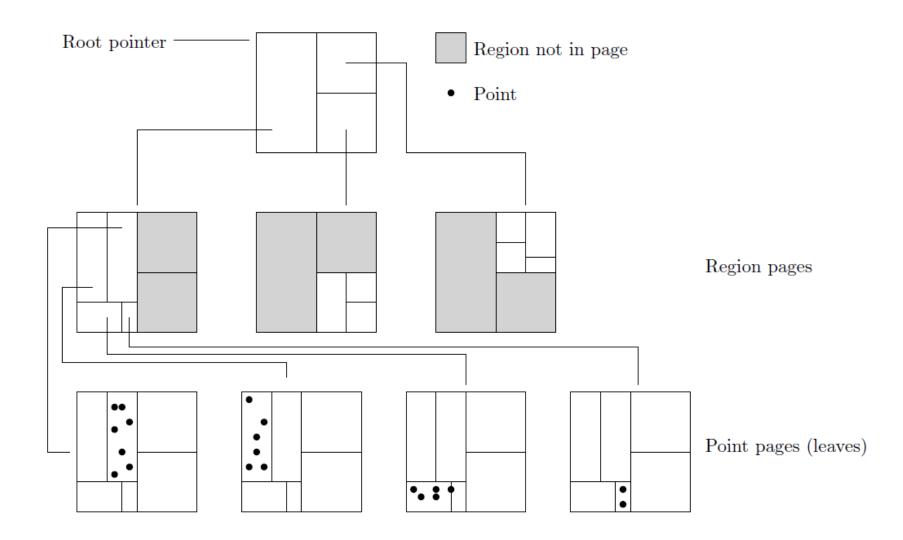


(k-dimensional B-tree)





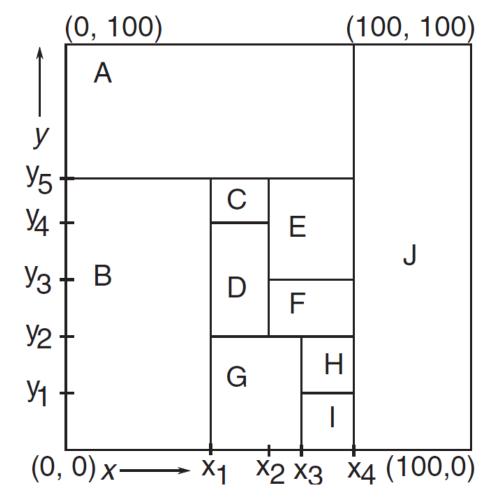
(k-dimensional B-tree)



- Todos los datos están en los nodos hoja.
- No podemos garantizar que cada *bucket* del k-d-B-*tree* esté lleno al 50%

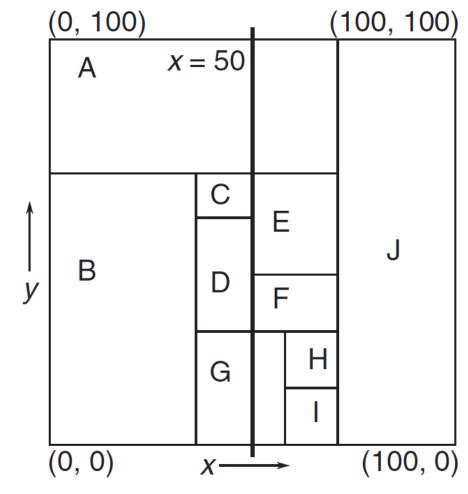


(k-dimensional B-tree)



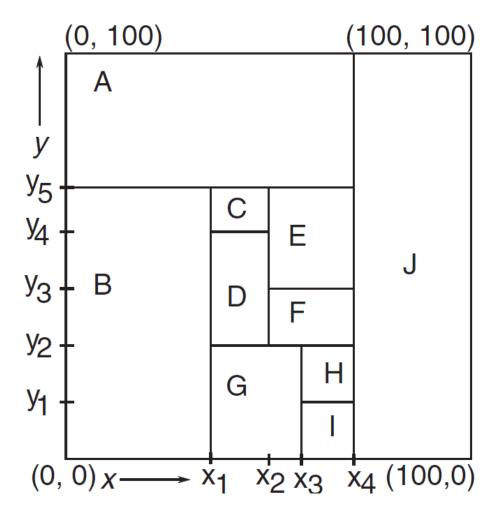
Desbordamiento del region page

#### Solución 1:



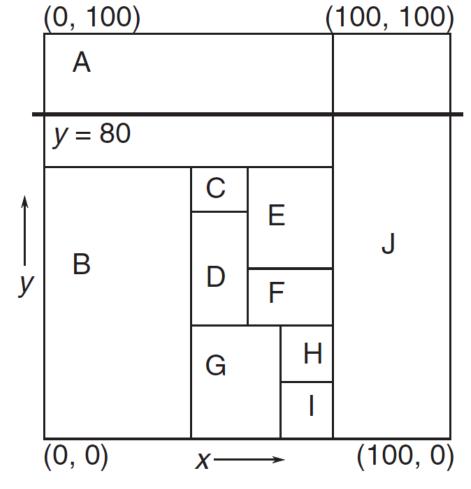
La división de propaga hacia abajo





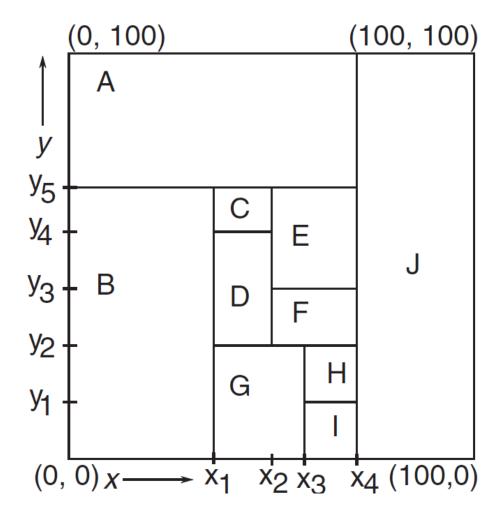
Desbordamiento del region page

Solución 2:



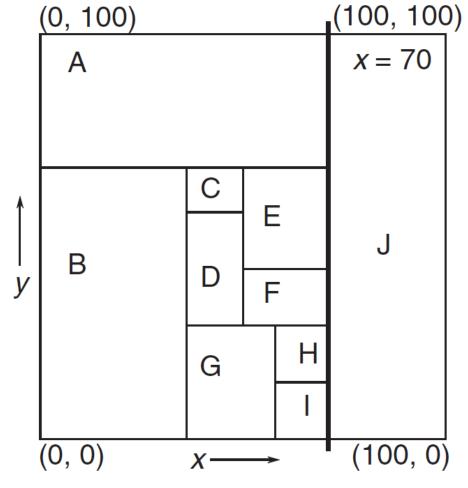
El desborde se mantiene





Desbordamiento del region page

Solución 3:



Sub-uso del region page.

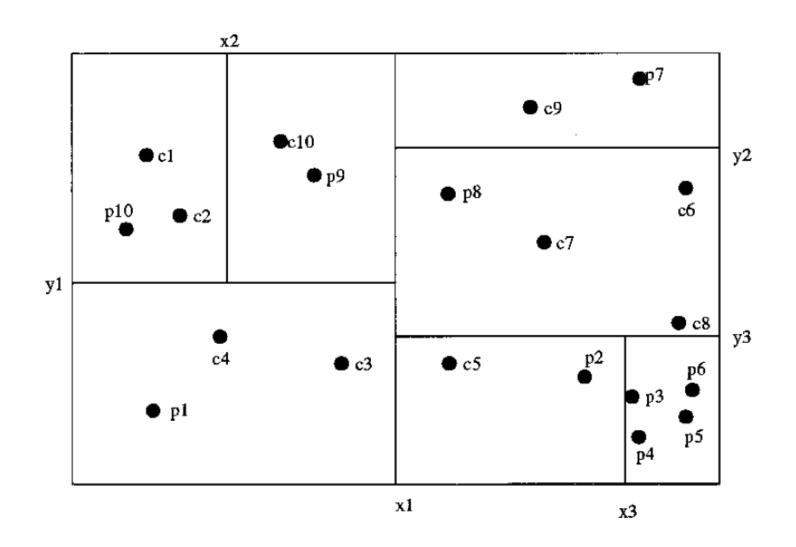


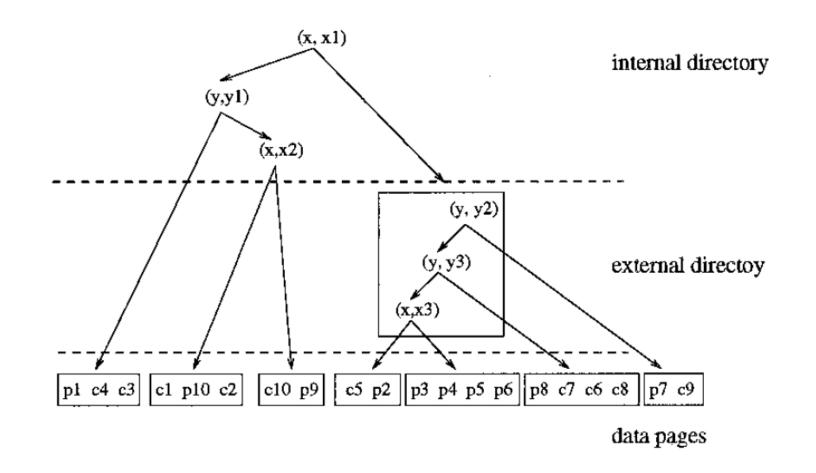
Entonces, que algoritmo podríamos utilizar?



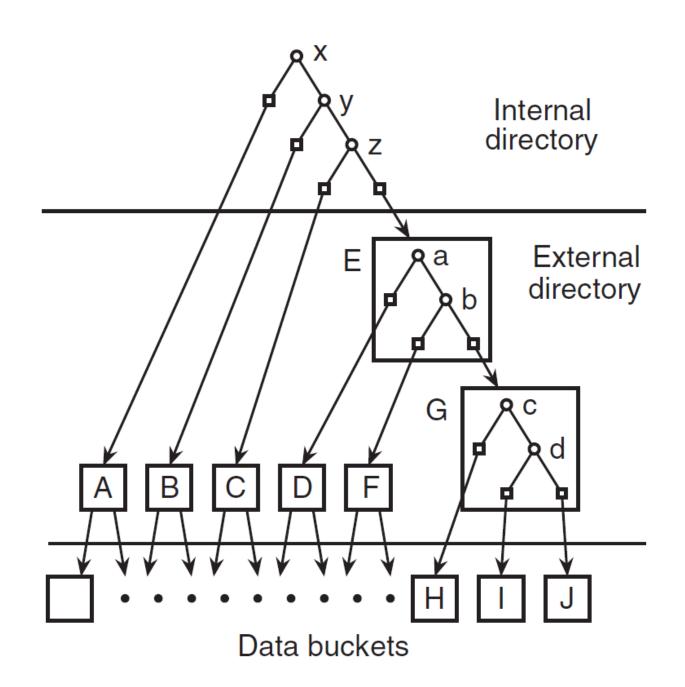


(Local Split Decision tree)

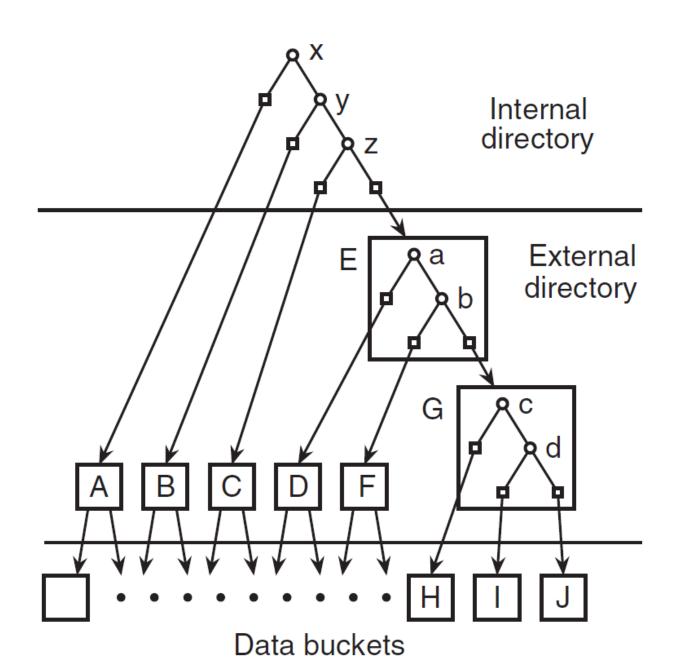


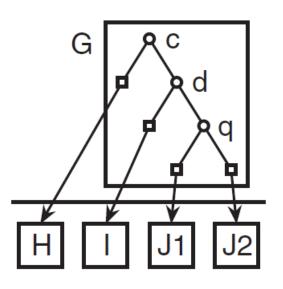




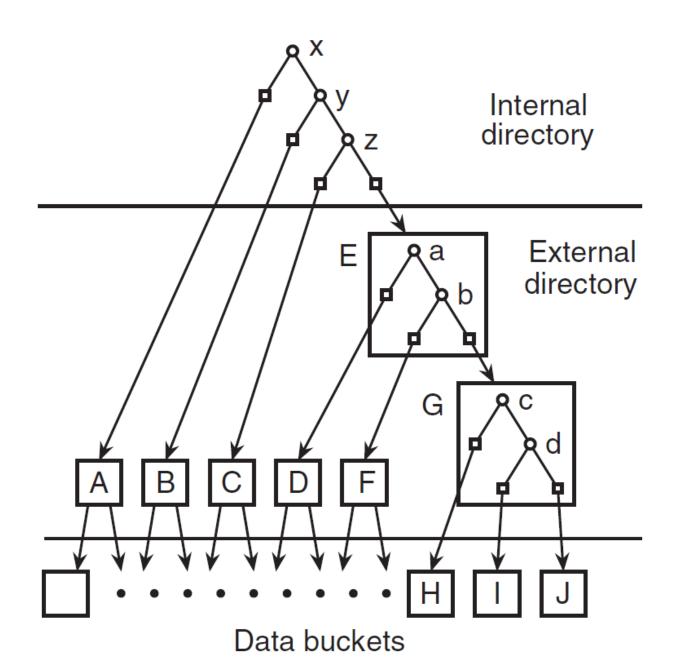


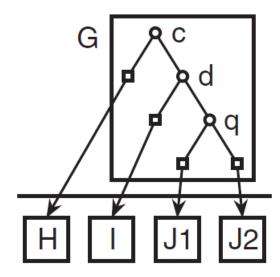


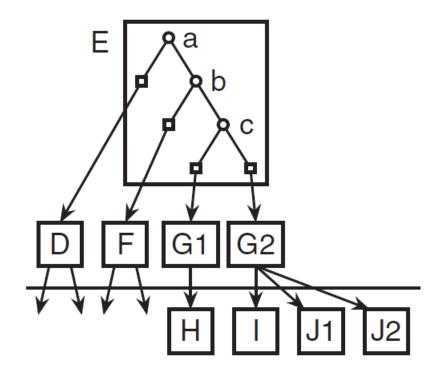




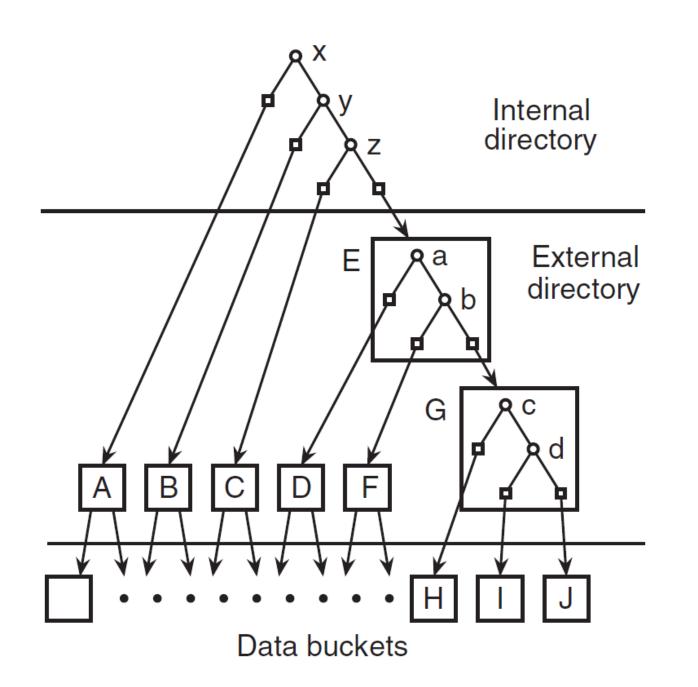


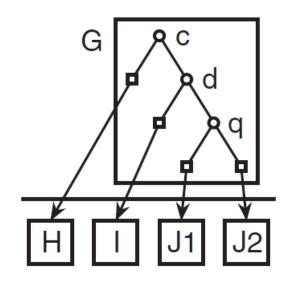


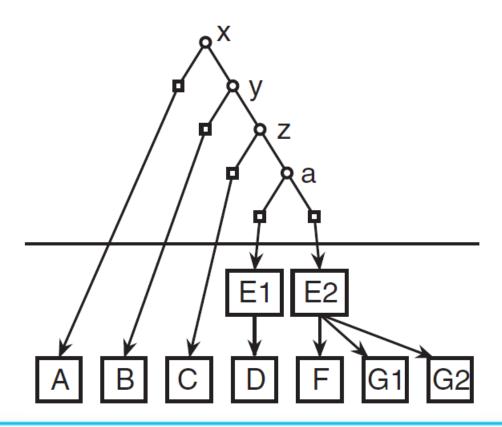


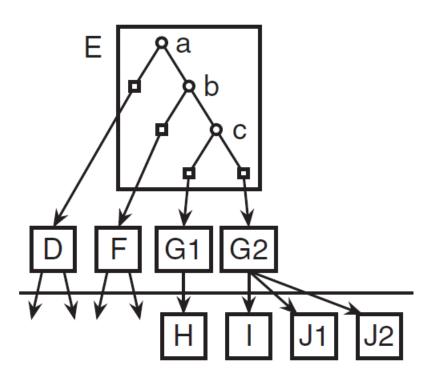




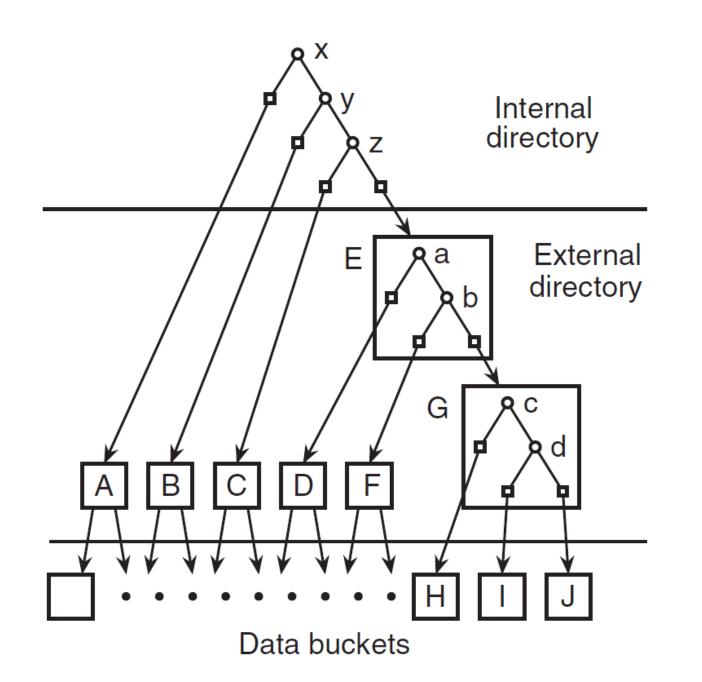


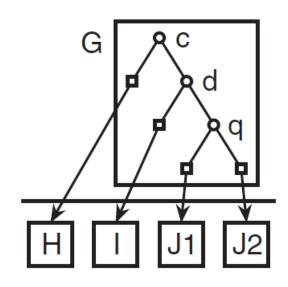


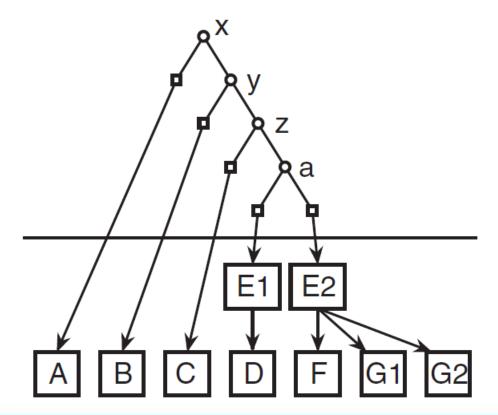


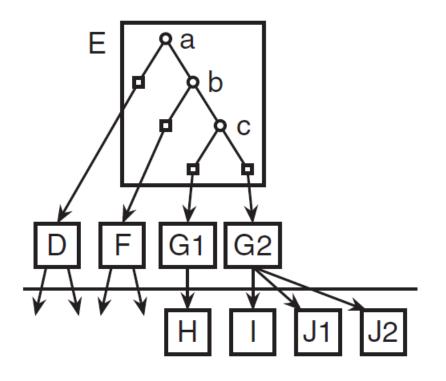


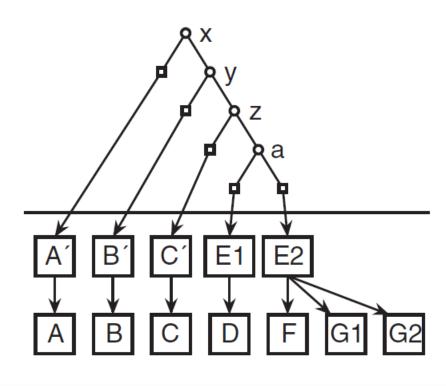




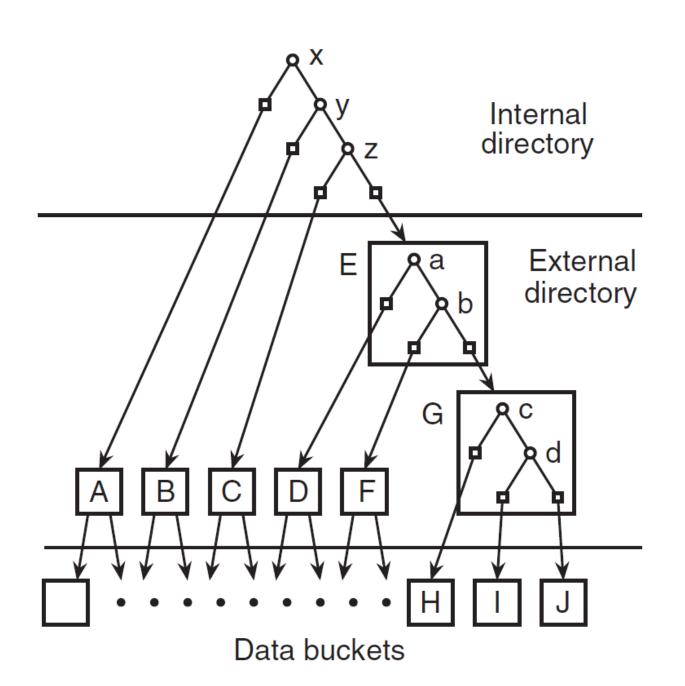


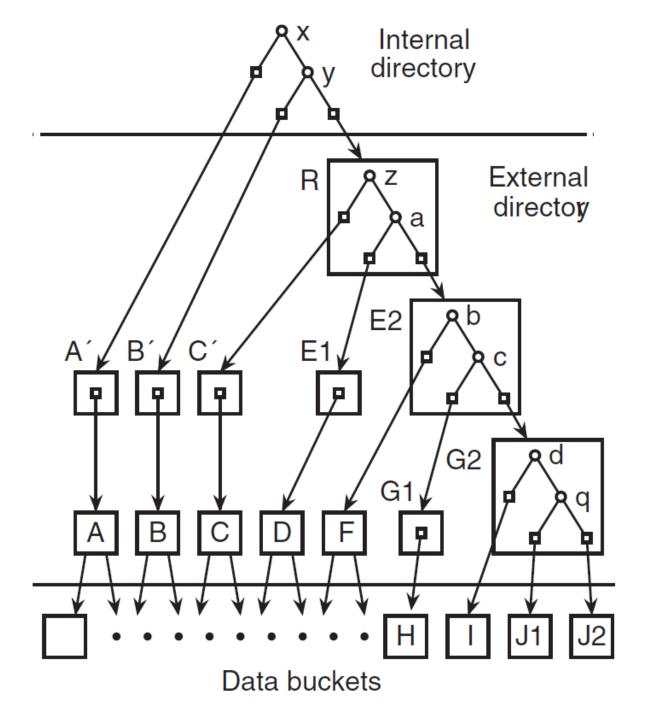








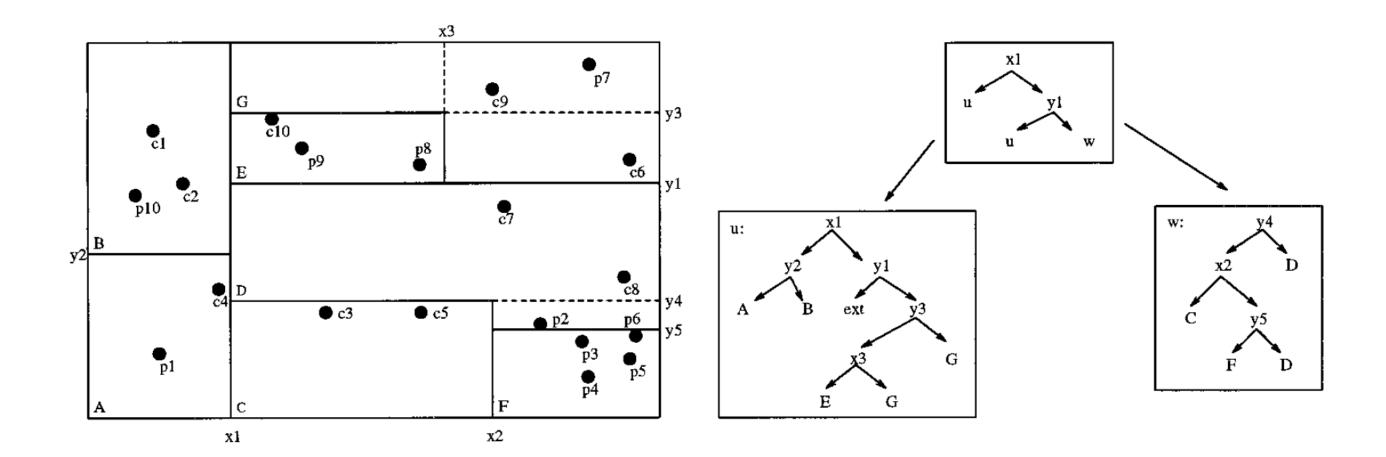




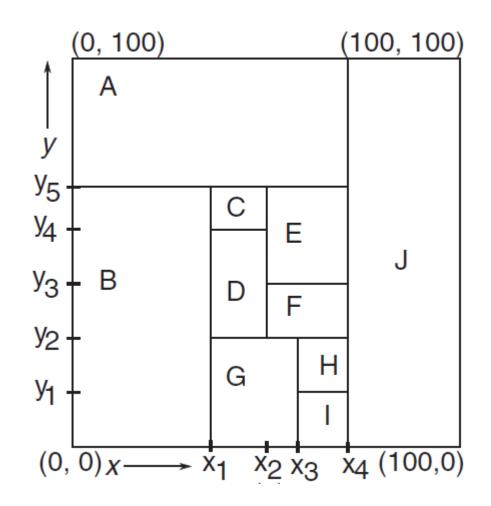


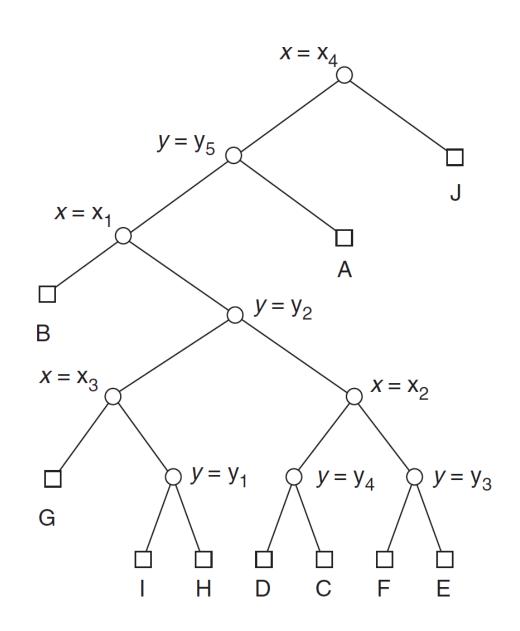


(Holey brick tree)

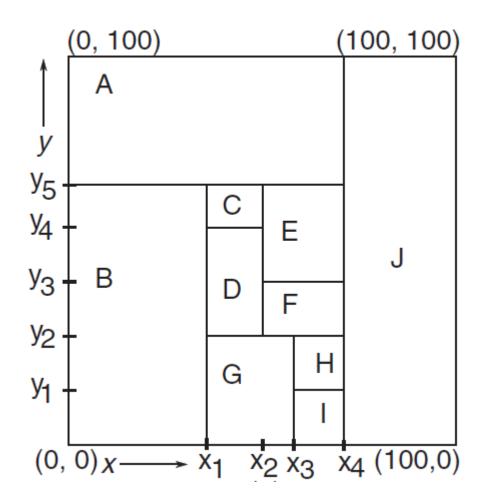


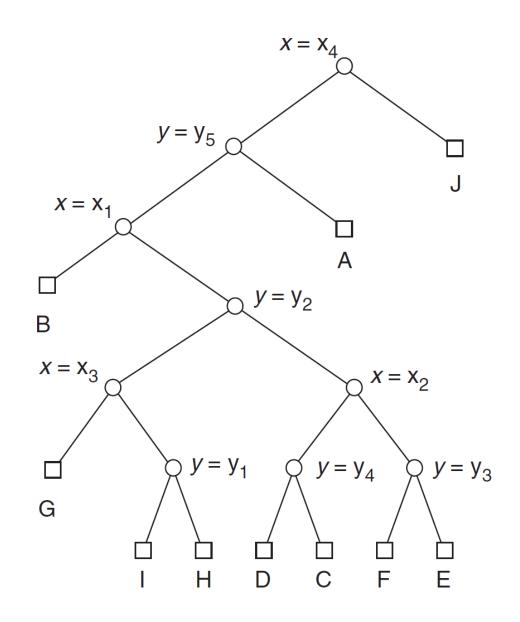












Recorremos el árbol descendiendo por elsubárbol con el mayor número de nodos de hoja, hasta obtener un tercio y dos tercios del número total de regiones de la página de regiones desbordadas



