Competitive Programming Algorithms

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1 Introduction and Miscellaneous

\overline{n}	Worst AC Algorithm	Comment
\leq [1011]	$O(n!), O(n^6)$	e.g., Enumerating permutations (Section 3.2)
$\leq [1719]$	$O(2^n \times n^2)$	e.g., DP TSP (Section 3.5.2)
$\leq [1822]$	$O(2^n \times n)$	e.g., DP with bitmask technique (Book 2)
$\leq [2426]$	$O(2^n)$	e.g., try 2^n possibilities with $O(1)$ check each
≤ 100	$O(n^4)$	e.g., DP with 3 dimensions + $O(n)$ loop, ${}_{n}C_{k=4}$
≤ 450	$O(n^3)$	e.g., Floyd-Warshall (Section 4.5)
$\leq 1.5K$	$O(n^{2.5})$	e.g., Hopcroft-Karp (Book 2)
$\leq 2.5K$	$O(n^2 \log n)$	e.g., 2 -nested loops $+$ a tree-related DS (Section 2.3)
$\leq 10K$	$O(n^2)$	e.g., Bubble/Selection/Insertion Sort (Section 2.2)
$\leq 200K$	$O(n^{1.5})$	e.g., Square Root Decomposition (Book 2)
$\leq 4.5M$	$O(n \log n)$	e.g., Merge Sort (Section 2.2)
$\leq 10M$	$O(n \log \log n)$	e.g., Sieve of Eratosthenes (Book 2)
$\leq 100M$	$O(n), O(\log n), O(1)$	Most contest problem have $n \leq 1M$ (I/O bottleneck)

1.1 Template

```
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;

int main() {
    cin.tie(0)->sync_with_stdio(0);

int n;
    cin >> n;
}
```

1.2 Codeforces Template

```
#include <bits/stdc++.h>
using namespace std;
4 typedef long long 11;
6 /* ======== */
8 void solve() {
    int n;
    cin >> n;
10
11 }
12
13 /* ======= */
15 int main() {
     cin.tie(0)->sync_with_stdio(0);
16
17
    int t = 1;
18
    cin >> t;
19
    while (t--) solve();
20
21 }
```

1.3 Generate Files

```
for i in {B..L}.cpp; do cp "A.cpp" "$i"; done
```

1.4 Binary Search

```
int binarysearch(function < bool(int) > f) {
    int lo = 0, hi = 100000, bestSoFar = -1;
    while (lo <= hi) {
        int mid = lo + (hi - lo) / 2;
        if (f(mid)) {
            bestSoFar = mid;
            hi = mid - 1;
        } else lo = mid + 1;
    }
    return bestSoFar;
}</pre>
```

1.5 Base 10 to Base m

```
char a[16] = {'0', '1', '2', '3', '4', '5', '6', '7', '8', '9', 'A', 'B', 'C', 'D', 'E', 'F'};

string tenToM(int n, int m) {
   int temp = n;
   string result = "";
   while (temp != 0) {
      result = a[temp % m] + result;
      temp /= m;
   }

return result;
}
```

1.6 Coordinate Compression

```
// coordinates -> (compressed coordinates).
map<int, int> coordMap;

void compress(vector<int>& values) {
    for (int v : values) coordMap[v] = 0;
    int cId = 0;
    for (auto it = coordMap.begin(); it != coordMap.end(); ++it) it->second = cId++;
    for (int& v : values) v = coordMap[v];
}
```

2 Data Structures

2.1 Order Statistic Tree (Set)

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>

using namespace __gnu_pbds;
using namespace std;

typedef tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>
ordered_set;
```

2.2 Order Statistic Tree (Map)

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>

using namespace __gnu_pbds;
using namespace std;

typedef tree<int, char, less<int>, rb_tree_tag, tree_order_statistics_node_update> ordered_map;
```

2.3 Union Find

```
1 const int N = 200010;
2 int parent[N];
3 int subtree_size[N];
5 void init(int n) {
      for (int i = 0; i < n; i++) {</pre>
6
           parent[i] = i;
           subtree_size[i] = 1;
8
      }
9
10 }
11
12 int root(int x) { return parent[x] == x ? x : parent[x] = root(parent[x]); }
13
void join(int x, int y) {
      x = root(x);
15
      y = root(y);
16
       if (x == y) return;
17
      if (subtree_size[y] < subtree_size[x]) swap(x, y);</pre>
18
19
       parent[x] = y;
       subtree_size[y] += subtree_size[x];
20
21 }
```

2.4 Sparse Table

```
1 const int N = 100000;
2 const int LOGN = 18;
4 int a[N];
5 // sparseTable[1][i] = max a[i..i+2^1)
6 int sparseTable[LOGN][N];
8 void precomp(int n) {
      // level 0 is the array itself
      for (int i = 0; i < n; i++) sparseTable[0][i] = a[i];</pre>
      for (int l = 1; l < LOGN; l++) { // inner loop does nothing if 2^2l > n
12
          int w = 1 << (1 - 1);</pre>
                                           // 2^(1-1)
13
14
           // a[i,i+2w) is made up of a[i,i+w) and a[i+w,i+2w)
15
           for (int i = 0; i + 2 * w <= n; i++)</pre>
16
               sparseTable[1][i] = max(sparseTable[1 - 1][i], sparseTable[1 - 1][i + w]);
17
      }
18
19 }
```

2.5 Range Tree

```
1 const int N = 100100;
2 int tree[1 << 18]; // 2^17 = 131,072</pre>
4 int n;
5
6 int query(int qL, int qR, int i = 1, int cL = 0, int cR = n) {
      if (cL == qL && cR == qR) return tree[i];
      int mid = (cL + cR) / 2;
      int ans = 0;
      if (qL < mid) ans += query(qL, min(qR, mid), i * 2, cL, mid);
10
      if (qR > mid) ans += query(max(qL, mid), qR, i * 2 + 1, mid, cR);
11
      return ans;
12
13 }
14
  void update(int p, int v, int i = 1, int cL = 0, int cR = n) {
15
       if (cR - cL == 1) {
16
17
           tree[i] = v;
           return;
18
      }
19
20
21
      int mid = (cL + cR) / 2;
      if (p < mid)
22
           update(p, v, i * 2, cL, mid);
23
24
           update(p, v, i * 2 + 1, mid, cR);
25
      tree[i] = tree[i * 2] + tree[i * 2 + 1];
26
27 }
28
void debug(int i = 1, int cL = 0, int cR = n) {
      cerr << "tree[" << cL << "," << cR << ") = " << tree[i];</pre>
30
31
      if (cR - cL > 1) {
32
          int mid = (cL + cR) / 2;
33
           debug(i * 2, cL, mid);
34
           debug(i * 2 + 1, mid, cR);
36
37 }
```

2.6 Range Tree on Trees

```
vector < int > children[N];

int indexInRangeTree[N], startRange[N], endRange[N];
int totId;

void compute_tree_ranges(int v) {
    indexInRangeTree[v] = startRange[v] = totId++;
    for (int w : children[v]) compute_tree_ranges(w);
    endRange[v] = totId;
}

void update_node(int id, int v) { update(indexInRangeTree[id], v); }

long long query_subtree(int id) { return query(startRange[id], endRange[id]); }
```

3 Dynamic Programming

3.1 Knapsack

```
int dp[N + 2][S + 1];

for (int i = N; i >= 1; --i) {
   for (int r = 0; r <= S; ++r) {
      int m = dp[i + 1][r];
      if (r - s[i] >= 0) m = max(m, dp[i + 1][r - s[i]] + v[i]);
      dp[i][r] = m;
}

}
```

3.2 Bitsets

3.3 Travelling Sales Person

```
const int N = 20;
const int INF = 1e9;
3 int n, adj[N][N]; // assume this is given.
4 int dp[1 << N][N]; // dp[x][i] is the shortest 0->i path visiting set bits of x
  int tsp(void) {
      for (int mask = 0; mask < (1 << n); mask++)</pre>
7
           for (int city = 0; city < n; city++) dp[mask][city] = INF;</pre>
8
      dp[1][0] = 0; // 1 represents seen set \{0\}
9
10
11
      int ans = INF;
      for (int mask = 1; mask < (1 << n); mask++) // for every subset of cities seen so far</pre>
12
           for (int cur = 0; cur < n; cur++)</pre>
               if (mask & (1 << cur)) {</pre>
                                                // cur must be one of the cities seen so far
14
                                               // distance travelled so far
                   int cdp = dp[mask][cur];
                   if (mask == (1 << n) - 1) // seen all cities, return to 0
16
                       // unlike the traditional TSP, we don't have to add adj[cur][0]
17
                       // to account for an edge back to vertex 0
                       ans = min(ans, cdp);
19
                   for (int nxt = 0; nxt < n; nxt++)</pre>
20
                       if (!(mask & (1 << nxt))) // try going to a new city
21
                           // new seen set is mask union {nxt}, and we will be at nxt
22
                            // distance incurred to get to this state is now no worse than
23
24
                            // cdp (current distance incurred) + edge from cur to nxt
25
                            dp[mask | (1 << nxt)][nxt] =
                                min(dp[mask | (1 << nxt)][nxt], cdp + adj[cur][nxt]);</pre>
26
               }
27
      return ans;
28
29 }
```

4 Graph Algorithms

4.1 Breath First Search

```
vector < int > edges [N];
  int dist[N];
3 int prev[N];
  void bfs(int start) {
      fill(dist, dist + N, -1);
6
      dist[start] = 0;
      prev[start] = -1;
      queue < int > q;
10
       q.push(start);
11
       while (!q.empty()) {
           int c = q.front();
           q.pop();
           for (int nxt : edges[c]) {
                if (dist[nxt] == -1) {
16
                    dist[nxt] = dist[c] + 1;
17
                    prev[nxt] = c;
18
                    q.push(nxt);
19
               }
20
           }
21
      }
22
23 }
```

4.2 Depth First Search

```
bool seen[N];

void dfs(int u) {
   if (seen[u]) return;
   seen[u] = true;
   for (int v : edges[u]) dfs(v);
}
```

4.3 Bridge Finding

```
vector < int > edges[N];
2 int preorder[N]; // initialise to -1
3 int T = 0;
  int reach[N];
5 vector<pair<int, int>> bridges;
  void dfs(int u, int from = -1) {
      preorder[u] = T++;
8
      reach[u] = preorder[u];
9
10
      for (int v : edges[u])
11
           if (v != from) {
               if (preorder[v] == -1) {
                   dfs(v, u);
14
                   if (reach[v] == preorder[v]) bridges.emplace_back(u, v);
15
16
               reach[u] = min(reach[u], reach[v]);
17
          }
18
19 }
```

4.4 Directed Cycle Detection

```
vector < int > edges [N];
1 int seen[N];
3 int active[N];
5 bool has_cycle(int u) {
      if (seen[u]) return false;
6
      seen[u] = true;
      active[u] = true;
      for (int v : edges[u]) {
9
           if (active[v] || has_cycle(v)) return true;
10
11
      active[u] = false;
12
13
      return false;
14 }
```

4.5 Tree Representation

```
1 \text{ const int } N = 1e6 + 5;
3 vector < int > edges[N];
                              // Parent. -1 for the root.
5 int par[N];
6 vector<int> children[N]; // Your children in the tree.
                              // As an example: size of each subtree.
7 int size[N];
  void constructTree(int c, int cPar = -1) {
      par[c] = cPar;
10
11
      size[c] = 1;
      for (int nxt : edges[c]) {
12
           if (nxt == par[c]) continue;
13
           constructTree(nxt, c);
14
           children[c].push_back(nxt);
15
           size[c] += size[nxt];
16
17
      }
18 }
```

4.6 Binary Lifting

```
const int N = 200010;
const int D = 30; // ceil(log2(10^9))
3 int parent[N][D];
  void precomp() {
      for (int i = 1; i <= n; ++i) cin >> parent[i][0];
      for (int j = 1; j < D; ++j) {
8
          for (int i = 1; i <= n; ++i) parent[i][j] = parent[parent[i][j - 1]][j - 1];</pre>
9
      }
10
11 }
12
int kth_parent(int x, int k) {
      for (int j = 0; j < D; ++j) {
14
           if (k & (1 << j)) x = parent[x][j];</pre>
15
      }
16
17 }
```

4.7 Kosaraju's Algorithm

```
int scc[N];
vector < int > edges[N], edges_r[N];
3 int n, m;
5 bool seen[N], seen_r[N];
6 int postorder[N];
7 int p = 0;
9 void dfs(int u) {
      if (seen[u]) return;
11
      seen[u] = true;
12
      for (int v : edges[u]) dfs(v);
13
      postorder[p++] = u;
14 }
15
void dfs_r(int u, int mark) {
17
      if (seen_r[u]) return;
       seen_r[u] = true;
18
       scc[u] = mark;
19
       for (int v : edges_r[u]) dfs_r(v, mark);
20
21 }
22
23
  int compute_sccs() {
24
       int sccs = 0;
25
      for (int i = 1; i <= n; i++)</pre>
           if (!seen[i]) dfs(i);
26
27
      for (int i = p - 1; i >= 0; i--) {
28
           int u = postorder[i];
29
           if (!seen_r[u]) dfs_r(u, sccs++);
31
32
       return sccs;
33 }
```

4.8 Topological Sort

```
set < int > dag[N]; // edges
2 bool seen_dag[N];
  void compute_topsort(int u, vector<int>& postorder) {
      if (seen_dag[u]) return;
5
      seen_dag[u] = true;
6
      for (int v : dag[u]) compute_topsort(v, postorder);
      postorder.push_back(u);
9 }
10
vector < int > topsort() {
      vector<int> res;
12
      for (int i = 0; i < nsccs; i++) compute_topsort(i, res);</pre>
13
      reverse(res.begin(), res.end());
14
      return res;
15
16 }
```

4.9 Compute SCC DAG

```
int main() {
    cin >> n >> m;
3
```

```
for (int i = 0; i < m; ++i) {</pre>
4
5
           int a, b;
           cin >> a >> b;
6
           edges[a].push_back(b);
           edges_r[b].push_back(a);
9
10
       int nsccs = compute_sccs();
11
       for (int i = 1; i <= n; ++i) {</pre>
12
           for (int j : edges[i]) {
13
                if (scc[i] != scc[j]) dag[scc[i]].insert(scc[j]);
14
15
       }
16
17
18
       vector < int > topo = topsort();
19 }
```

4.10 2-SAT

```
struct TwoSatSolver {
      int n_vars;
3
      int n_vertices;
       vector < int >> adj, adj_t;
      vector < bool > used;
      vector<int> order, comp;
6
      vector < bool > assignment;
8
9
      TwoSatSolver(int _n_vars)
          : n_vars(_n_vars),
10
             n_vertices(2 * n_vars),
             adj(n_vertices),
             adj_t(n_vertices),
             used(n_vertices),
14
             order(),
15
             comp(n_vertices, -1),
16
             assignment(n_vars) {
17
           order.reserve(n_vertices);
19
       void dfs1(int v) {
20
           used[v] = true;
21
           for (int u : adj[v]) {
22
               if (!used[u]) dfs1(u);
23
24
25
           order.push_back(v);
26
27
      void dfs2(int v, int cl) {
28
           comp[v] = cl;
29
           for (int u : adj_t[v]) {
30
31
               if (comp[u] == -1) dfs2(u, cl);
32
33
34
       bool solve_2SAT() {
35
           order.clear();
36
           used.assign(n_vertices, false);
37
           for (int i = 0; i < n_vertices; ++i) {</pre>
38
39
               if (!used[i]) dfs1(i);
           }
40
41
           comp.assign(n_vertices, -1);
42
           for (int i = 0, j = 0; i < n_vertices; ++i) {</pre>
43
```

```
int v = order[n_vertices - i - 1];
44
45
               if (comp[v] == -1) dfs2(v, j++);
46
47
48
           assignment.assign(n_vars, false);
           for (int i = 0; i < n_vertices; i += 2) {</pre>
49
               if (comp[i] == comp[i + 1]) return false;
50
               assignment[i / 2] = comp[i] > comp[i + 1];
51
           return true;
53
      }
54
      void add_disjunction(int a, bool na, int b, bool nb) {
56
           // na and nb signify whether a and b are to be negated
57
58
           a = 2 * a ^ na;
           b = 2 * b ^nb;
           int neg_a = a ^ 1;
60
           int neg_b = b ^ 1;
61
           adj[neg_a].push_back(b);
62
63
           adj[neg_b].push_back(a);
           adj_t[b].push_back(neg_a);
64
           adj_t[a].push_back(neg_b);
65
      }
66
67 };
```

4.11 Kruskal's Algorithm

```
struct edge {
     int u, v, w;
3 };
4 bool operator < (const edge & a, const edge & b) { return a.w < b.w; }
6 edge edges[N];
                             // union-find root with path compression
7 int root(int u);
8 void join(int u, int v); // union-find join with size heuristic
  int mst() {
      sort(edges, edges + m); // sort by increasing weight
12
      int total_weight = 0;
      for (int i = 0; i < m; i++) {</pre>
13
           edge& e = edges[i];
14
          if (root(e.u) != root(e.v)) {
15
16
               total_weight += e.w;
               join(e.u, e.v);
          }
18
19
      }
      return total_weight;
20
21 }
```

4.12 Prim's Algorithm

```
typedef pair < int, int > ii;

vector < ii > edges[N]; // pairs of (weight, v)
bool in_tree[N];
priority_queue < ii, vector < ii > pq;

int mst() {
   int total_weight = 0;
   in_tree[0] = true;
```

```
for (auto edge : edges[0]) pq.emplace(edge.first, edge.second);
10
11
      while (!pq.empty()) {
          auto edge = pq.top();
12
          pq.pop();
          if (in_tree[edge.second]) continue;
14
           in_tree[edge.second] = true;
15
           total_weight += edge.first;
16
           for (auto edge : edges[edge.second]) pq.emplace(edge.first, edge.second);
17
18
      return total_weight;
19
20 }
```

4.13 Shortest Path Algorithms

4.13.1 Dijkstra's Algorithm

```
#include <bits/stdc++.h>
2 using namespace std;
3
4 typedef long long ll;
5 typedef pair<11, int> edge; // (distance, vertex)
6 const int N = 100100;
8 vector < edge > edges [N];
9 11 dist[N];
10 bool seen[N];
priority_queue <edge, vector <edge>, greater <edge>> pq;
  void dijkstra(int s) {
13
      fill(seen, seen + N, false);
14
      pq.push(edge(0, s));
      while (!pq.empty()) {
16
17
           edge cur = pq.top();
18
          pq.pop();
          int v = cur.second;
19
          11 d = cur.first;
20
           if (seen[v]) continue;
21
22
           dist[v] = d;
23
24
           seen[v] = true;
25
           for (edge nxt : edges[v]) {
26
               int u = nxt.second;
27
               ll weight = nxt.first;
28
               if (!seen[u]) pq.push(edge(d + weight, u));
29
           }
30
      }
31
32 }
```

4.13.2 Bellman Ford

```
const long long INF = 1e9 + 7;

struct edge {
   int u, v, w; // u -> v of weight w
   edge(int _u, int _v, int _w) : u(_u), v(_v), w(_w) {}
};

int n, m, cycleStart;
vector<long long> dist;
```

```
10 vector<int> parent;
vector<edge> edges;
12 set <int > visited;
14 bool relax() {
      bool relaxed = false;
      for (edge e : edges) {
16
           if (dist[e.u] != INF && dist[e.v] > dist[e.u] + e.w) {
17
               relaxed = true;
18
               dist[e.v] = dist[e.u] + e.w;
19
               parent[e.v] = e.u;
20
               cycleStart = e.v;
21
22
               visited.insert(e.u);
           }
23
      }
25
      return relaxed;
26 }
27
  bool bellman_ford(int start) {
28
      fill(dist.begin(), dist.end(), INF);
29
      fill(parent.begin(), parent.end(), -1);
30
      dist[start] = 0;
31
      for (int i = 0; i < n - 1; i++)</pre>
32
           if (!relax()) break;
33
34
35
      return relax();
36 }
```

4.13.3 Finding Negative Cycles

```
int main() {
       cin >> n >> m;
       for (int i = 0; i < m; ++i) {</pre>
3
           int a, b, c;
4
           cin >> a >> b >> c;
           edges.push_back({a, b, c});
6
       dist.resize(n);
9
10
       parent.resize(n);
11
       bool res = false;
12
       for (int i = 0; i < n; ++i) {</pre>
13
14
           if (visited.find(i) == visited.end() && bellman_ford((i))) {
15
                res = true;
                break;
16
           }
17
18
19
       if (!res) cout << "NO\n";
20
21
       else {
           cout << "YES\n";</pre>
22
23
           for (int i = 0; i < n; ++i) cycleStart = parent[cycleStart];</pre>
24
25
           vector < int > cycle;
26
27
           for (int v = cycleStart;; v = parent[v]) {
28
                cycle.push_back(v);
                if (v == cycleStart && cycle.size() > 1) break;
29
30
31
           reverse(cycle.begin(), cycle.end());
32
```

4.13.4 Floyd Warshall

```
for (int u = 0; u < n; ++u)
    for (int v = 0; v < n; ++v) dist[u][v] = INF;

for (edge e : edges) dist[e.u][e.v] = e.w;

for (int u = 0; u < n; ++u) dist[u][u] = 0;

for (int i = 0; i < n; i++)
    for (int u = 0; u < n; u++)
    for (int v = 0; v < n; v++) dist[u][v] = min(dist[u][v], dist[u][i] + dist[i][v]);</pre>
```

5 Flow Networks

5.1 Dinic's Algorithm

```
1 typedef long long 11;
3 const int INF = 1e9 + 7;
  struct FlowNetwork {
      int n;
      vector < vector < ll >> adjMat, adjList;
      // level[v] stores dist from s to v
      // uptochild[v] stores first non-useless child.
9
      vector < int > level, uptochild;
11
      FlowNetwork(int _n) : n(_n) {
12
13
           // adjacency matrix is zero-initialised
           adjMat.resize(n);
14
           for (int i = 0; i < n; i++) adjMat[i].resize(n);</pre>
           adjList.resize(n);
16
           level.resize(n);
17
           uptochild.resize(n);
18
      }
19
20
       void add_edge(int u, int v, ll c) {
21
           // add to any existing edge without overwriting
22
           adjMat[u][v] += c;
23
           adjList[u].push_back(v);
24
           adjList[v].push_back(u);
25
26
27
      void flow_edge(int u, int v, ll c) {
28
           adjMat[u][v] -= c;
29
           adjMat[v][u] += c;
30
31
32
33
       // constructs the level graph and returns whether the sink is still reachable
       bool bfs(int s, int t) {
34
           fill(level.begin(), level.end(), -1);
35
           queue < int > q;
36
           q.push(s);
37
           level[s] = 0;
38
           while (!q.empty()) {
```

```
int u = q.front();
41
               q.pop();
               uptochild[u] = 0; // reset uptochild
42
               for (int v : adjList[u])
43
                   if (adjMat[u][v] > 0) {
44
                        if (level[v] != -1) // already seen
45
                            continue;
46
                        level[v] = level[u] + 1;
47
                        q.push(v);
48
49
           }
50
           return level[t] != -1;
51
      }
52
54
      // finds an augmenting path with up to f flow.
55
      11 augment(int u, int t, ll f) {
           if (u == t) return f; // base case.
56
           // note the reference here! we increment uptochild[u], i.e. walk through u's neighbours
57
           // until we find a child that we can flow through
58
           for (int& i = uptochild[u]; i < adjList[u].size(); i++) {</pre>
59
               int v = adjList[u][i];
60
               if (adjMat[u][v] > 0) {
61
                   // ignore edges not in the BFS tree.
62
                   if (level[v] != level[u] + 1) continue;
63
                   // revised flow is constrained also by this edge
64
                   ll rf = augment(v, t, min(f, adjMat[u][v]));
65
                   // found a child we can flow through!
66
67
                   if (rf > 0) {
68
                        flow_edge(u, v, rf);
                        return rf;
69
                   }
70
               }
71
           }
72
           level[u] = -1;
73
74
           return 0;
75
76
      11 dinic(int s, int t) {
77
78
           11 \text{ res} = 0;
79
           while (bfs(s, t))
               for (11 x = augment(s, t, INF); x; x = augment(s, t, INF)) res += x;
81
           return res;
      }
82
83 };
```

5.2 Min-cut

```
void check_reach(int u, vector < bool > & seen) {
      if (seen[u]) return;
      seen[u] = true;
      for (int v : adjList[u])
4
           if (adjMat[u][v] > 0) check_reach(v, seen);
5
6 }
8 vector<pair<int, int>> min_cut(int s, int t) {
9
      ll value = dinic(s, t);
10
11
      vector < bool > seen(n, false);
      check_reach(s, seen);
13
      vector<pair<int, int>> ans;
14
      for (int u = 0; u < n; u++) {</pre>
15
```

```
if (!seen[u]) continue;
for (int v : adjList[u])

if (!seen[v] && !is_virtual[u][v]) // need to record this in add_edge()

ans.emplace_back(u, v);

return ans;
}
```

6 Mathematics

6.1 Fast Exponentiation

```
const int MOD = 1e9 + 7;
typedef long long ll;

ll modpow(ll x, ll n, int m) {
    if (n == 0) return 1;

ll a = modpow(x, n / 2, m);
    a = a * a % m;
    if (n % 2 == 1) a = a * x % m;
    return a;
}
```

6.2 Primality Testing

```
bool isprime(int x) {
   if (x < 2) return false;

for (int f = 2; f * f <= x; f++)
        if (x % f == 0) return false;

return true;
}</pre>
```

6.3 Prime Factorisation

6.4 Sieve of Eratosthenes

```
bool marked[N + 1];
vector<int> primefactorization[N + 1];

for (int i = 2; i <= N; i++) {
   if (!marked[i]) {</pre>
```

```
primefactorization[i].push_back(i);
           for (int j = 2 * i; j <= N; j += i) {</pre>
                marked[j] = true;
               int tmp = j;
9
                while (tmp % i == 0) {
10
                    primefactorization[j].push_back(i);
11
                    tmp /= i;
12
                }
13
           }
14
       }
15
16 }
```

6.5 GCD

```
int gcd(int a, int b) { return b ? gcd(b, a % b) : a; }
```

6.6 LCM

```
int lcm(int a, int b) { return a * b / gcd(a, b); }
```

6.7 Extended Euclidean Algorithm

```
int euclidean_algorithm(int a, int b, int& x, int& y) {
      if (a == 0) {
2
          x = 0;
3
          y = 1;
          return b;
      }
6
      int x1, y1;
      int d = euclidean_algorithm(b % a, a, x1, y1);
      x = y1 - (b / a) * x1;
9
      y = x1;
10
11
      return d;
12 }
```

6.8 Matrices

```
struct Matrix {
       vector < vector < long long >> v;
3
4
       \texttt{Matrix(int } \_\texttt{n)} \; : \; \texttt{n(\_n)} \; \{
5
            v.resize(n);
6
            for (int i = 0; i < n; i++)</pre>
                 for (int j = 0; j < n; j++) v[i].push_back(0);</pre>
9
10
       Matrix operator*(const Matrix &o) const {
11
            Matrix res(n);
            for (int i = 0; i < n; i++)</pre>
13
                 for (int j = 0; j < n; j++)</pre>
14
                      for (int k = 0; k < n; k++) res.v[i][j] += v[i][k] * o.v[k][j];
15
            return res;
16
17
18
       static Matrix getIdentity(int n) {
19
            Matrix res(n);
```

```
for (int i = 0; i < n; i++) res.v[i][i] = 1;</pre>
21
22
           return res;
23
24
       Matrix operator^(long long k) const {
25
           Matrix res = Matrix::getIdentity(n);
26
           Matrix a = *this;
27
           while (k) {
28
                if (k & 1) res = res * a;
29
                a = a * a;
30
               k /= 2;
31
           }
32
           return res;
       }
34
35 };
```

6.9 Combinations

```
1 typedef long long 11;
3 const int N = 1001001;
4 const int MOD = 1e9 + 7;
5 ll f[N + 1];
6 ll inv[N + 1];
7 ll modpow(ll a, ll b, int c); // as earlier
9 ll choose(ll n, ll r) { return ((f[n] * inv[r]) % MOD * inv[n - r]) % MOD; }
10
int main() {
      f[0] = 1;
12
      for (int i = 1; i < N; i++) f[i] = (i * f[i - 1]) % MOD;</pre>
13
14
15
      inv[N] = modpow(f[N], MOD - 2, MOD);
      for (int i = N; i >= 1; --i) inv[i - 1] = (inv[i] * i) % MOD;
16
17 }
```

7 Computational Geometry

7.1 Cross Product

```
const double EPS = 1e-8;
typedef pair < double, double > pt;

#define x first

#define y second

pt operator - (pt a, pt b) { return pt(a.x - b.x, a.y - b.y); }

bool zero(double x) { return fabs(x) <= EPS; }

double cross(pt a, pt b) { return a.x * b.y - a.y * b.x; }

// true if left or straight
// sometimes useful to instead return an int
// -1, 0 or 1: the sign of the cross product
bool ccw(pt a, pt b, pt c) { return cross(b - a, c - a) >= 0; }
```

7.2 Three Points Collinear

7.3 Segment-Segment Intersection

```
typedef pair<pt, pt> seg;
2
3 bool collinear(seg ab, seg cd) { // all four points collinear
      pt a = ab.first, b = ab.second, c = cd.first, d = cd.second;
      return zero(cross(b - a, c - a)) && zero(cross(b - a, d - a));
6 }
8 double sq(double t) { return t * t; }
  double dist(pt p, pt q) { return sqrt(sq(p.x - q.x) + sq(p.y - q.y)); }
11
12 bool intersect(seg ab, seg cd) {
      pt a = ab.first, b = ab.second, c = cd.first, d = cd.second;
13
14
      if (collinear(ab, cd)) {
15
          double maxDist =
16
              max({dist(a, b), dist(a, c), dist(a, d), dist(b, c), dist(b, d), dist(c, d)});
17
          return maxDist < dist(a, b) + dist(c, d) + EPS;</pre>
18
19
20
      return ccw(a, b, c) != ccw(a, b, d) && ccw(c, d, a) != ccw(c, d, b);
21
22 }
```

7.4 Polygon Area (Trapezoidal Rule)

```
double area(vector<pt> pts) {
    double res = 0;
    int n = pts.size();
    for (int i = 0; i < n; i++) {
        res += (pts[i].y + pts[(i + 1) % n].y) * (pts[(i + 1) % n].x - pts[i].x);
    }
    return res / 2.0;
}</pre>
```

7.5 Polygon Area (Cross Product)

```
double area(vector < pt > pts) {
    double res = 0;
    int n = pts.size();
    for (int i = 1; i < n - 1; i++) {
        // i = 0 and i = n-1 are degenerate triangles, OK to omit
        // e.g. if i = 1 is ABC, and i = 2 is ACD, then i = 0 is AAB
        res += cross(pts[i] - pts[0], pts[i + 1] - pts[0]);
    }
    return res / 2.0;
}</pre>
```

7.6 Convex Hull

```
vector<pt> half_hull(vector<pt> pts) {
      vector <pt> res;
2
3
      for (int i = 0; i < pts.size(); i++) {</pre>
           // ccw means we have a left turn; we don't want that
5
           while (res.size() >= 2 && ccw(pts[i], res[res.size() - 1], res[res.size() - 2])) {
6
               res.pop_back();
           }
          res.push_back(pts[i]);
8
      }
9
10
      return res;
11
  }
12
  vector<pt> convex_hull(vector<pt> pts) {
13
      sort(pts.begin(), pts.end());
14
      vector < pt > top = half_hull(pts);
16
17
      reverse(pts.begin(), pts.end());
      vector <pt> bottom = half_hull(pts);
18
19
      top.pop_back();
20
      bottom.pop_back();
21
      vector<pt> res(top.begin(), top.end());
22
23
      res.insert(res.end(), bottom.begin(), bottom.end());
24
      return res;
25 }
```

7.7 Half Plane Intersection

```
1 typedef pair < double , double > pt;
3 struct line {
      double a, b, c;
4
5 };
7 struct half_plane {
      line 1;
      bool neg; // is the inequality <= or >=
9
10 };
11
12 const double EPS = 1e-8;
14 pt intersect(line f, line g) {
      double d = f.a * g.b - f.b * g.a;
15
      double y = (f.a * g.c - f.c * g.a) / (f.b * g.a - f.a * g.b);
16
      double x = (f.c * g.b - f.b * g.c) / (f.b * g.a - f.a * g.b);
17
18
      return pt(x, y);
19 }
20
21 bool in_half_plane(half_plane hp, pt q) {
      if (hp.neg) return hp.l.a * q.x + hp.l.b * q.y + hp.l.c <= EPS;</pre>
22
      else return hp.l.a * q.x + hp.l.b * q.y + hp.l.c >= -EPS;
23
24 }
25
26
  vector<pt> intersect_half_planes(vector<half_plane> half_planes) {
      int n = half_planes.size();
27
      vector <pt> pts;
28
      for (int i = 0; i < n; i++) {</pre>
29
           for (int j = i + 1; j < n; j++) {
30
               pt p = intersect(half_planes[i].1, half_planes[j].1);
31
               bool fail = false;
```

```
for (int k = 0; k < n; k++)</pre>
33
                      if (!in_half_plane(half_planes[k], p)) fail = true;
34
                 if (!fail) pts.push_back(p);
35
           }
36
       }
37
38
       vector < pt > res = pts;
if (pts.size() > 2) pts = convex_hull(res);
39
40
       return pts;
41
42 }
```