

WAECO Intro to AECs talk part 1

Based largely on material learned in Rami Grossberg's Model Theory II course.

1st order logic: All familiar with theory T in language L ,
study $(\text{Mod}(T), \leq)$
elementary substructure.

Call these elementary classes.

Lots of useful tools:

- compactness
- LS theorems (\uparrow & \downarrow)
- Chain theorems ($\langle M_i : i < \delta \rangle$ cts line,
then $\bigcup M_i \in \text{Mod}(T)$,
 $M_j \not\leq \bigcup M_i$, and
 $M_j \leq N \vee j \Rightarrow \bigcup M_i \not\leq N$)

But there are other logics & classes of interest, e.g.

~~(*)~~ Infinitary logic: Models of some $\varphi \in L_{\omega, \omega}$ perhaps
 $(\text{Mod}(\varphi), \leq_F)$ where $\varphi \in L_{\lambda^+, \omega}$

(i.e. φ allows \wedge and \vee of size $< \lambda^+$), $|L| \leq \lambda$,
and F a fragment of $L_{\lambda^+, \omega}$ containing φ with $|F| \leq \lambda$.

(i.e. $F \subseteq \text{Form}(L_{\lambda^+, \omega})$ where

- F closed under f.o. connectives $\wedge, \vee, \neg, \forall x, \exists y$
- If $\varphi(\underline{x}) \in F$, $\underline{\tau}(\underline{y}) = \langle \tau_1(\underline{y}), \dots, \tau_k(\underline{y}) \rangle$
where τ_i terms, then $\varphi(\underline{\tau}(\underline{y})) \in F$
- F is closed under subformulas.

so 'F is well enough closed'

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$$\textcircled{X} \quad \text{ECC}(\Gamma, \Gamma) = (\{\text{M} \models \Gamma : \text{Models all } q \in \Gamma\}, \leq).$$

In these, compactness fails - so if we move to a more general framework that also covers these examples, we need to weaken our assumptions.

Definition: An abstract class (AC) is a pair $\underline{K} = (K, \leq_K)$ where

- K is a class of models in some language $L(K)$
- \leq_K is a partial ordering on K refining submodel, i.e.

$$M \leq_K N \Rightarrow M \subseteq N$$

- 1st iso axiom: K is closed under isomorphism: i.e.

if $M \in K$, $N \cong M$, then $N \in K$.

- 2nd iso axiom: \leq_K respects isomorphism, i.e.

given $M_1 \leq_K M_2$ and $M_2 \stackrel{f}{\cong} N$

then $f[M_1] \in K$ and $f[M_1] \leq_K N$

$$\begin{array}{ccc} M_2 & \xrightarrow{\cong f} & N \\ \vee^K & & \vee^K \\ M_1 & \xrightarrow{f} & f[M_1] \end{array}$$

These assumptions are still too weak to capture much of the theory of the last two examples. So we have ...

Definition: An abstract elementary class (AEC) is an AC $\underline{K} = (K, \leq_K)$ that also satisfies

- Cohesiveness: If $M_1, M_2, M_3 \in K$ and $M_1 \leq_K M_3$,
 $M_2 \leq_K M_3$, and $M_1 \subseteq M_2$, then $M_1 \leq_K M_2$.

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- TV chain axioms: If $\langle M_i : i < \delta \rangle$ a \leq_K -inc cts \leq_K -chain of models in K where δ is limit, then
 - $\bigcup_{i < \delta} M_i \in K$
 - $M_j \leq_K \bigcup_{i < \delta} M_i$ for all $j < \delta$
 - If $M_j \leq_K N$ for all j , then $\bigcup_{i < \delta} M_i \leq_K N$.
- LS Axion: There is an infinite cardinal $LS(K)$ such that $\forall M \in K \forall A \subseteq |M| \exists N \in K$ s.t. $N \leq_K M$, $A \subseteq |N|$, and $|N| \leq |A| + LS(K)$.
We denote the least such cardinal by $LS(K)$.
- (sometimes) $\forall M \in K |M| \geq LS(K)$.

Examples:

- $(Mod(T), \leq)$, $LS(K) = |L(T)| + \aleph_0$
- $(Mod(\varphi), \leq_F)$ where $\varphi \in L_{\lambda^+, \omega}$, $|L|, |F| \leq \lambda$
then $LS(K) = \lambda$
- $EC(T, T)$, $LS(K) = |L(T)| + \aleph_0$
- Lots of specific algebraic examples:
 - groups with $\leq_K = \subseteq$
or $\leq_K =$ 'pure subgroup'
$$\left(\text{def } H \leq_K G \iff \forall n \in \omega \forall a \in H \left[G \models \exists x (a = x^n) \Rightarrow H \models \exists x (a = x^n) \right] \right)$$
 - locally finite groups
 - nil rings ($\forall n \in \mathbb{N} \exists a \in R \forall x \in R^n = 0$)
 - modules over a ring with $\leq_K = \subseteq$, $LS(K) = |R| + \aleph_0$
 - Well orders of type $\leq \lambda^+$ with $\leq_K =$ 'initial segment', $LS(K) = \lambda$

- modules over a ring with $\leq_K =$ 'pure submodel', $LS(K) = |R| + \aleph_0$

- Non-examples:
- well orders with $\leq_K = \leq$ (chain breaks \dots)
 - well orders with \leq_K initial segment (LS axiom breaks)
not well founded

Notation:

- Given λ cardinal, $K_\lambda = \{M \in K : |M| = \lambda\}$
- $I(K, \lambda) = |K_\lambda|_{\leq} | (\leq 2^{|L(T)| + \aleph_0 + \lambda})$

Question: What can $I(K, \lambda)$ look like?

Note $\{ \lambda \geq LS(K) : I(K, \lambda) \neq 0 \} = [LS(K), *)$ for some $* \in [LS(K), \infty)$
or $* = \infty$
by the LS axiom.

$* < \infty$ is possible, but

Fact (Shelah): If $* \geq \beth_{(2^{LS(K)})^+}$, then $* = \infty$.

Fact: There are examples of $* = \beth_\alpha$ for all $\alpha < (2^{LS(K)})^+$.

Considering Morley's / Shelah's categoricity theorems,

Question: what can $\{ \lambda : I(K, \lambda) = 1 \}$ look like?

Big open problem!

Shelah's categoricity conjecture: If $\mu_0 = \beth_{(2^{LS(K)})^+}$ and $\exists \lambda_0 \geq \mu_0$

such that K is categorical in λ_0 , then K is categorical in all $\lambda \geq \mu_0$.

or maybe use a different bound for μ_0 in terms of $LS(K)$.

Lifting basic notions to an AEC setting:

Rather than elementary embeddings ...

Definition: Given $M, N \in K$, $f: M \rightarrow N$ is a K -embedding iff
 $\exists N^* \leq_K N$ such that $f: M \cong N^*$ (i.e. $f: M \cong f[M] \leq_K N$)

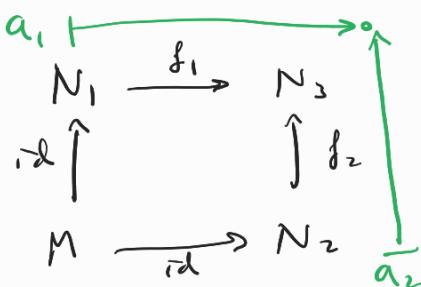
Rather than types ...

Definition: Kan AEC.

- $K^3 = \{(M, N, a) : M, N \in K, a \in N, M \leq N\}$.

- Given $(M_\ell, N_\ell, a_\ell) \in K^3$ for $\ell = 1, 2$

$(M_1, N_1, a_1) \sim (M_2, N_2, a_2)$ iff $M_1 = M_2$ and
 $\exists N_3 \in K \exists f_\ell: N_\ell \xrightarrow{M_1} N_3$ K -embeddings such that
 $f_1(a_1) = f_2(a_2)$



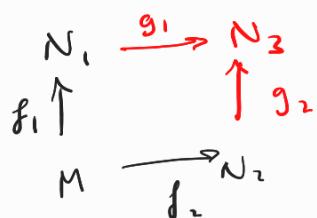
- \approx is the transitive closure of \sim

- Given $(M, N, a) \in K^3$, $\text{gfp}(a/M, N) = [(M, N, a)]_{\approx}$.

In many cases, $\sim = \approx$.

Definition: Suppose K is an AC. We say K has the amalgamation property if whenever we have $f_\ell: M \rightarrow N_\ell$ for $\ell = 1, 2$,

$\exists N_3 \exists g_\ell: N_\ell \rightarrow N_3$ such that $g_1 \circ f_1 = g_2 \circ f_2$

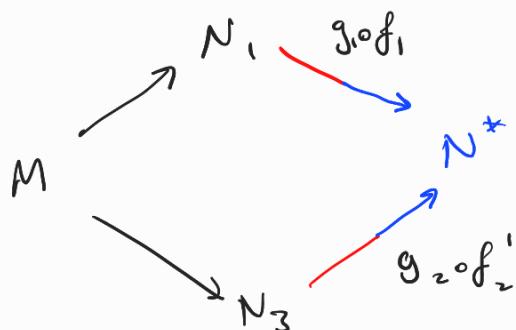
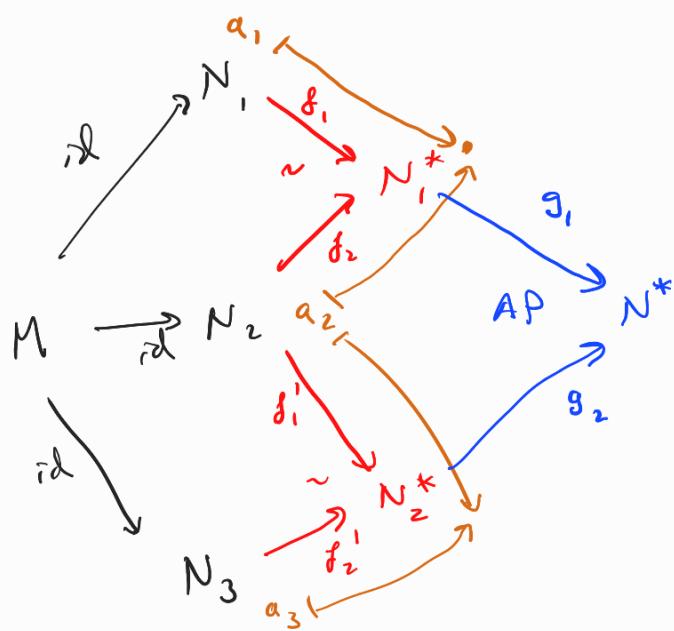


Proposition: If K has AP, then $\sim = \approx$.

proof: Suppose $(M_1, N_1, a_1) \sim (M_2, N_2, a_2) \sim (M_3, N_3, a_3)$.

Let $M = M_1 = M_2 = M_3$.

sketch:



$$g_1 \circ f_1(a_1) = g_1 \circ f_2(a_2) = g_2 \circ f_1'(a_2) = g_2 \circ f_2'(a_3)$$

$\sim \text{def}$ AP $\sim \text{def}$

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Definition: Let K be an $A \models C$, $M \in K$.

$$S(M) = \{g_{tp}(a/M, N) : M \subseteq N, a \in |N|\}$$

Fact: $|S(M)| \leq 2^{|M|}$

(wlog $|N| = |M|$ by LS Ax & def of \sim ,

so $|M|$ choices for $|N| \setminus |M|$,
 $\leq 2^{|M|}$ choices for interpreting $L(M)$ in N ,)
 $\leq |M|$ choices for A .

Definition: If $\rho = \text{gfp}(\alpha / M, N)$, we say ρ is realised in N .

Definition: M is λ -saturated if $\forall M_0 \subseteq M$ st. $\|M_0\| < \lambda$,
 $\forall \rho \in S(M_0)$, ρ is realised in M .

M is saturated if $\|M\|$ -saturated.

These still behave as we expect.

Proposition: Suppose K is an AEC, $\lambda > LS(K)$, $K_{<\lambda}$ has AP,
and $M \in K_\lambda$. Then

M is saturated $\iff \forall N \subseteq M \ \forall N' \geq_K N$ st. $\|N'\| < \|M\|$
 $\exists f: N' \xrightarrow[N]{} M$
 M is model homogeneous.

proof ideas: Like 1st order, but be careful.

$m\text{-h} \iff \forall \mu < \lambda \quad M \text{ is } m\text{-h for } \|N\| = \|N'\| = \mu$
sat
is the 1st order,
but there is no notion of (M, N) -maps in AECs,
(all sets are models), so we have all maps between
'larger' models.

Other important notions:

Definition: We say an AC K has the joint embedding property (JEP) if
 $\forall M_1, M_2 \in K \ \exists N \in K \ \exists f_i: M_i \rightarrow N$.



Fact: $(\text{Mod}(T), \leq)$ has JEP $\Leftrightarrow T$ is complete.

(\Rightarrow if not complete, get $M_1 \not\models M_2$

\Leftarrow use monster model, or large saturated model, or N).

Definition: We say an AC K has no maximal models (NMM) if $\forall M \in K \exists N \in K \quad M \not\subseteq_K N$.

Fact: If K has AP, JEP, NMM, then K has monster models similar to 1st order logic, i.e. $\forall \mu > \text{LS}(K) \quad \exists M^* \in K$ st.

- All models of card $< \mu$ embed into M^*
- M^* is μ -m.h. (i.e. m-h holds for $\|N\| = \|N'\| < \mu$)
- $\forall N, N' \subseteq M^*$ st. $\|N\| = \|N'\| < \mu, \forall f: N \cong N'$
 f extends to an automorphism of M^* .

Note: Existence of such M^* recovers AP, JEP, NMM

proof: JEP, NMM easy

AP:

$$\begin{array}{ccc} N_1 & & N_2 \\ \downarrow id & M & \uparrow f \\ \end{array}$$

$M, N_1 \subseteq M^*$ wlog

get $g: N_2 \cong N_2$ extending f . By m-h,

$$\exists h: N_2 \xrightarrow{M} M^*$$

then

$$N_1 \xrightarrow{id} M^*$$

$$\uparrow id \quad \uparrow h \circ g^{-1}$$

$$M \xrightarrow{f} N_2$$



Note: - With a monster model, types $gtp(a/M, N)$ over small sets become orbits under automorphisms fixing M

- The notion of types agrees with the first order notion: if $a \in N \models_K M \vdash$,
 $tp(a/M, N) = tp(b/M_1, N_1) \Leftrightarrow gtp(a_1/M_1, N_1) = gtp(b_1/M_2, N_2)$.

END OF PART I