

## Optimal Waiting Time

Consider two lanes 1 and 2, where at the current time  $t = 0$ , some vehicle  $j = 0$  has just been completed at lane 1. We know that another vehicle  $j = 2$  will arrive to lane somewhere in the future at  $r = r_1$ , which is a random variable. Furthermore, some other vehicle  $j = 2$  is available (so  $r_2 \leq 0$ ) for service on lane 2. Assume that all processing times are  $p_j = p = 1$ . We need to choose how long we wait for vehicle  $j = 1$ , before switching to lane 2. Let  $y_j$  denote the scheduled starting (crossing) time of the vehicles. The goal is to minimize

$$\sum_j y_j. \quad (1)$$

Let the waiting time be denoted as  $x \geq 0$ , then our problem may be stated as

$$\min_{x \geq 0} \mathbb{E} [\Pi(x, r)], \quad (2)$$

where the value function  $\Pi$  is given by

$$\Pi(x, r) = \begin{cases} r + (r + 1 + s) & \text{if } x \geq r, \\ \max(x, s) + \max(r, \max(x, s) + 1 + s) & \text{if } x < r. \end{cases} \quad (3)$$

We calculate

$$\begin{aligned} \mathbb{E}[\Pi(x, R)] &= \int_0^\infty \Pi(x, r) dF(r) \\ &= \int_0^x 2r + s + 1 dF(r) + \int_x^\infty \max(r, \max(x, s) + 1 + s) dF(r) \\ &= (s + 1)F(x) + 2 \int_0^x r dF(r) \\ &\quad + \max(x, s)(1 - F(x)) \\ &\quad + (\max(x, s) + s + 1) (F(\max(x, s) + s + 1) - F(x)) \\ &\quad + \int_{\max(x, s) + s + 1}^\infty r dF(r), \end{aligned}$$

where the part involving  $\max(x, s)$  can be verified by considering the cases  $x < s$  and  $x > s$  separately.

## Exponential interarrival times

We can now optimize this expression as function of  $x$ . For example, assume that  $R \sim \text{Exp}(\lambda)$ , with  $F(r) = 1 - e^{-\lambda r}$  for  $r \geq 0$ . Using the fact that

$$\int_A^B r dF(r) = \left(-B - \frac{1}{\lambda}\right) e^{-\lambda B} - \left(-A - \frac{1}{\lambda}\right) e^{-\lambda A}, \quad (4)$$

we obtain the explicit expression

$$\begin{aligned}
\mathbb{E}[\Pi(x, R)] &= (s+1)(1 - e^{-\lambda x}) \\
&\quad + 2 \left( \frac{1}{\lambda} - \left( x + \frac{1}{\lambda} \right) e^{-\lambda x} \right) \\
&\quad + \max(x, s) e^{-\lambda x} \\
&\quad + (\max(x, s) + s + 1) \left( -e^{-\lambda x} + e^{-\lambda(\max(x, s) + s + 1)} \right) \\
&\quad + \left( \max(x, s) + s + 1 + \frac{1}{\lambda} \right) e^{-\lambda(\max(x, s) + s + 1)} \\
&= s + 1 + \frac{2}{\lambda} \\
&\quad + \left( \max(x, s) - s - 1 - \frac{2}{\lambda} \right) e^{-\lambda x} \\
&\quad - 2x e^{-\lambda x} \\
&\quad + \left( 2(\max(x, s) + s + 1) + \frac{1}{\lambda} \right) e^{-\lambda(\max(x, s) + s + 1)}.
\end{aligned}$$

For  $x < s$ , this simplifies to

$$\mathbb{E}[\Pi(x, R)] = s + 1 + \frac{2}{\lambda} - \left( 1 + \frac{2}{\lambda} \right) e^{-\lambda x} - 2x e^{-\lambda x} + \left( 4s + 2 + \frac{1}{\lambda} \right) e^{-\lambda(2s+1)},$$

with derivative

$$(\lambda + 2\lambda x) e^{-\lambda x}.$$

For  $x > s$ , this simplifies to

$$\mathbb{E}[\Pi(x, R)] = s + 1 + \frac{2}{\lambda} + \left( x - s - 1 - \frac{2}{\lambda} \right) e^{-\lambda x} - 2x e^{-\lambda x} + \left( 2x + 2s + 2 + \frac{1}{\lambda} \right) e^{-\lambda(x+s+1)},$$

with derivative

$$(\lambda x + \lambda s + \lambda - 3) e^{-\lambda x} + (1 - 2\lambda x - 2\lambda s - 2\lambda) e^{-\lambda(x+s+1)}.$$

We can probably show that these two derivatives are always negative and positive, respectively, so that the minimum is always achieved at  $x = s$ .

## Simulation

Simulation verifies that  $x = s$  is indeed optimal, see `wait.py`.