

Formal Languages

An introduction

ALPHABETS, STRINGS AND LANGUAGES

An alphabet is a set of symbols.

ex: $\{ a, b, c, d, \dots, z \}$
 $\{ 0, 1 \}$

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- A string (word) is a **concatenation** of symbols of an alphabet.

ex: abc

 jfhI

 rnl

- The length of a string is the number of symbols in the string.

ALPHABETS, STRINGS AND LANGUAGES

A (formal) language is a set of strings.

ALPHABETS, STRINGS AND LANGUAGES

- The empty string, λ , is the string with no symbols
- A Capital Letter is used to denote an alphabet.

ex: $A = \{ a, b, c, d, \dots, z \}$

$B = \{ 0, 1 \}$

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- **Some Notations**

- W notation for the set of all possible strings that can be formed out of a given alphabet.
- W_k represents a set of strings with length k , where length is equal to the number of symbols.

ex: $P = \{ 0, 1, a, b, c, d, \text{else}, \text{for} \}$

$\text{length}(\text{elsefor}) = 2$

- Given an alphabet:

$$W_k = \{ \varphi \mid \varphi \text{ is a string with length} = k \}$$

$$W_0 = \{ \lambda \}$$

$$W = \bigcup W_k = W_0 \cup W_1 \cup W_2 \cup W_3 \cup \dots$$

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- **CONCATENATION**

(“multiplication” of strings)

- The concatenation of $\omega = \omega_1\omega_2\omega_3 \dots \omega_n$ and $\varphi = \varphi_1\varphi_2\varphi_3 \dots \varphi_n$ is

$$\omega\varphi = \omega_1\omega_2\omega_3 \dots \omega_n \varphi_1 \varphi_2 \varphi_3 \dots \varphi_n = \omega\varphi$$

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- **OPERATIONS WITH LANGUAGES**

1. Union
2. Intersection
3. Concatenation
4. Complement (with reference to W)

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Exercise:

Let $L1 = \{ a, ab \}$ be a language

Let $L2 = \{ c, bc \}$ be a language

Union: $L1 \cup L2 = \{ a, ab, c, bc \}$

Intersection: $L1 \cap L2 = \{ \} = \phi$

ALPHABETS, STRINGS AND LANGUAGES

Exercise:

Let $L1 = \{ a, ab \}$ be a language

Let $L2 = \{ c, bc \}$ be a language

Concatenation:

$$L1 \cdot L2 = \{ ac, abc, abbc \}$$

$$L2 \cdot L1 = \{ ca, cab, bca, bcab \}$$

$$L1 \cdot \lambda = \{ a\lambda, ab\lambda \} = \{ a, ab \}$$

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Exercise:

Let $L1 = \{ a, ab \}$ be a language

Let $L2 = \{ c, bc \}$ be a language

Note: $\lambda \neq \phi$;

$$\lambda = \{ \lambda \}$$

$$\phi = \{ \}$$

$$L2 \cdot \phi = \phi$$

$$L2 \cdot \lambda = \{ c\lambda \cdot bc\lambda \} = \{ c, bc \}$$

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- KLEENE-STAR (*)

Suppose $A = \{0,1\}$, the kleene-star of A , $A^* = \{\lambda, 0, 1, 00, 01, 10, 11, \dots\}$ is the set of strings that can be formed out of the elements of the alphabet/set A including the null string.

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Suppose $L_1 = \{ a, ab \}$

The kleene-star of L_1 ,

$L_1^* = \{ \epsilon, a, ab, aab, \dots \}$ = the set of all strings that can be formed out of the strings of L_1 .

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- Exercise:

Let $A = \{ a, ab \}$ $B = \{ c, bc \}$ $C = \{ c \}$

Verify that:

- a) concatenation of languages is associative
- b) concatenation distributes over union

Solution

$$\begin{aligned} \text{a) } A \cdot (B \cdot C) &= (A \cdot B) \cdot C \\ \{a, ab\} \cdot (\{c, bc\} \cdot \{c\}) &= (\{a, ab\} \cdot \{c, bc\}) \cdot \{c\} \\ \{a, ab\} \cdot \{cc, bcc\} &= \{ac, abc, abbc\} \cdot \{c\} \\ \{acc, abcc, abbcc\} &= \{acc, abcc, abbcc\} \end{aligned}$$

Therefore, concatenation is associative

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$$b) A \cdot (B \cup C) = A \cdot B \cup A \cdot C$$

$$\{a, ab\} \cdot (\{c, bc\} \cup \{c\}) = (\{a, ab\} \cdot \{c, bc\}) \cup (\{a, ab\} \cdot \{c\})$$

$$\{a, ab\} \cdot \{c, bc\} = \{ac, abc, abbc\} \cup \{ac, abc\}$$

$$\{ac, abc, abbc\} = \{ac, abc, abbc\}$$

Therefore, concatenation distributes over union

ALPHABETS, STRINGS AND LANGUAGES

Note: $A^3 = A . A . A$

Let $A = \{a, ab\}$

$A^3 = ?$

$A^3 = (\{a, ab\} . \{a, ab\} . \{a.ab\})$

$A^3 = \{aa, aab, aba, abab\} . \{a, ab\}$

$A^3 = \{aaa, aaab, aaba, aabab, abaa, abaab, ababa, ababab\}$

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Note: ω^R = Reverse of ω

$$\omega = ab \quad \omega^R = ba$$

Let $A = \{\omega \mid \omega \text{ is a string}\} = \{a, ab\}$

Determine the strings in B if $B = \{\omega\omega^R \mid \omega \in A\}$

ALPHABETS, STRINGS AND LANGUAGES

Note: ω^R = Reverse of ω

$$\omega = ab \quad \omega^R = ba$$

Let $A = \{a, ab\}$

Determine the strings in B if $B = \{\omega\omega^R \mid \omega \in A\}$

Answer

$B = \{aa, abba\}$ = mirror image language

ALPHABETS, STRINGS AND LANGUAGES

SINGLETON SET NOTATION

The singleton set notation for $\{0\}$ is the boldfaced **0**.

Similarly, $\{a\} = \mathbf{a}$

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- Exercise:

Let $V = \{0, 1\}$. Using singleton sets and operators \cup , $.$, and $*$, find an expression

for the language that has

- All $\omega \in V^*$ such that the string 101 appears as a substring of ω .
- All $\omega \in V^*$ such that each 0 appearing in ω is immediately followed by at least two 1's.
- All $\omega \in V^*$ such that each 1 appearing in ω is immediately followed by the substring 10.

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- Exercise:

Let $v = \{0, 1\}$. Using singleton sets and operators \cup , $.$, and $*$, find an expression for all $\omega \in v^*$ such that the string 101 appears as substring of ω

Answer: $(0 \cup 1)^* . 101 . (0 \cup 1)^*$