Paving the Way for NFV: Simplifying Middlebox Modifications with StateAlyzr

Junaid Khalid, Aaron Gember-Jacobson, Roney Michael, Archie Abhashkumar, Aditya Akella



Perform sophisticated operations on network traffic

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Firewall



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Intrusion detection system (IDS)



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Caching proxy



Perform sophisticated operations on network traffic

Firewall



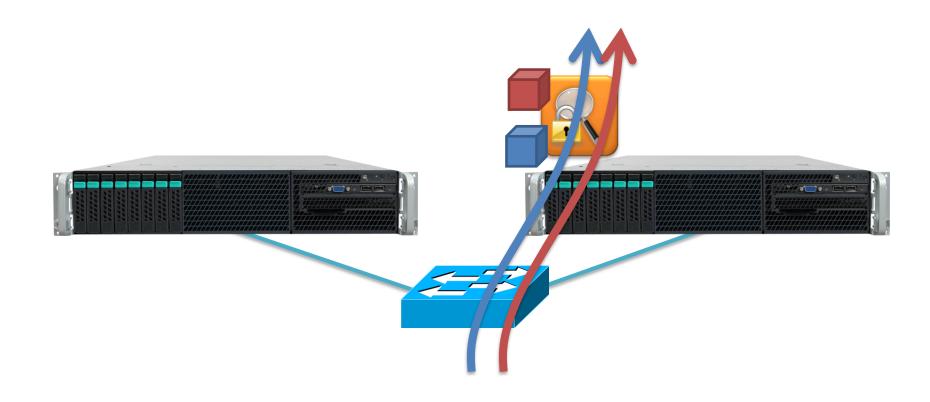
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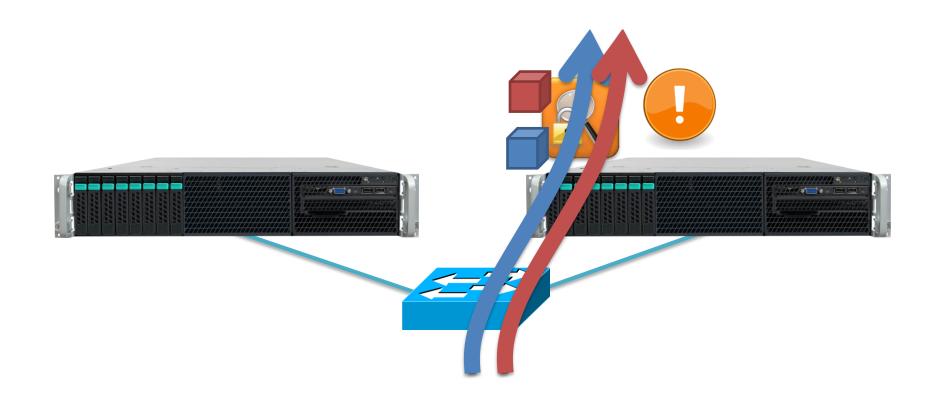


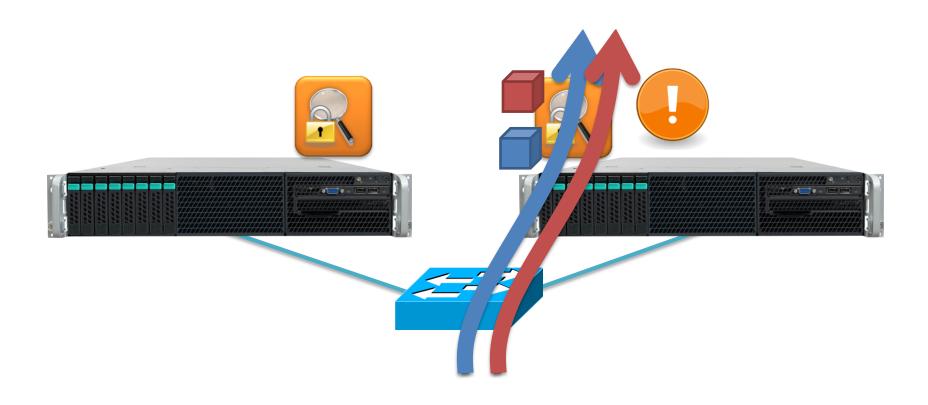
Caching proxy



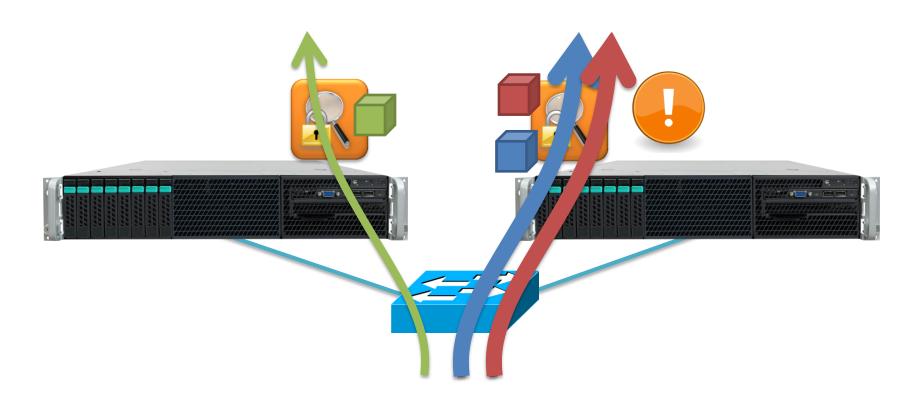
Maintain state about connections and hosts





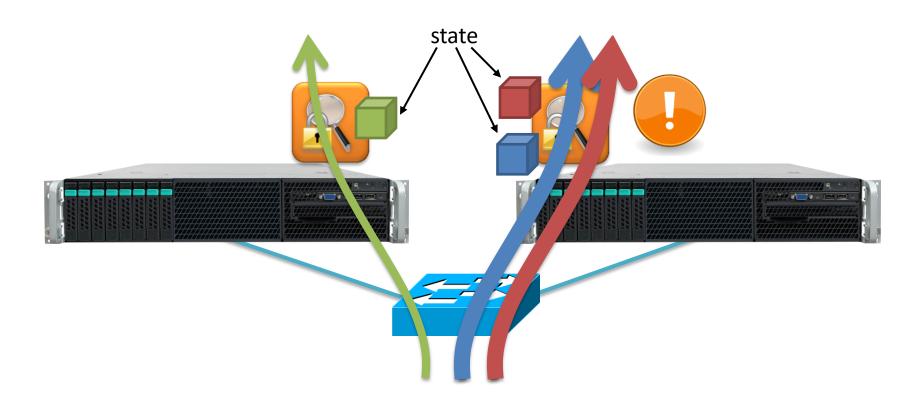


NFV enables *elastic scaling* and *high availability*

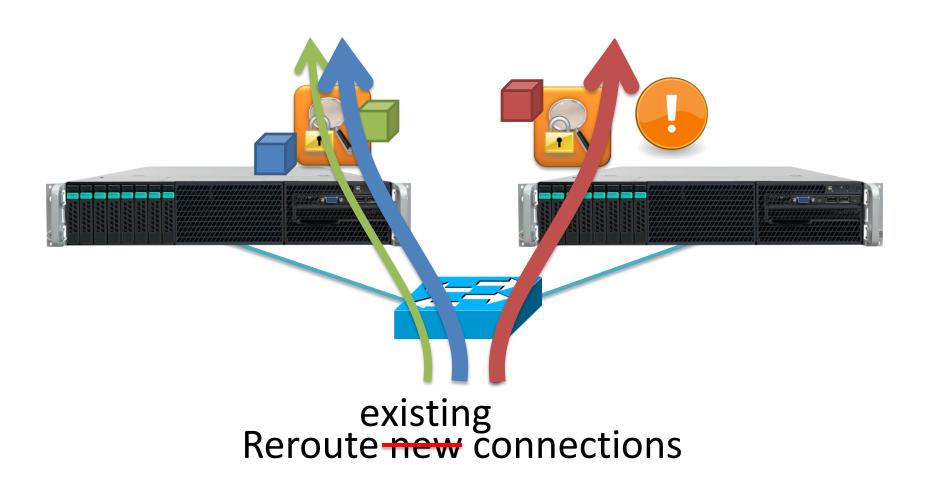


Reroute new connections

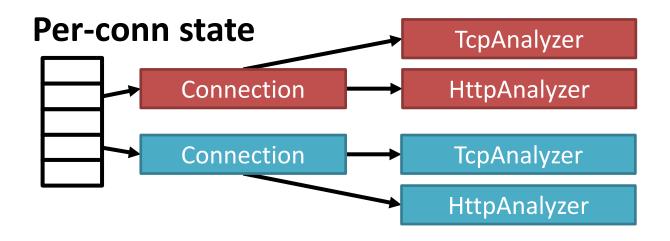
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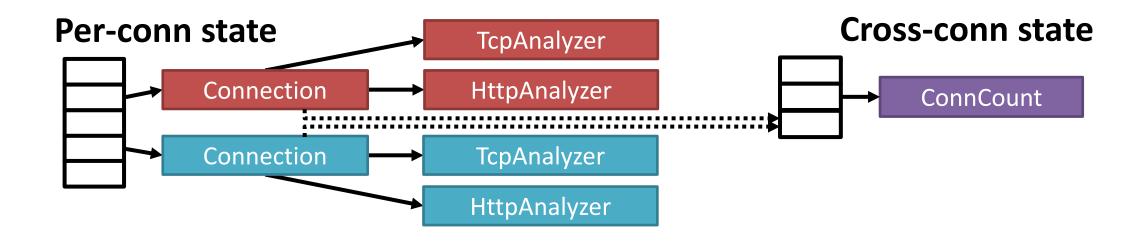
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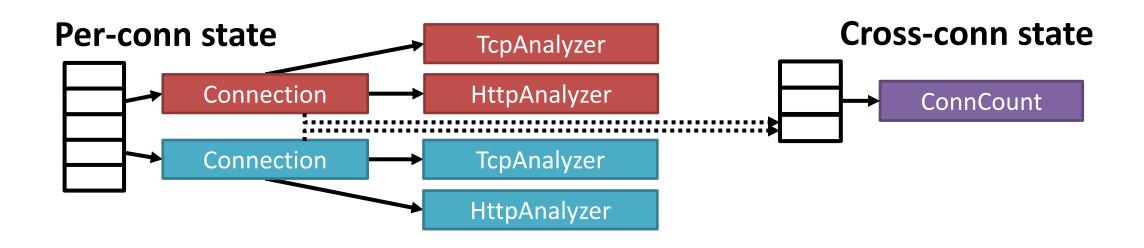








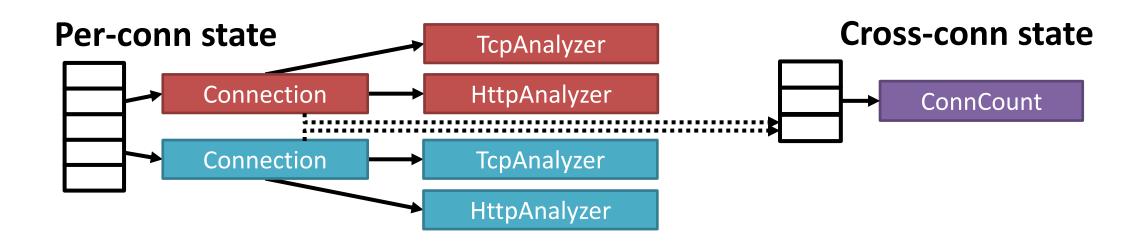
State created or updated by a middlebox applies to either a single connection or a set of connections





All-conns state

Statistics

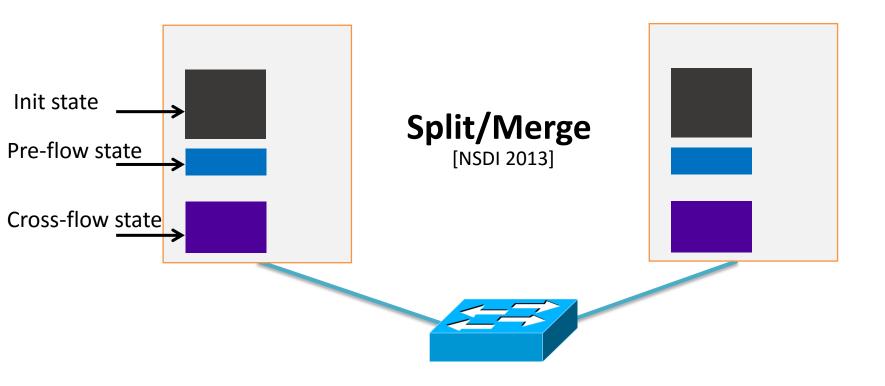




Frameworks for transferring, or sharing live middlebox state

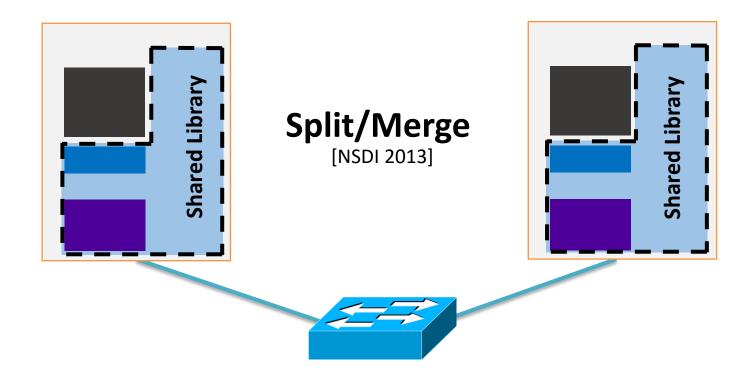
Frameworks for transferring, or sharing live middlebox state

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• Require modifications or annotation to middlebox code

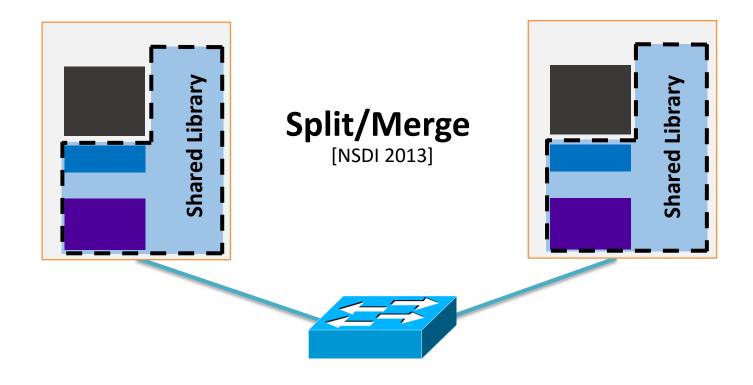


Required modifications:

1. State allocation

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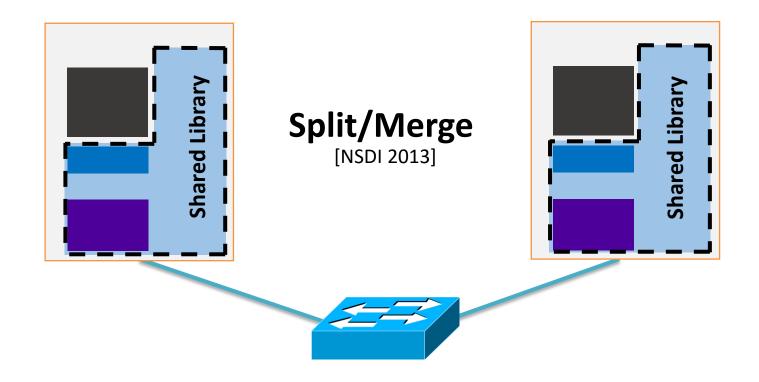


Required modifications:

- 1. State allocation
- 2. State access

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Required modifications:

- 1. State allocation
- State access
- 3. State merge

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Framework	State Allocation	State Access	Serialization	Merge State
Split/Merge [NSDI 2013]	✓	✓		✓

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Stateless NF [HotMiddlebox 2015]	✓	✓		

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Middleboxes are complex

MB	LOC (C/C++)	Classes/ Structs	Level of pointers	Number of Procedures
PRADS	10K	40	4	297
OpenVPN	62K	194	2	2023
HAProxy	63K	191	8	2560
Bro IDS	97K	1798	-	3034
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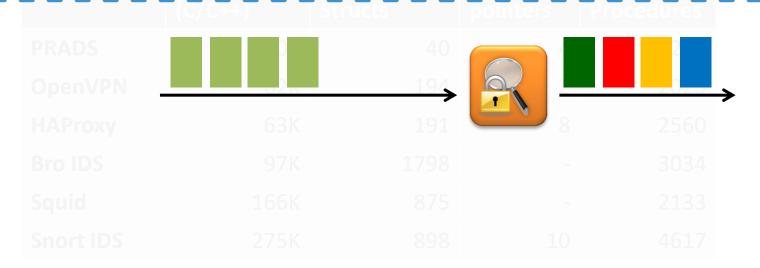
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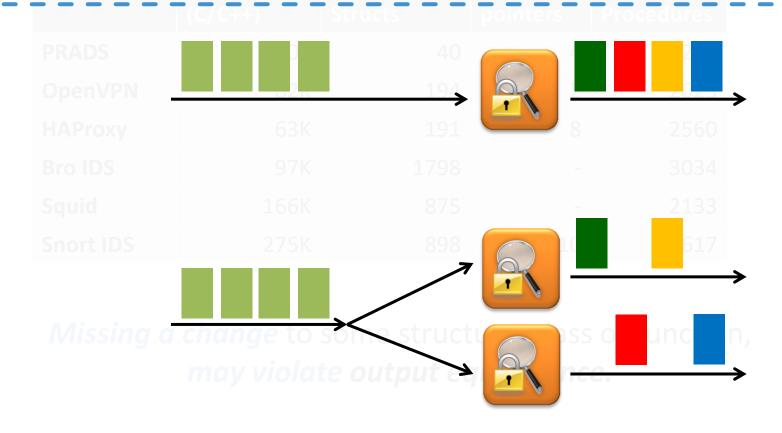
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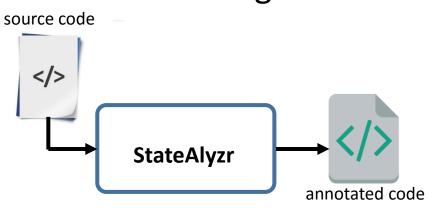
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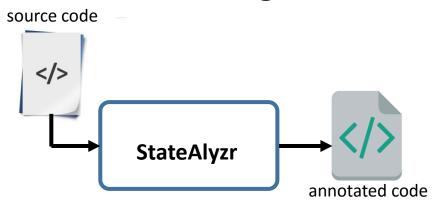
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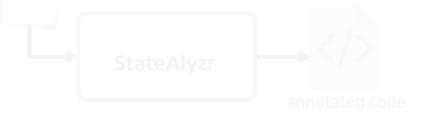
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StateAlyzr annotated code

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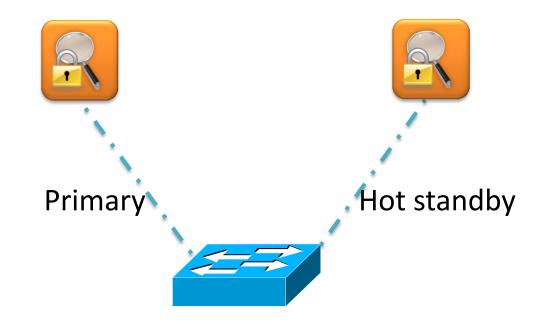
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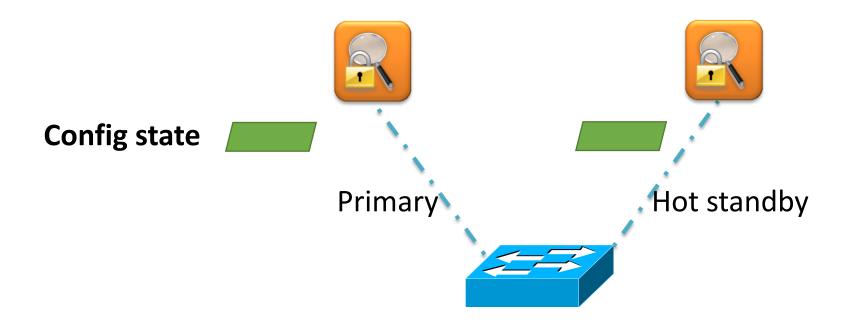
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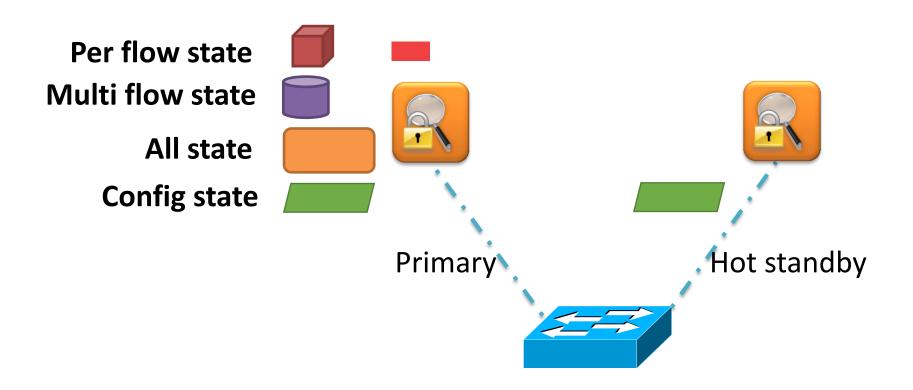
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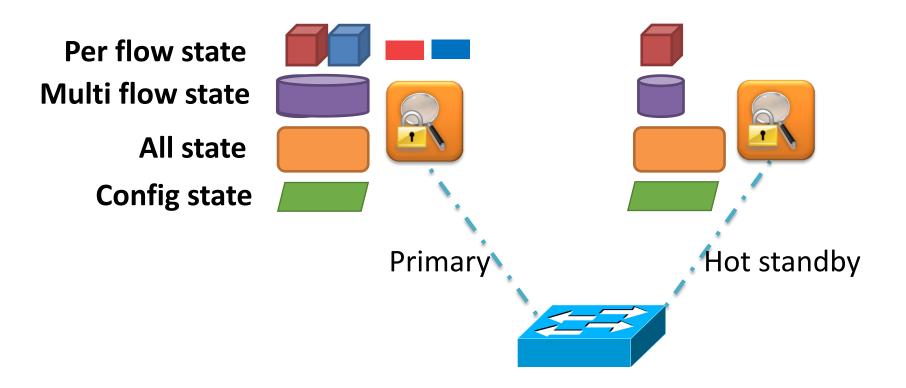
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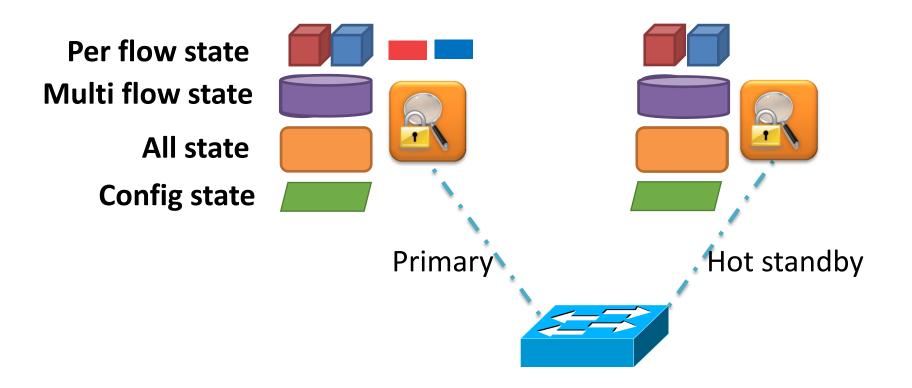






updated

The *primary* sends a copy of the state to the *hot standby* after each packet

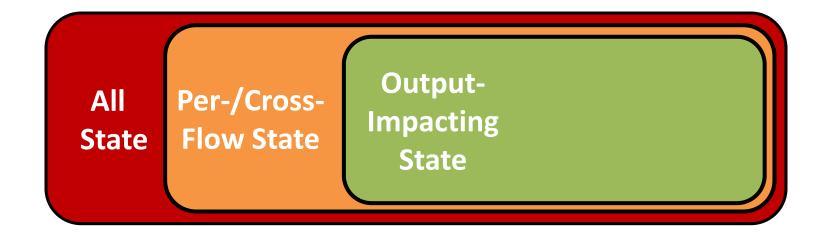


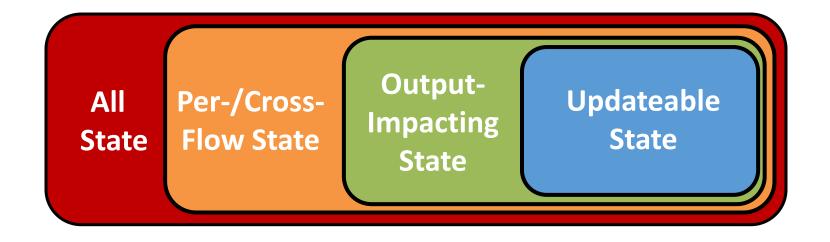
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All State

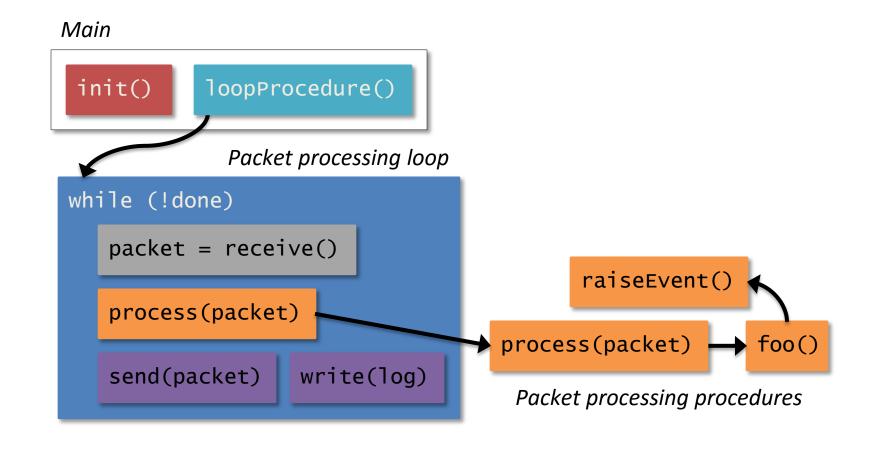






Main init()

Main init() loopProcedure()



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  int count = 0;
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    struct pcap_pkthdr pcapHdr;
    char *pkt = pcap_next(extPcap, &pcapHdr);
    .
```

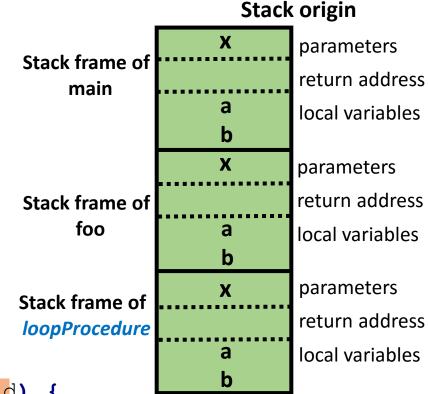
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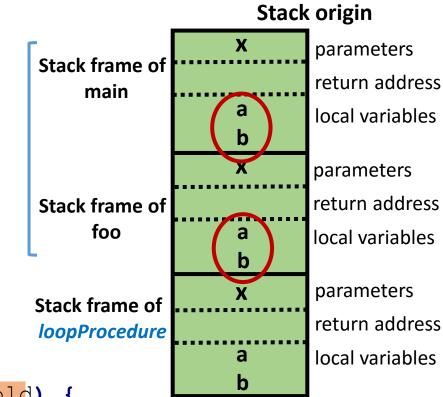


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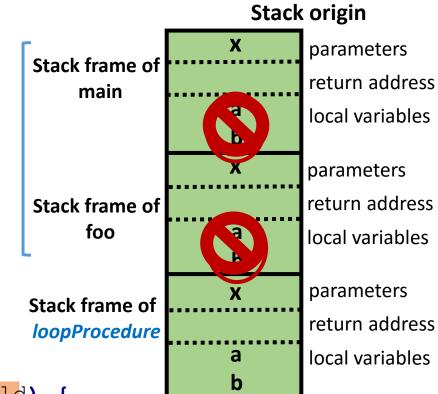


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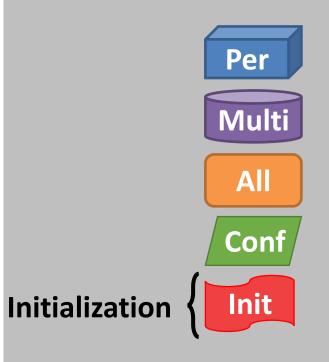
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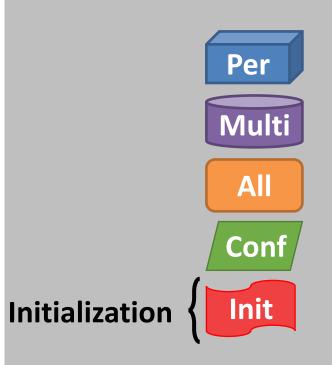


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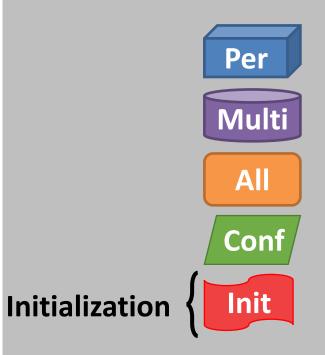


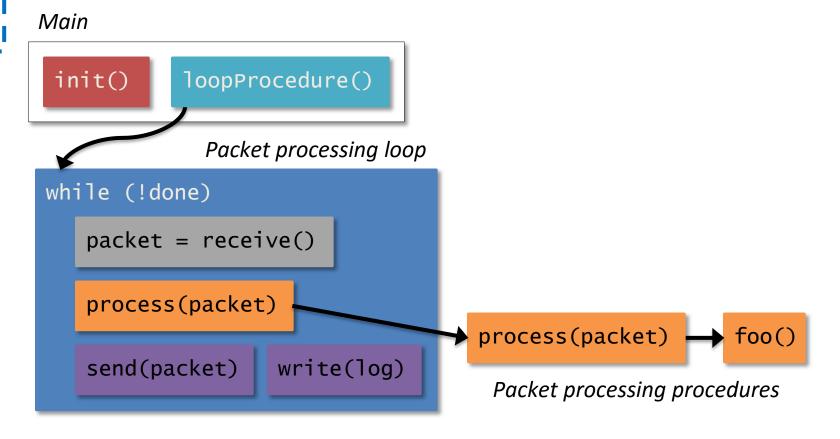
Improve precision by considering variables which are used in packet processing code

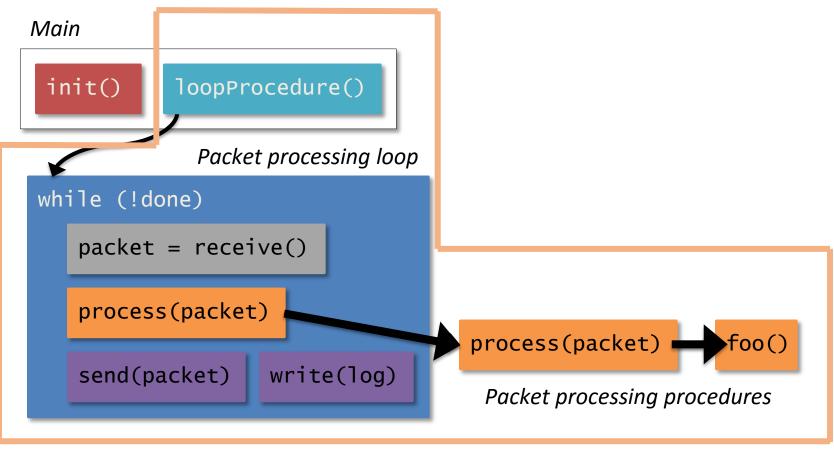


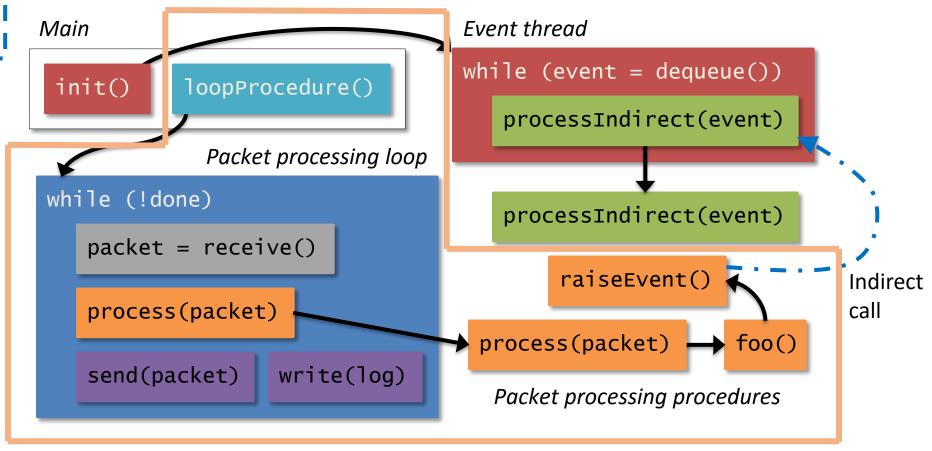
1. Per-/cross-flow state identification

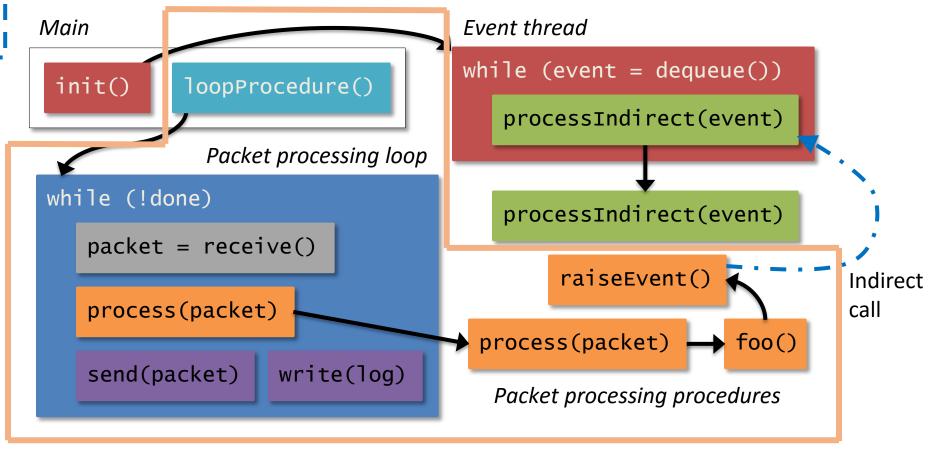
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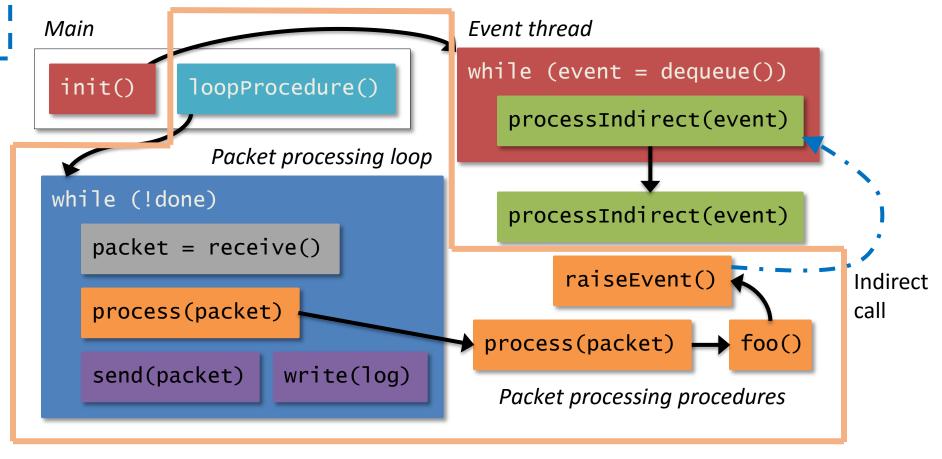




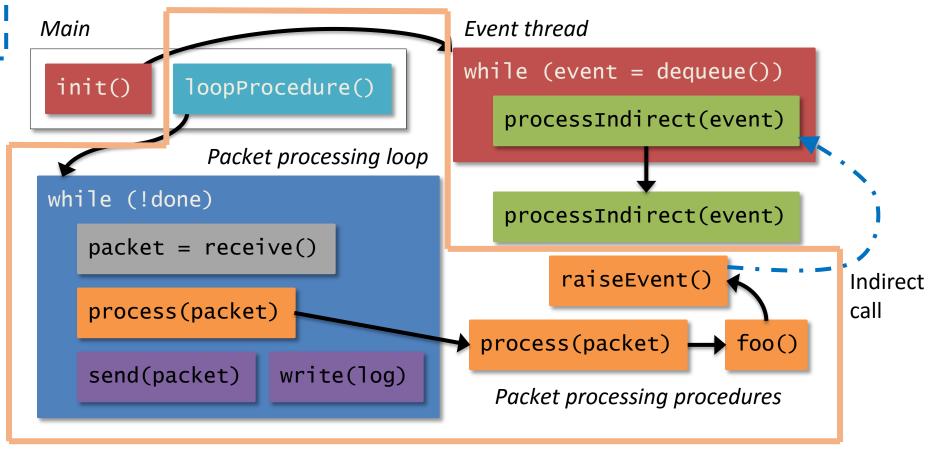




```
struct pktHdr *pkt = recv(extPcap);
src_ip = pkt->ip_src_addr;
packet_count ++;
index = src ip + offset
```

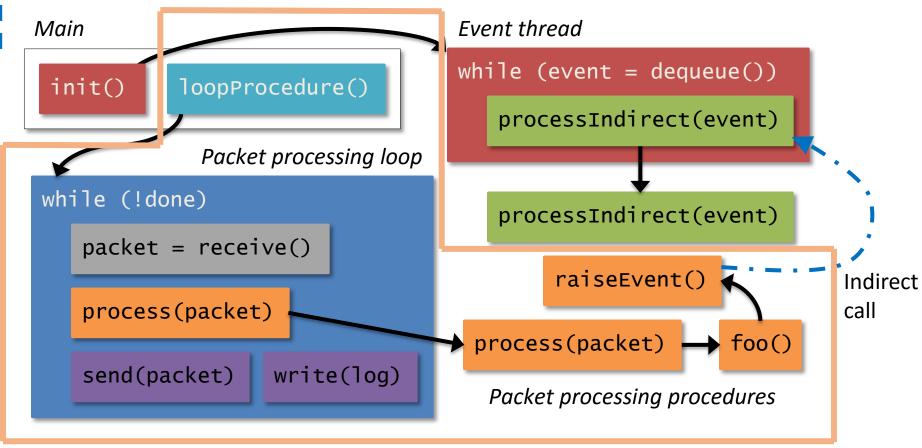


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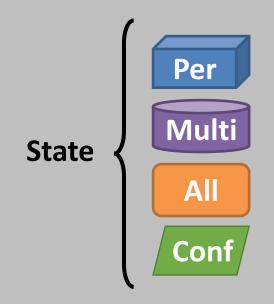


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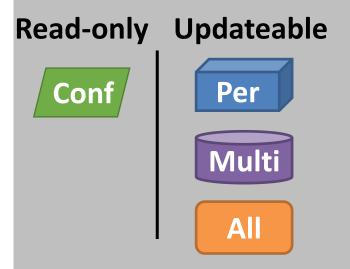
Computes a *forward slice* from packet recv function. Any procedure appearing *in the slice* is considered as *packet processing procedure*.





Read-only Updateable

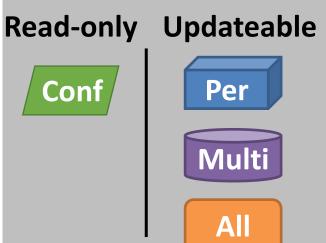
Whether the state is updated while processing the packet?



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- Strawman approach
 - Identify top-level variable on the left-hand-side(LHS) of assignment statement

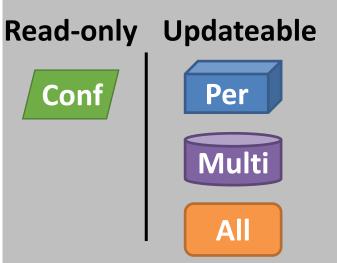
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per-/cross-flow var
in_port = pkt.src_port;
```



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Falls short due to aliasing



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Falls short due to *aliasing*

```
per-/cross-flow var
int *index = &tail;
*index = (*index + 1)%100;
```

Read-only Updateable



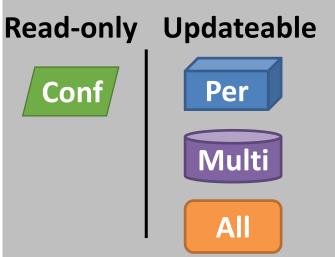


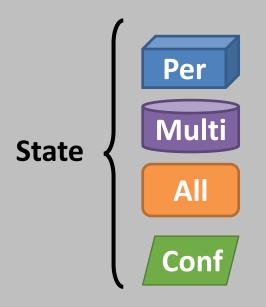


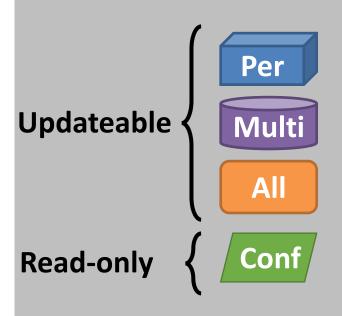
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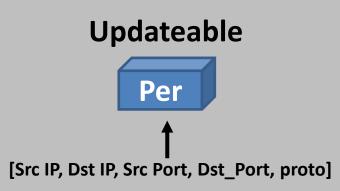
StateAlyzr employs flow-, context-, and field-insensitive *pointer analysis* to identify updateable variables

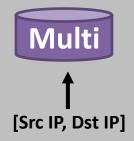




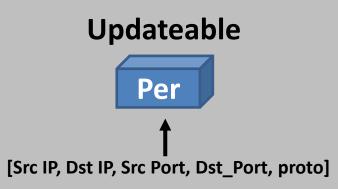


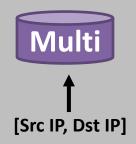
Identify a set of *packet header fields* that delineate the subset of traffic that relates to the state



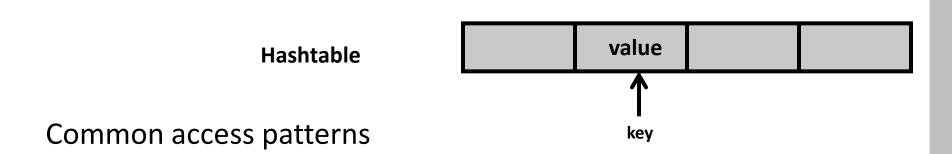


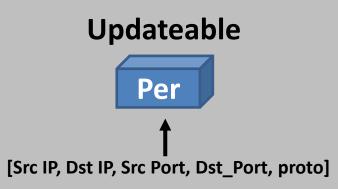
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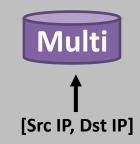




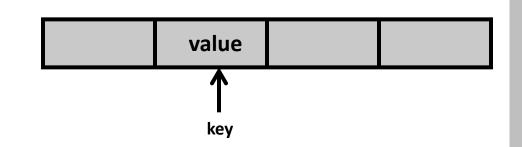
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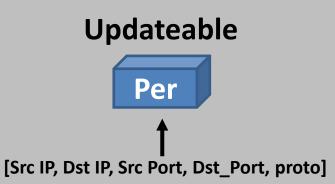


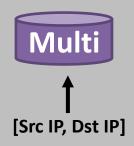
Hashtable

Common access patterns

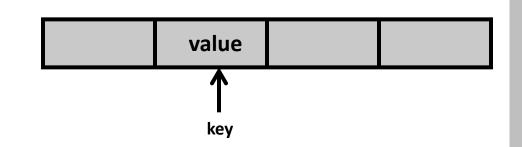
1. Square brackets

```
entry = table[index];
```





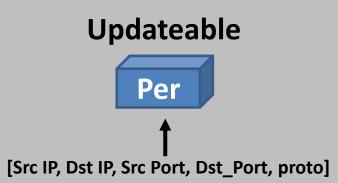
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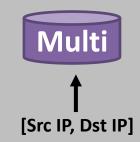


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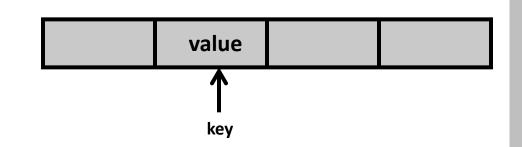
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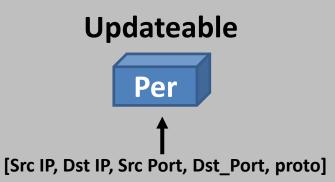
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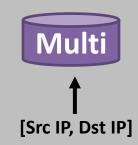


Hashtable

- 1. Square brackets
- 2. Pointer arithmetic

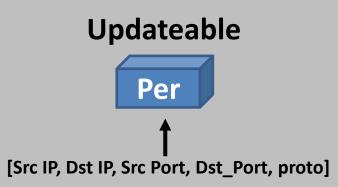
```
entry = head + offset;
```

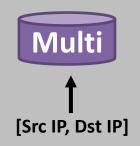




Identify a set of *packet header fields* that delineate the subset of traffic that relates to the state

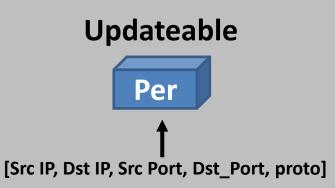
- 1. Square brackets
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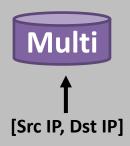




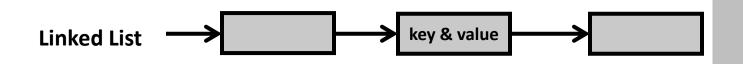
Identify a set of *packet header fields* that delineate the subset of traffic that relates to the state

- 1. Square brackets
- 2. Pointer arithmetic
- 3. Iteration



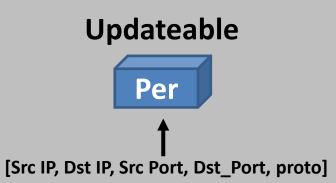


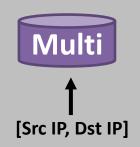
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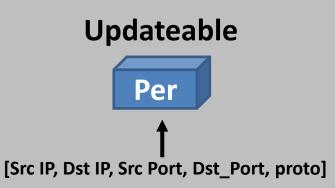
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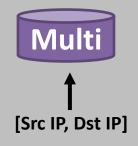
```
struct host *lookup(uint ip) {
  struct host *curr = hosts;
  while (curr != NULL) {
    if (curr->ip == ip)
      return curr;
    curr = curr->next;
  }
}
```





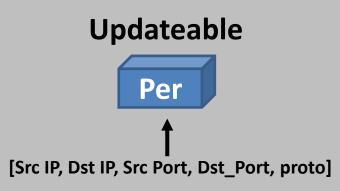
Identify a set of *packet header fields* that delineate the subset of traffic that relates to the state

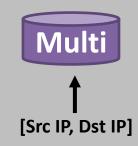




Identify a set of *packet header fields* that delineate the subset of traffic that relates to the state

Program *chopping* to determine relevant *header fields*





```
struct pktHdr *pkt = recv(extPcap);
```

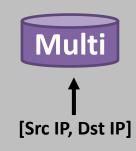
3. Identify states' flowspace dimensions

Identify a set of *packet header fields* that delineate the subset of traffic that relates to the state

Program *chopping* to determine relevant *header fields*

```
Updateable
Per

[Src IP, Dst IP, Src Port, Dst_Port, proto]
```



```
struct pktHdr *pkt = recv(extPcap);
src_ip = pkt->ip_src_addr;
packet_count ++;
index = src_ip + offset
entry = host_map[index]
```

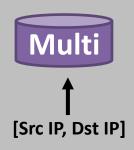
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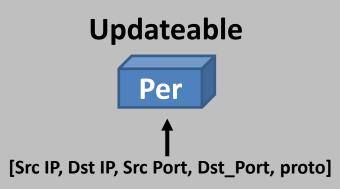


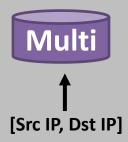
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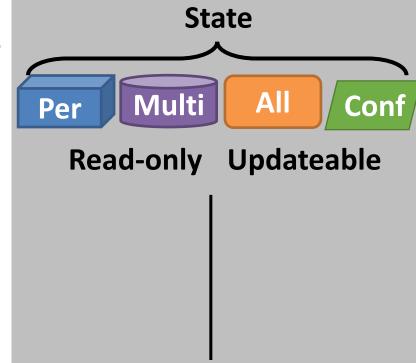


1. Identify Per-/Cross-flow state

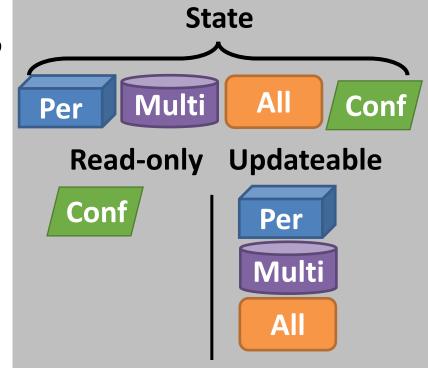
State
Per Multi All Conf

1. Identify Per-/Cross-flow state

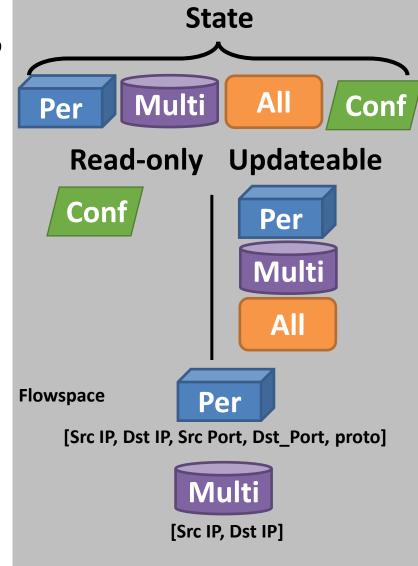
- 1. Identify Per-/Cross-flow state
- 2. Identify Updateable State



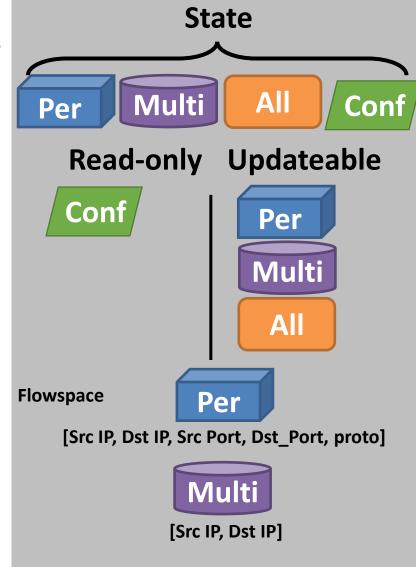
- 1. Identify Per-/Cross-flow state
- 2. Identify Updateable State



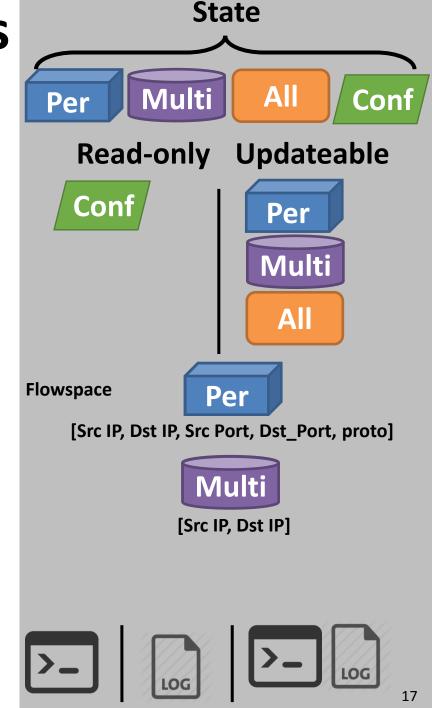
- 1. Identify Per-/Cross-flow state
- 2. Identify Updateable State
- 3. Identify States' Flowspace Dimensions



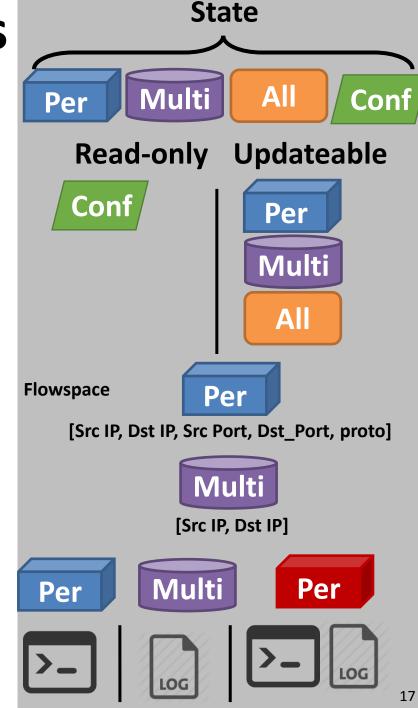
- 1. Identify Per-/Cross-flow state
- 2. Identify Updateable State
- 3. Identify States' Flowspace Dimensions
- 4. Output Impacting State



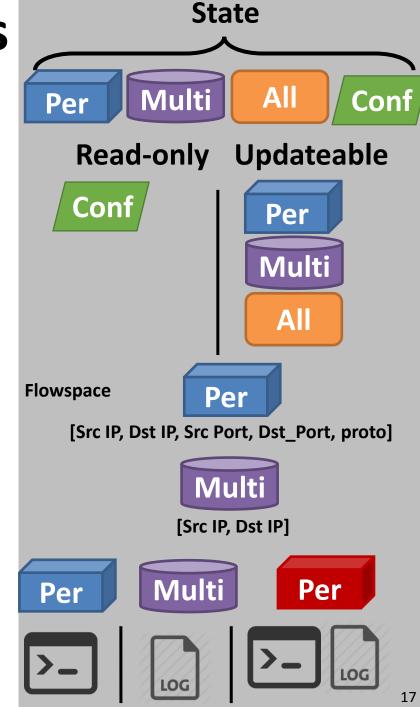
- 1. Identify Per-/Cross-flow state
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- 3. Identify States' Flowspace Dimensions
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 - Identify the type of output (log or packet) that updateable state affects



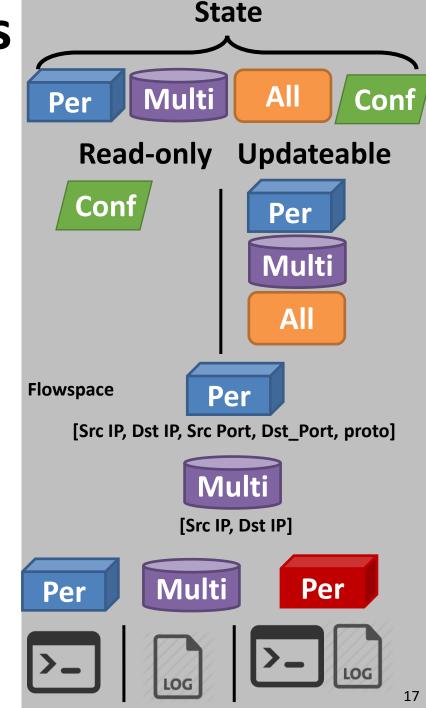
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- 4. Output Impacting State
 - Identify the type of output (log or packet) that updateable state affects
- 5. Tracking Run-time Update



- 1. Identify Per-/Cross-flow state
- 2. Identify Updateable State
- 3. Identify States' Flowspace Dimensions
- 4. Output Impacting State
 - Identify the type of output (log or packet) that updateable state affects
- 5. Tracking Run-time Update
 - Insert statements to do run time monitoring to track whether a variable is updated



Used CodeSurfer to implement StateAlyzr

CodeSurfer has built-in support for

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 - Control flow graph construction

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 - Flow and context-insensitive pointer analysis

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Analyzed four open-source middleboxes

1. PRADS – a monitoring middlebox

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- 1. PRADS a monitoring middlebox
- 2. Snort an IDS

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- 1. PRADS a monitoring middlebox
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- 3. HAProxy a load balancing proxy
- 4. OpenVPN a VPN gateway

Evaluation

Evaluation

Precision

Evaluation

- Precision
- Performance benefits at run time

	Step 0		Step 1	Step 2
МВ	All variables	Persistent variables	per-/cross- flow variables	Updateable variables
PRADS	1529	61	29	10
Snort IDS	18393	507	333	148
HAproxy	7876	272	176	115
OpenVPN	8704	156	131	106

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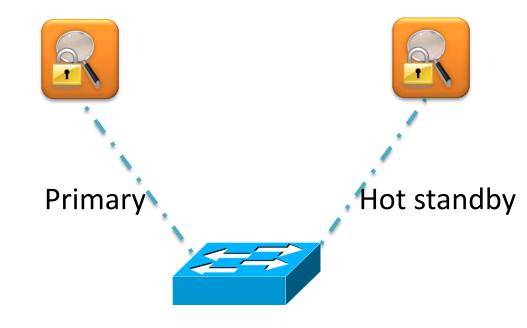
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				A

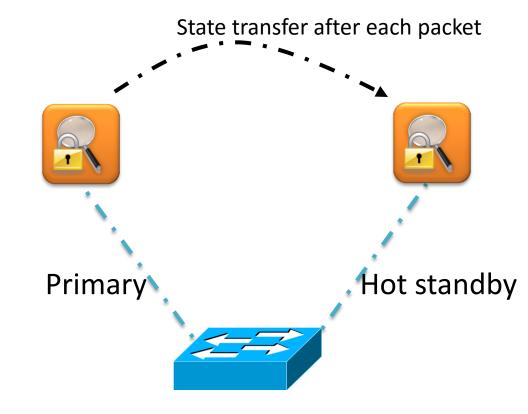
StateAlyzr offers useful *improvements* in *precision*

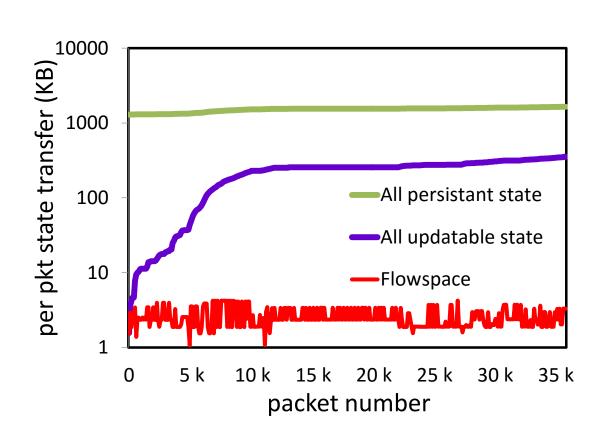
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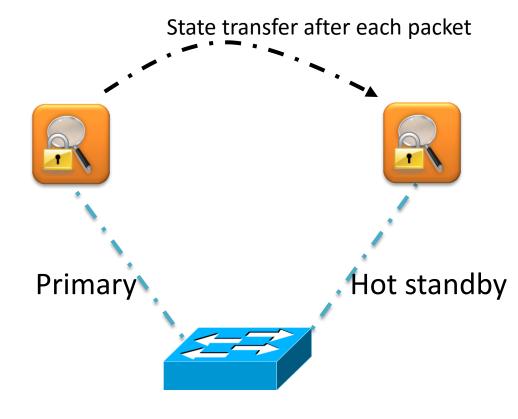
StateAlyzr offers useful *improvements* in *precision*

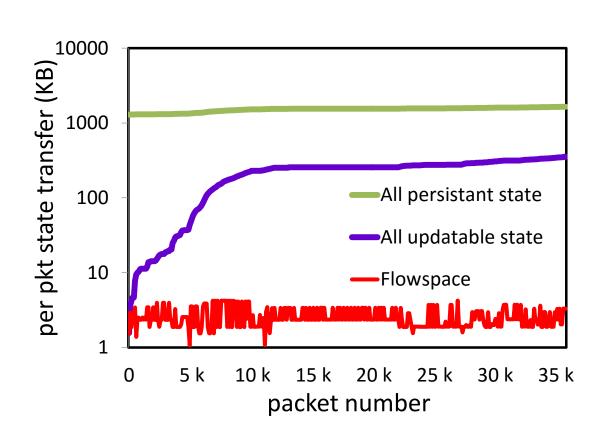
Theoretically *proved* the *soundness* of our algorithms

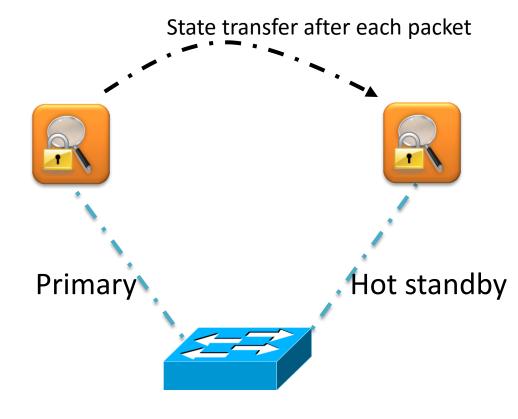


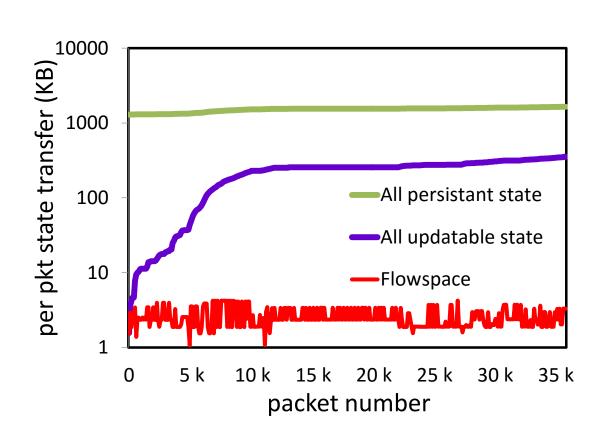


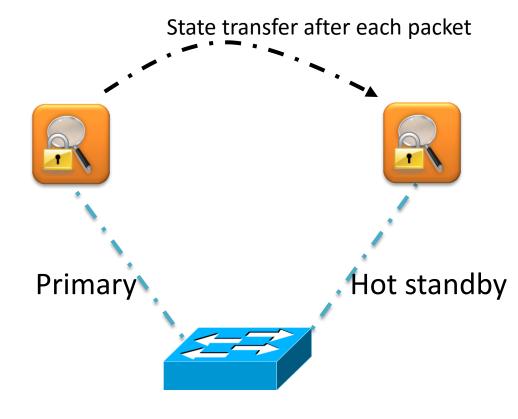


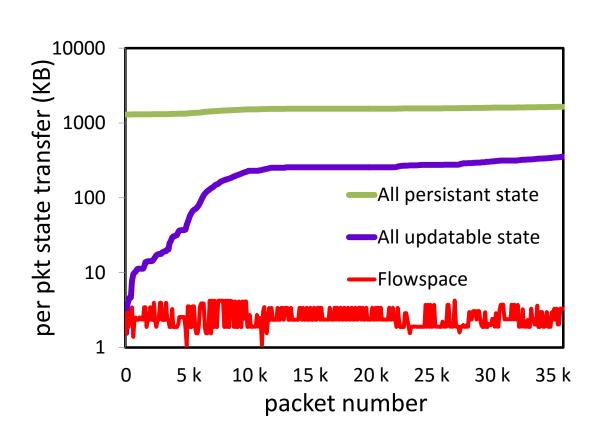




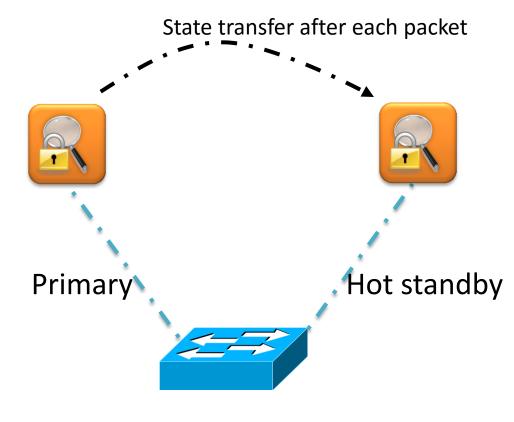


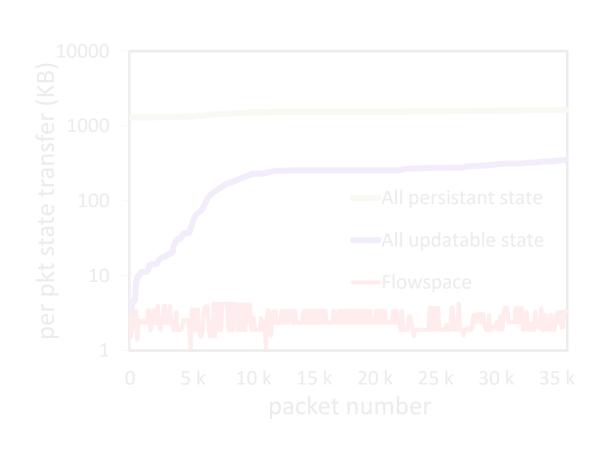




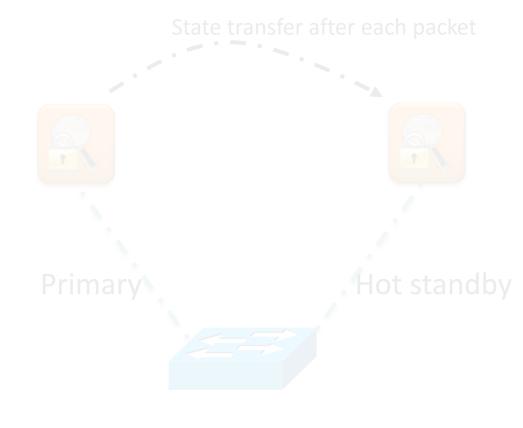


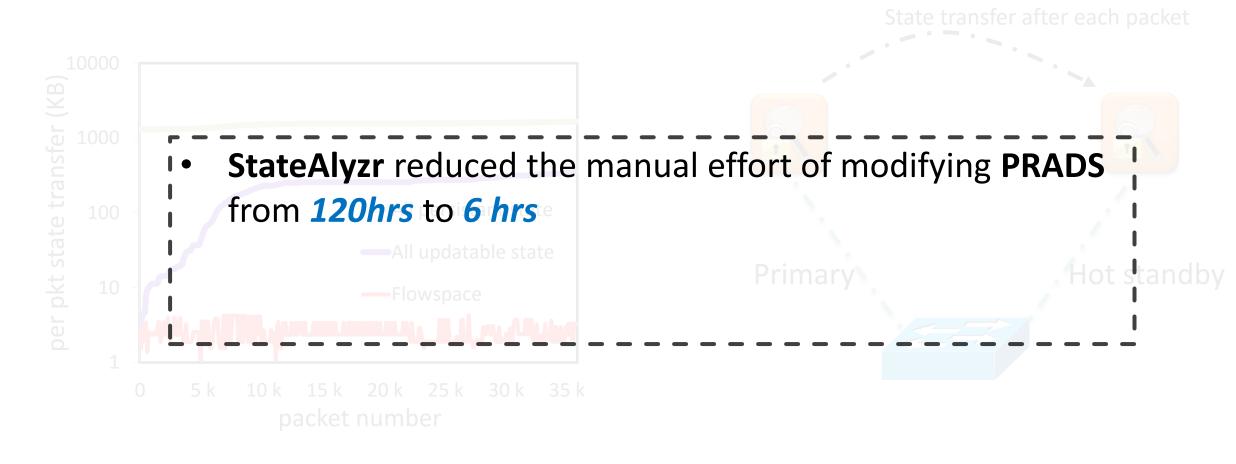
Reduction in the state transfer by 305x



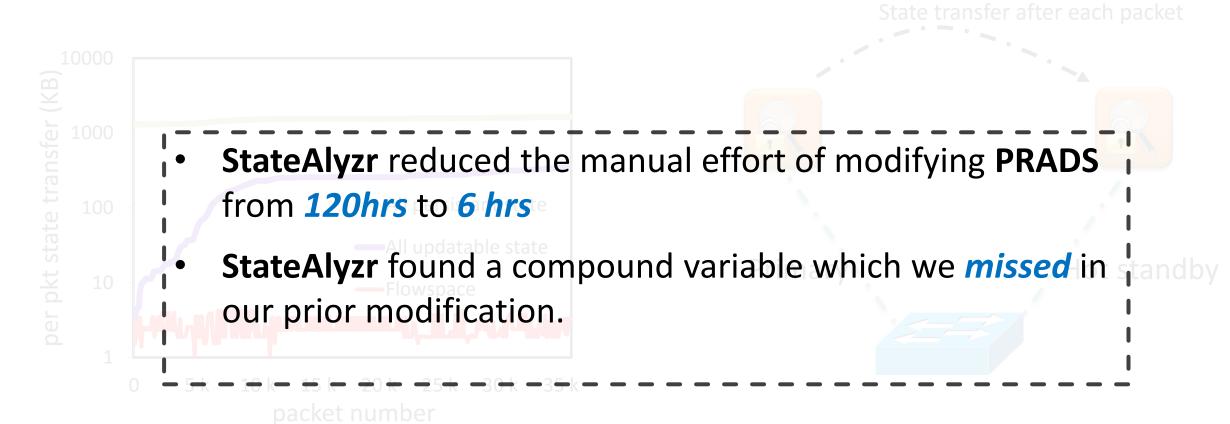


Reduction in the state transfer by 305x





Reduction in the state transfer by 305x



Reduction in the state transfer by 305x

 Goal is to aid middlebox developers to identify state objects that need explicit handling

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- Novel state characterization algorithms that adapt standard program analysis tools
- Ensure soundness and high precision
- Ultimate goal is to fully automate the process