Natueal Algorithm

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Linear System - Matrix Tree Theorem

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In which we discuss the connection between graph Laplacian and spanning tree.

1 Introduction and Notation

Given a connected graph G = (V, E), we fix the following notations:

- L is the graph Laplacian of G.
- $\tau(G)$ denote the number of spanning trees in graph G.
- Let $e \in E$, $G \setminus e$ is the graph deleting edge e.
- Let $e \in E$, G/e is the graph contracting edge e, i.e., contracting the endpoints of e.
- Let M be a square matrix and S is a subset of its indexes, M[S] is the submatrix of M deleting the rows and columns in S.
- Let M be a square matrix, denote M(i, j) be the submatrix of M deleting the i-th row and the j-th column.
- The *ij*-cofactor of M is $(-1)^{i+j} \det M(i,j)$.
- The adjugate of M, denoted by adj M has ij-entry being the ji-cofactor of M.

2 Theorems and Lemmas

2.1 Number of spanning trees

Theorem 1 For any $u \in V$, $\det L[u] = \tau(G)$.

PROOF: We prove this theorem with four steps:

- 1. For any edge $e \in E$, $\tau(G) = \tau(G \setminus e) + \tau(G/e)$. Observe that $\tau(G \setminus e)$ = the number of spanning trees without e and $\tau(G/e)$ = the number of spanning trees containing e.
- 2. Let $e = (u, v) \in E$, $L[u] = L(G \setminus e)[u] + \mathbf{1}_{i=j=v}$. This is because taking away e decrease the degree of v by 1.

3.
$$\det L[u] = \det \left(L(G \backslash e)[u] \right) + \det \left(L(G \backslash e)[u, v] \right).$$

$$\det L[u] = \det \left(L(G \backslash e)[u] + \mathbf{1}_{i=j=v} \right)$$

$$= \sum_{w \in V} (-1)^{w+v} \left(L(G \backslash e)[u] + \mathbf{1}_{i=j=v} \right)_{v,w} \cdot \det \left(L(G \backslash e)[u, v, w] \right)$$

$$= \sum_{w \in V} (-1)^{w+v} \left(L(G \backslash e)[u] \right)_{v,w} \cdot \det \left(L(G \backslash e)[u, v, w] \right)$$

$$+ (-1)^{v+v} \left(\mathbf{1}_{i=j=v} \right) \cdot \det \left(\det L(G \backslash e)[u, v] \right)$$

$$= \det \left(L(G \backslash e)[u] \right) + \det \left(L(G \backslash e)[u, v] \right)$$

4. Observe that $L(G \setminus e)[u, v] = L(G/e)[u, v]$, then by induction, the theorem holds.

Lemma 2 For any square matrix M, Madj(M) = det(M)I.

PROOF: Consider the *ij*-entry of the left hand side.

$$\left(M\operatorname{adj}(M)\right) = \sum_{k} M_{i,k}\operatorname{adj}(M)_{k,j} = \sum_{k} M_{i,k} \cdot (-1)^{j+k} \cdot \det(M(j,k)) \tag{1}$$

When i = j, this is exactly $\det(M)$ and when $i \neq j$, this will be the determinant of replacing the j-th row of M by the i-th row of M, which is 0. \square

The next theorem connects the graph Laplacian to the number of spanning tree via the matrix adjugate.

Theorem 3 We have $adj(L) = \tau(G)J$, where J is the all 1 matrix.

PROOF: By Lemma 2, we have $Ladj(L) = \det(L)I = 0$. That is, the columns of adj(L) are in the null space of L. If G is disconnected, then the entries of adj(L) are all zero and $\tau(G) = 0$, the theorem holds. If G is connected, then the null space of L is spanned by the uniform vector, i.e., the columns of adj(L) are all uniform vector. As adj(L) is symmetric, we know that every entries of adj(L) are the same. Finally, by Theorem 1, we know that the diagonal of adj(L) is $\tau(G)$ and thus the theorem holds. \square

Theorem 4 (matrix tree theorem) Let $0 = \lambda_1 \leq \lambda_2 \leq \cdots \leq \lambda_n$ be the eigenvalues of the graph Laplacian L, then $\tau(G) = \frac{1}{n} \prod_{i=2}^{n} \lambda_i$.

PROOF: If G is disconnected, then $\lambda_2 = 0$ and the theorem holds. If G is connected, let's consider the characteristic polynomial $\phi(t) = \det(tI - L) = \prod_{i=1}^{n} (t - \lambda_i)$ of L. The linear term of $\phi(t)$ is

$$(-1)^{n-1} \sum_{i=1}^{n} \prod_{j \in [n] \setminus \{i\}} \lambda_j = (-1)^{n-1} \prod_{i=2}^{n} \lambda_i$$
 (2)

On the other hand, if we directly expand det(tI - L) using Laplace's expansion, we will have

$$\det(tI - L) = \sum_{i=0}^{n-1} (-1)^{n-i} t^i \sum_{S \subseteq [n], |S| = i} \det(L[S])$$
(3)

Thus, the coefficient of the linear term in $\phi(t)$ is also $(-1)^{n-1} \sum_{u \in [n]} \det(L[u]) = n\tau(G)$ from Theorem 1. As a result, the theorem holds. \square

References

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