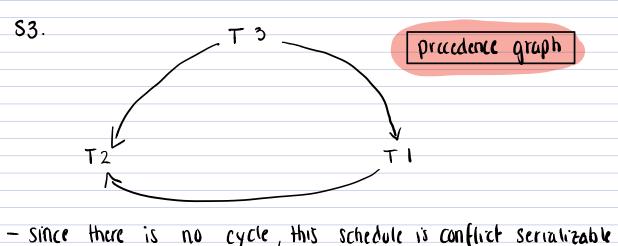
Jerry Perez

301.				
31.	T	T2	Т 3	schedule
	read A	rcad Ar Write B Write C	write (
	write C			
T3				Precedence graph
T2 T1 - this series is conflict-serializable since there is no cycle. A possible serial schedule is:				
		T3, T1, T2		
S2.	TI	Τ2	T 3	T 4
Schedule	readA	wate C	write A	
	write B read C	rcad A		read A
		read O		write 0
TI Freedence graph				
T 2 — T 4				
- the schedule is conflict serializable since Here is no cycle. A possible serial schedule would be:				



- since there is no cycle, this schedule is conflict serializable a possible serial schedule would be:

T3, T1, T2