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April 10, 2017

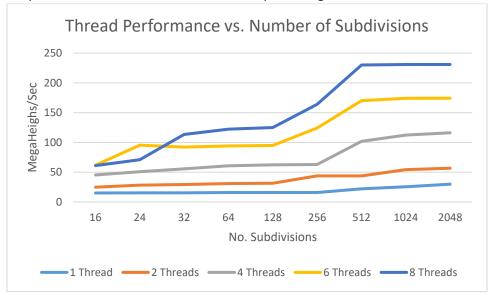
CS 475 – Parallel Programming

Project 1 – Written Commentary

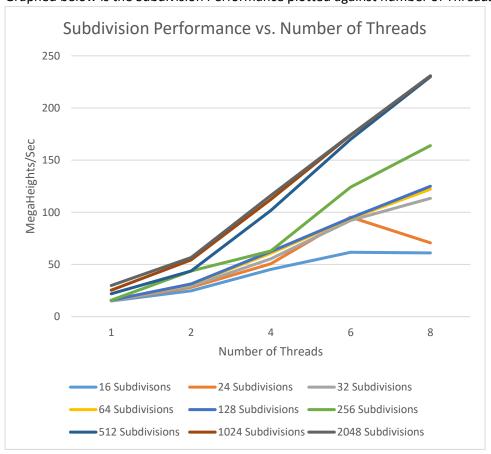
- 1. My program was run on the OSU Flip server at 6:47 PM Pacific on April 10, 2017.
- 2. My program suggests the actual volume of the surface is around 25.312 units³. This is based on the results of computing the volume using 2048 Subdivisions for both X and Y.
- 3. The average megaheights/sec over 100 runs is listed in the table below.

	16	24	32	64	128	256	512	1024	2048
	Subdivisions								
1	15.059	15.357	15.236	15.697	15.769	15.79	21.938	25.393	29.845
Thread									
2	24.776	28.133	29.19	30.883	31.367	43.79	43.789	54.378	56.665
Threads									
4	45.284	50.678	55.469	60.797	62.264	62.959	101.702	112.253	116.171
Threads									
6	61.716	95.394	92.242	94.148	94.736	124.096	169.993	174.005	174.226
Threads									
8	61.044	70.818	113.449	122.141	124.994	163.942	229.979	230.777	230.881
Threads									

3 cont.
Graphed below is the Thread Performance plotted against number of subdivisions.



Graphed below is the Subdivision Performance plotted against number of Threads.



- 4. From the graphs above, we can see that Performance generally increases by using additional threads, at least when comparing 1, 2, 4, 6, and 8 threads.

 Furthermore, we see that performance generally increases by creating additional subdivisions, but this levels off after a sufficiently large amount of subdivisions have been created.
- 5. I think the performance gain that is found by increasing the amount of threads is resulting from the additional threads running in parallel on their own cores. This seems to be the case because the performance from using 1, 2, 4, 6, and 8 threads seems to be multiples of 1, 2, 4, 6, and 8 respectively. When computing average (mean) S values for the different runs, I found that the S value for 2 threads is 0.504653, for 4 threads was 0.263481, for 6 threads was 0.164097, and for 8 threads is 0.142613. Inverting these suggest that 2 threads is 1.981559 times faster than 1 thread, 4 threads is 3.795347 times faster than 1 thread, 6 threads is 6.093975 times faster than 1 thread, and 8 threads is 7.011975 times faster than 1 thread.

The value that doesn't make sense here is 6.093975 for 6 threads, which shouldn't theoretically exceed 6. This may be occurring due to variations in CPU load as the testing script was completing its runs, causing extremely good average values for 6 threads, or extremely bad average values for 1 thread.

6. Since we know the average speedup values (as calculated in number 5) we can use the inverse Amdahl's law formula to help compute Fp. I will be computing Fp from each of the S values, and taking the mean of these as an approximation for Fp.

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Fp(n) = (n/(n-1))*(1-1/S(n)), where S(n) is the speedup for n threads.
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 $Fp(2) = (2/1)^* (1-.504653) \approx .9906$

 $Fp(4) = (4/3)^* (1-.263481) \approx .9820$

 $Fp(6) = (6/5)* (1-.164097) \approx 1.0031$

 $Fp(8) = (8/7)^* (1-.142613) \approx .9799$

FpAverage ≈ .9889. This suggests that around 98.89% of the code in this test is parallelizable.

7. If we use the parallel fraction .9889, we can compute the maximum speed-up we could ever get by using a limit (theoretically infinite amount of processor).

Speedup(n) = 1/((Fp/n) + (1-Fp)) as n goes to infinity simplifies to

Speedup(Max) = $1/(1-\text{Fp}) \approx 1/(1-.9889) \approx 90.09$ times theoretical max speedup.