

HW3

A1 Formalization in FOL

- 1) $R(x)$: x is a red thing
 $B(x)$: x is in the box

a) All red things are in the box.
 $\forall x (R(x) \rightarrow B(x))$

b) Only red things are in the box
 $\forall x (B(x) \rightarrow R(x))$

c) No animal is both a cat and a dog.
 $A(x)$: x is an animal
 $C(x)$: x is a cat
 $D(x)$: x is a dog

$$\forall x (A(x) \rightarrow \neg (C(x) \wedge D(x)))$$
$$\forall x (A(x) \rightarrow (\neg C(x) \vee \neg D(x)))$$

d) Every prize was won by a boy.
 $P(y)$: y is a prize
 $B(x)$: x is a boy
 $W(x, y)$: x won y

(Adam won the lottery)

$$\forall y (P(y) \rightarrow \exists x (B(x) \wedge W(x, y)))$$

#1 Formalization in FOL

1) e) A boy won every prize.

$P(y)$: y is a prize

$B(x)$: x is a boy

$W(x,y)$: x won y

$$\exists x \forall y (B(x) \wedge P(y) \rightarrow W(x,y))$$

2) a) $\forall x \forall y (x=y)$ (where all quantifiers have the same domain)

- Every pair of elements consists of two equal elements.

b) $\exists x \exists y (x \neq y)$ (where all quantifiers have the same nonempty domain)

- There exists at least one pair of elements such that the elements are not equal.

c) $\exists x \exists y (x \neq y \wedge \forall z (z \neq x \rightarrow z=y))$ (where all quantifiers have the same nonempty domain)

- There exists a model with only two distinct elements.

		Note: $z \neq x$	$z=x$	$z \neq x \rightarrow z=y$	$x \neq y \wedge \forall z (z \neq x \rightarrow z=y)$
2 distinct elements	$(x, y=z, y=2)$	T	T	T	T
3 distinct elements	(x, y, z)	T	F	F	F
1 distinct element	$(x=z, y=2)$	F	T	T	F
2 distinct elements	$(x=2, y=2)$	F	F	T	T

#1 Formalization in FOL

- 3) If there are any tax payers, then all politicians are tax payers. If there are any philanthropists, then all tax payers are philanthropists. Therefore, if there are any tax-paying philanthropists, then all politicians are philanthropists.

$T(x)$: x is a tax payer

$P(x)$: x is a politician

$H(x)$: x is a philanthropist

$$\exists x (T(x) \rightarrow (\forall x (P(x) \rightarrow T(x))))$$

$$\exists x (H(x) \rightarrow \forall x (T(x) \rightarrow H(x))),$$

$$\vdash \exists x (T(x) \wedge H(x)) \rightarrow \forall x (P(x) \rightarrow H(x))$$

#1 Formalizing in FOL

- 4) It is possible to specify a formula for models with a specified concrete number of elements; for example a model with exactly three distinct elements or even at most three distinct elements.

When not specifying a concrete bound as in the previous two examples, it is possible to specify an upper bound with the " $\bigwedge_{i=1}^k$ " operator. Within this operator, the upper bound is 'k', where 'k' is an arbitrary natural number. A subset bounded by k is a subset of infinitely many elements.

A subset of an infinite set is still an infinite set. Therefore, there can be no formula in first-order logic that is true only for models with finitely many distinct elements.