

## HW 2.b — Jessup Barrueco

**Q1** D-H,  $q=71$ ,  $\alpha=7$

**a** Given A has private key  $x_A=5$ , what is pub. key  $Y_A$ ?

$$Y_A = \alpha^{x_A} \bmod q = 7^5 \bmod 71 = \boxed{51}$$

**b** B has  $x_B=12$ , find  $Y_B$ :

$$Y_B = 7^{12} \bmod 71 = \boxed{4}$$

**c** Shared secret key:

$$4^5 \bmod 71 = 51^{12} \bmod 71 = \boxed{30}$$

**d** If participants exchanged  $x^\alpha \bmod q$  instead of  $\alpha^x \bmod q$ , an attacker would be able to easily recover the original value of  $x$ . Because  $\alpha$  is a publicly known value,  $x^\alpha \bmod q$  can be reduced  $\alpha$ -times to recover  $x \bmod q$ .

Q2

X appends 64-bit hash code, encrypted w/X's private key.

(a) The Birthday Attack is carried out as follows. Attacker generates  $2^{m/2}$  variations of a valid msg in msg-space  $m$  ( $m=64$  in given question params), these msgs all have a similar (valid) meaning. Then, the attacker generates  $2^{m/2}$  variations of a desired fraudulent msg. According to the Birthday Paradox, the probability of a fraudulent msg having the same hash value as a valid msg is  $>0.5$ . The user then signs the valid msg, so the attacker is able to use that valid signature with the fraudulent msg that has the same hash value.

(b) For an  $M$ -bit message, the attacker would need  $2^{m/2}$  memory. If  $m=64$ , the attacker would need  $2^{32}$  memory. If  $m=128$ , memory needed =  $2^{64}$  (d)

(c)  $2^{20}$  hash/second,  $2^{\frac{m}{2}+1}$  hashes in total.  
 $\frac{2^{33}}{2^{20}} = 2^{13}$  seconds for  $m=64$

$\frac{2^{65}}{2^{20}} = 2^{45}$  seconds for  $m=128$  (d)

Q3

Trapdoor One-Way Fcn

plaintext  $P = 01010111$

private key  $S = \{5, 9, 21, 45, 103, 215, 450, 946\}$

note: set  $S$  is superincreasing

$a = 1019$        $p = 1999$

step 1: generate public key  $\beta$ :  $\beta_i = a s_i \bmod p$

$\beta = \{1097, 1175, 1409, 1877, 1009, 1194, 779, 456\}$   
          0      1      0      1      0      1      1      1

step 2: add elements of  $\beta \bmod p$  according to plaintext

$$C = 1175 + 1877 + 1194 + 779 + 456 = 5481 \equiv 1483 \bmod 1999$$

step 3: to decrypt the ciphertext  $C$  back to obtain the plaintext,

first find  $a^{-1} \bmod p$ . From Fermat's Little Thm, see

$$a^{-1} \bmod p = a^{p-2} \bmod p = 1019^{1997} \bmod p = 1589$$

$$C \cdot a^{-1} = 1483 \cdot 1589 \equiv 1665 \bmod 1999$$

step 4: verify 1665 is correct plaintext:

$S = \{5, 9, 21, 45, 103, 215, 450, 946\}$   
          0      1      0      1      0      1      1      1

$$9 + 45 + 215 + 450 + 946 = 1665$$

$\Rightarrow$  process successful