Lecture 5 - Tree (Part 1)

CPE112 - Programming with Data Structures 26 February 2025

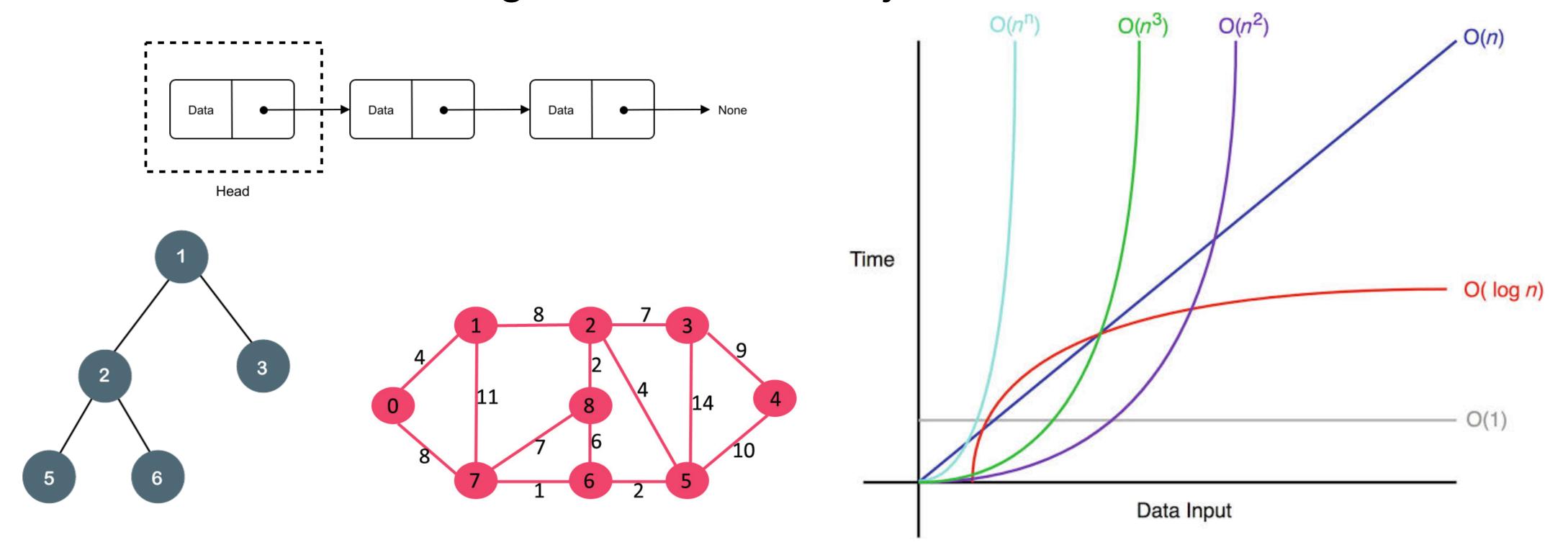
Dr. Piyanit Wepulanon & Dr. Taweechai Nuntawisuttiwong

Department of Computer Engineering KMUTT



Review

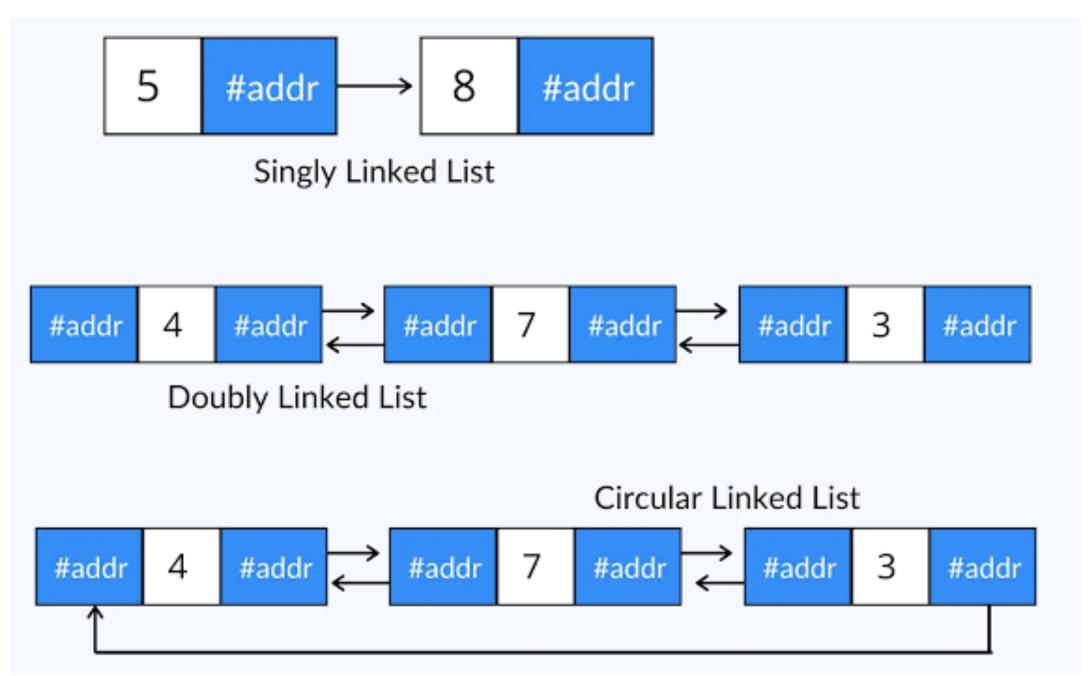
- Introduction to Programming with Data Structures
 - O Data structures, Algorithms, Efficiency



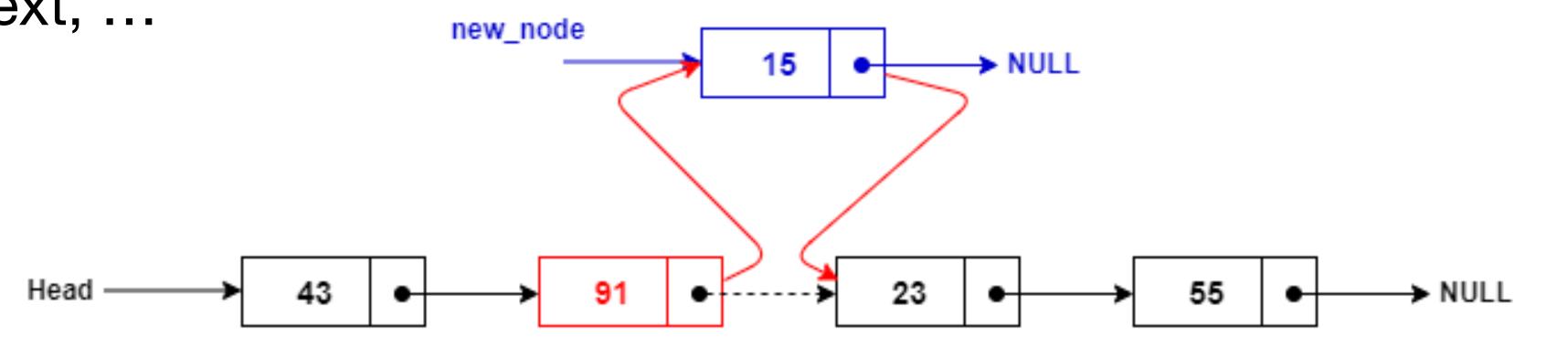
Review

- List (Array & Linked List)
 - Pros & cons, capacity, accessibility, insert/remove elements
 - Structure

Types of Linked List

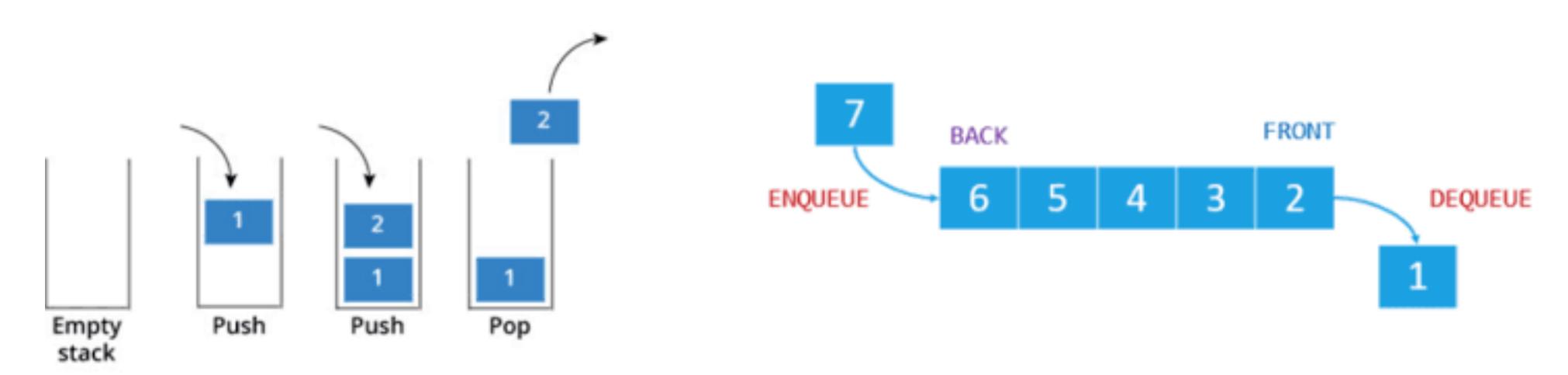


 Main functions: Create, Destroy, Size, InsertAt, Append, RemoveAt, Get, Reset, GetNext, ...



Review

- Stack & Queue
 - Structure
 - Main functions (Create, Destroy, Size, Push, Pop, Enqueue, Dequeue)



Linear Data Structure

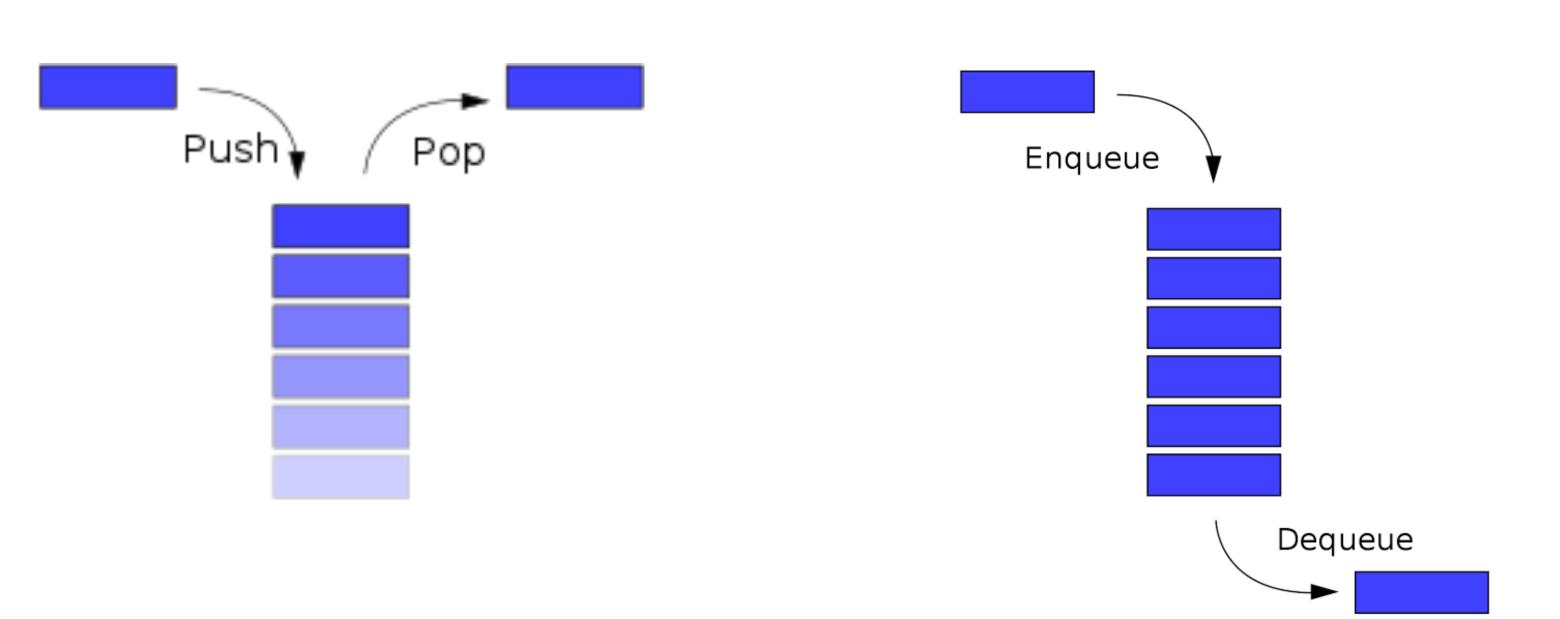
Data structure

 100
 101
 102
 103
 104
 105
 106

 A
 B
 C
 D
 E
 F
 G

 0
 1
 2
 3
 4
 5
 6

- Sequentially arrange
- Start/End of structure
- Previous/Next element

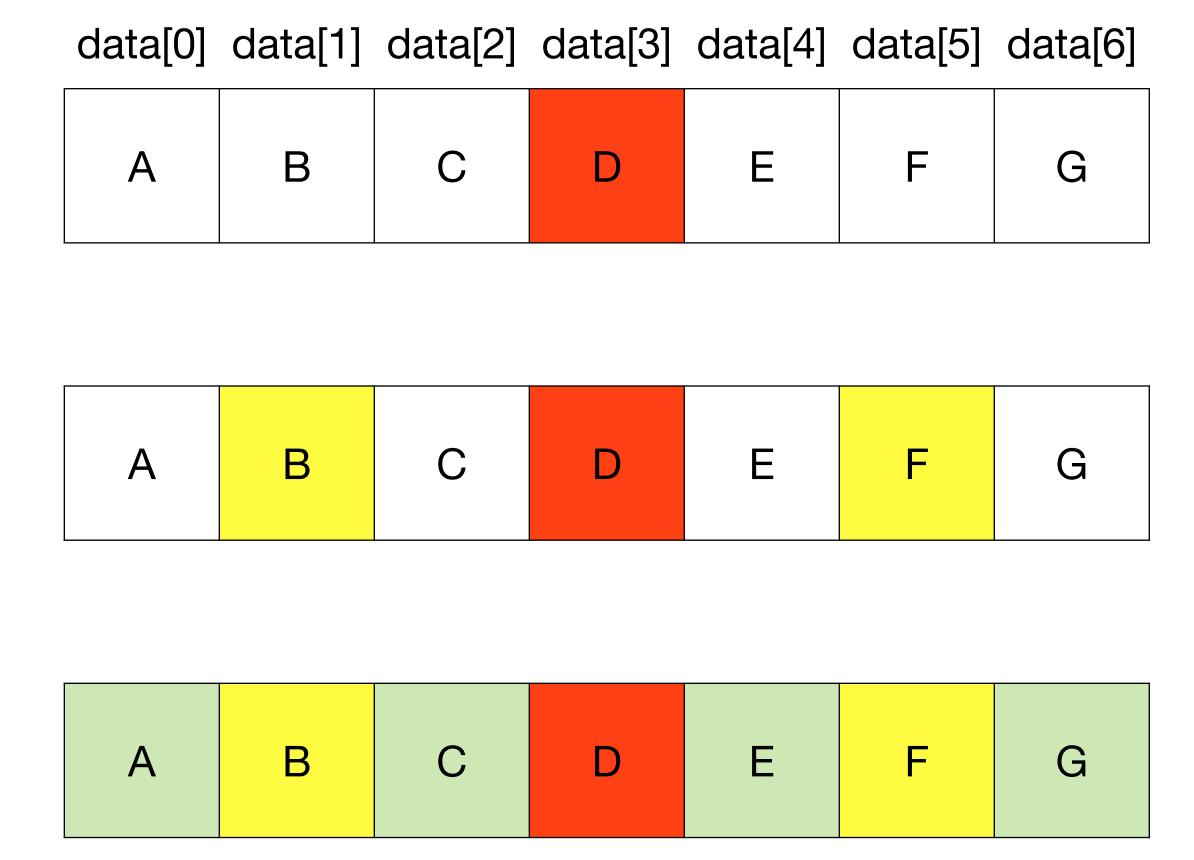




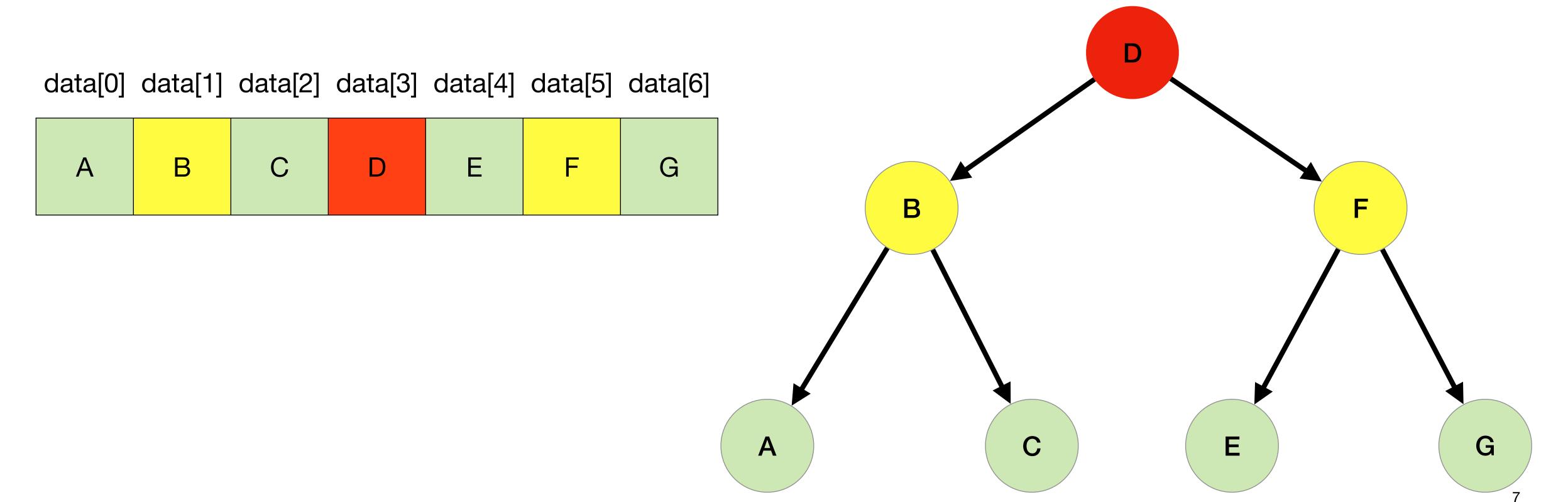
Linear Data Structure

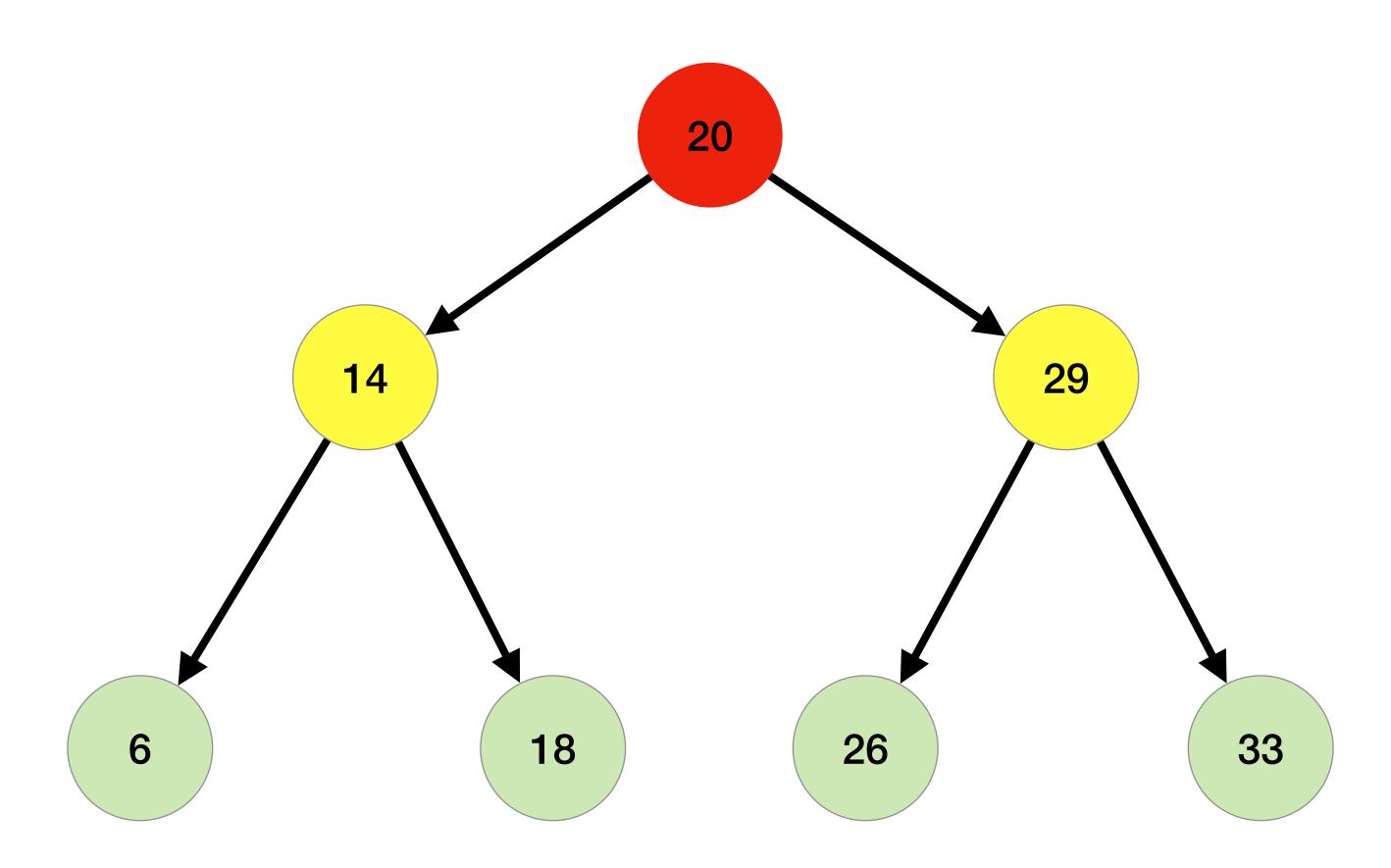
Data structure

Example binary search



Example binary search

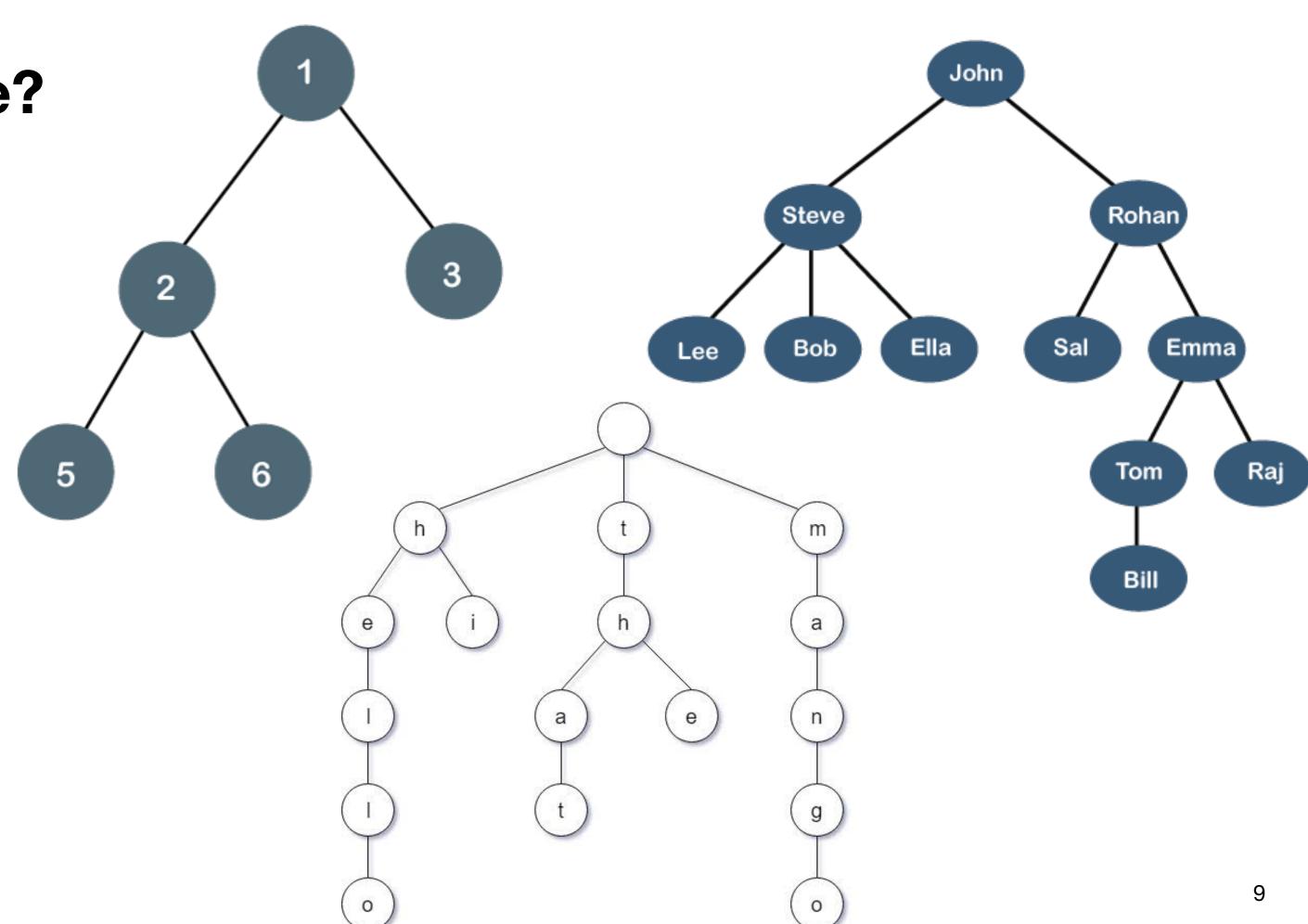




Today's Goal

What is "tree" data structure?

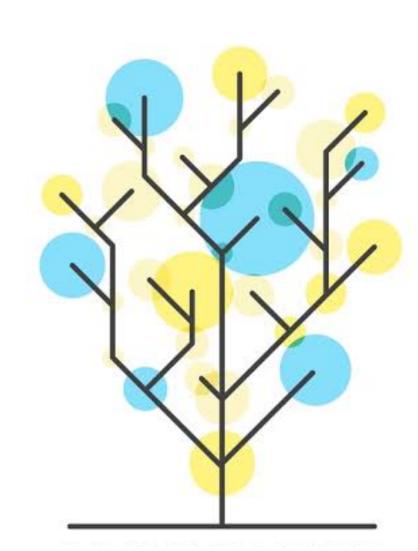
- Terminologies
- Binary tree
- Tree traversal
- Lab practice



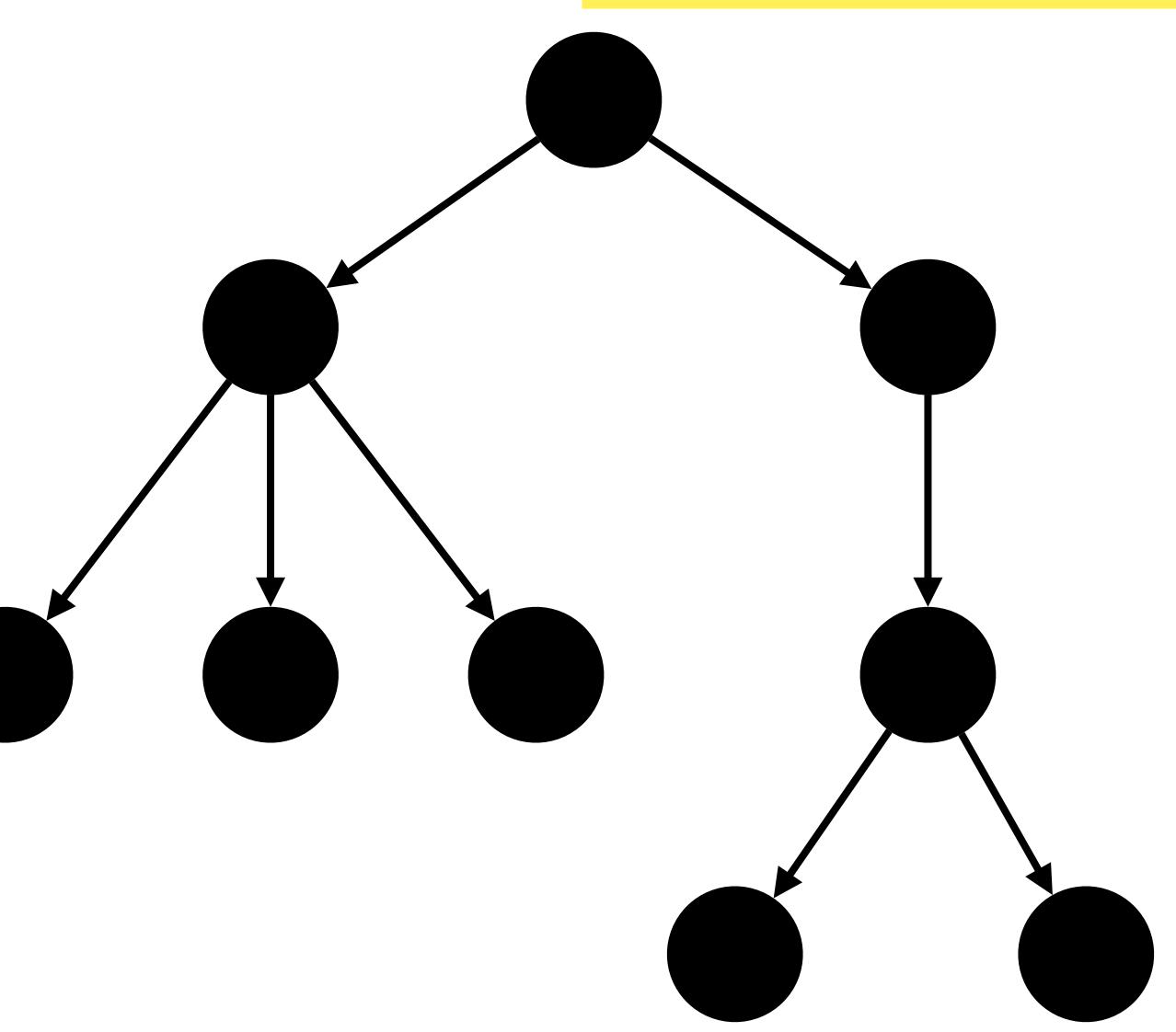
Tree data structure

Define with a logical model

- Node/Edge/Path
- Parent/Children/Sibling
- Ancestor/Descendant
- Root/Leaf



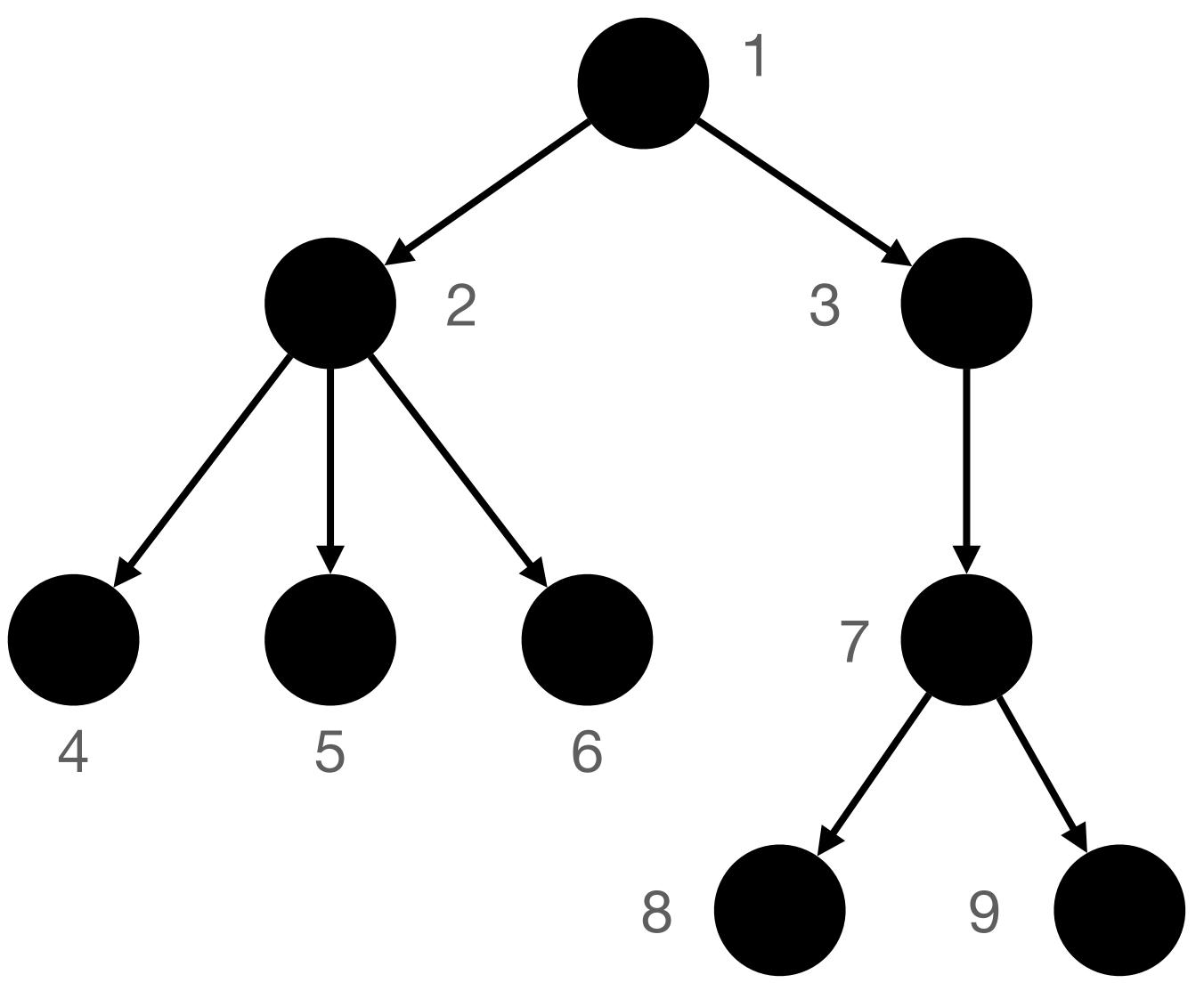
N nodes -> N-1 edges
(1 incoming edge for each node)



Tree data structure

Parent-child relationship

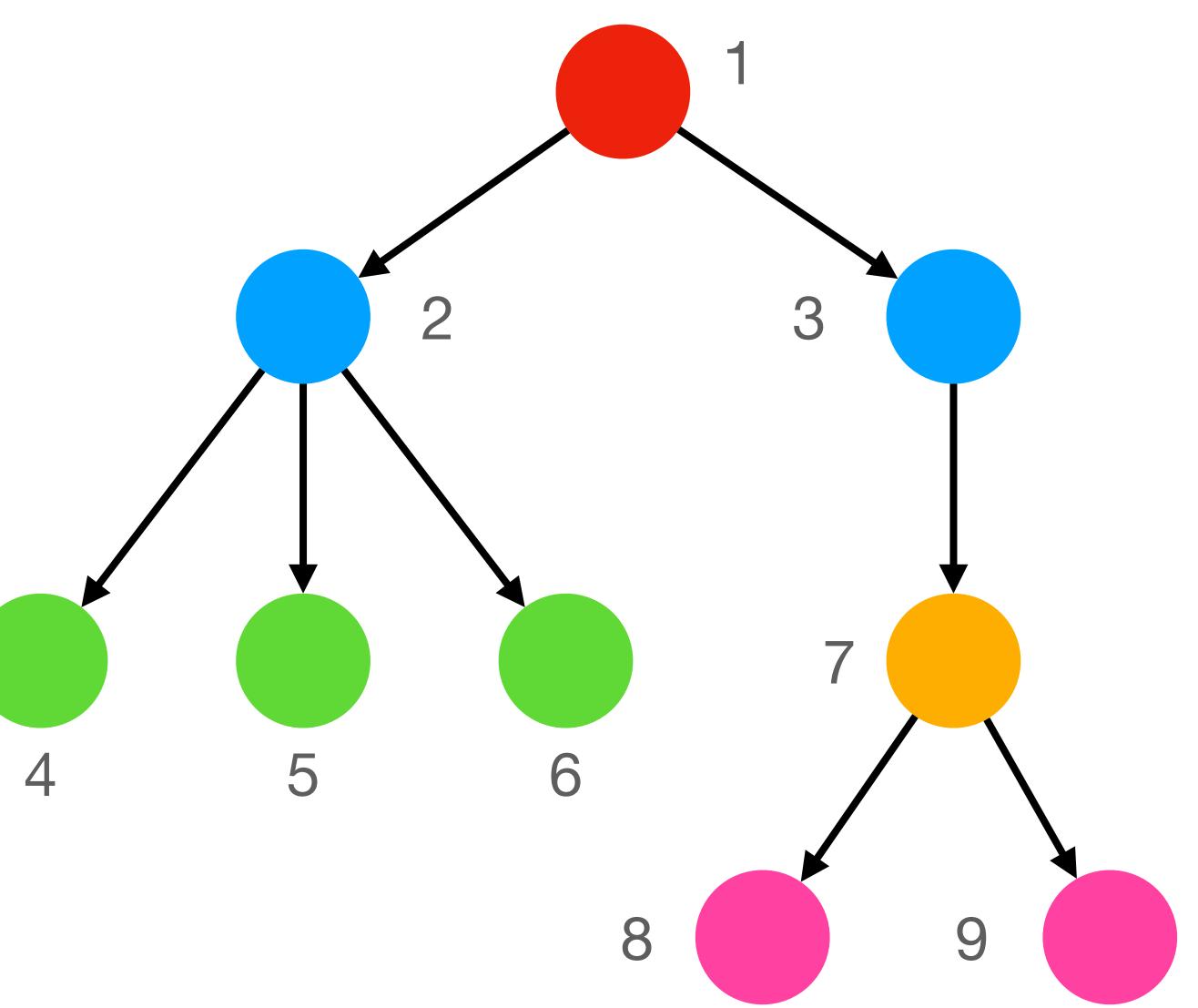
Node	Parent	Children
2	1	4,5,6
7		
3		
1		
5		



Tree data structure

Parent-child relationship

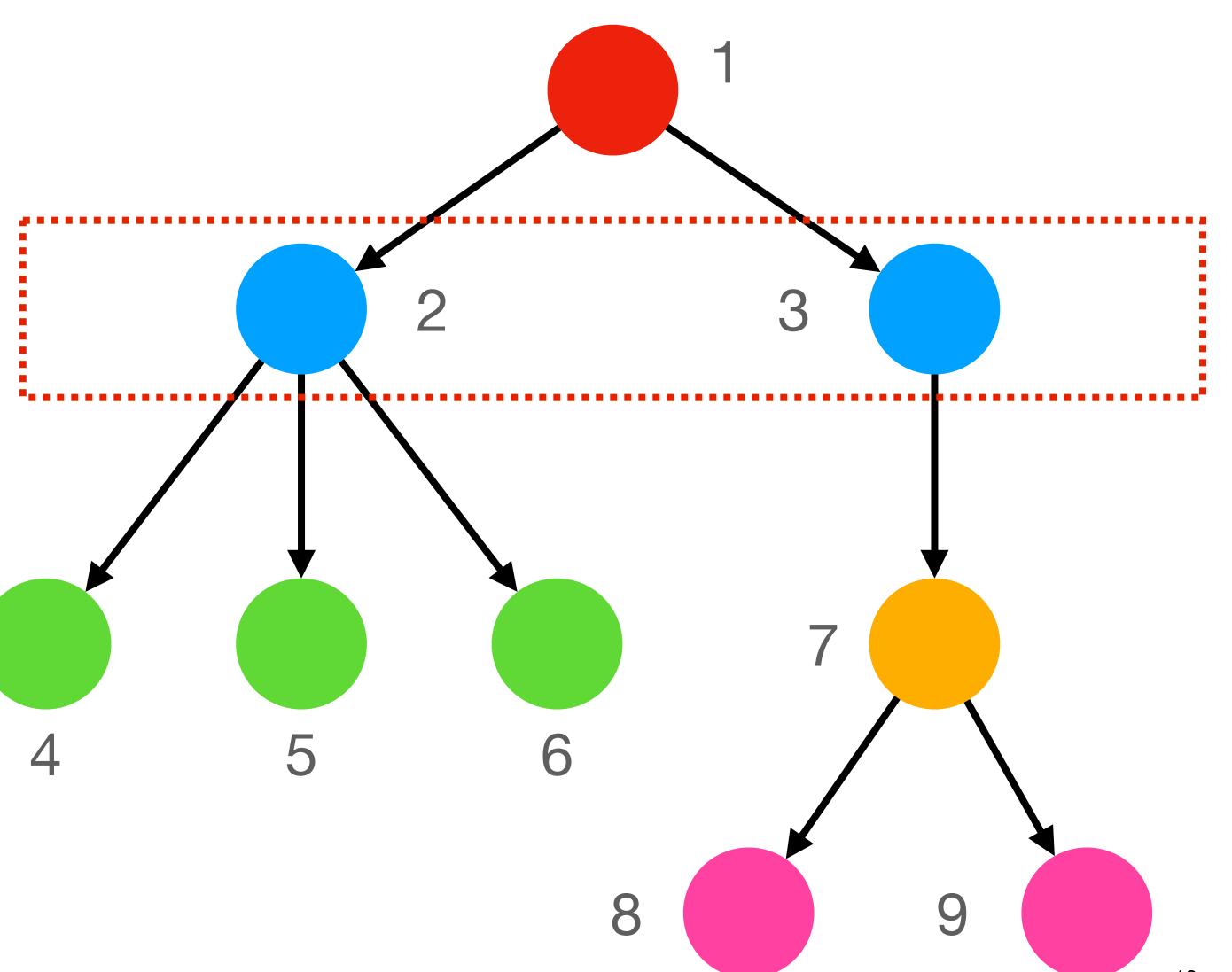
Node	Parent	Children
2	1	4,5,6
7	3	8,9
3	1	7
1	_	2,3
5	2	



Tree data structure

Sibling

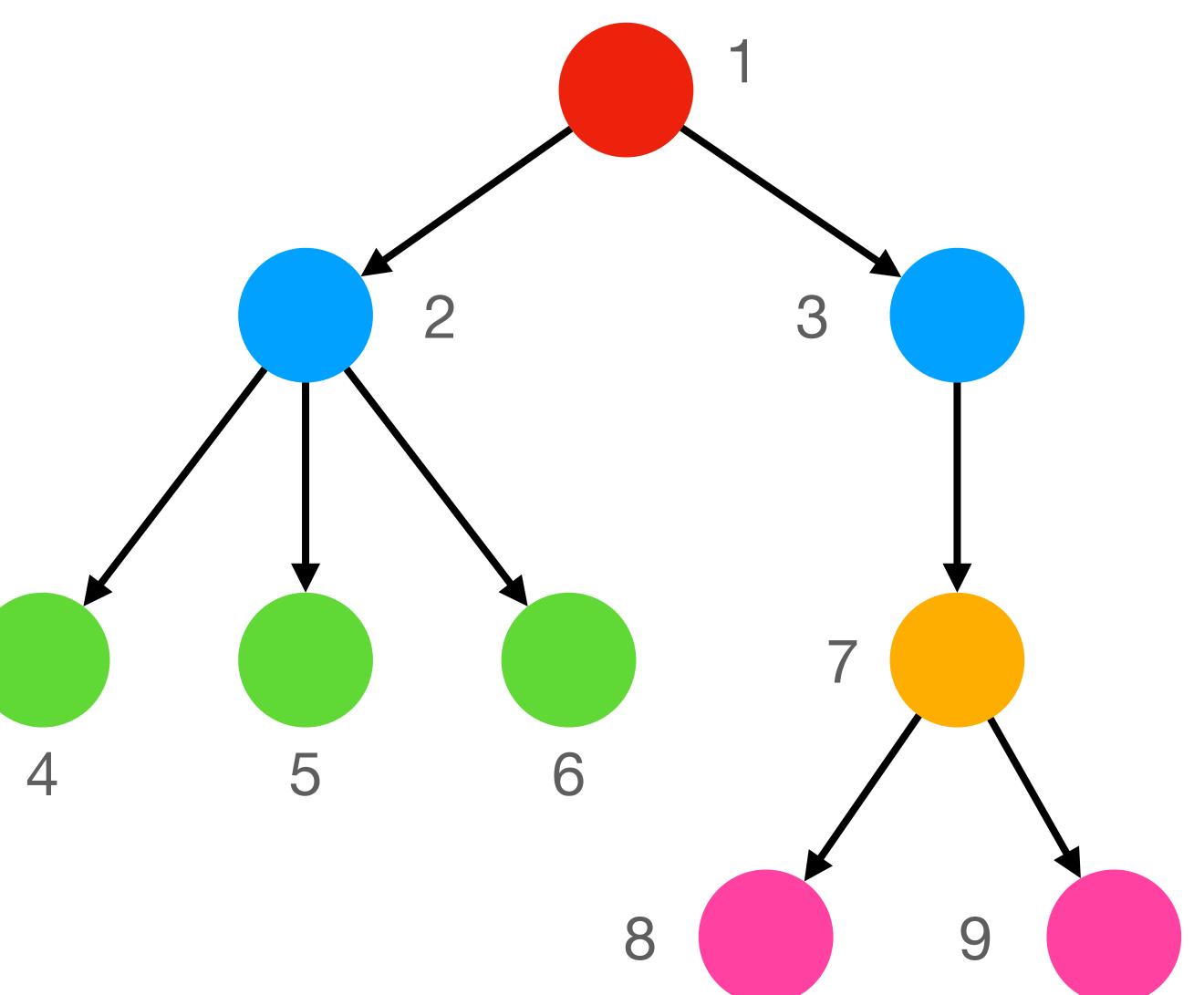
Node	Parent	Children
2	1	4,5,6
7	3	8,9
3	1	7
1	_	2,3
5	2	_



Tree data structure

Root/Leaf

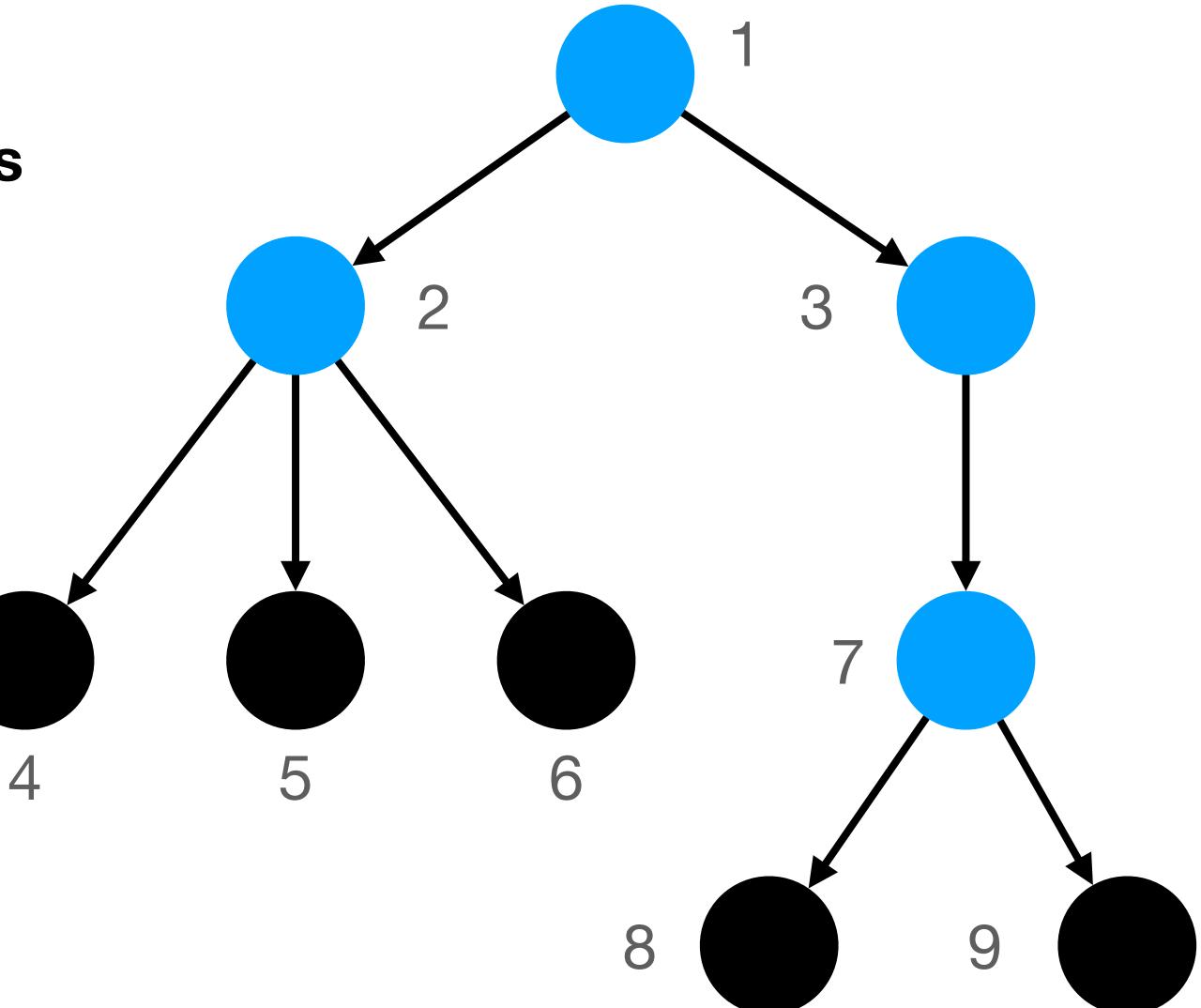
Node	Parent	Children
2	1	4,5,6
7	3	8,9
3	1	7
1	_	2,3
5	2	_



Tree data structure

Internal or non-terminal nodes

External or terminal nodes

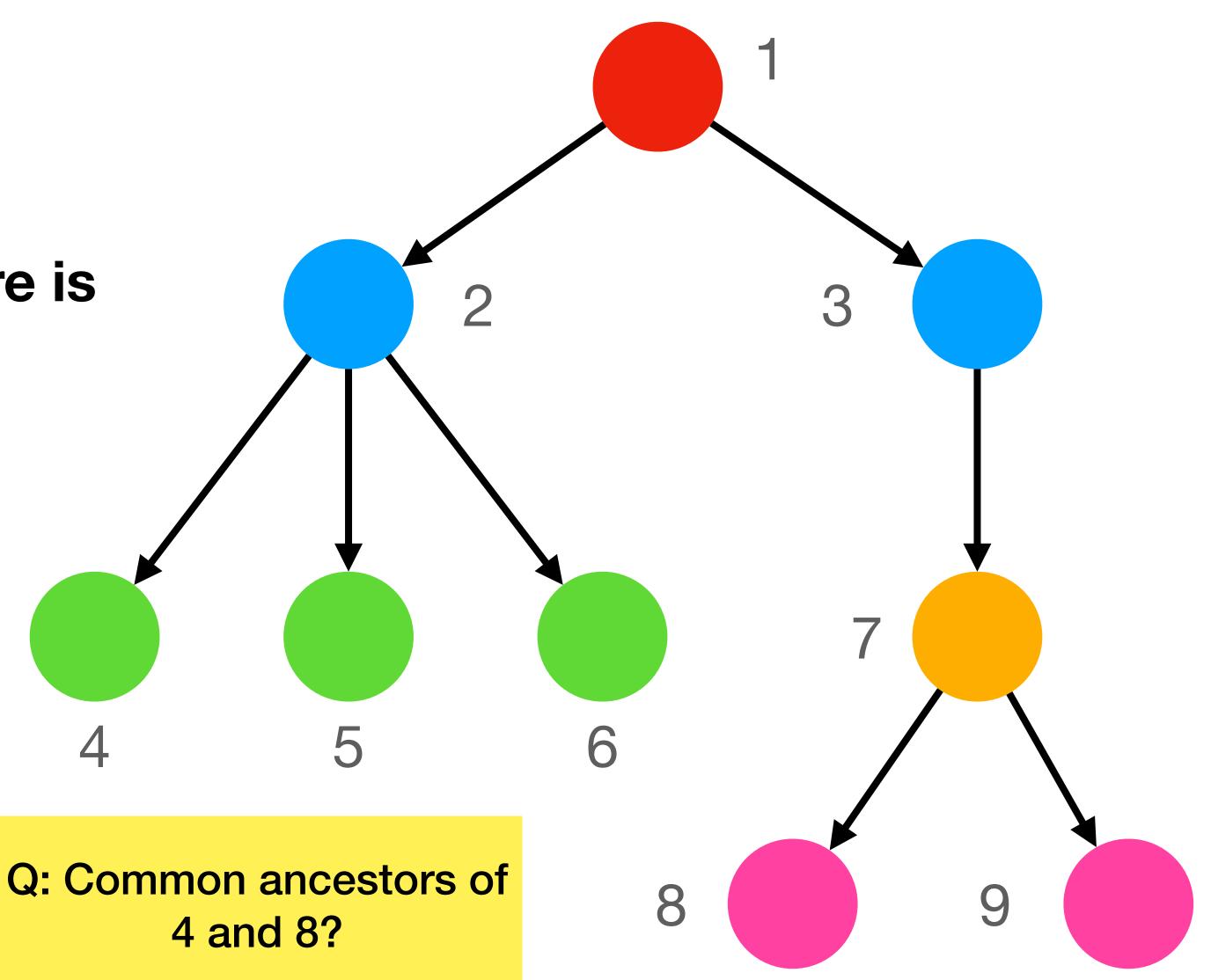


Tree data structure

Ancestor/Descendant

 A is an ancestor of B if there is a path from A to B

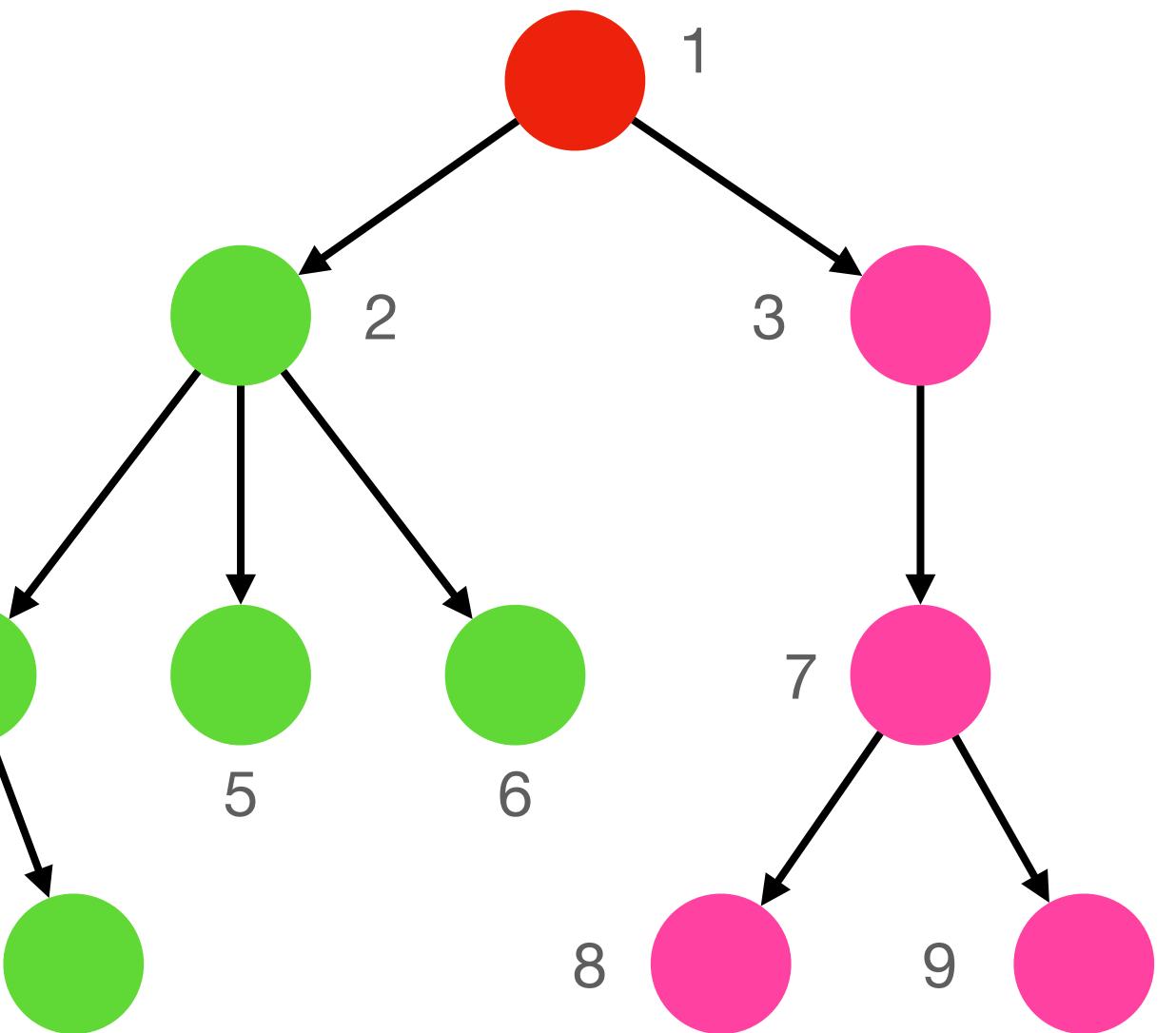
Node	Ancestors
4	1,2
7	
8	



Recursive data structure

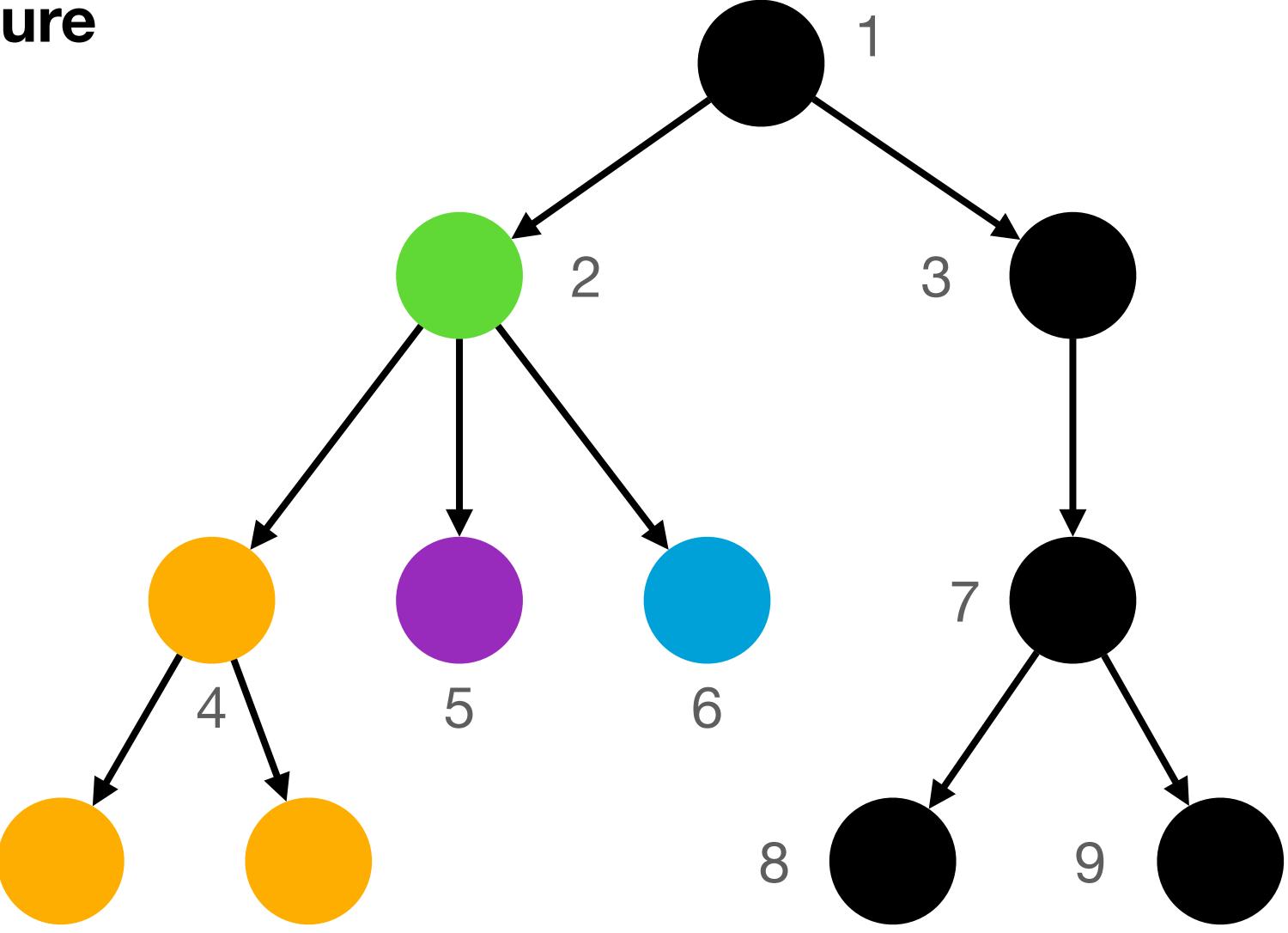
Sub-tree

• Forest: There are 2 trees if node 1 is removed.



Recursive data structure

Sub-tree



node 2 = ?

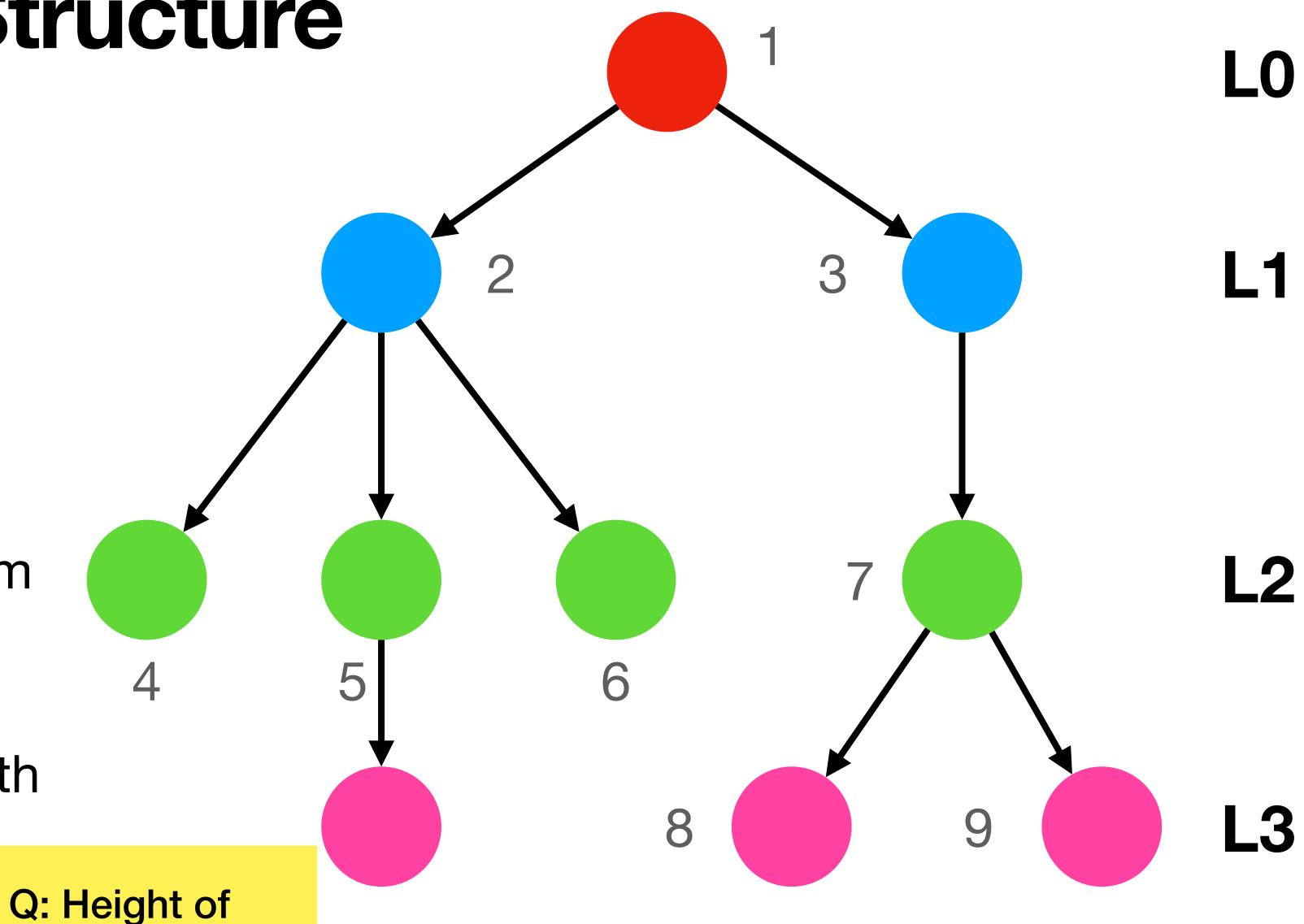
Tree data structure

• Level: 4 levels

Path length: #edges
 between 2 nodes

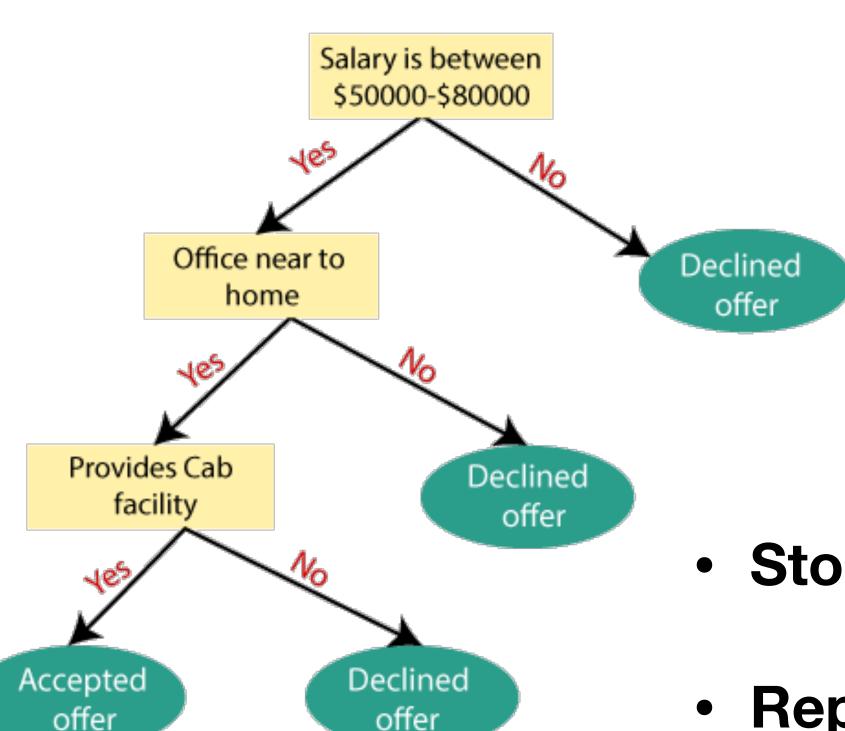
 Depth: path length from the root (= level, L)

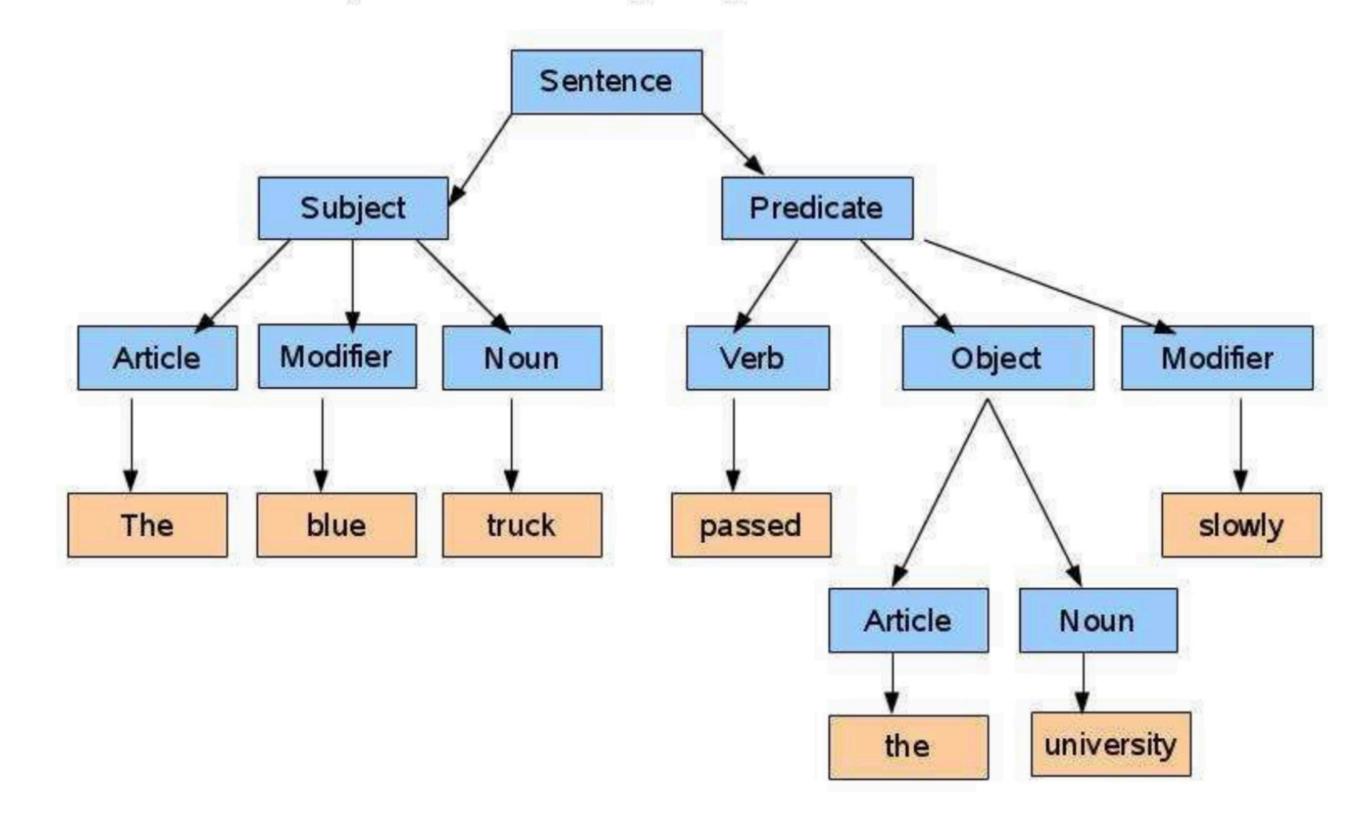
 Height: the longest path length to the leaf



Applications

Tree data structure





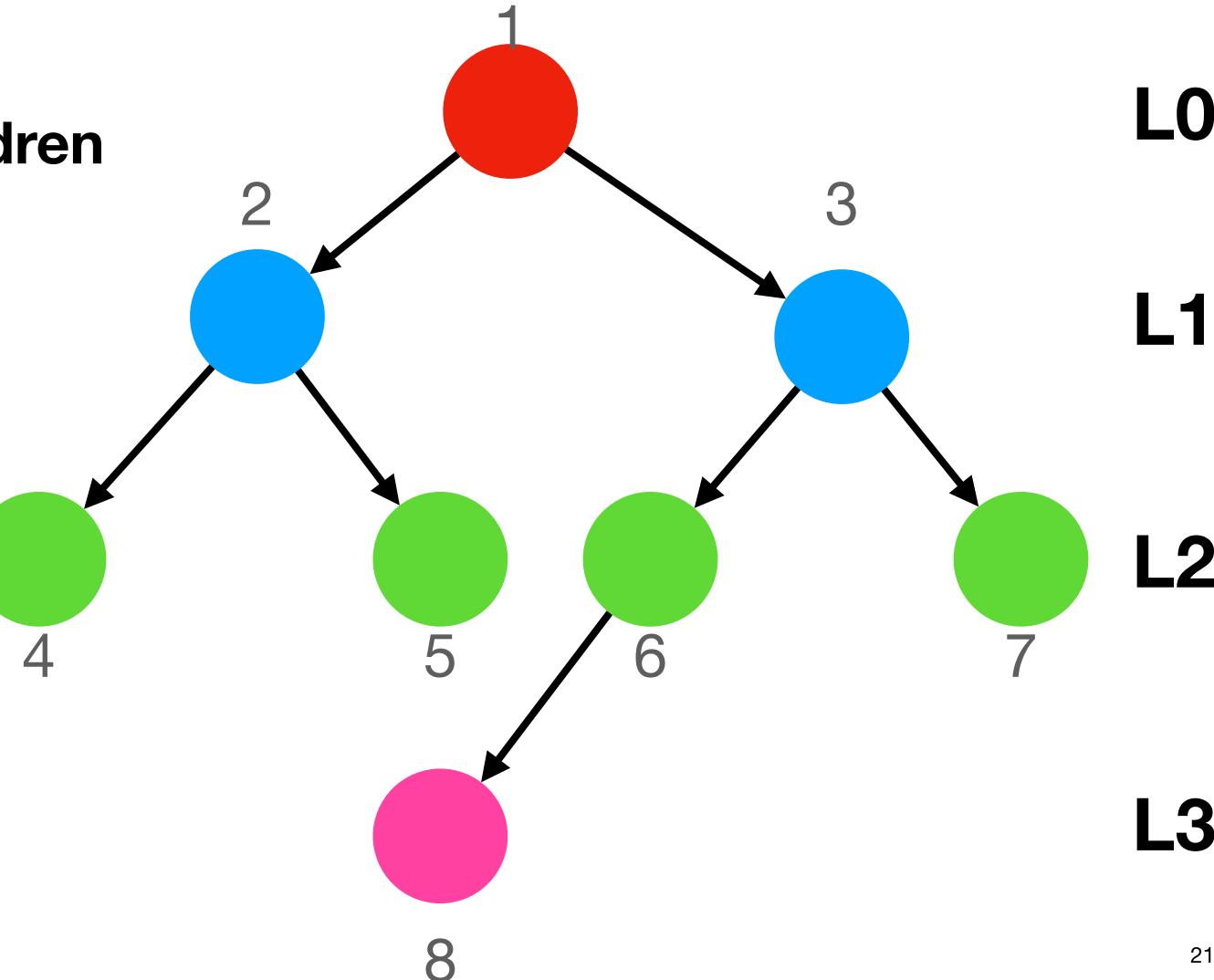
- Storing hierarchical data: file system, family tree, decision tree
- Represent structure of a sentence in natural language
- Routing information

Binary tree

Any node have at most 2 children

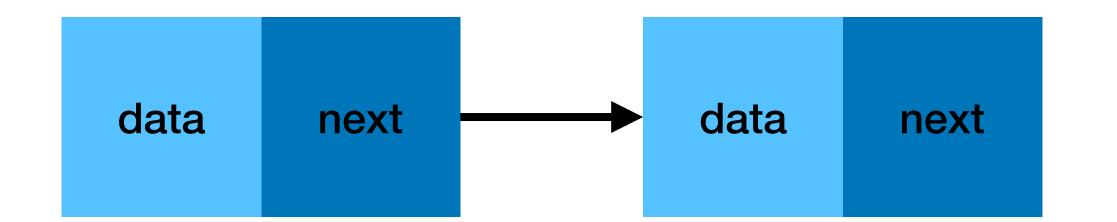
Left Subtree & Right Subtree

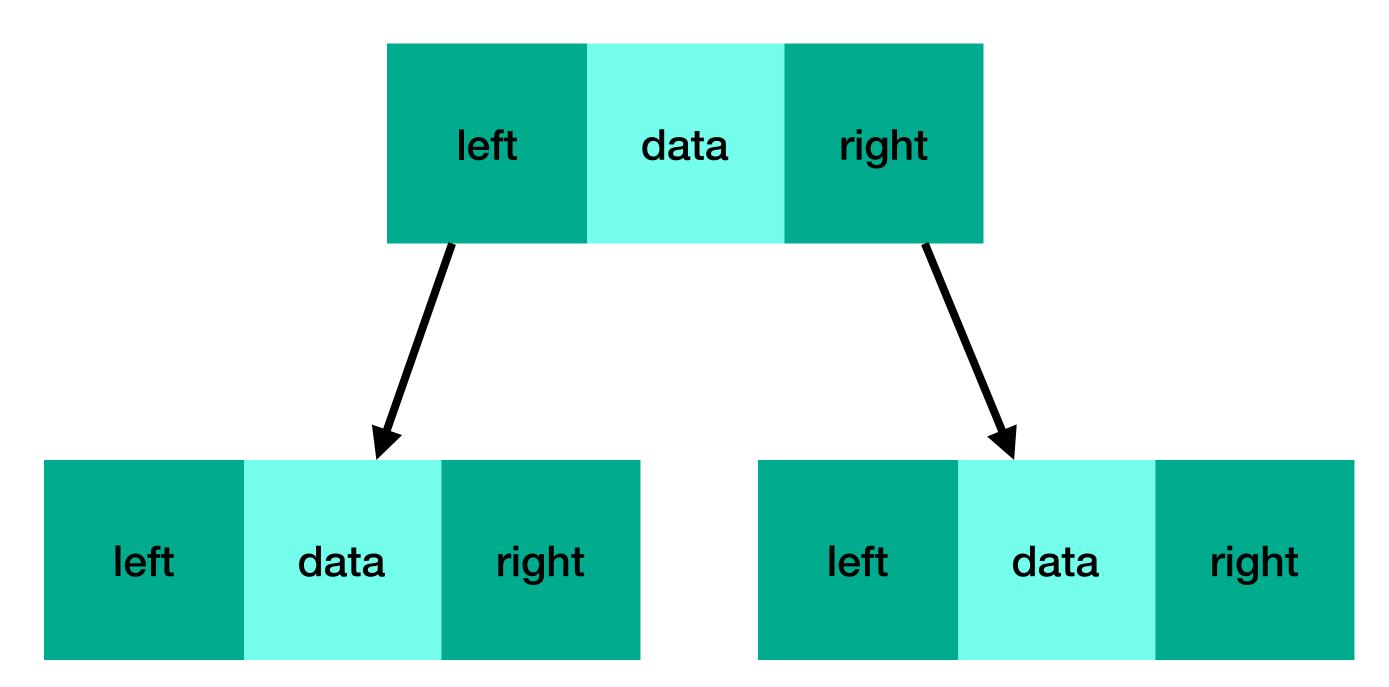
Left Child & Right Child



Node Structure

Linked List vs Tree





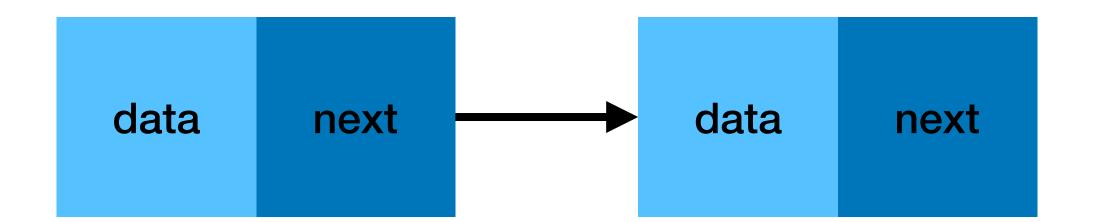
```
typedef struct _listnode
{
  int data;
  struct _listnode* next;
} LISTNODE_T;
```

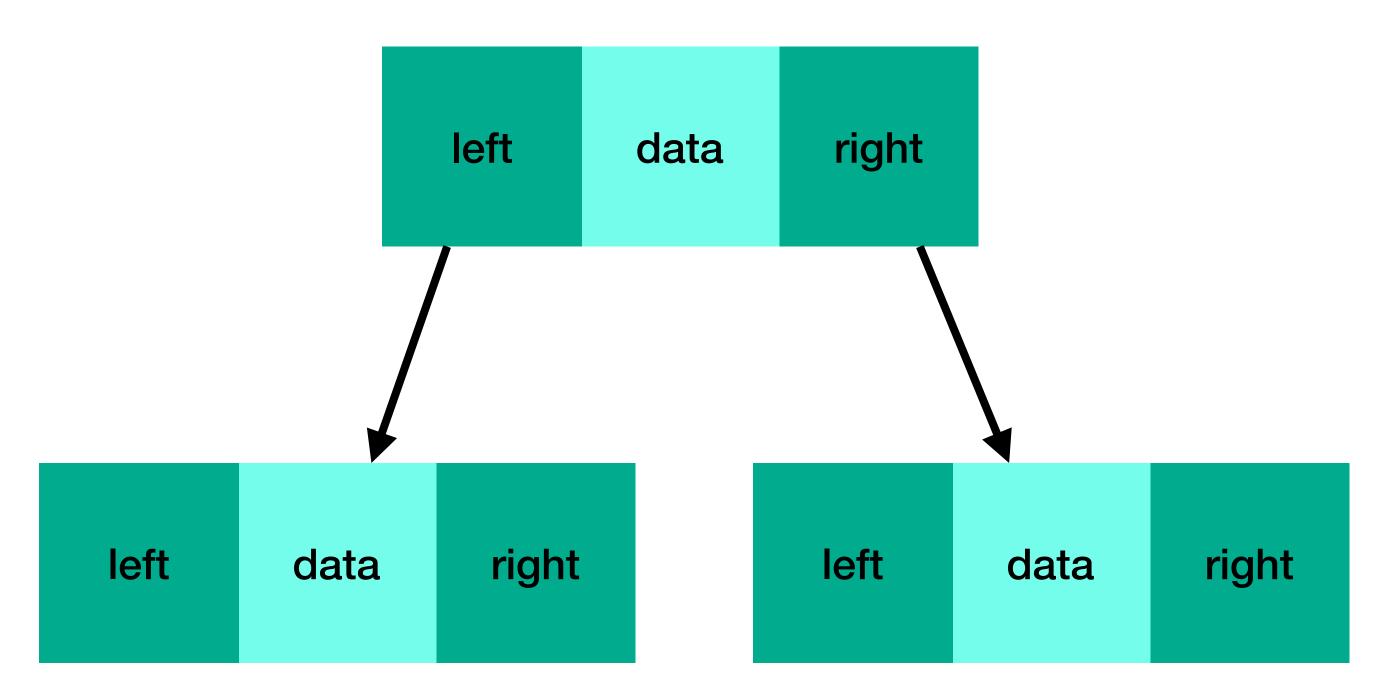
```
typedef struct _treenode
{

TREENODE_T;
```

Node Structure

Linked List vs Tree





```
typedef struct _listnode
{
   int data;
   struct _listnode* next;
} LISTNODE_T;
```

```
typedef struct _treenode
{
    int data;
    struct _treenode* left;
    struct _treenode* right;
} TREENODE_T;
```

Representation of binary trees

```
typedef struct _treenode
{
  int data;
  struct _treenode* left;
  struct _treenode* right;
} TREENODE_T;
```

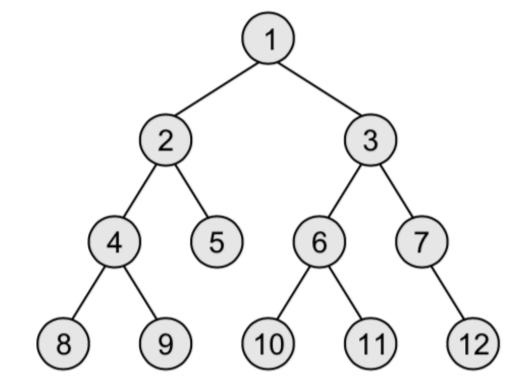


Figure 9.10 Binary tree T

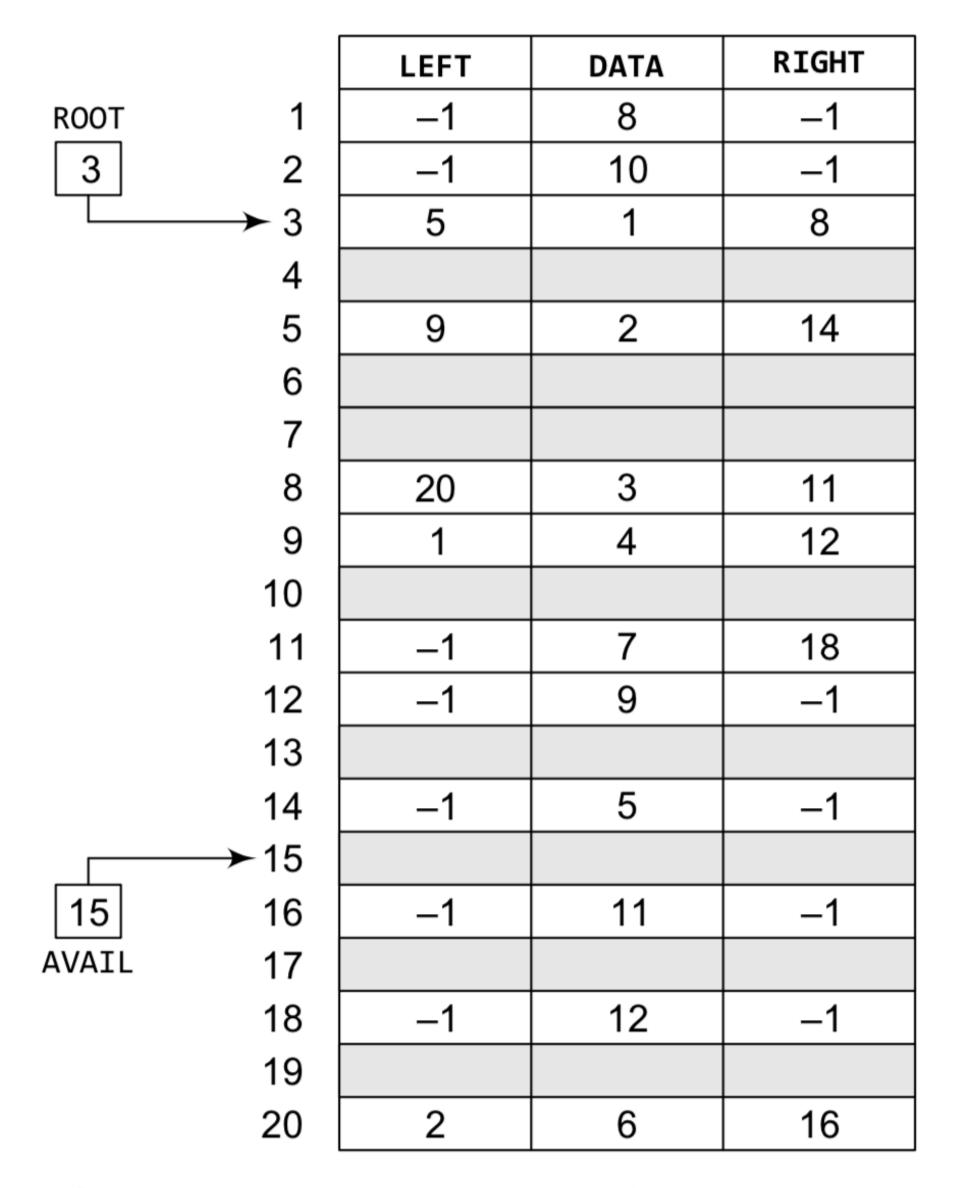


Figure 9.11 Linked representation of binary tree T

Representation of binary trees

Janak

Ridhiman

Pallav

Rudraksh

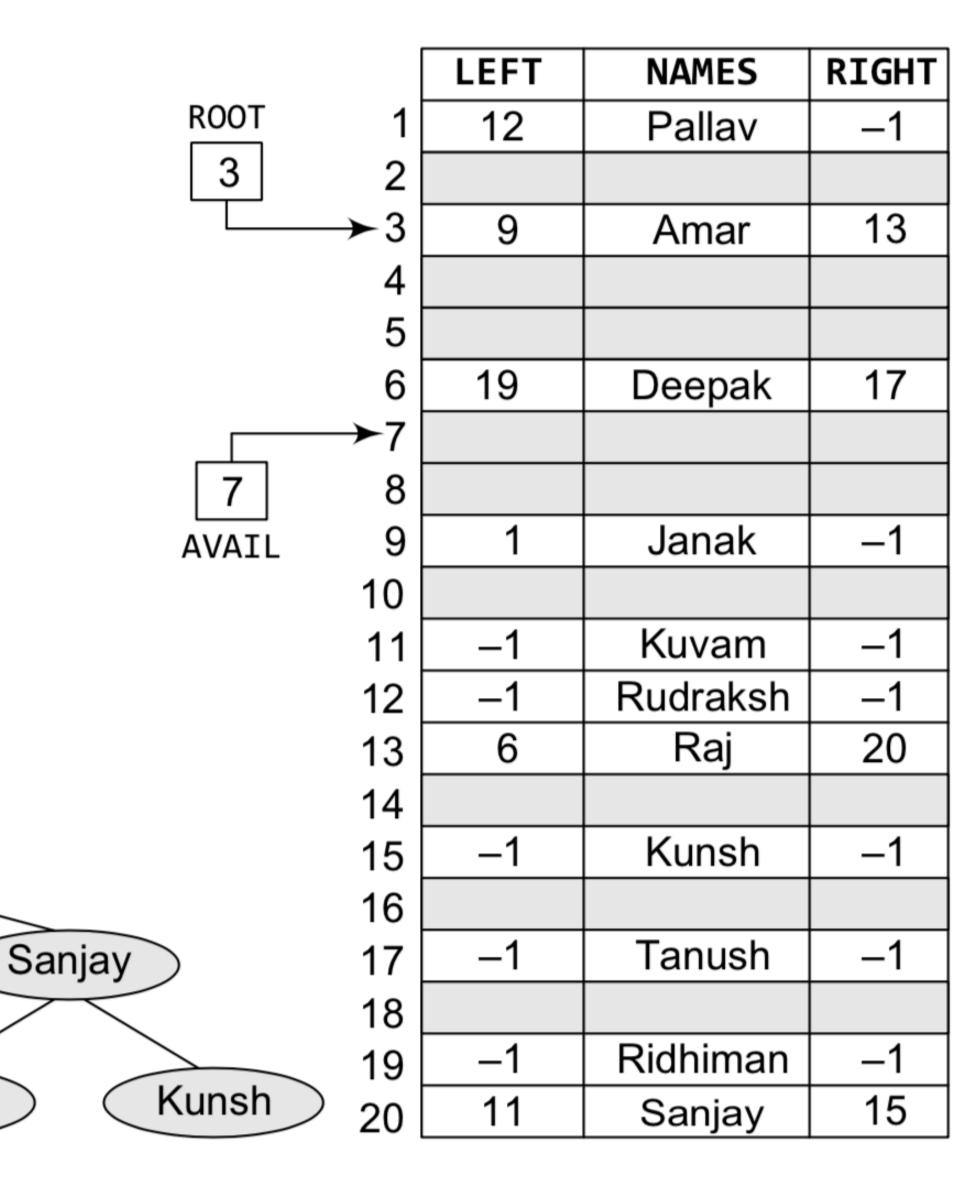
Amar

Deepak

Raj

Kuvam

Tanush



Custom - Structure & New node

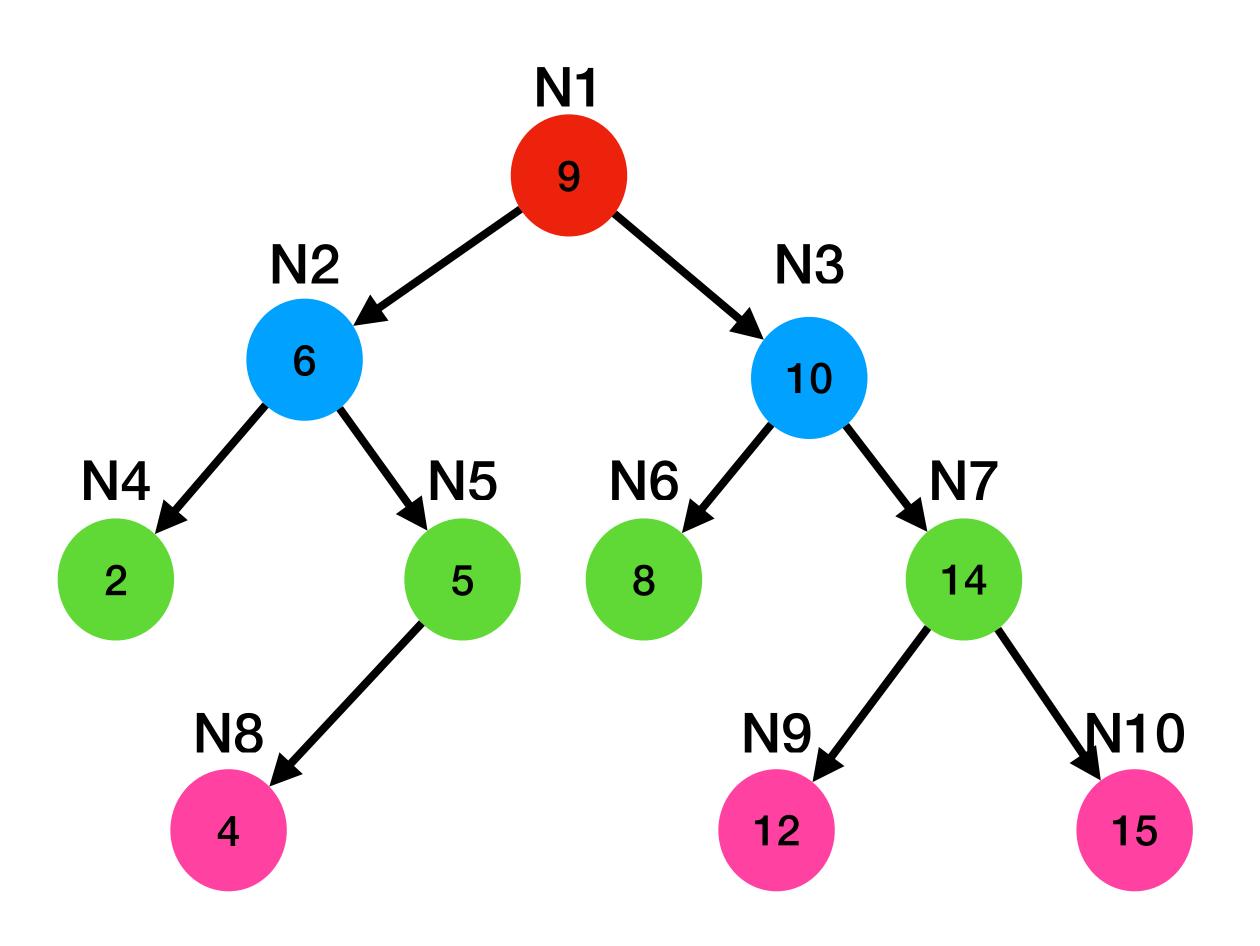
tree.h

```
typedef struct _treenode
{
  int data;
  struct _treenode* left;
  struct _treenode* right;
} TREENODE_T;
```

tree.c

```
TREENODE_T *newNodeCreate(int item)
{
   TREENODE_T* new_node;
   new_node = (TREENODE_T*)
        calloc(1, sizeof(TREENODE_T));
   new_node->data = item;
   new_node->left = NULL;
   new_node->right = NULL;
   return new_node;
}
```

Custom - Create new node



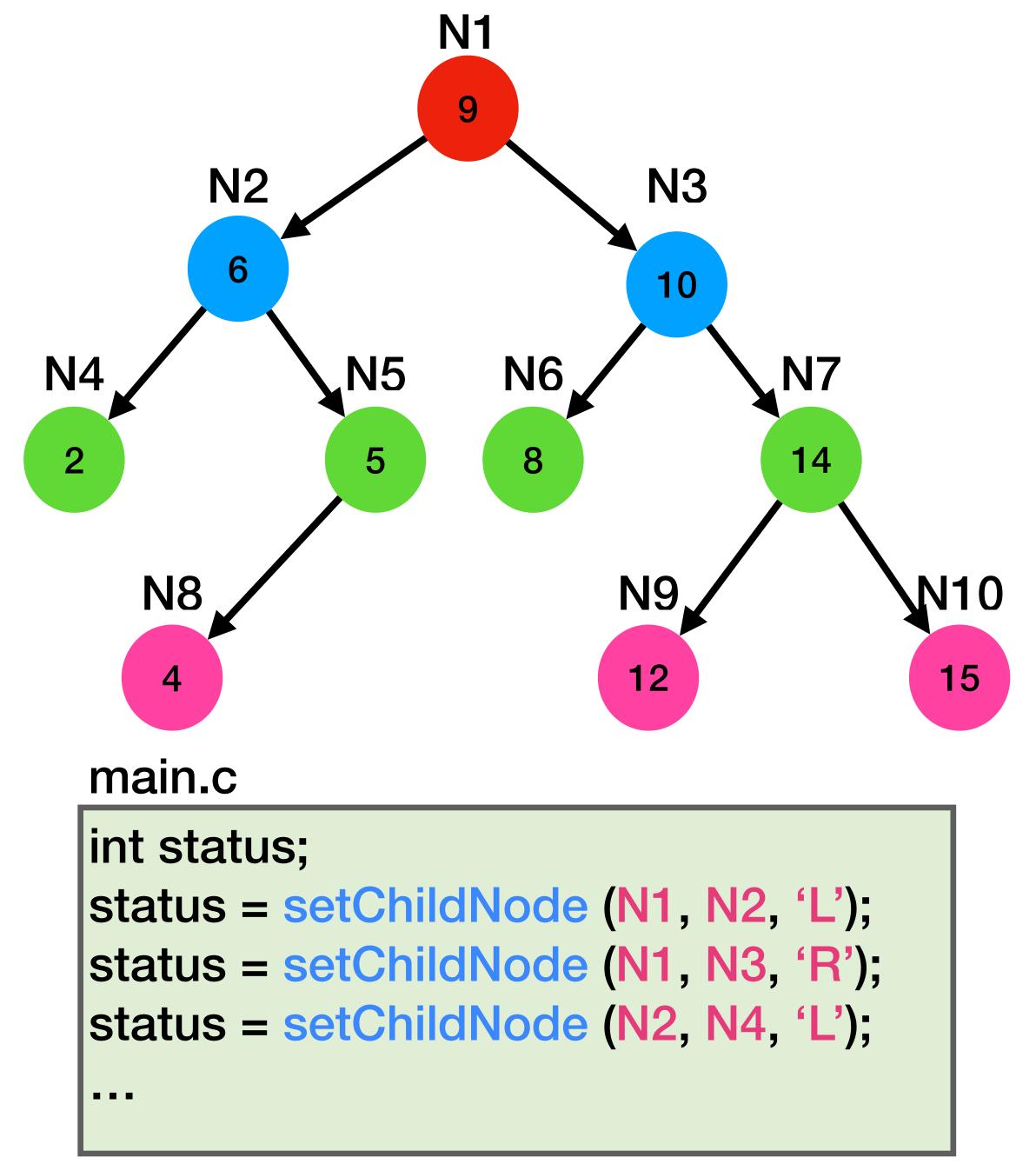
main.c

```
TREENODE_T* N1 = newNodeCreate(9);
TREENODE_T* N2 = newNodeCreate(6);
TREENODE_T* N3 = newNodeCreate(10);
TREENODE_T* N4 = newNodeCreate(2);
...
```

Custom - Child node

tree.c

```
int setChildNode (TREENODE_T* parent,
TREENODE_T* child, char direction)
 int success = 1;
  //check null node or other conditions?
 if (direction == 'L')
    parent->left = child;
 else if (direction == 'R')
    parent->right = child;
  else
   success = 0;
 return success;
```

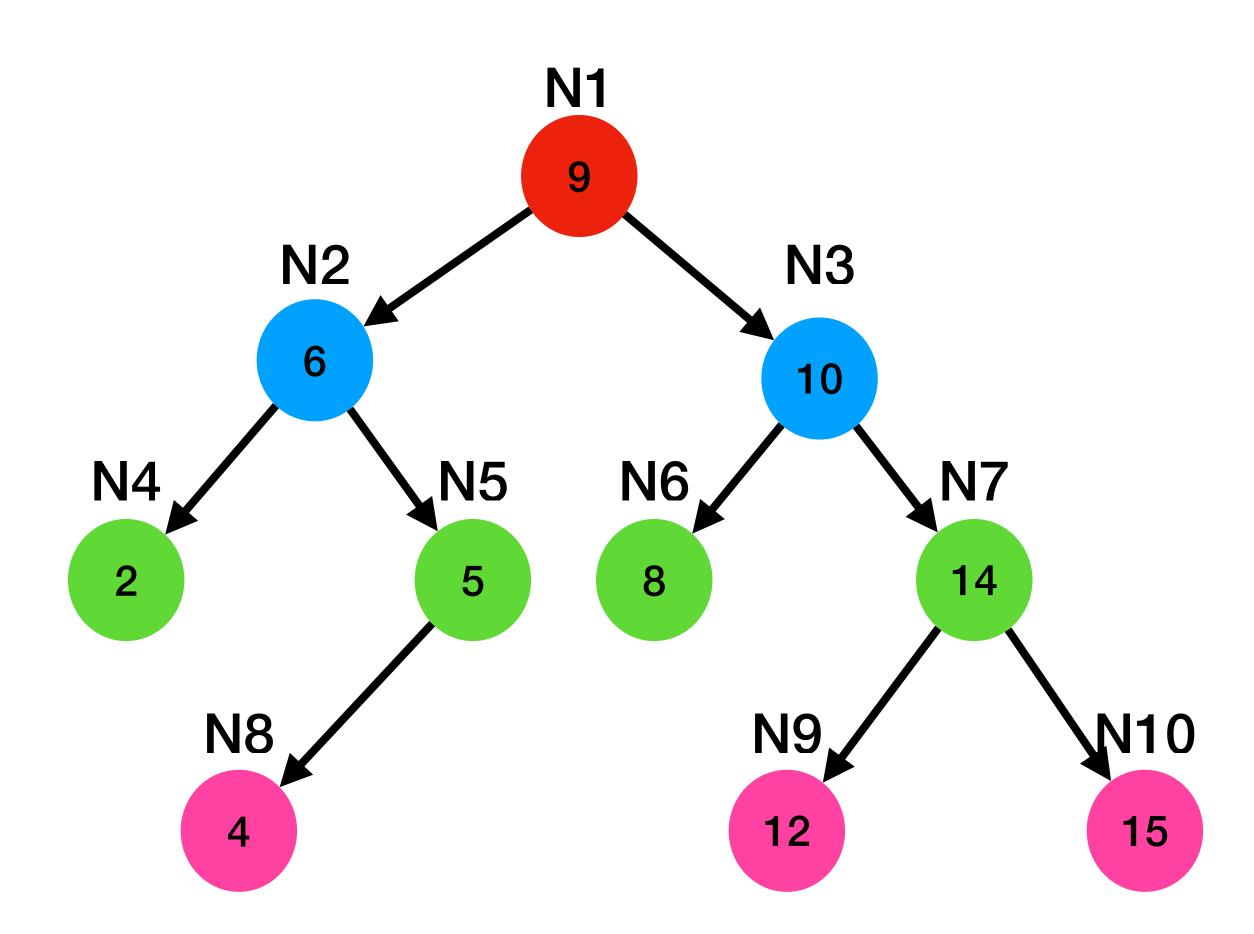


Custom - Free node

main.c

```
free (N10);
N7->right = NULL;
free (N9);
free (N7);
...
```

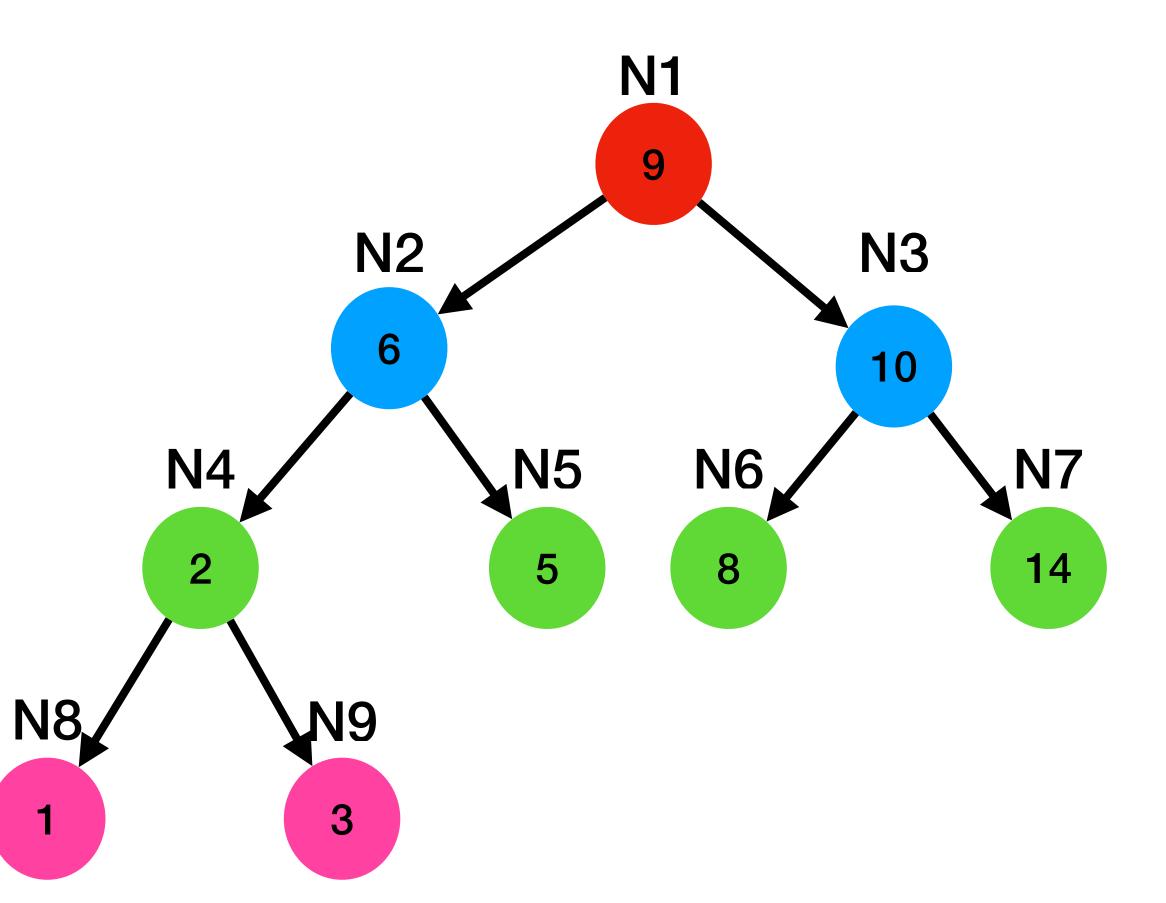
Print to check data before free EX: N1->left->data



Complete Binary Trees

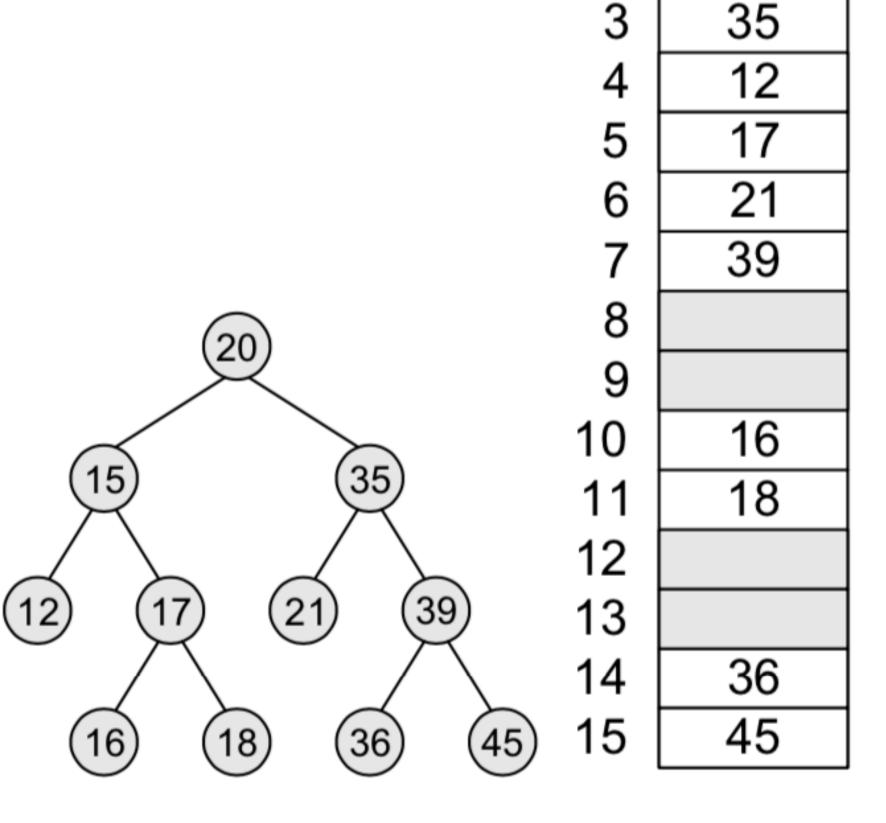
Properties

- Every level except possibly the last is completely filled.
- All nodes appear as far left as possible.
- A level r can have at most 2^r nodes



Sequential representation of binary trees

- 1D Array
- The root will be stored in the first location.
- The children of a node in location K will be stored in locations (2 x K) and (2 x K + 1).

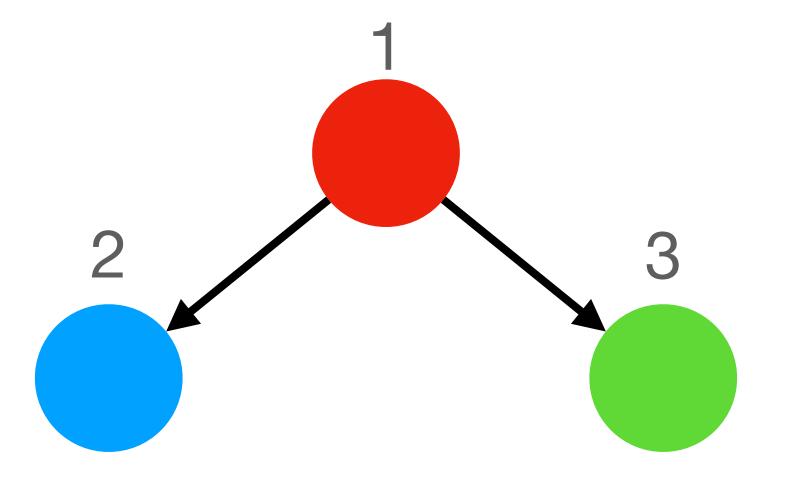


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Tree Traversal

Depth-first



R: Root Node,

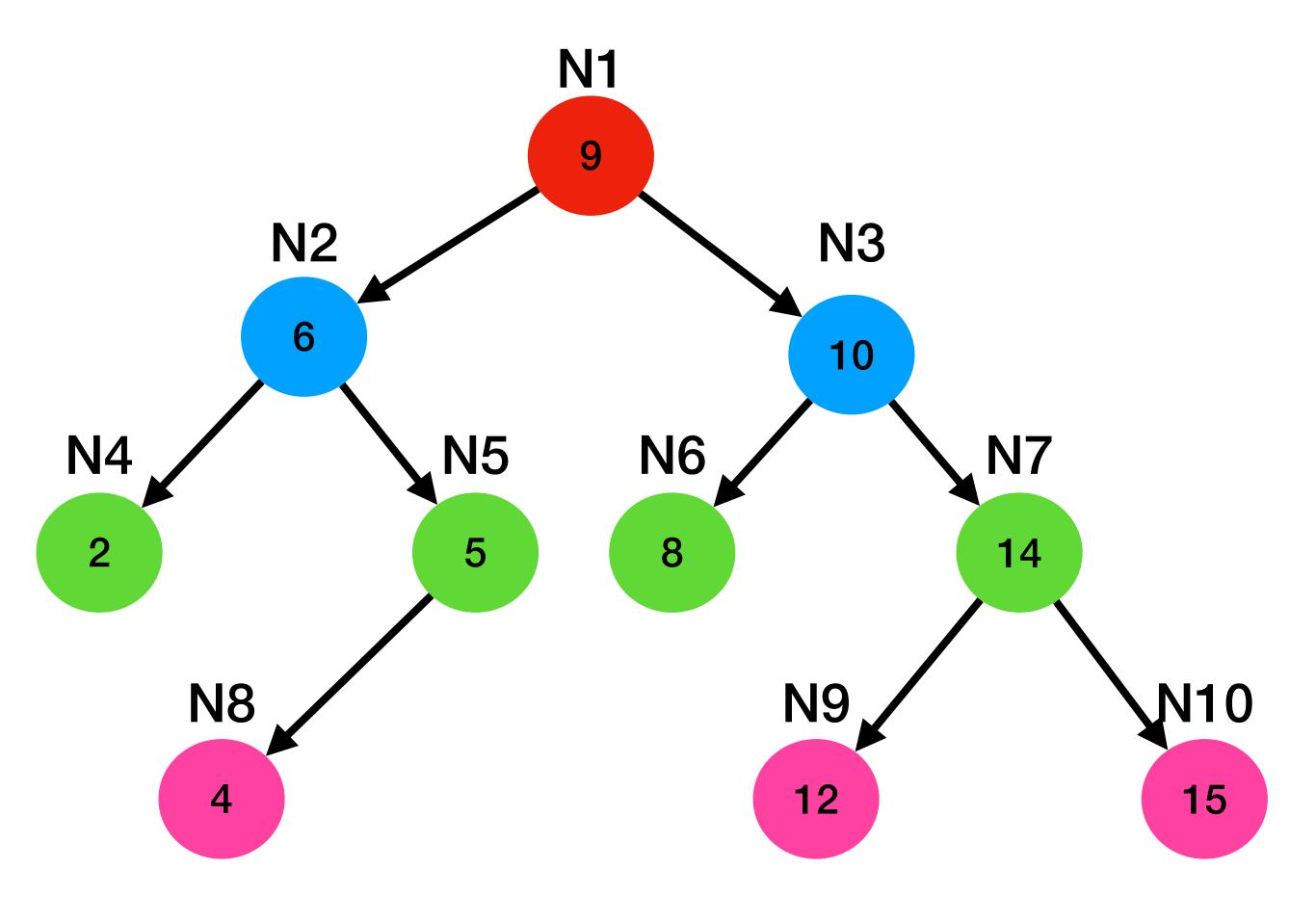
LT: Left sub-tree

RT: Right sub-tree

Order	(1)	(2)	(3)
Pre-order	R	LT Recursion	RT Recursion
In-order	LT Recursion	R	RT Recursion
Post-order	LT Recursion	RT Recursion	R

Tree Traversal

Depth-first

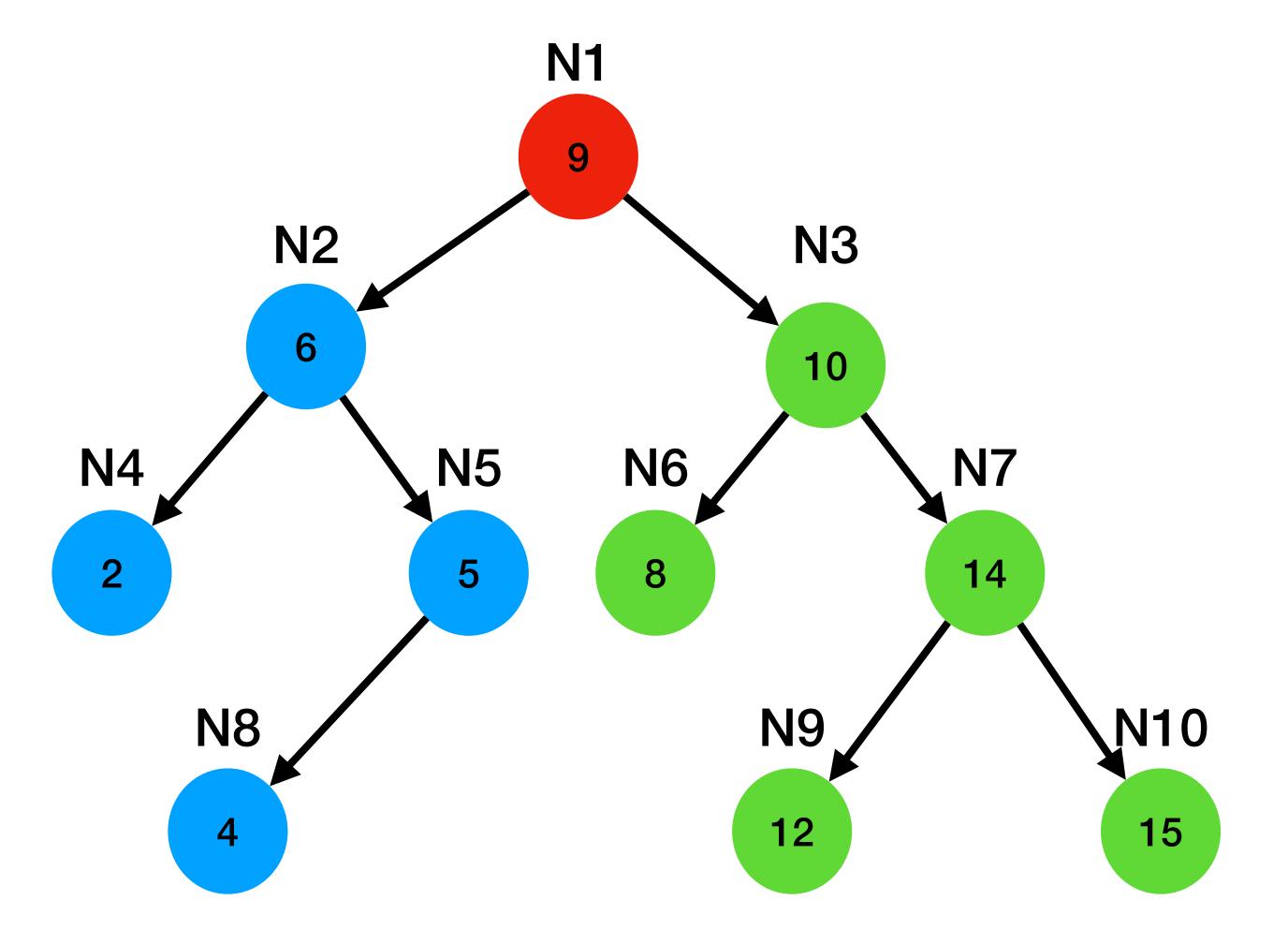


R-LT-RT LT-R-RT LT-RT-R

Pre-order	In-order	Post-order
N1		
N2		
N4		
N5		
N8		
N3		
N6		
N7		
N9		
N10		

Pre-order traversal

Root -> Left sub-tree -> Right sub-tree



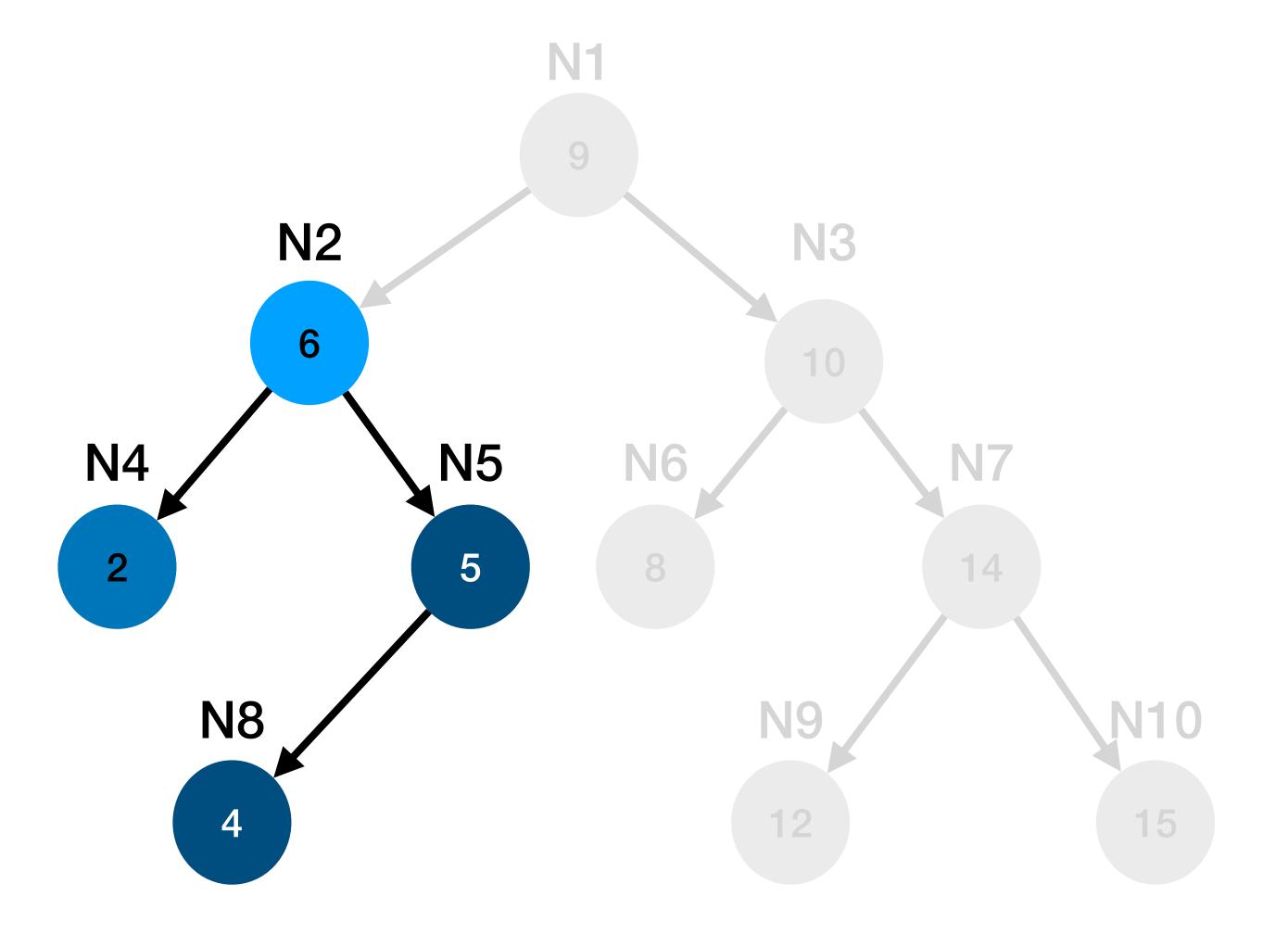
Consider the full tree

Step	Result
Find root	N1
Find Left Sub-tree	Blue Group
Find Right Sub-tree	Green Group

Order:			
N1			

Pre-order traversal

Root -> Left sub-tree -> Right sub-tree



Consider Left Sub-tree (Blue Group)

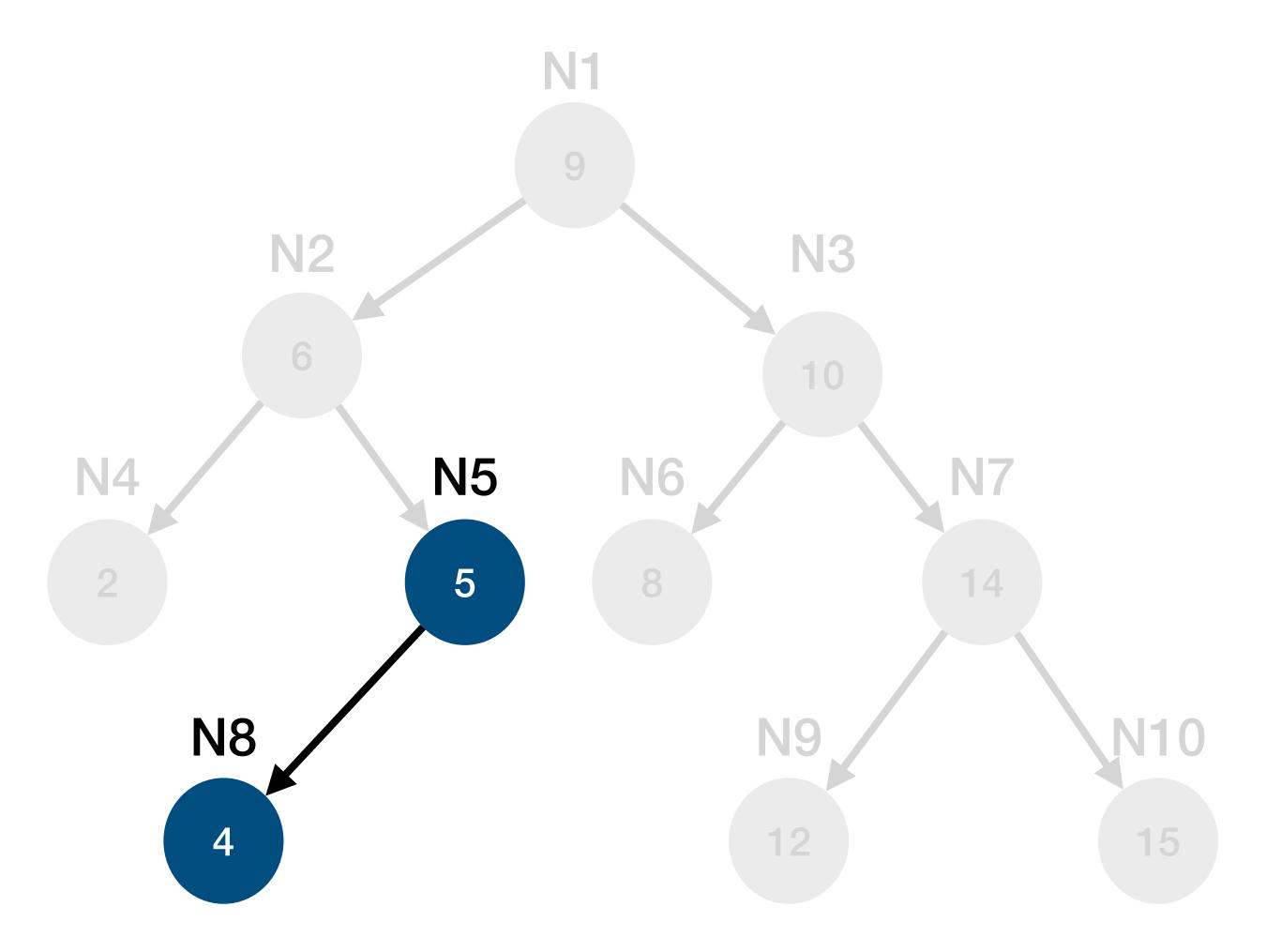
Step	Result
Find root	N2
Find Left Sub-tree	N4
Find Right Sub-tree	Dark blue group

Order:

N1 -> N2 -> N4

Pre-order traversal

Root -> Left sub-tree -> Right sub-tree



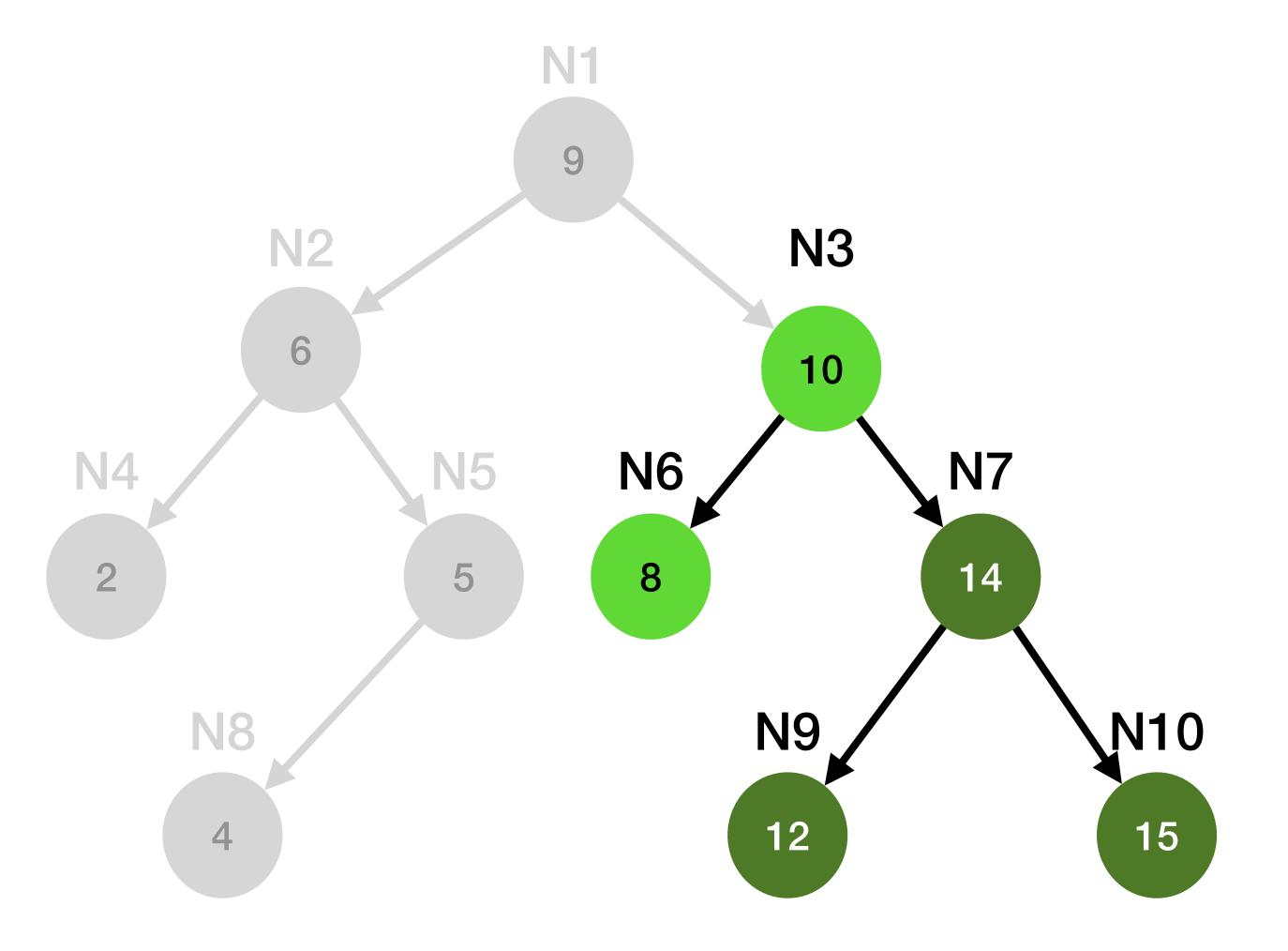
Consider Right Sub-tree (Dark blue group)

Step	Result
Find root	N5
Find Left Sub-tree	N8
Find Right Sub-tree	_

Order:

Pre-order traversal

Root -> Left sub-tree -> Right sub-tree



Consider Right Sub-tree (Green Group)

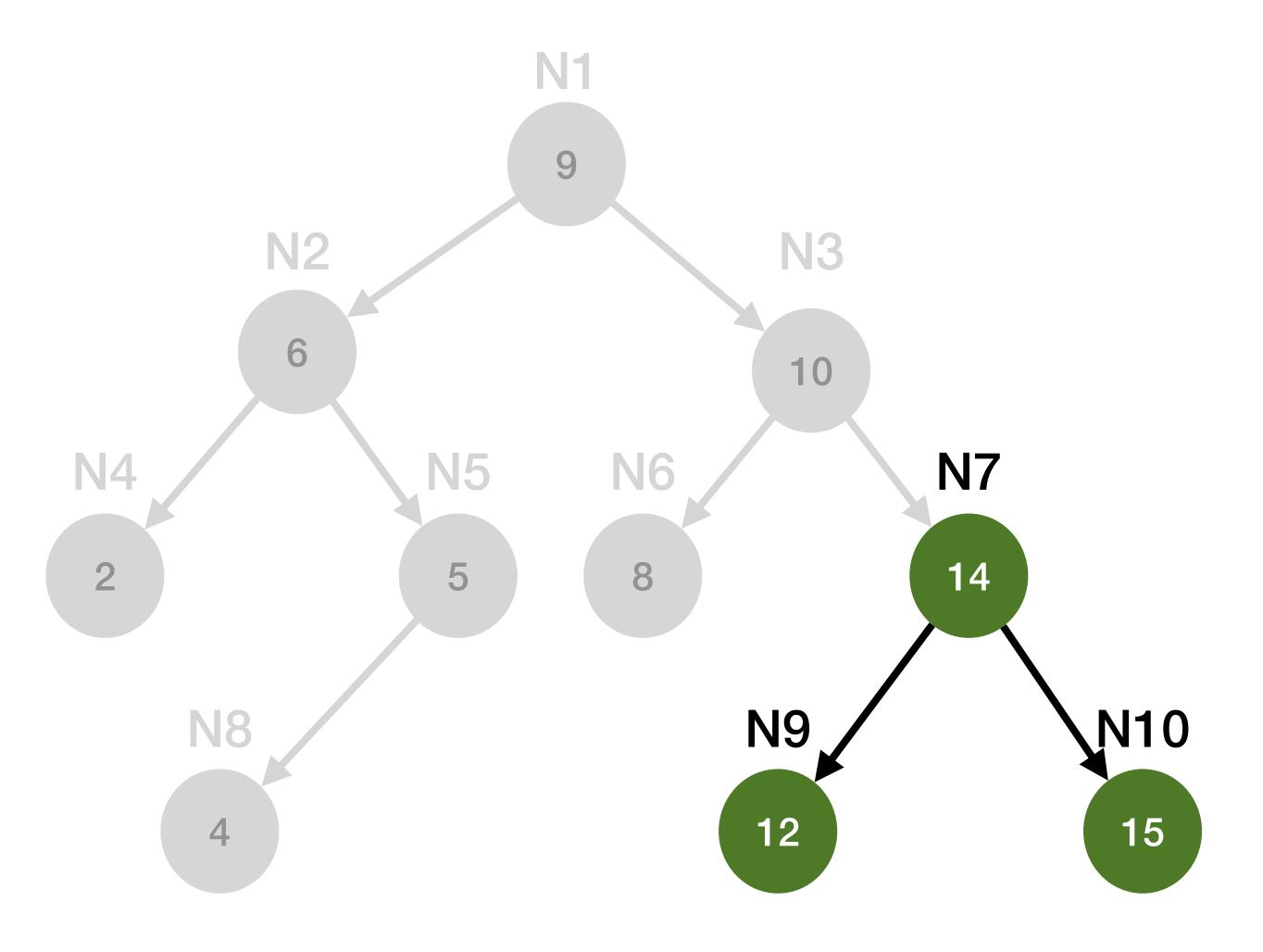
Step	Result
Find root	N3
Find Left Sub-tree	N6
Find Right Sub-tree	Dark Green Group

Order:

N1 -> N2 -> N4 -> N5 -> N8 -> N3-> N6

Pre-order traversal

Root -> Left sub-tree -> Right sub-tree

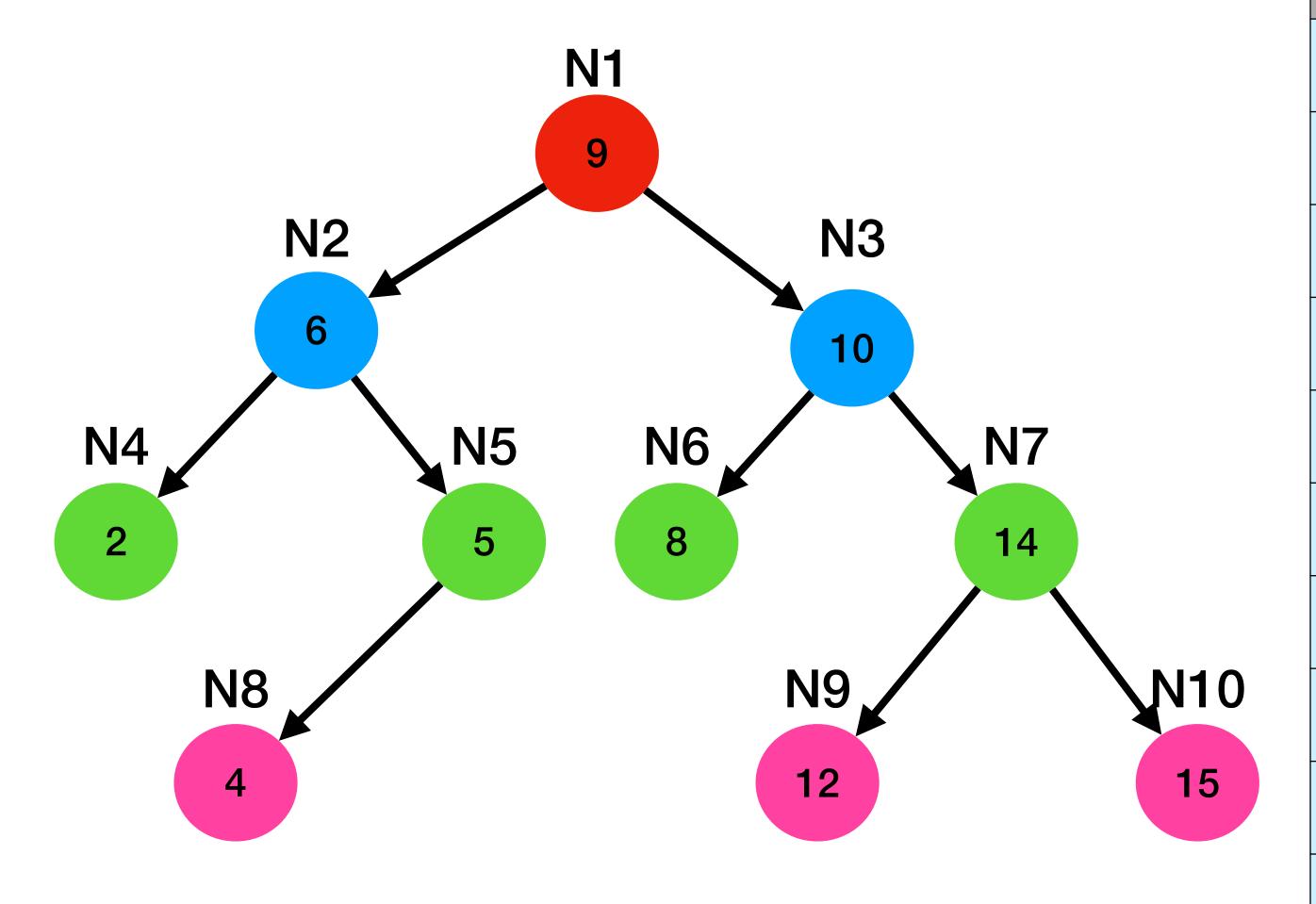


Consider Right Sub-tree (Dark Green Group)

Step	Result
Find root	N7
Find Left Sub-tree	N9
Find Right Sub-tree	N10

Order:

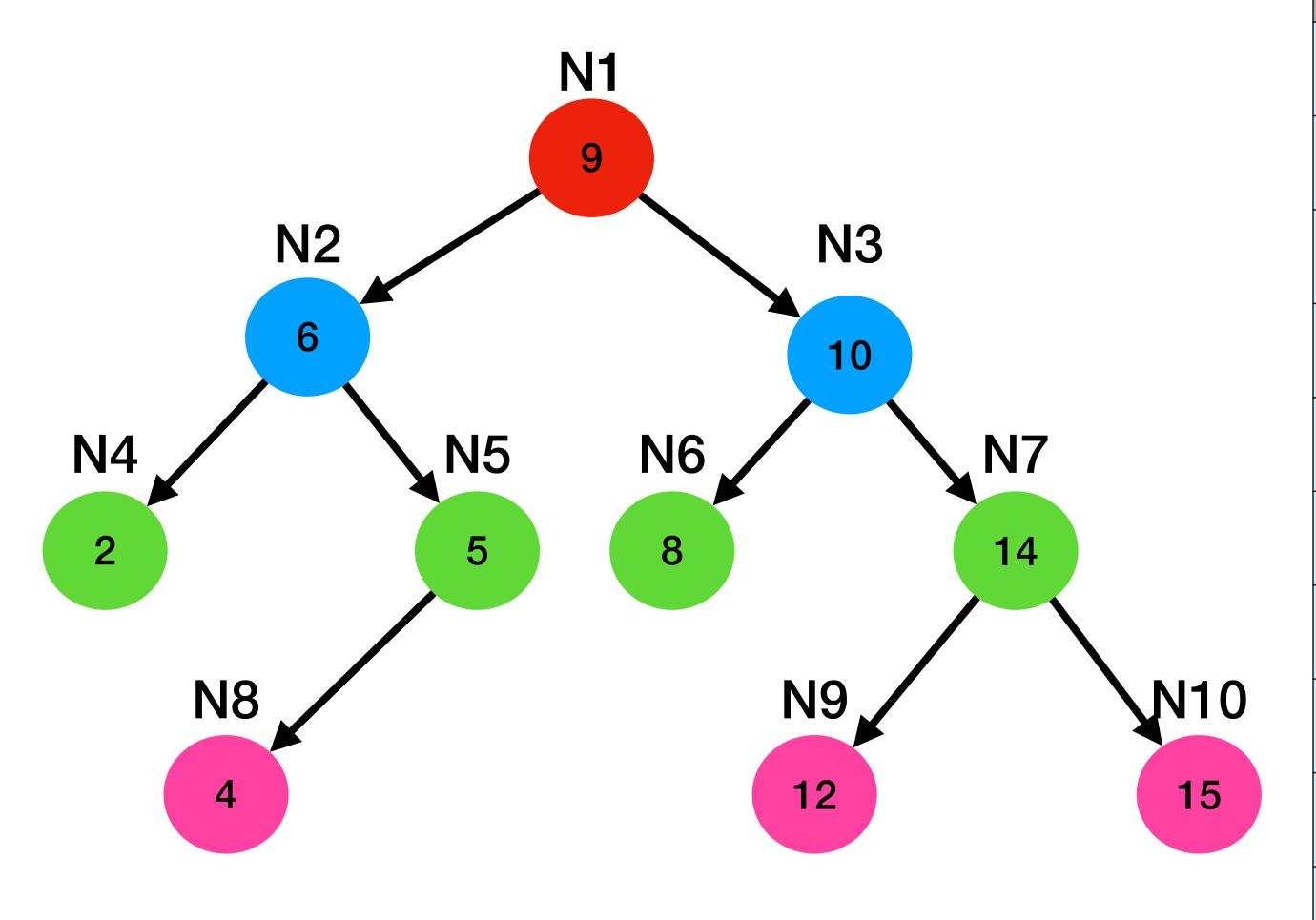
Depth-first



R-LT-RT LT-RT-RT-R

Pre-order	In-order	Post-order
N1	N4	
N2	N2	
N4	N8	
N5	N5	
N8	N1	
N3	N6	
N6	N3	
N7	N9	
N9	N7	
N10	N10	

Depth-first



R-LT-RT LT-R-RT LT-RT-R

Pre-order	In-order	Post-order
N1	N4	N4
N2	N2	N8
N4	N8	N5
N5	N5	N2
N8	N1	N6
N3	N6	N9
N6	N3	N10
N7	N9	N7
N9	N7	N3
N10	N10	N1

Depth-first

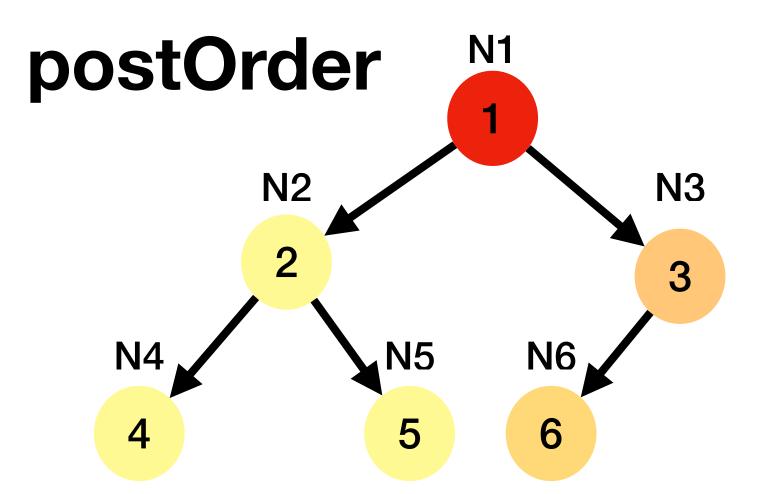
tree.c

```
void postOrder (TREENODE_T* node)
 if (node->left != NULL)
   postOrder(node->left);
 if (node->right != NULL)
   postOrder(node->right);
 printf("access %d\n", node->data);
```

main.c

```
TREENODE_T* root = N1;
postOrder(root);
```

Recursive Function

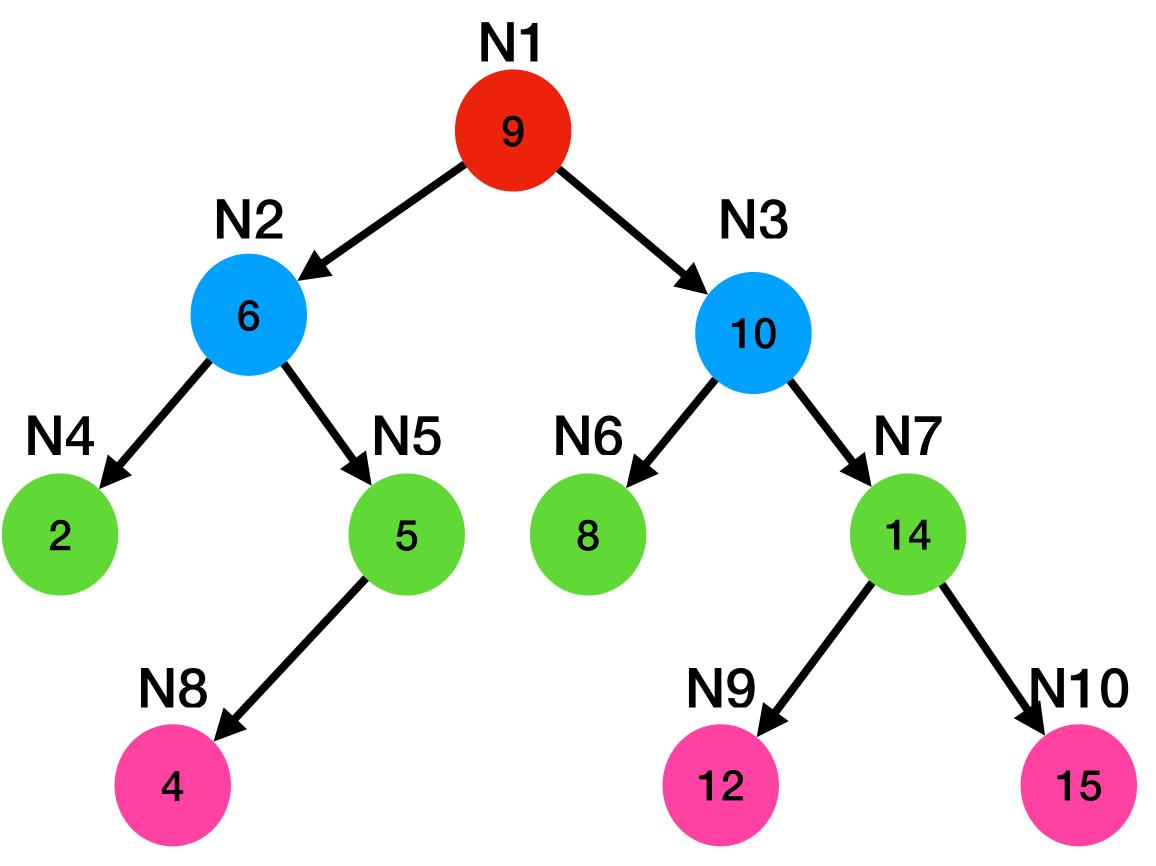


```
1 void postOrder (TREENODE_T* node)
2 {
3 if (node->left != NULL)
4  postOrder(node->left);
5 if (node->right != NULL)
6  postOrder(node->right);
7 printf("access %d\n", node->data);
8 }
```

```
TREENODE_T* root = N1;
postOrder(N1);
```

Г					
	Step	Current Node	Action	Pass	End function
	1	N1	Recursion#1 at Line 4	N1->left (N2)	No
	2	N2	Recursion#2 at Line 4	N2->left (N4)	No
	3	N4	Print access 4	_	Back to Step2 Line5
	4	N2	Recursion#3 at Line 6	N2->right (N5)	No
	5	N5	Print access 5	_	Back to Step2 Line7
	6	N2	Print access 2	_	Back to Step1 Line5
	7	N1	Recursion#4 at Line 6	N1->right (N3)	No
	8	N3	Recursion#5 at Line 4	N3->left (N6)	No
	9	N6	Print access 6	_	Back to Step8 Line5
	10	N3	Print access 3	_	Back to Step7 Line7
	11	N1	Print access 1	_	Back to Main 42

Breadth-first



Root -> Leaf L -> Leaf R

Dequeue	PTR	Enqueue	Queue
Start	N1	<u>N2, N3</u>	N2, N3
N2	N2	<u>N4, N5</u>	N3, <u>N4, N5</u>
N3	N3	<u>N6, N7</u>	N4, N5, <u>N6, N7</u>
N4	N4	_	N5, N6, N7
N5	N5	<u>N8</u>	N6, N7, <u>N8</u>
N6	N6	_	N7, N8

Construct a binary tree from traversal results

- Require at least 2 traversal results
 - o In-order > left & right child node
 - Pre-order or post-order > root node

In-order: DBEAFCG

Pre-order: A B D E C F G

Construct a binary tree from traversal results

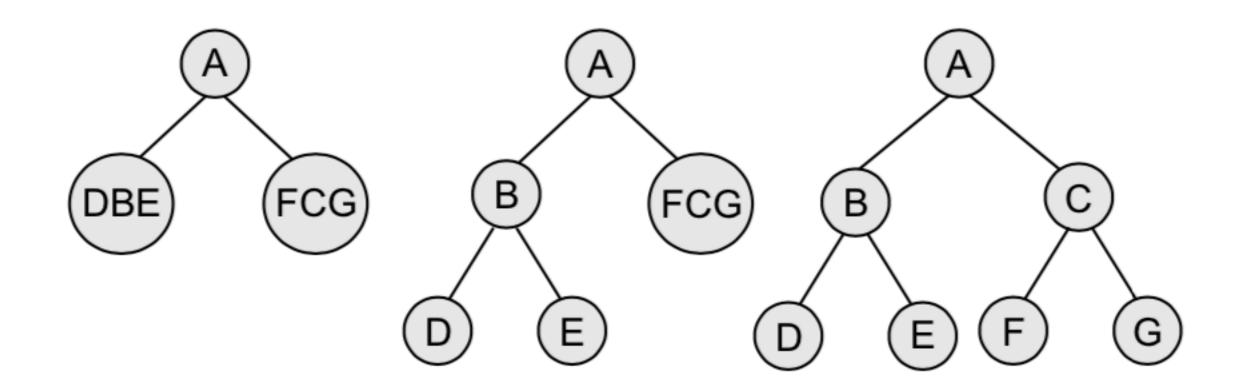
STEP1 - Use pre-order/post-order to determine the root (first/last).

STEP2 - Use in-order to determine the left & right subtree of the root.

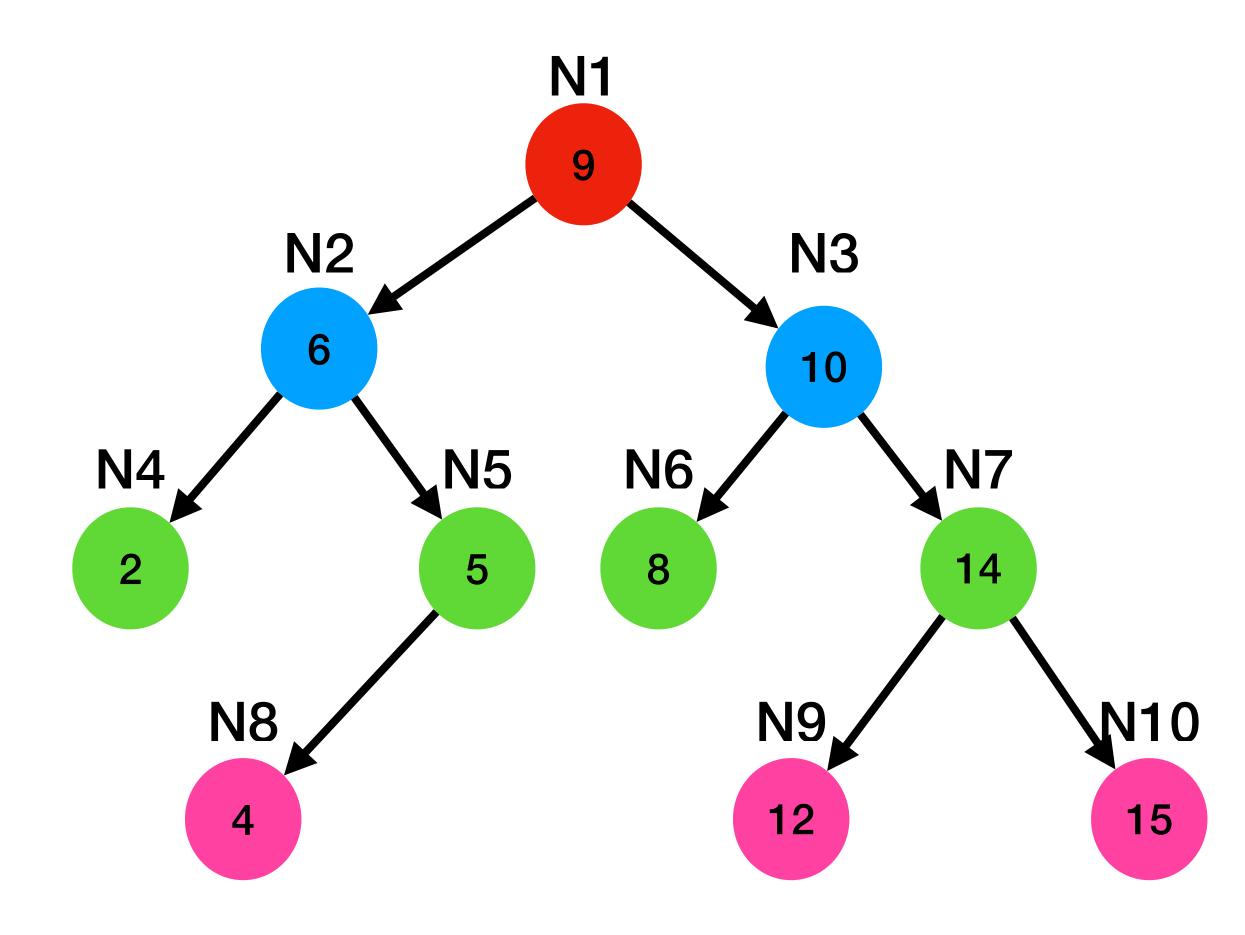
STEP3 - Recursively select each element from STEP1 to be the root of the subtree.

In-order: DBEAFCG

Pre-order: A B D E C F G



```
int CountNodes (node *n)
       int nodes = 0;
       if (n==NULL)
5
          nodes = 0;
6
       else
          nodes = CountNodes(n->left) +
8.1
                  CountNodes(n->right) + 1;
8.2
       return nodes;
10
```



Step	Current Node	Action	Pass	End function
1	N1	Recursion#1 at Line 8.1	N1->left (N2)	No
2	N2	Recursion#2 at Line 8.1	N2->left (N4)	No
3	N4	Recursion#3 at Line 8.1	N4->left (NULL)	No
4	NULL	Return 0	_	Back to Step3 Line 8.1 (0 + N4->right + 1)
5	N4	Recursion#4 at Line 8.2	N4->right (NULL)	No
6	NULL	Return 0	_	Back to Step5 Line 8.2 (0 + 0 + 1)
7	N4	Return 0+0+1 = 1	_	Back to Step2 Line 8.1 (1 + N2->right + 1)
8	N2	Recursion#5 at Line 8.2	N2->right (N5)	No
9	N5	Recursion#6 at Line 8.1	N5->left (N8)	No
10	N8	Recursion#7 at Line 8.1	N8->left (NULL)	No

```
void AddRoot (int bt[], int rootValue)
   bt[0] = rootValue;
void AddLeftChild (int bt[], int parentIndex, int value)
   bt[2 * parentlndex + 1] = value;
void AddLeftChild (int bt[], int parentIndex, int value)
   bt[2 * parentlndex + 2] = value;
```

```
int CountNodes (int bt[], int size)
   int i = 0; count = 0;
   for (i=0; i<size; i++)
      count++;
      if (bt[i+1]==0 \&\& bt[i+2]==0)
      break;
   return count;
```

Wrap up

- Hierarchical data structure
- Terminologies
- Binary tree
- Tree traversal
 - O Depth-first search: pre-order, in-order, post-order
 - Breadth-first search