# Lecture 1 - Introduction

**CPE112 - Programming with Data Structures 15 January 2025** 

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## What to learn?

#### **Programming with Data Structures**



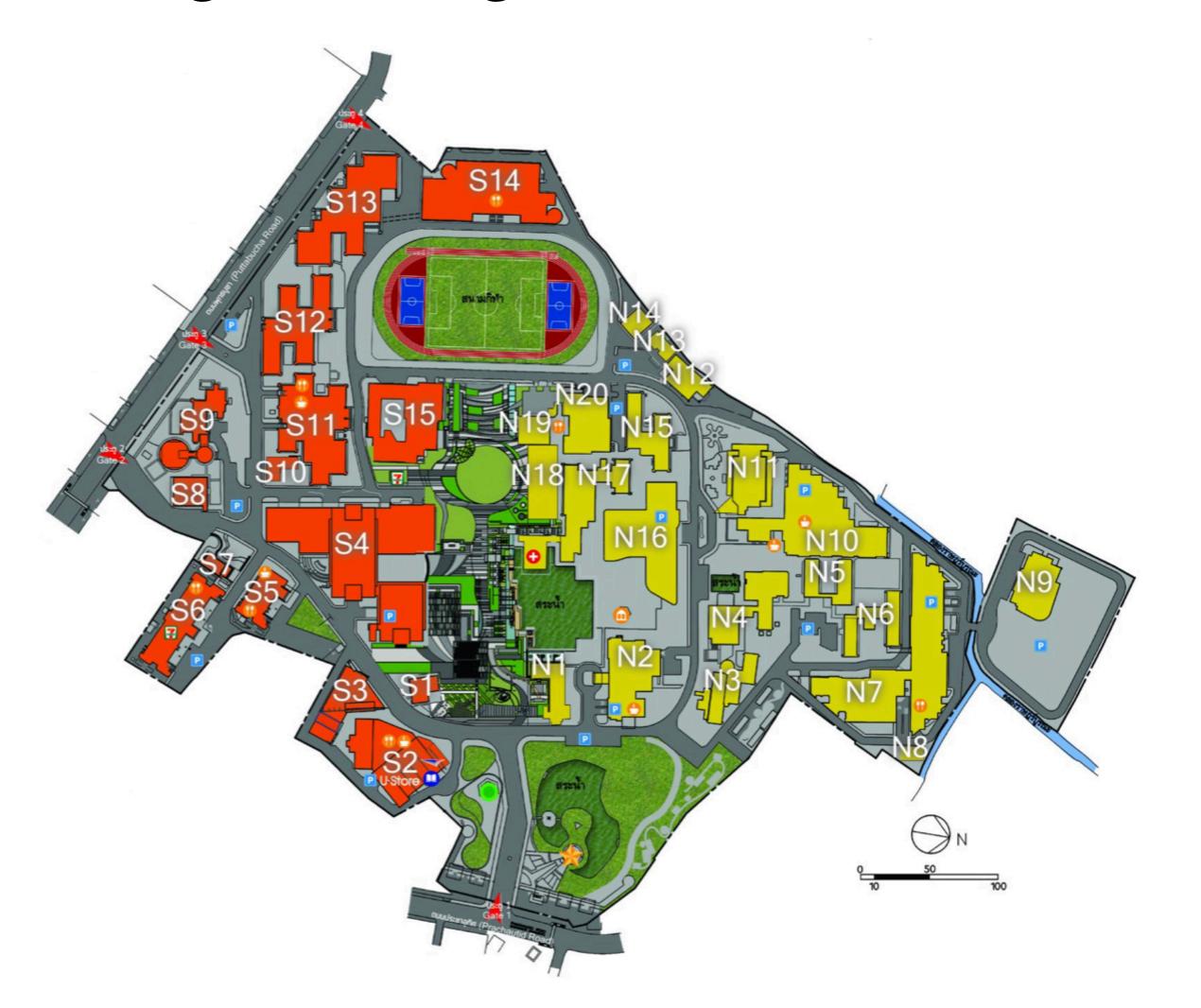
- Programs manipulate data to solve problems
- Ways to organize & store information
  - Insert/delete/sort numbers: array
  - Student information: structure

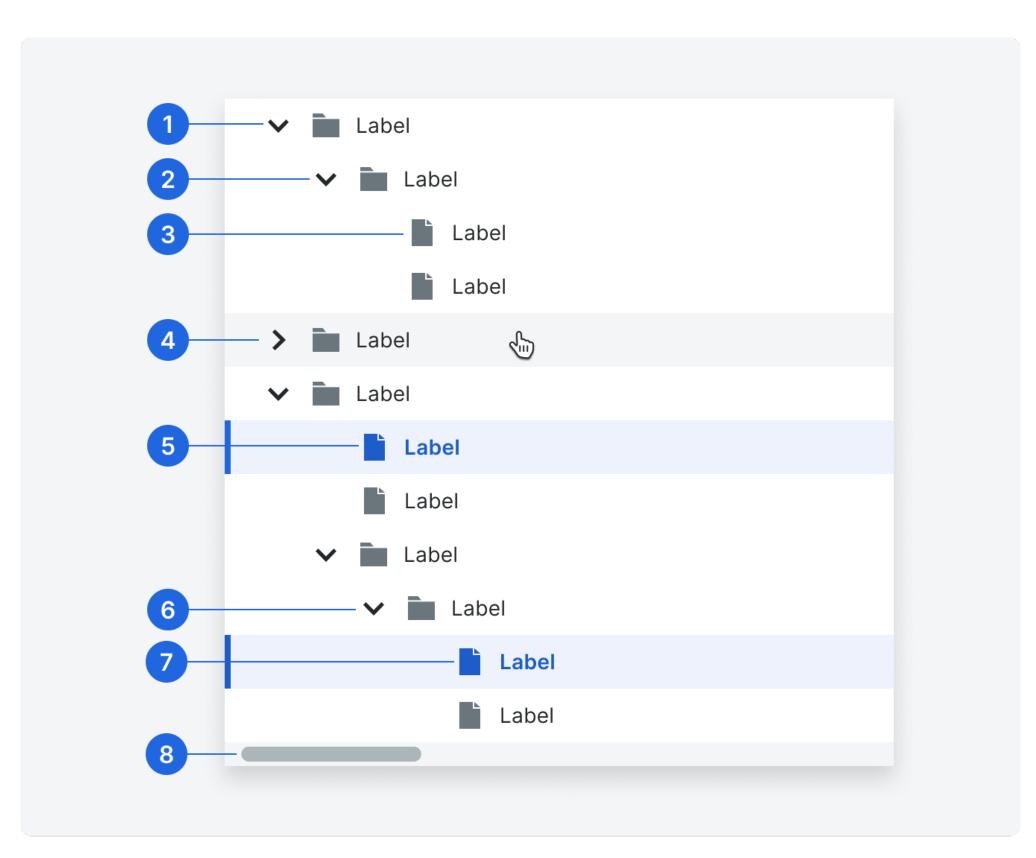
Score
67
78
87
90
82
65
68

Name	Midterm	Final
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## What to learn?

### **Programming with Data Structures**

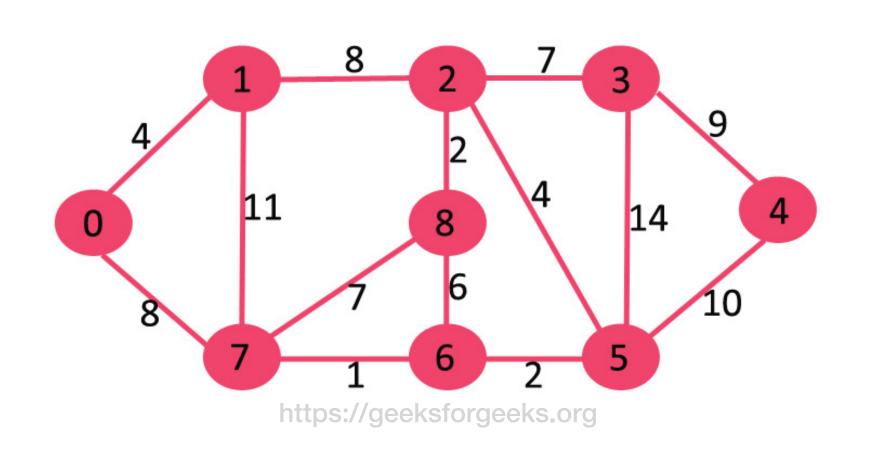


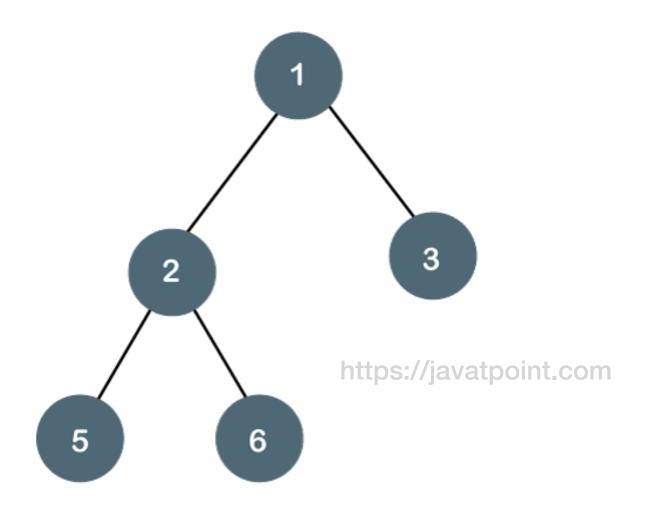


- 1. Expanded Folder (1st level)
- 2. Expanded Folder (2nd level)
- 3. Unselected File (2nd level)
- 4. Collapsed Folder (1st level) / Hover
- 5. Selected File (1st level)
- 6. Expanded Folder (3rd level)
- 7. Selected File (3rd level)
- 8. Horizontal Scrollbar

## What to learn?

#### **Programming with Data Structures**





- Important data structures strategies for organizing information in a program
- Important algorithms well-known methods for accomplishing particular kinds of tasks
- Efficiency of different data structure algorithms

## Data Structure & Algorithm



Find the first recurring character

"ABCA" -> A
"BCABA" -> B
"ABC" -> NULL

## Data Structure & Algorithm



Find the first recurring character

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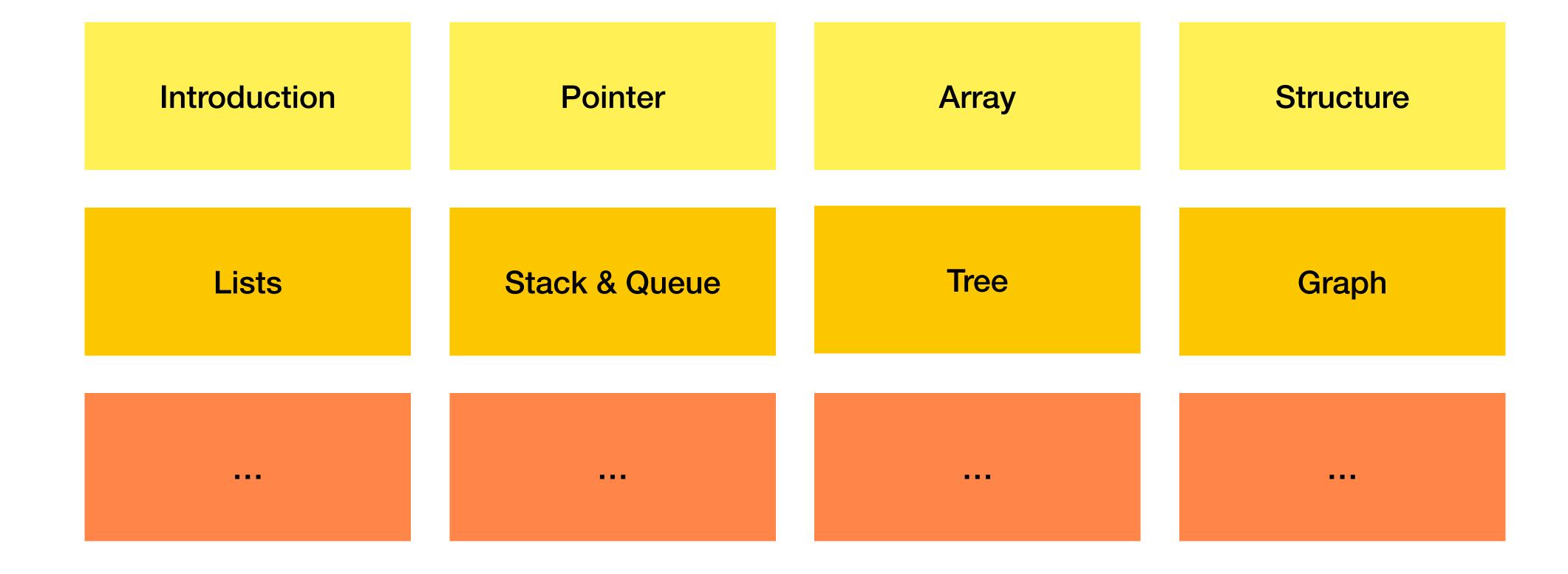
Your algorithm will be evaluated.

Correctness (work well)

Efficiency (minimum time & memory)

Style (easy to understand & modify)

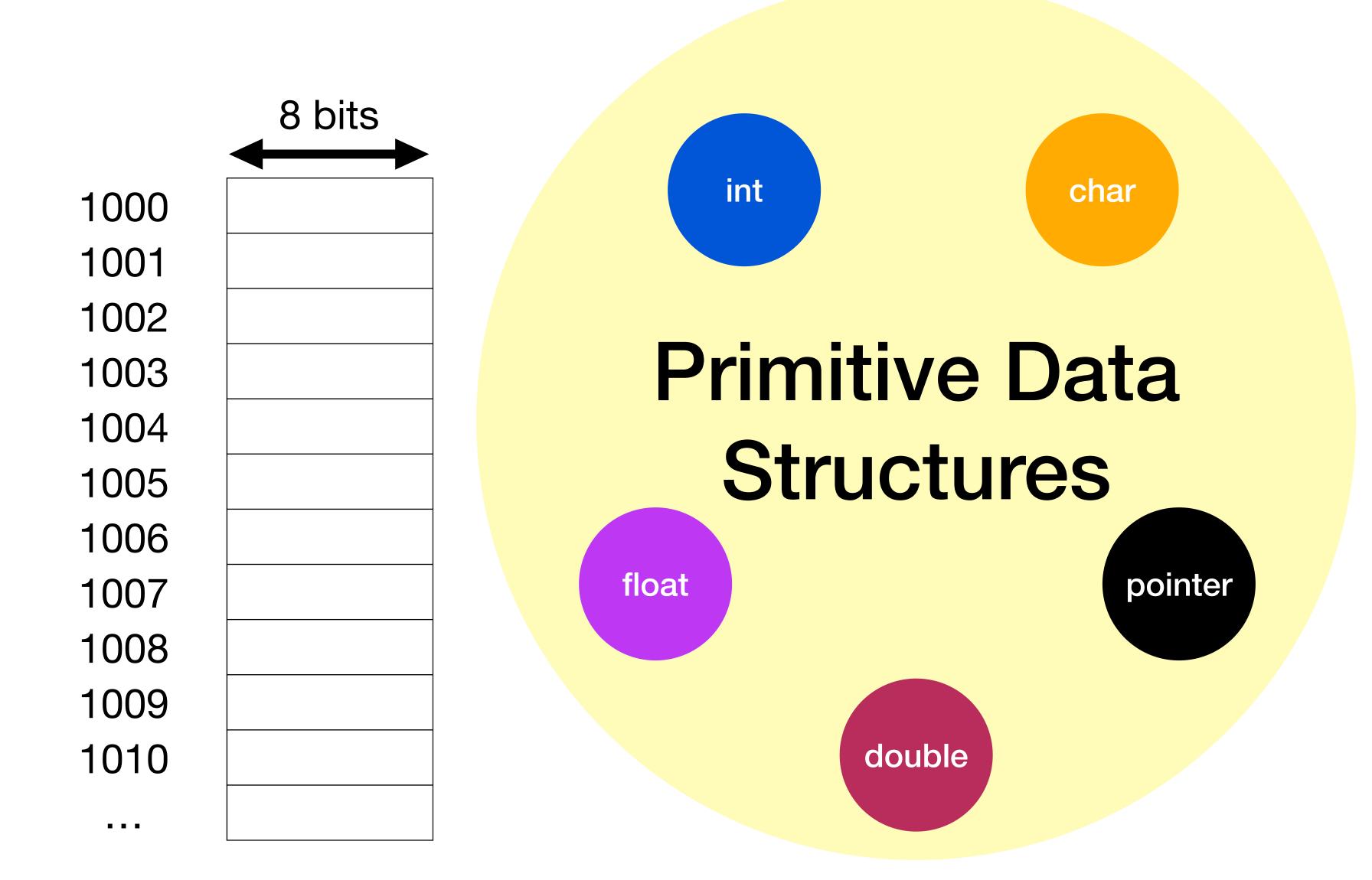
# Major Contents



#### Why we learn CPE100?

- <u>Programmers</u> need to translate <u>problems</u> and/or <u>user requirements</u> into <u>source</u> codes that a computer can understand.
- Two questions that every programmers should solve
  - How we store data in our computer memory? => Data Structure (CPE112)
  - Our How can process the data efficiently?
    - Procedures to solve common (programming) problems => Algorithm (CPE223)
    - Domain specific algorithms e.g. speech recognition, image processing, data analytics => Elective courses (CPE3xx, CPE4xx)

# Review Memory



#### **Pointer**

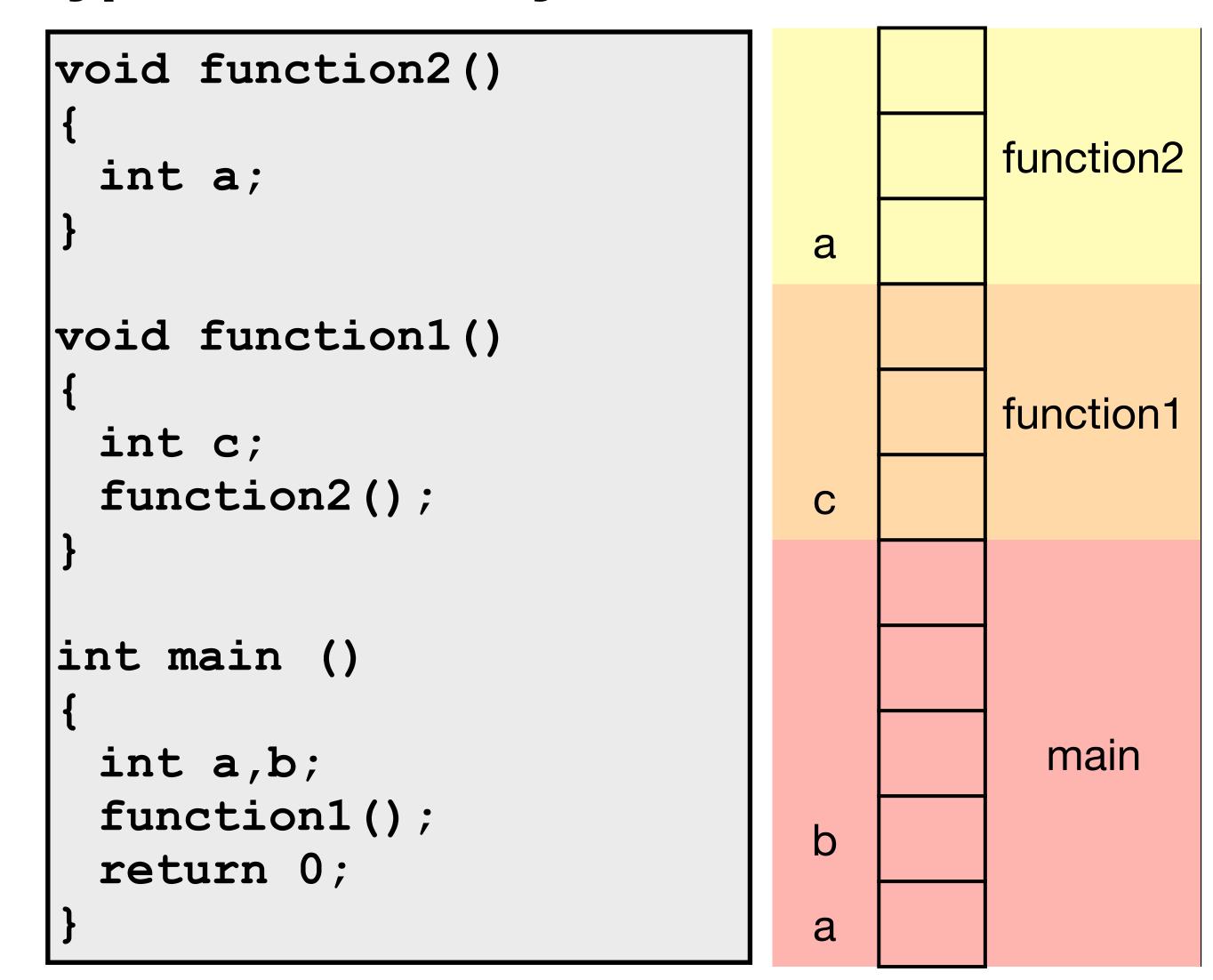
- A variable type store the address of some value
  - & Get address
  - \* Look into address

```
data-type *pointer-variable-name;

char c='K';
char *cp=&c;
char **cpp=&cp;
```

1001	K	С
1002		
1003		00
1004		ср
1005		
1006		
1007		onn
1008		срр
1009		
1010		

#### **Types of Memory - Stack**



#### **Types of Memory - Heap**

```
void function1(int* x)
 x[0] = /*...*/;
int main ()
 int *a = malloc(4);
 int *b = malloc(6);
 function1(a);
 free(a);
 free(b);
 return 0;
```

X	100	function1
		main
b	104	
a	100	

# Review Types of Memory

	(How & when they are allocated/deallocated)	(Where they can be accessed)
Stack	Allocate - program enters the function that they are declared <b>Deallocate</b> - program leaves the function	Can be accessed and will have valid value within their own block
Static	Allocate - program starts running Deallocate - program exits	Can be accessed by <u>any code in C</u> <u>module</u>
Dynamic	Allocate - call calloc Deallocate - call free	Can be accessed and will have valid values anywhere in the program, by passing around variables that hold the start address

Scope

Life cycle

#### **Dynamic Memory Allocation**

#### **Stack Memory**

1000	Н
1001	i
1002	\0
1003	С
1004	Р
1005	E
1006	1
1007	1
1008	2
1009	\0
1010	

#### Example:

char word1[3] = "Hi"; char word2[10] = "CPE112"

"Hi" -> "Hello"

#### Heap Memory

### calloc free

#### Example:

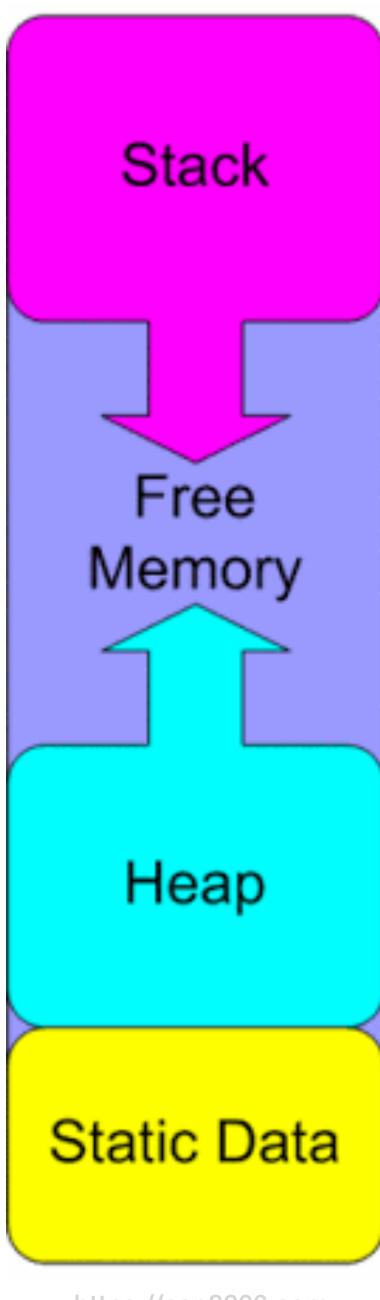
#define ARRAYSIZE 300 int\* myValues = calloc(ARRAYSIZE,sizeof(int));



# **Review**Dynamic Memory Allocation

calloc but not free

Memory crash!!!



#### Structure

 Create own custom data type that is a composition of multiple other data types

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```
typedef struct {
   char name[100];
   float midterm;
   int final;
} student_info;
```

#### Structure

```
void selectionSort (student info data[], int count)
 int i, j, minPos;
                                                void swap (student info *x,
 for (i = 0, i < count-1; i++)
                                                 student info *y)
   minPos = i;
                                                  student info tmp;
   for (j=i+1; j<count; j++)
                                                  tmp = *x; *x=*y; *y=tmp;
     if (data[j].final < data[minPos].final)</pre>
      minPos = j;
   swap(&data[i], &data[minPos]);
```

## Data Structure

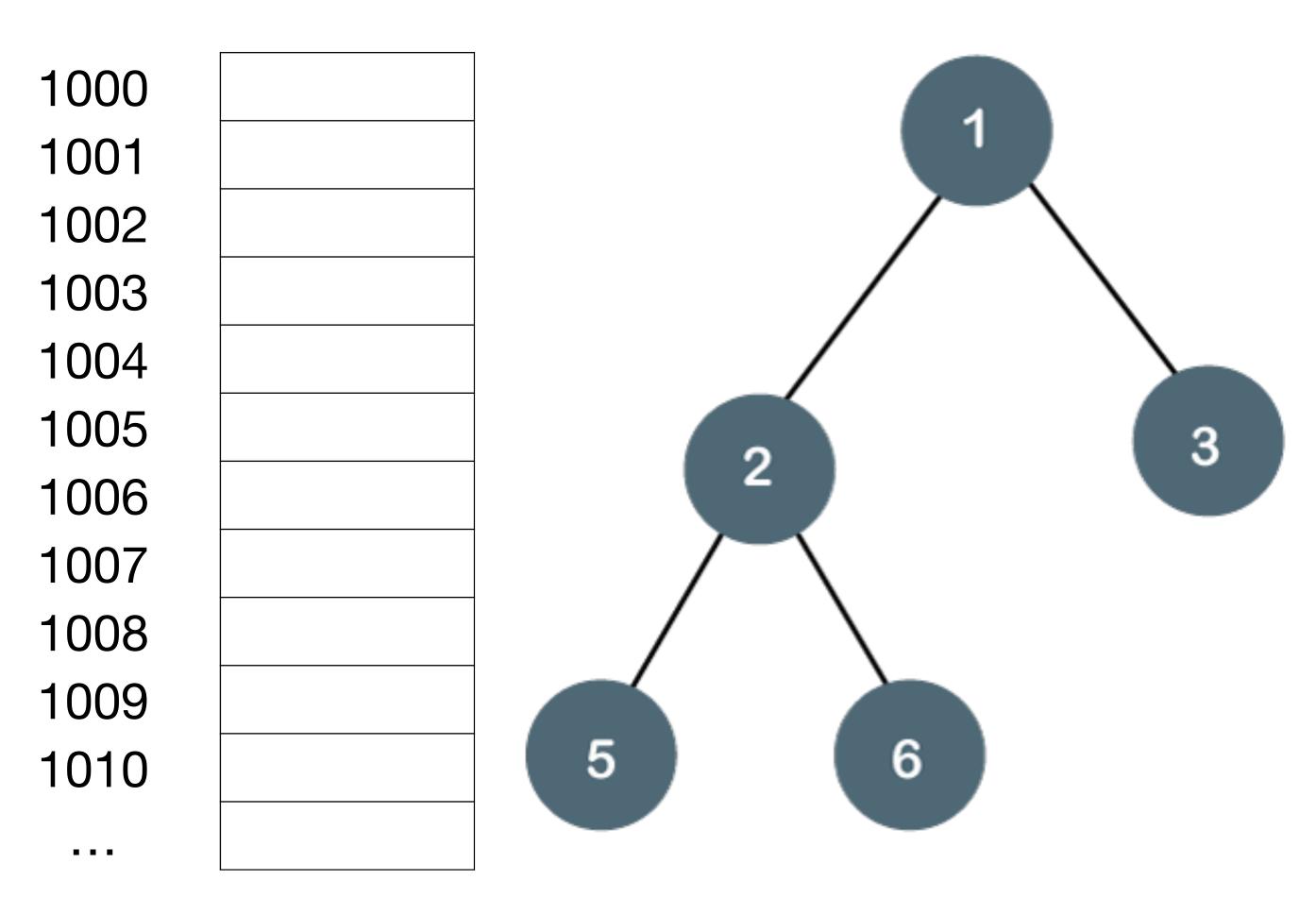
- A data structure a way to organize information stored in computer memory.
- Different kinds of data structures works better for different sorts of problems.
- Appropriate structure -> Good solution

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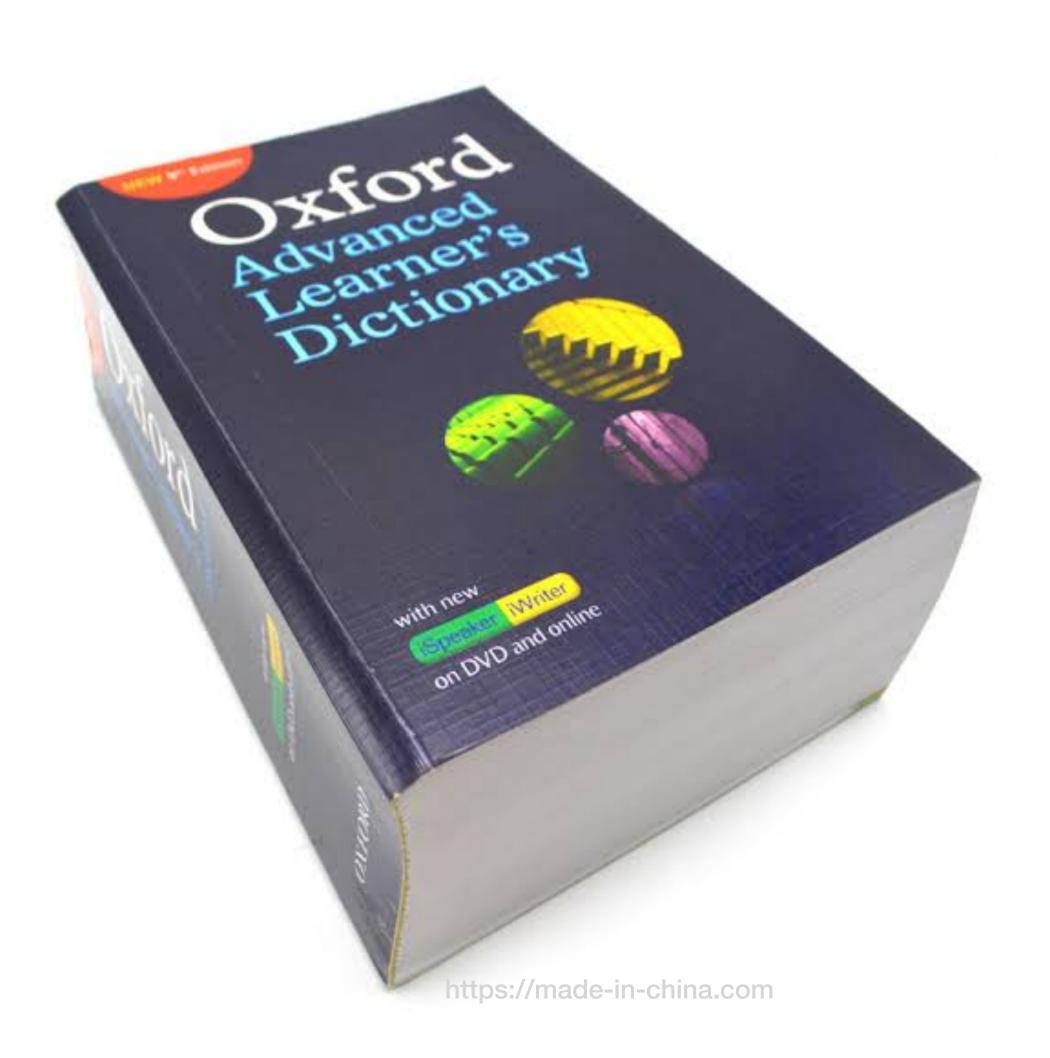
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## Data Structure

- Creating a data structure
  - Ability to refer to different memory locations
  - Stitch them together into a custom shape



# Data Structure & Algorithm Example

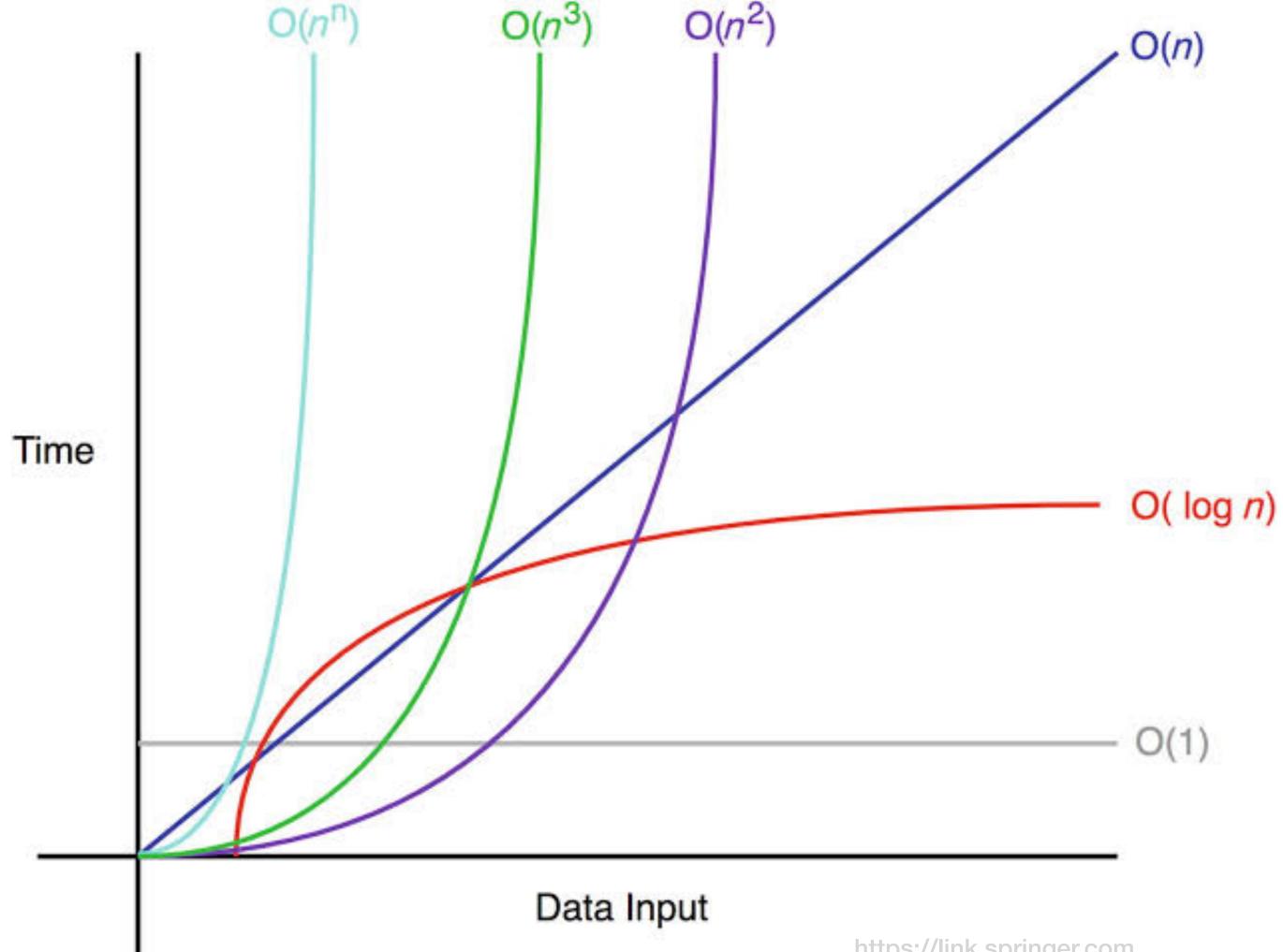


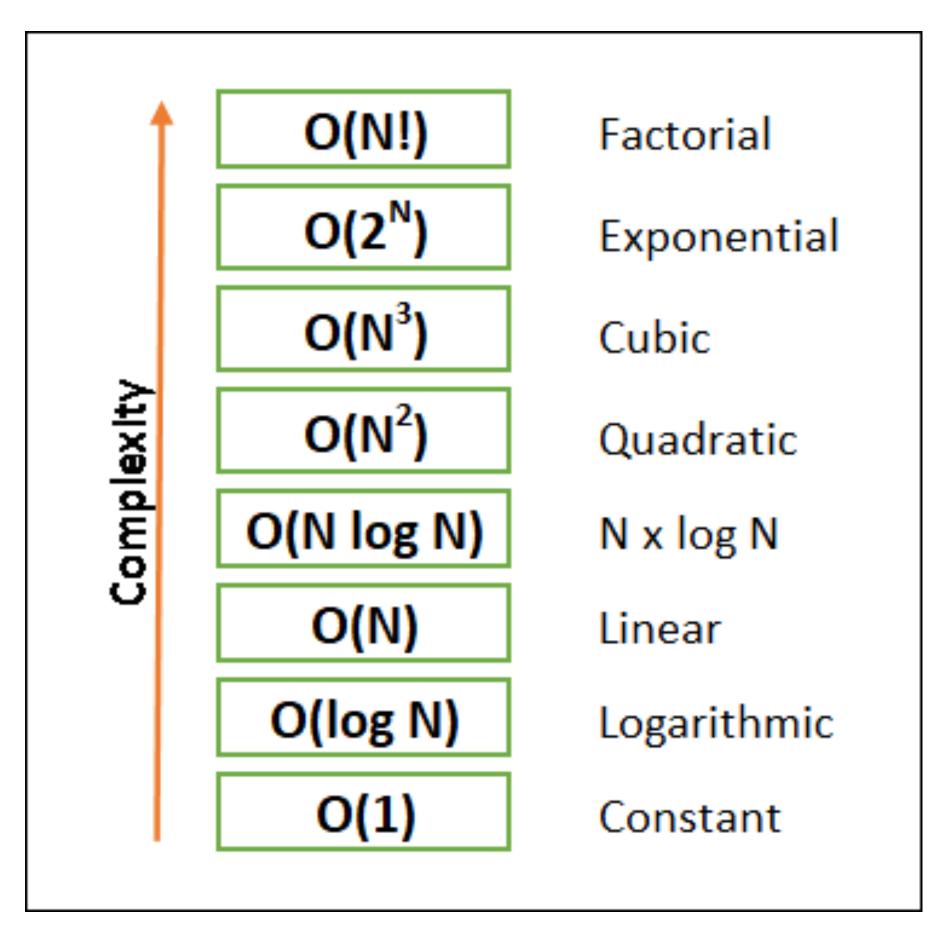
How to find "Structure"?

	#operations for 1 page	#operations for 100 page	#operations for 1000 page	
Method#1				
Method#2				

Time to solve

### **Big O Notation**





https://towardsdatascience.com

# Algorithm Efficiency Big O Notation

- 1. Understand how the algorithm works
- 2. Identify a basic unit of the algorithm to count
- 3. Map growth of count from step 2 to appropriate Big O class

## Big O - Examples

1	x = 5 + (5 * 5);	O()
2	for (i = 0; i < N; i++) printf("%d", i);	O()
3	for (i = 0; i < N; i++) for (j = 0; i < N; i++) printf("%Id", i*j);	O()
4	x = 5 + (5 * 5); for (i = 0; i < N; i++) printf("%d", i); for (i = 0; i < N; i++) for (j = 0; i < N; i++) printf("%ld", i*j);	?

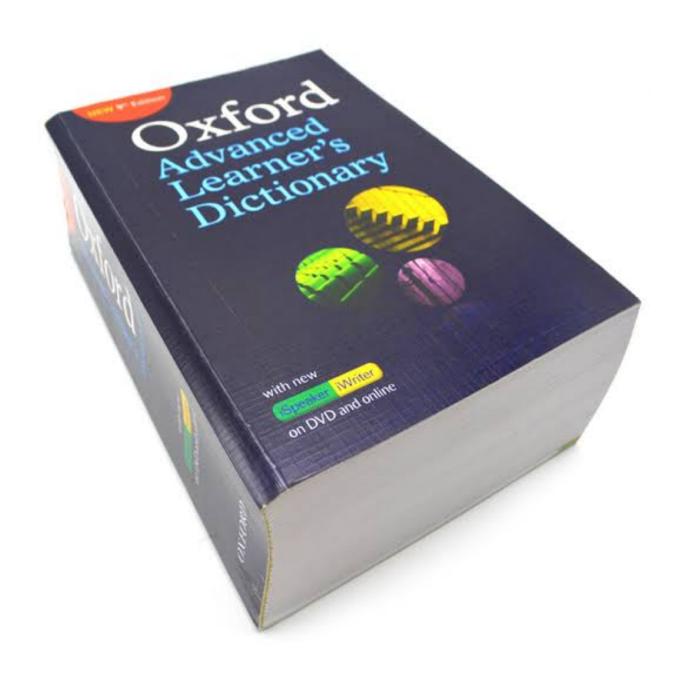
#### **Big O Definition**

Let f(n) and g(n) be functions from positive integers to positive reals. f(n) = O(g(n)) if there is a constant c > 0 such that  $f(n) \le c*g(n)$  for large n

นิยาม 1: f(n) และ g(n) เป็นฟังก์ชันซึ่งมีโดเมนเป็นเลขจำนวนเต็มบวก และมี ranges เป็นจำนวนจริงบวก <u>สามารถเขียนได้เป็น</u>

f(n) = O(g(n)) ถ้ามีค่าคงที่บวก c ที่ทำให้ f(n) ≤ c\*g(n) เป็นจริงทุกค่าสำหรับ x ที่เป็นจำนวนเต็มบวก

Big O - Examples



Find "Ant", "List", "Zombie", ...

#### Cases

### Find character 'A' in a given string length n

**Best case** 

$$C(n) = 1$$

Worst case

$$C(n) = n$$

Average case

$$C(n) = (1+2+3+...+n)1/n$$
  
=  $[n(n+1)/2] * [1/n]$   
=  $(n+1)/2$ 

#### Complexity Summary of Array Sorting Algorithms

Algorithm	Time Complexity			
	Best	Average	Worst	
Quicksort	$\Omega(n \log(n))$	θ(n log(n))	0(n^2)	
Mergesort	$\Omega(n \log(n))$	Θ(n log(n))	O(n log(n))	
<u>Timsort</u>	<u>Ω(n)</u>	θ(n log(n))	O(n log(n))	
<u>Heapsort</u>	$\Omega(n \log(n))$	θ(n log(n))	O(n log(n))	
<b>Bubble Sort</b>	<u>Ω(n)</u>	Θ(n^2)	0(n^2)	
Insertion Sort	<u>Ω(n)</u>	Θ(n^2)	0(n^2)	
Selection Sort	Ω(n^2)	Θ(n^2)	0(n^2)	
Tree Sort	$\Omega(n \log(n))$	θ(n log(n))	0(n^2)	
Shell Sort	$\Omega(n \log(n))$	θ(n(log(n))^2)	0(n(log(n))^2)	
<b>Bucket Sort</b>	$\Omega(n+k)$	Θ(n+k)	0(n^2)	
Radix Sort	Ω(nk)	Θ(nk)	0(nk)	
Counting Sort	$\Omega(n+k)$	$\theta(n+k)$	0(n+k)	
Cubesort	<u>Ω(n)</u>	θ(n log(n))	O(n log(n))	

## Designing an Algorithm

- Perform operations on the stored data contained in data structures.
- Modularization process a complex algorithm is often divided into smaller units called modules.
  - Top-down approach dividing the complex algorithm into one or more modules
  - Bottom-up approach designing the most basic or concrete modules and then proceed towards designing higher level modules

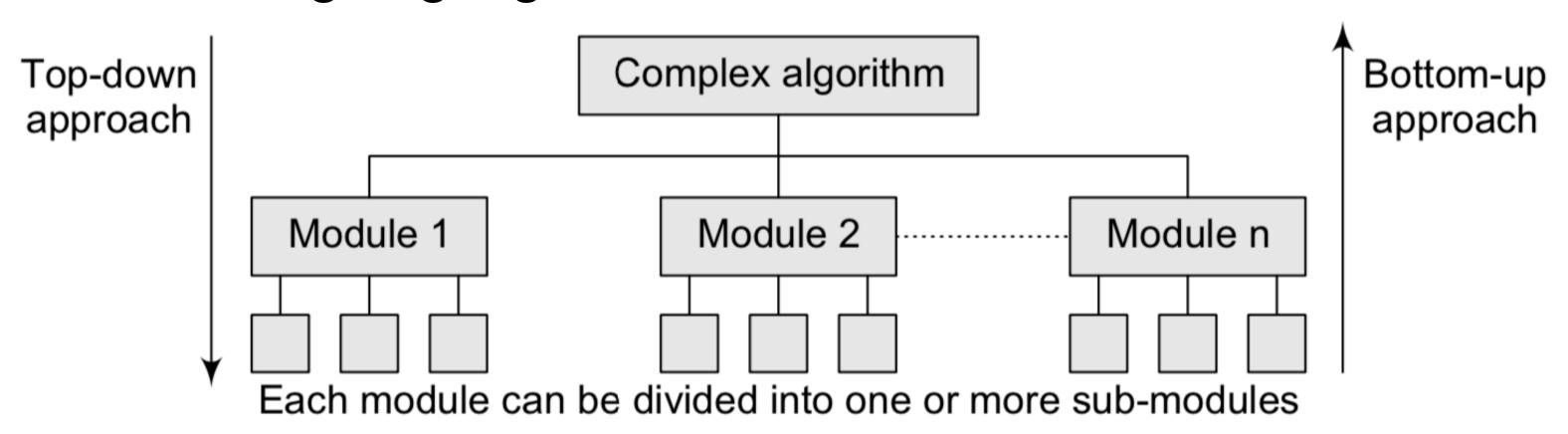


Figure 2.9 Different approaches of designing an algorithm

## Wrap up

- Course introduction
- What to learn?
- Data structures & algorithms
- Review CPE100
- Big O