

Nuts and Bolts of Deploying Real Cloudproxy Applications

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Overview

Cloudproxy is a software system that provides *authenticated* isolation, confidentiality and integrity of program code and data for programs even when these programs run on remote computers operated by powerful, and potentially malicious system administrators. Cloudproxy defends against observation or modification of program keys, program code and program data by persons (including system administrators), other programs or networking infrastructure. In the case of the cloud computing model, we would describe this as protection from co-tenants and data center insiders. To achieve this, Cloudproxy uses two components: a “Host System” (raw hardware, Virtual Machine Manager, Operating System) which provides capabilities described below to the protected program or “Hosted System” (VM, Application, Container).

This document focuses on installation, security and deployment tradeoffs for Cloudproxy applications. In addition, we explain at length several options for providing key management and program operation continuity as programs are upgraded or new programs are introduced. Programmers new to Cloudproxy often worry that program upgrade is impossible or difficult but it is not. We also describe successful key management strategies which support frequent key rotation which is strongly recommended to ensure ongoing security as standard practice as well as providing resilience in the face of errors or omissions.

We hope these instructions allow safe and rapid installation and configuration of Cloudproxy nodes without extensive training or preparation. We assume readers are familiar with [4] which explains Cloudproxy basic concept like measurement and principal names and explains how to build Cloudproxy applications.

Readers can consult [1] for a fuller description. Source code for Cloudproxy as well as all the samples and documentation referenced here is in [2]. We assume, in this document, that the reader is familiar with [5] which is in [2] along with the latest version of this document.

Downloading and compiling Cloudproxy

First, you should download the Cloudproxy repository from [2]. To do this, assuming you have git repository support, type

```
git clone https://github.com/jlmucb/cloudproxy,
```

or,

```
go get github.com/jlmucb/cloudproxy/... .
```

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This latter command will also install the needed go libraries. You can also download a zipped repository from github. You should probably install this in `~/src/github.com/jlmucb` (which we refer to as `$CLOUDPROXYDIR`) to save go compilation problems later; the go get command above does this automatically if your GOPATH environment variable is set to `$HOME`. It's a good idea to put go binaries in `~/bin` as is common in Go. Follow the installation instructions in `$CLOUDPROXYDIR/Doc`; that directory also contains [1] and an up to date version of [4].

You must also install the Go development tools (and C++ development tools if you use the C++ version) as well as protobuf, gtest and gflags as described in the Go documentation.

We will continue to use the `simpleexample` application described in [4].

Security Model

Any deployment of a security system must consider the security or threat model that a system seeks to provide. For the purpose of this Deployment Guide, we consider you are running software for your clients in a cloud data center (or other shared facility) and storing your data there.

Our “Security Goal” is to provide the principled, verifiable *confidentiality and integrity* for programs and the data they read and store from unrestricted attacks from outsiders (e.g. – other clients of the cloud provider) as well as attacks from individual insiders in the cloud data center².

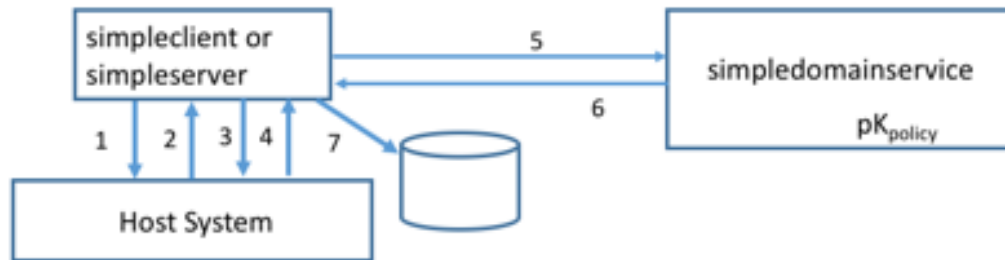
We do assume that data center insiders do not have physical access to computers running your software but only while that software is running. If insiders have physical access to computers while they execute programs, they could, for example, attach memory probes to buses and read and write data at will. Fortunately, restricting physical access to computers can be ensured by techniques commonly in use in data centers, for example by enclosing the computers running client software in locked cages that prevent physical access while the computers are running Cloudproxy programs, and for a few minutes afterwards (to protect against insiders harvesting memory remnants of running programs). An extra measure to ensure compliance might include cameras monitoring the caged systems with automated monitoring to detect breaches. Achieving our security model, however, *does not* require that all computers or networking equipment be in such cages. In fact, storage systems and networking equipment can remain unprotected without violating the foregoing security guarantee.

We also require that the software you provide (or use) not have remotely exploitable vulnerabilities but we *do not* limit what software you can run or significantly change the common programming model.

² Often we only require the integrity of the executing code. If code changes, it can leak or misuse data so verifiable code integrity is a basic security requirement which cannot be violated. Cloudproxy can also preserve code confidentiality by loading encrypted code as it would encrypted data but this is not always a requirement and none of our examples emphasize this equally achievable goal.

Overview of installation on Linux

The figure below should be familiar from [4]. We will refer to this throughout the document.



Initialization

1. simpleclient (or simpleserver) generates public/private key pair $PK_{\text{simpleclient}}$, $pK_{\text{simpleclient}}$. simpleclient requests Host System attest $PK_{\text{simpleclient}}$.
2. Host System returns attestation
3. simpleclient generates additional symmetric keys and request Host System seal symmetric keys and $pK_{\text{simpleclient}}$.
4. Host System returns sealed blobs.
5. simpleclient connects to simplifiedomainservice and transmits attestation.
6. simplifiedomainservice returns signed Program Certificate.
7. simpleclient stores sealed blobs and Program Certificate for later activations.

Initialization



Operation

1. Simpleserver reads previous sealed blobs and Program Certificate.
2. Simpleserver requests Host System unseal blobs yielding symmetric keys and private program key.
3. Host system returns unsealed blobs.
4. Simpleclient reads previous sealed blobs and Program Certificate.
5. Simpleclient requests Host System unseal blobs yielding symmetric keys and private program key.
6. Host system returns unsealed blobs.
7. Simpleclient and simpleserver open encrypted, integrity protected channel using their program keys and certificates.
8. Simpleclient transmits a request to retrieve secret.
9. Simpleserver retruns secret.

Operation

All the hardware enforced mechanisms used by the current version of Cloudproxy rely on the primary system software being Linux based. This includes a native Linux booted on an Intel SMX/TPM enabled platforms. Many of the currently popular hypervisors (Xen and KVM, for example) are Linux based so they too can run Cloudproxy “out of the box”. We use KVM as our hypervisor and provide the necessary support. A hypervisor that is hardware protected can extend, in Cloudproxy model, the confidentiality, integrity and other assurances to almost any layer of software without additional hardware. For example, we describe how Linux VM’s, each with independent Cloudproxy security assurances can be stacked on a Cloudproxy enabled KVM. Similarly, Docker based containers running within a Linux OS can be protected by Cloudproxy whether the Linux instance is booted on hardware directly or runs as a virtual machine on a KVM hypervisor.

We refer to a Cloudproxy Host booted on hardware as a Base Cloudproxy Host. Thus as the foregoing indicates one might have a KVM Base Host or a native Linux Based Host. With fairly minor changes, one could also make an BSD based platform a Base Cloudproxy Host or for that matter a Base Host employing a non-Linux based hypervisor or OS.

Hardware hosted Linux Base Cloudproxy systems employ a multi-boot sequence mediated by Grub to provide the required security properties, namely, the measurement of the Base software

and the Hardware mediated cryptographic services that ensure that secrets are bound to software. The booting Linux environment consists of:

1. *Grub*, which loads all the other booting software into memory.
2. *Tboot*, which mediates the authenticated boot using SMX and the TPM. Tboot plays the critical role of establishing that all the booted foundational software is measured and can use the TPM to store secrets and attest to the actual software booted.
3. The Linux kernel with TPM support.
4. A custom in memory file system, *initramfs* which is measured as part of the boot. *Initramfs* contains the compiled Cloudproxy components, and all security critical libraries and configuration files required by the Base Cloudproxy Host. Unlike, as is common with the normal Linux boot sequence, *initramfs* remains the root file system throughout the base system's lifetime so the integrity of its contents is assured. All other file systems are mounted under it. All such mounted file systems (as well as any network connected storage) are treated as untrusted by Cloudproxy programs, and files are encrypted and integrity protected for reading and writing.
5. *Dmccrypt* a common Linux component is used to encrypt the page-file and swap space³ used by the Cloudproxy Base system software, so no unencrypted data is stored on the paging disk.

We describe the role and general installation procedure for each of these configurations below. For each, there is also an appendix with detailed installation instructions. In addition, we describe preparing, measuring and running Linux based VM's with Cloudproxy support over a Cloudproxy enabled KVM. This configuration (or one using Xen) is the one that would commonly run in a large “*Infrastructure as a Service*” cloud provider facility and so Cloudproxy could thus provide the “security guarantees” above to traditional cloud customers.

The overall procedure for preparing and installing a Cloudproxy Base system is:

1. Install Linux. TPM2.0 requires a Version 4 kernel or later. The kernel should contain the appropriate TPM driver and *dmccrypt*. We remark that, for security reasons, the TCB of the root Cloudproxy Base system should be as small as possible. A carefully constrained kernel is a critical element of security. For example, our Linux base Cloudproxy systems either disallow loading dynamic modules or subject proposed loads to a small whitelist of signed modules.
2. Enable SMX, IOMMU and VT extensions in BIOS. Activate the TPM using the BIOS in the case of TPM 2.0 or the Cloudproxy TPM 1.2 utilities (*tpm-takeownership*) in the case of TPM 1.2.
3. Take ownership of the TPM and generate the appropriate keys.
4. Build TBOOT. Retrieve the required Authenticated Code Module (ACM) required by *Tboot* and place it and the *Tboot* image in */boot* as described below.

³ A subtle point: *Dmccrypt* at the time of this writing, encrypts but does not provide integrity protection. As a result, the “paging disk” (which might be flash or spinning media) should be in the “physical protection barrier” described in the Security Considerations section of this document.

5. Build the Cloudproxy binaries. Then build a custom *initramfs* which includes these binaries and any security critical configuration information as described below. Put the newly created *initramfs* in */boot/*
6. Configure Grub. You will boot a grub configuration which names Tboot, the ACM, your Linux kernel and your custom *Initramfs*. Grub configuration is handled automatically by later versions of *Tboot*.
7. Boot!

The “measurement” of a TPM mediated Base system will be stored in PCR 17 and 18 of your TPM. You will need this measurement for your *domain service* and the easiest way to get it is to have Cloudproxy report it after it boots on your test hardware. Cloudproxy does the rest! After you have successfully carried out these steps, you can build applications as described in [4].

Next, we describe each of these steps in more detail. After that, we discuss secure, scalable deployment of Cloudproxy applications including key management techniques followed by a brief description of Security Considerations which should be thought through for your particular deployment needs.

Hardware root of trust

Because of a lack of direct control over the hardware, maintaining our security model requires the use of a hardware root of trust that measures and reports the base system software that boots and ensure that only software trusted by you can access basic system secrets like signing keys. Without this, insiders could boot any privileged software and violate the security model you seek to enforce. Cloudproxy, as mentioned above can employ TPM1.2 or TPM 2.0 as “off the shelf” mechanisms hardware roots of trust.

The TPM is a dedicated microprocessor designed to offer facilities for generation of cryptographic keys, random number generation, remote attestation and data sealing. In this document, we use the TPM as the hardware root of trust of our Cloudproxy system. The Trusted Computing Group developed and maintains security standards for TPM design, implementation and provisioning and there are many compliant TPM manufacturers.

To use the TPM we need to enable it, and take ownership of it. Taking ownership of a TPM is the process of setting up an authorization value (e.g. – a password), which is subsequently required for certain kinds of TPM operations.

Taking ownership of a TPM is the process of setting up authorization values (eg. passwords) for the *ownerAuth*, *endorsementAuth*, and *lockoutAuth* types of authorization on TPM 1.2. TPM 2.0 employs similar concepts although TPM 2.0 is considerably easier to use than TPM 1.2. Many new Intel chip sets have firmware based TPM 2.0's.

Here are instructions for using TPM 1.2 and TPM 2.0 with Cloudproxy.

TPM 1.2

There are several keys in a TPM 1.2:

1. TPM Endorsement Key (EK): This key is created by the manufacturer and cannot be removed. Sometimes it can be changed by the owner of the computer.
2. TPM Storage Root Key (SRK): Is the 2048 bit RSA key created when configuring the ownership. This key is stored inside the chip and can be removed. This key is also used to seal/unseal data.
3. TPM Attestation Identity Key (AIK): This key is used by the TPM to sign attestations (to the state of the PCRs). It is encrypted and integrity protected by the SRK.

Cloudproxy uses the SRK to seal/unseal data and the AIK to sign attestations. Following are instructions on how to configure a TPM 1.2 for use by Cloudproxy.

1. First, we install tools needed to talk to the TPM device. We have a utility to do this. To obtain it, type:

```
go get -u -v github.com/google/go-tpm/...
```

2. Reset your TPM device as follows.

Disable the TPM device on your BIOS screen and reboot the system **after powering it off** to re-enable the TPM device in BIOS, and select the option to clear the device if you see one.

TPM devices require an assertion of physical presence to be cleared, and for most machines powering off and on does the trick. How presence is asserted depends on your machine, so consult your BIOS manual for detailed instructions. Also, note that the TPM 1.2 devices come disabled by default.

3. Take ownership of the device. This requires root privileges. The current Cloudproxy implementation assumes these authorizations are not set, when talking to the TPM.

```
unset TPM_OWNER_AUTH
unset TPM_SRK_AUTH
sudo su --preserve-environment
tpm-takeownership
```

4. Generate an *aikblob* and set permissions so that Cloudproxy may read it (note that this blob is encrypted and integrity protected by the TPM so it can be made public). To do this run

```
genaik
chmod 777 aikblob
```

Some troubleshooting tips for TPM 1.2 devices can be found in an appendix.

TPM 2.0

Enabling and taking ownership of and clearing a TPM 2.0 device is done entirely through BIOS and is typically relatively simple once you find the appropriate vendor specific instructions. Consult your BIOS manual for detailed steps. Detailed instructions for the Intel NUC 5i5MYHE can be found in an appendix. We have used both NUC based physical TPMs and firmware with uniform success.

Tboot

Trusted Boot (*tboot*) is a pre-kernel/VMM module that uses Intel Trusted Execution Technology (Intel TXT) to perform a measured and verified launch of an OS kernel/VMM. It uses the TPM for this and when successfully run, it stores the kernel/VMM measurement in the TPM PCRs, which is then read by the root Cloudproxy Tao.

Tboot installation overview: After making sure your TPM device is enabled (see previous sections), download the tboot source tarball from sourceforge.net, and build and install using `make` and `sudo make install`. This builds a `tboot.gz` image and places it in `/boot` for your bootloader to use when booting up. You will also need a SINIT AC module for your type of processor that you must manually place in `/boot`. This module enables TXT to load tboot. Tboot does the necessary grub configuration. Detailed instructions for tboot compilation and installation are in an appendix.

Some miscellaneous but probably useful notes on *tboot* follow.

txt-stat: The `txt-stat` tool installed by *tboot* is a Linux application that reads some of the TXT registers and will display the *tboot* boot log if *tboot* was run with '`logging=memory`'. It is useful to confirm that *tboot* booted the Base System successfully.

Launch control policy (LCP) and Verified Launch Policy (VLP): These are policies that govern which measured launch environments (MLEs, for eg. *tboot*) and kernels TXT and *tboot* are allowed to launch respectively. Tboot comes with a set of tools to install LCPs and VLPs, however they are not needed for Cloudproxy.

PCR values after trusted boot:

PCR 17 will contain a cryptographic hash as specified in Intel MLE Developers Manual including a SHA-1 hash of the tboot policy control value (4 bytes) | SHA-1 hash of tboot policy (20 bytes),

where the hash of the *tboot* policy will be 0s if TB_POLCTL_EXTEND_PCR17 is clear
PCR 18 will contain a cryptographic hash of the following (in this order):

- A SHA-1 hash of *tboot* (as calculated by *lcp_mlehash*)

- A SHA-1 hash of first module in *grub.conf* (e.g. Xen or Linux kernel)

- A SHA-1 hash of subsequent modules (including *initramfs*) specified in the multiboot process.

Detailed copy and paste instructions are in an appendix.

Preparing the Linux or KVM image

Linux systems need to be properly configured to meet our security goals. First, on most images, there should be no way for someone to log on and become root (for example, via SSH). Becoming root in a Cloudproxy image means the person doing so can violate policy since they can run any program or inspect any memory. On most Linux systems there should be no interactive users at all and any administrative tasks should be carried out using constrained, authenticated interfaces. The kernel should be as small as possible and no one should be able to load kernel modules --- i.e. any device driver should be examined and compiled into the original kernel image. Linux images should be carefully configured, maintained, curated and inspected.

We also need to ensure the measurement, authentication and integrity of any privileged code which certainly includes Cloudproxy components (like *linux_host*) as well as any dynamically linked libraries used by privileged applications. Cloudproxy applications should never trust “spinning” disks. All data written or read from these devices should be encrypted and integrity protected using keys protected by Cloudproxy. Finally, the disk to which images are swapped or paged should be encrypted (and either be physically secured, for example, by being in the same secure physical enclosure as the CPU and memory or cryptographically integrity protected).

We achieve these goals using two Linux features: *initramfs* and *dmccrypt*. *Initramfs* is a ram based storage system which is loaded when the kernel is loaded and measured by *tboot*. We put all security critical code including *linux_host* and dynamic libraries used by Cloudproxy applications in *initramfs*; often we also include the actual Cloudproxy programs (although this is not strictly necessary). *Initramfs* is mounted as the root file system during boot. On most Linux systems, it is dismounted and replaced by the “system disk.” On Cloudproxy Linux, *initramfs* not dismounted and serves as the permanent system disk; we do this by changing the shell script in *initramfs* that normally dismounts *initramfs* and mounts the “real” disk; instead, we simply mount the “real disk” under *initramfs*. All other disks are mounted under *initramfs* and should not be trusted. This procedure provides the complete control over system components and libraries we need. In addition, we encrypt the page file and swap devices using *dmccrypt*. *Dmccrypt* uses keys randomly generated at boot and we often require this “swap disk” be inside the secure physical cage protecting the computer (CPU and memory) running Cloudproxy.

Building InitRamfs

Everything upon which the security of your software relies must be measured, including basic system software and libraries you rely on. On Linux, we can protect and measure all that requires protection by incorporating it into a small, in memory, filesystem provided with the kernel boot image.

Detailed copy and paste instructions for `initramfs` are in an appendix.

Using Dmccrypt

Since most Linux systems employ virtual memory systems which temporarily store memory pages on disk, we must prevent secrets from being saved to a disk in a way that could be modified as programs are running or read after programs stop running. This would violate our security model. To do this, we employ `dmccrypt` which encrypts the disk used for paging. We also assume this disk is contained within the cage housing the Cloudproxy programs (since `dmccrypt` does not provide integrity protection). However, we *do not* require that the paging disk be protected before or after being used.

Detailed copy and paste instructions for configuring Linux to use `dmccrypt` are in an appendix.

Grub configuration

Grub is used to load the ACM, `tboot`, the OS (or hypervisor) and `initramfs` into main memory so that the measured boot can be carried out. Note that one can employ a grub configuration that boots non Cloudproxy enabled images without harm.

Grub is automatically configured by Tboot. Detailed copy and paste instructions for grub setup are in an appendix.

Installing and configuring a Cloudproxy supported KVM instance

We use KVM as the Cloudproxy hypervisor that can host isolated, measured Cloudproxy VMs. KVM Cloudproxy requires that the `libvirt-dev` and `libtspi-dev` be installed.

Detailed copy and paste instructions for KVM setup are in an appendix.

Installing and configuring a Cloudproxy supported TPM hosted Linux

In addition to a hypervisor rooted Cloudproxy stack, one could simply boot a Cloudproxy enabled Linux.

Detailed copy and paste instructions for Linux root setup are in an appendix.

Installing and configuring a Linux stacked on KVM

When using a hypervisor rooted Cloudproxy system, we need to provide Cloudproxy enabled VMs if those VMs need Cloudproxy protection (and they should!). Cloudproxy can support “stacked” systems that extend root protection to any layer of a software stack.

Detailed copy and paste instructions for configuring a stacked Linux Tao are in an appendix.

Installing and configuring Docker containers

Another stacked system protected by Cloudproxy is a Docker container within a Linux host (or stacked Linux host).

We have not completely documented this feature yet.

Host Certificates

Your TPM must have a signed certificate for its the root TPM key. This key is called an AIK in the case of TPM 1.2 and an Endorsement Key in the case of TPM 2.0. In many cases, an Endorsement Certificate, signed by the manufacturer, is included in NvRam for TPM 2.0; however, we usually produce certificates for these keys signed by the “policy-key” in our examples.

Domains

Domains and domain initialization was covered in [5] but when booting under TPM 1.2 and TPM 2.0, it is important to specify the proper PCR’s and, in the case of TPM 1.2, make sure the AIK was signed by the policy key using *aiksigner* and, in the case of TPM 2.0, the endorsement certificate is signed by *Endorsement*. In the TPM 1.2 case, the AIK serves as the attestation key. In the TPM 2.0 case, the attestation key is called the “Quote Key.” The quote key is certified using the endorsement key; this is done by the *QuoteServer*, provided in the sample code, so you need too make sure an instance of *QuoteServer* is available.

SimpleExample with TPM Tao’s

Complete instructions on running *simpleexample* under a TPM 1.2 Tao and a TPM 2.0 Tao appear in an appendix.

Considerations in deploying Cloudproxy

Key management

Key management supports the storage, distribution and critically, the rotation of symmetric content keys as well as establishing the “Program Key” for Cloudproxy programs (and hosts) and rotating the Program Key as keys expire or Cloudproxy programs are changed.

In most of our programs, keys and content types (like files) are objects and have domain-wide hierarchical universal names and epochs. Epochs are monotonically increased. For example, a file used by several Cloudproxy programs may contain customer information like customer names, addresses, etc. Since these are stored objects, they are encrypted and integrity protected and typically stored redundantly at several network accessible locations. The file `/jlm/file/us-zone/customer-file-location1`, epoch 2, type: file refers to the file with the indicated name; the epoch indicates the version of the file, different versions of the file are usually encrypted with different keys. These encryption keys also have domain-wide hierarchical universal names and epochs. Encryption keys protect other objects like keys or files. To decrypt a file, we need a chain of keys starting at the sealing key for the Cloudproxy program on a given node and terminating in the desired object (say a file). Each node of the chain consists of the name of the protector object, the protected objects and an encrypted blob consisting of the protected object encrypted (and integrity protected) by the protector key.

For example

```
/jlm/program-name/master-sealing-key, epoch 5, type: aes-128-gcm  
/jlm/key/us-zone/customer-folder-keys, epoch 2, type: aes-128-gcm  
/jlm/key/us-zone/customer-file-key, epoch 3, type: aes-128-gcm  
/jlm/file/us-zone/customer-file-location1, epoch 2, type: file
```

Often domains will want to partition keys to provide resilience by *protection zones*. One common way to do this is to have different keys for different geographic locations. This is easily accomplished by reserving a level of the hierarchical object namespace for the zone name as is done above.

Although the foregoing name, epoch based mechanisms facilitate key rotation, it is not a complete solution. Let’s first consider, given this framework, how we might rotate keys.

There are two common circumstances where new keys are introduced into a Cloudproxy environment. First, keys should be rotated regularly as part of good cryptographic hygiene and second, when programs change, their measurements change which means they can no longer access sealed data for earlier versions of the program. In both cases, similar techniques can be used to carry out the key rotation.

Key Rotation and software upgrades

We have provided sample code to maintain keys in `go/support/support_libraries/protected-objects` and `go/support/support_libraries/rotation_support`, although you can, of course, use

other software or customize our software; these libraries provide routines to store, retrieve, encrypt and decrypt content keys and file. They can build chains of keys like:

Protector Object: /jlm/program-name/master-sealing-key, epoch 5, type: aes-128-gcm

Protected Object: /jlm/key/us-zone-key, epoch 2, type: aes-128-gcm

Encrypted Protected Object

Protector Object: /jlm/key/us-zone-key, epoch 2

Protected Object: /jlm/key/us-zone-key/customer-folder-key, epoch 3, type: aes-128-gcm

Encrypted Protected Object

Protector Object: /jlm/key/us-zone-key/ customer-folder-key, epoch 3

Protected Object: /jlm/key/us-zone-key/customer-folder-key/customer-file-key, epoch 2, type: aes-128-gcm

Encrypted Protected Object

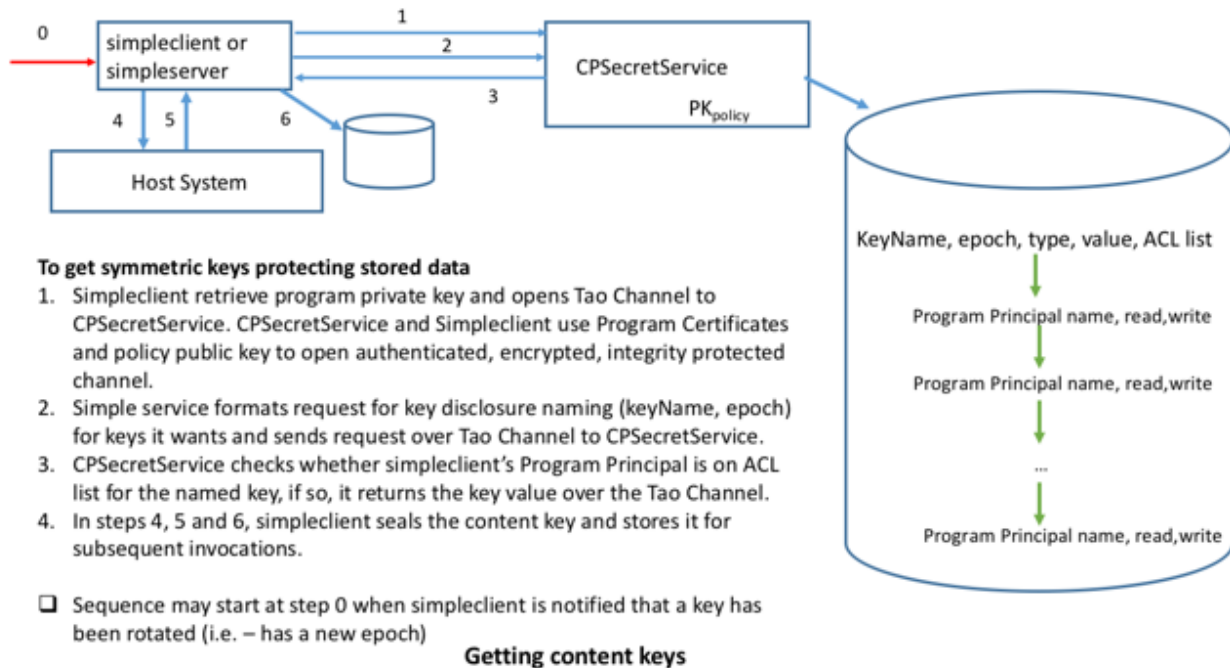
Protector Object: /jlm/key/us-zone/customer-file-key, epoch 3, type: aes-128-gcm

Protected Object: /jlm/file/us-zone/customer-file-location1, epoch 2, type: file

Using a keystore

The first, and maybe the simplest way to maintain keys and provide for program upgrade is by using a keystore. A sample keystore with the required functionality is implemented in `$CLOUDPROXY/go/support/infrastructure_support/secret_service`. Needless to say, this program itself would likely be implemented as a Cloudproxy program.

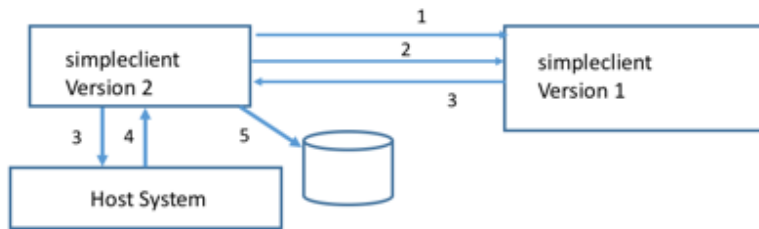
When you use a keystore, once a Cloudproxy program has obtained a Program Certificate, it simply opens a Tao Channel to the keystore to request other keys. The keys are “acl’ed” to Tao program names. When programs are upgraded, the new Tao Principal names are simple added to the acl list for the appropriate keys. That’s all there is to it! Of course, once a program obtains its keys, it can seal them and obtain them locally during subsequent invocations.



Since keys can be fetched any time after startup and no pre-existing state is required (except for the Program key), key rotation is easy and is focused at the Secret Service. The secret service can also publish alerts when new keys are available.

Using policy key based key disclosure

An alternative key provisioning mechanism when programs are upgraded is policy based provisioning. This is also quite simple. When programs are created, the private policy key holder simply signs a rule saying “Programs with Tao Principal Name $x_{old-program}$ can disclose keys K_1, K_2, \dots to the program with Tao Principal Name $x_{new-program}$. When the new program starts up for the first time, it opens a Tao Channel with an older version of the program that has access to the named keys, and supplied the signed rule. The old version of the program can authenticate the new version (using the Tao Channel), verify the rule applies and provide the named keys.



To get symmetric keys protecting stored data from another program or version

1. Simpleclient, version 2 retrieves program private key and opens Tao Channel to Simpleclient, version 2. Simpleclient, version 2 and Simpleclient, version 1 use Program Certificates and policy public key to open authenticated, encrypted, integrity protected channel.
2. Simple service formats request for key disclosure naming (keyName, epoch) for keys it wants and sends request over Tao Channel to CPSecretService accompanied with policy-key signed certificate stating policykey says ProgramPrincipalName (simpleclient, version 2) can read /jlm/zone-us/customer-folder-key, epoch 2.
3. Simpleclient, version 1 transmits keys over tao channel.
4. 4, 5, 6, Simpleclient, version 2 seals retrieved key for subsequent use.

Getting content keys

See `$CLOUDPROXY/go/support/support_libraries/secret_disclosure_support` for sample code to construct and verify a policy-key signed secret-disclosure statement.

Cloudproxy guards provide (see [5]) rather general support for building secure authorization policies for key disclosure to Cloudproxy authenticated programs. Key disclosure rules can easily be written (and verified) allowing groups of keys to be managed as a unit; for example, consider a simple rule like:

```
Policy-key says Tao-Principal-Name can-read /jlm/key/us-zone/customer-
folder-key, epoch 2
```

Other key upgrade and rotation mechanisms

While either the keystore mechanism or policy based secret disclosure mechanism can be used to protect any key, Cloudproxy allows several “peer programs” to collaborate to provide secrets and upgrade policy to authenticated peers, so almost any current key rotation or key provisioning mechanism can be adapted for Cloudproxy use.

Scale and deployment considerations

The instructions in this deployment guide are adequate for modest deployment scale (hundreds of servers). Large scale operation of Cloudproxy domains requires the same care as is needed for non-Cloudproxy platforms and infrastructures. While Cloudproxy generally imposes no new restrictions or barriers neither does it automatically provide complete support for scale deployment.

If you deploy Cloudproxy “at scale” be sure to consider the following:

1. Redundancy: To provide reliability and resilience as well as immunity to one (or too few) points of failure, services should be redundantly provisioned and data should be redundantly stored. Be careful to avoid unnoticed “correlated failures” that may make you believe you have greater resilience than you actually do.
2. Backup: The word should be enough.
3. Audit and logging: At scale it is impossible to detect problems as they occur without persistent, automatic and generous real time audit and logging. With Cloudproxy, you will want to encrypt and integrity protect your logs and audit trails and make sure the logs are also stored in a way that prevents loss in the face of failure or tampering.
4. Dashboards, analytics, predictive maintenance, and forensics: Once you have logs, you should automatically analyze them (again in real time) to spot potential attacks or emerging failures. Sometimes a “dashboard” is enough but often more attention will pay off in the end.
5. Key management and insider protection: While Cloudproxy frees you from the customary “untrustworthy CA” key management problem, you must safeguard critical keys (like the policy key) and put processes in place to avoid being at risk of having the entire security of your deployment compromised (deliberately or accidentally) by one or two “insiders” with regular access.
6. Break glass procedures: Make sure that when Black Swan failures occur you can respond immediately with documented and practiced processes. These practices should enable rapid “cold restart” while offering adequate insight into potential compromised data without extensive post incident investigation.

There is an emerging body of practice on running services at scale that is worth learning.

Just write your applications properly: simple, right?

While the Cloudproxy security model makes security reliance transparent and easily manageable, you still have to write your programs so that adversaries cannot exploit flaws in the programs you write (or the author of the VMM or BIOS wrote). This is not a trivial task but more programming tools, better programming languages and new techniques (such as proof-carrying-code) present plausible models for “much safer” software you can trust.

Oops: What to do if there’s a gigantic breach

Although it is hopefully a rare event, a large scale failure (maybe caused by flaws in critical software) can be remedied simply by redistributing the application with a new policy key (after fixing the security flaws in this case).

Bugs and suggestions

Please post bugs and suggestions to the Github repository. Those suggestions will be treated as having been licensed under the Cloudproxy license.

Acknowledgements

Tom Roeder, Fred Schneider, Kevin Walsh and Sid Telang.

References

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- [2] **CloudProxy Source code**, <http://github.com/jlmucb/cloudproxy>. Kevin Walsh and Tom Roeder were principal authors of the Go version.
- [3] **TCG, TPM specs**, http://www.trustedcomputinggroup.org/resources/tpm_library_specification
- [4] **Beekman, Manferdelli, Wagner**, Attestation Transparency: Building secure Internet services for legacy clients. AsiaCCS, 2016.
- [5] **Manferdelli**, CloudProxy Nuts and Bolts.
- [6] **Intel**, The MLE Development Guide

Appendix 1 - Supported Hardware

Intel or AMD CPU's and chipsets with secure extensions that support SKINIT. In addition, there must be a physical TPM on the motherboard or in firmware. Most new Intel CPU's have firmware TPM's so most new hardware is "automatically" Cloudproxy enabled. Most OEM advertising makes it very difficult to determine whether a model is "SMX" or "Trusted Technology" enabled and this together with a TPM, preferably 2.0 is what we need. Most "VPro" branded machines have the right hardware.

We've used the following:

1. Dell Omniplex
2. HP Elite laptops
3. Lenovo laptop (but be careful, we've had some models of Lenovo fail even when they should have worked)
4. Some Intel "NUC" machines have SMX support but make sure before you buy (only some i5 and i7 models).
5. We've also verified our software works on some new Dell machines with firmware TPMs.

Appendix 2 - TPM 1.2 Initialization and Operation

Troubleshooting TPM 1.2

We no longer use the trousers initialization routines but the ones from *go-tpm*. If you have Trousers running on your machine, you will first need to kill the *tcsd* daemon since it owns the TPM. You may not need to do all the following, but it works.

1. Clear the TPM in BIOS.
2. Turn off the machine.
3. Enable the TPM in BIOS.
4. Turn off the machine.
5. Boot normally and run *tpm-takeownership* after building Cloudproxy. This will be in the standard binary directory.
6. Follow the examples in *simpleexample* scripts.

Note that the TPM can become unusable if there are too many keys loaded. Run *tpm-keys* to see what keys are loaded. To unload existing keys run

```
tpm-keys -close
```

This is useful during testing since you may kill some malfunctioning programs before they unload tpm keys.

tpm-takeownership fails: *tpm-takeownership* assumes the owner of the device is not set. Reset the TPM through BIOS as explained in the main TPM 1.2 section to reset ownership.

tpm-takeownership not found: Make sure `$GOPATH/bin` is in your PATH environment variable.

Using tpm-clear: is NOT recommended for resetting the device or clearing ownership, and it is known to leave the TPM in a mucked state from where *tpm-takeownership* does not work. Reset the TPM through BIOS as explained in the main TPM 1.2 section to reset the device/ownership.

The following is obsolete:

Trousers supports some TPM initialization operations but we have tried to structure installation so that you do not need to use Trousers in simple Cloudproxy deployments. As a result, this section is informational.

Clearing and taking ownership of TPM 1.2 with Trousers

```
# tpm_clear --force
```

Tspi_TPM_ClearOwner failed: 0x00000007 - layer=tpm, code=0007 (7), TPM is disabled
We can see that the TPM is disabled, which is why we can't clear it. This can happen if we forget to actually enable the TPM in BIOS. The first thing to do would be to actually enable the TPM in BIOS. But if the TPM has been initialized before, we would receive the output that can be seen below:

```
# tpm_clear --force
```

TPM Successfully Cleared. You need to reboot to complete this operation. After reboot the TPM will be in the default state: unowned, disabled and inactive.

This would require us to reboot the computer for changes to take effect. When clearing the TPM we'll return it to the default state, which is unowned, disabled and inactive, as already mentioned. To enable the TPM afterwards, we need the owner password. But since the TPM owner has been cleared, there is no owner password and we can set a new one without entering the old one. We can also receive an error like the following:

```
# tpm_clear --force
```

Tspi_TPM_ClearOwner failed: 0x0000002d - layer=tpm, code=002d (45), Bad physical presence value

DO NOT SET ANY PASSWORD for the TPM.

```
# tpm_takeownership -z -y
```

If we later want to change either of the commands, we can do it with the `tpm_changeownerauth` command. If we pass the `--owner` argument to the `tpm_changeownerauth` command we'll be changing the administration password and if we pass the `--srk` into the `tpm_changeownerauth` command we'll be changing the SRK password. We can see the example of both commands in the output below:

```
# tpm_changeownerauth --owner
```

Enter owner password:

Enter new owner password:

Confirm password:

```
# tpm_changeownerauth --srk
```

Enter owner password:

Enter new SRK password:

Confirm password:

There are 5 keys in TPM:

TPM Endorsement Key (EK): This key is created by the manufacturer and cannot be removed. Sometimes it can be changed by the owner of the computer.

TPM Storage Key (SRK): Is the 2048 bit RSA key created when configuring the ownership.

This key is stored inside the chip and can be removed. The key is used to encrypt the Storage Key (SK) and Attestation Identity Key (AIK).# tpm_setenable --enable

Enter owner password:

Disabled status: false

tpm_setactive

Enter owner password:

Persistent Deactivated Status: false

Volatile Deactivated Status: false

There are usually two Endorsement Keys (EK): the public and private one. The private key is always stored at the TPM and cannot even be seen by anyone, while the public key can be displayed with the tpm_getpubek command.

tpm_getpubek

Tspi_TPM_GetPubEndorsementKey failed: 0x00000008 - layer=tpm, code=0008 (8), The TPM target command has been disabled

Enter owner password:

Public Endorsement Key:

Version: 01010000

Usage: 0x0002 (Unknown)

Flags: 0x00000000 (!VOLATILE, !MIGRATABLE, !REDIRECTION)

AuthUsage: 0x00 (Never)

Algorithm: 0x00000020 (Unknown)

Encryption Scheme: 0x00000012 (Unknown)

Signature Scheme: 0x00000010 (Unknown)

Public Key:

```
a350b3a3 3edddc30 06248f4f 5d3eb80a 34fcbea0 83dde002 8dffa703 e116f8b0
eb1962ee a65998b3 384aeb6e 85486be9 0316a6ca a189a5ba 2217b2a2 9da014db
dfbe7731 fb675e7a 438c4775 deea54fb 0c75de5d ba961950 3eda4555 d27a9a30
e94d39d0 a4ea314d a70eaf08 e49dd354 d57ed34d 234220d9 604471a9 86173050
9ff9b0e5 b65cb4b5 5f46a7f9 4378bd7e 8c61b91b ad312974 fef5d70f 84f4484f
e5c95300 0eef76f2 1667443f dc2fa82e 351d945e 6b5f75e8 828d010f 61541552
```

[...]

Appendix 3 - Tboot Launch Policy for TPM 1.2 (Optional)

You may build a Launch Control Policy (LCP) which controls what can be booted. If you have no LCP, the default policy will be applied and this should work, so this appendix, like the last is informational only.

Trousers and trousers-devel packages must already be installed in order to build lcpools.

Next you must modify grub so that you can choose the tbooted linux from the boot window.

If the verified launch policy is non-empty, it is extended into PCR 17 by default. Subsequent loaded modules are extended into PCR 18 by default (Check this!). This behavior can be changed with the VLP.

The Grub(2) configuration file is usually in /boot/grub and is called grub.cfg. It is updated automatically when a kernel is updated or when you run update-grub.

Tboot module must be added as the 'kernel' in the grub.conf file.

Note that the Lenovo T410 is known not to work with the Intel IOMMU; it will not successfully boot with TBOOT.

The final grub configuration file will look something like:

```
menuentry 'Ubuntu Linux 3.0.0-16-generic with TXT' --class ubuntu --class gnu-linux --
class gnu --class os {
    recordfail
    set gfxpayload=text
    insmod gzio
    insmod part_msdos
    insmod ext2
    set root='(hd0,msdos5)'
    search --no-floppy --fs-uuid --set=root 8ab78657-8561-4fa8-af57-bff736275cc6
    echo 'Multiboot'
    multiboot /boot/tboot.gz /boot/tboot.gz logging=vga,memory,serial
    echo 'Linux'
    module /boot/vmlinuz-3.0.0-16-generic /boot/vmlinuz-3.0.0-16-generic
    root=UUID=8ab78657-8561-4fa8-af57-bff736275cc6 ro splash vt.handoff=7 intel_iommu=on
    echo 'initrd'
    module /boot/initrd.img-3.0.0-16-generic /boot/initrd.img-3.0.0-16-generic
    echo 'sinit'
    module /boot/sinit51.bin /boot/sinit51.bin
```

Modify grub.conf to load the policy data file:

Edit grub.conf and add the following:

```
module /list.data
```

where you should use the path to this file.
Copy sinit into /boot and change run grub.conf then run update-grub.
Check the /boot directory to make sure tboot.gz is there.

The utility txt-stat in the utils subdirectory of tboot, can be used to view the DRTM boot log.

Reported problems with TPM 1.2

The TPM, driver/char/tpm/tpm.c, depends on TPM chip to report timeout values for timeout_a, b, and d. The Atmel TPM in Dell latitude 6430u, reports wrong timeout values (10 ms each), instead of TCG specified (750ms, 2000ms, 750ms, 750ms for timeout_a/b/c/d respectively). You can fix the driver code and it will work.

A permissive Tboot policy is required for the non-hypervisor solutions. In this case, the system will boot whatever it finds, but it will still perform measured launch, and clients interacting with the system will still be able to check that the right software was booted by checking the PCRs. Once the hypervisor is written, a more restrictive launch control policy is possible and maybe desirable.

Put 11_tboot in /etc/grub.d

Information on PCR Usage on TPM 1.2

PCR 17 will be extended with the following values (in this order):

- The values as documented in the MLE Developers Manual
- SHA-1 hash of: tboot policy control value (4 bytes) SHA-1 hash of tboot policy (20 bytes)
: where the hash of the tboot policy will be 0s if TB_POLCTL_EXTEND_PCR17 is clear

PCR 18 It will be extended with the following values (in this order):

- SHA-1 hash of tboot (as calculated by lcp_mlehash)
- SHA-1 hash of first module in grub.conf (e.g. Xen or Linux kernel)

PCR *: tboot policy may specify modules' measurements to be extended into PCRs specified in the policy

The default tboot policy will extend, in order, the SHA-1 hashes of all modules (other than 0) into PCR 19.

The LCP consists of policy elements.

To specify launch policy via a list of hashes:

```
0. sudo bash
1. lcp_mlehash -c "command line for tboot from grub.cfg" /boot/tboot.gz > mle_hash
   - the command line in this case is the string from /boot/grub/grub.cfg
   - after multiboot tboot.gz, e.g., "logging=vga,memory,serial"
   - copy and paste:
./lcp_mlehash -c "logging=vga" /boot/tboot.gz>mle_hash
2. lcp_crtpolelt --create --type mle --ctrl 0x00 --minver 17 --out mle.elc mle_hash
   - copy and paste
./lcp_crtpolelt --create --type mle --ctrl 0x00 --minver 17 --out mle.elc mle_hash
```

```

3. lcp_crtpollist --create --out list_unsig.lst mle.elc pconf.elc
   - copy and paste
./lcp_crtpollist --create --out list_unsig.lst mle.elc pconf.elc
4. lcp_crtpol2 --create --type list --pol list.pol --data list.data list_unsig.lst
   - copy and paste
./lcp_crtpol2 --create --type list --pol list.pol --data list.data list_unsig.lst
5. cp list.pol /boot
   - copy and paste
cp owner_list.data /boot

```

Next create a verified Launch policy:

```

1. tb_polgen/tb_polgen --create --type nonfatal vl.pol
   - copy and paste
./tb_polgen/tb_polgen --create --type nonfatal vl.pol
2. tb_polgen/tb_polgen --add --num 0 --pcr 18 --hash image
   --cmdline "the command line from linux in grub.conf"
   --image /boot/vmlinuz-2.6.18.8
   vl.pol
   - copy and paste
./tb_polgen/tb_polgen --add --num 0 --pcr 18 --hash image --cmdline "root=UUID=cf6ae6b5-
abb5-4d5d-b823-bd798a0621de ro quiet splash $vt_handoff" --image /boot/vmlinuz-3.5.0-23-
generic vl.pol
3. tb_polgen/tb_polgen --add --num 1 --pcr 19 --hash image
   --cmdline ""
   --image /boot/initrd-2.6.18.8
   vl.pol
   - copy and paste
./tb_polgen/tb_polgen --add --num 1 --pcr 19 --hash image --image /boot/initrd-3.5.0-23-
generic vl.pol

```

If any mle can be launched:

```
./lcp_crtpol2 --create --type any --pol any.pol
```

Define tboot error TPM NV index:

```

1. lcptools/tpmnv_defindex -i 0x20000002 -s 8 -pv 0 -rl 0x07 -wl 0x07
   -p TPM-password
   - copy and paste
./lcptools/tpmnv_defindex -i 0x20000002 -s 8 -pv 0 -rl 0x07 -wl 0x07

```

Define LCP and Verified Launch policy indices:

```

1. lcptools/tpmnv_defindex -i owner -s 0x36 -p TPM-owner-password
   - copy and paste
./lcptools/tpmnv_defindex -i owner -s 0x36 -p TPM-owner-password
2. lcptools/tpmnv_defindex -i 0x20000001 -s 256 -pv 0x02 -p TPM-owner-password
   - copy and paste
./lcptools/tpmnv_defindex -i 0x20000001 -s 256 -pv 0x02 -p TPM-owner-password

```

Write LCP and Verified Launch policies to TPM:

```

(modprobe tpm_tis; tcscd;)
1. lcptools/lcp_writepol -i owner -f [any.pol|list.pol] -p TPM-password
   - copy and paste
./lcp_writepol -i owner -f list.pol -p <ownerauth password>
2. If there is a verified launch policy:
   lcptools/lcp_writepol -i 0x20000001 -f vl.pol -p TPM-password
   - copy and paste
./lcptools/lcp_writepol -i 0x20000001 -f vl.pol -p TPM-password

```

Use lcp_crtpollist to sign the list:

1. openssl genrsa -out privkey.pem 2048

2. `openssl rsa -pubout -in privkey.pem -out pubkey.pem`
3. `cp list_unsig.lst list_sig.lst`
4. `lcp_crtpollist --sign --pub pubkey.pem --priv privkey.pem --out list_sig.lst`

Use openssl to sign the list:

1. `openssl rsa -pubout -in privkey.pem -out pubkey.pem`
2. `cp list_unsig.lst list_sig.lst`
3. `lcp_crtpollist --sign --pub pubkey.pem --nosig --out list_sig.lst`
4. `openssl genrsa -out privkey.pem 2048`
5. `openssl dgst -sha1 -sign privkey.pem -out list.sig list_sig.lst`
6. `lcp_crtpollist --addsig --sig list.sig --out list_sig.lst`

Appendix 4 - TPM 2.0 Installation and Operation

The interface to TPM 2.) is via BIOS, unlike TPM 1.2. TPM 2's have pretty severe limitations on the number of open handles.

Troubleshooting TPM 2.0

Too many open handles: Most TPM 2.0's can support only a small number of open handles (as little as 4). To clear handles, run

```
tpm2_util.exe -command=Flushall
```

You must compile the utility to do this by:

```
cd $CLOUDPROXY/src/tpm2
make -f tpm2.mak
```

TPM2 Error code 0x9a2:

This error code corresponds to 'bad authorization' and is returned when a command has incorrect authorization values. If you encounter this when using CloudProxy, it means your TPM is not provisioned properly, and you would need to clear the TPM.

TPM2 Error code 0x902:

This error code is returned when the TPM is out of memory for creating new objects. Run the TPM2_FlushContext command for all open handles (which can be obtained by running the TPM2_GetCapability command).

TPM2_Clear on the Intel NUC 5i5MYHE:

Clearing the TPM has the effect of setting all authorization values to *emptyAuth* (i.e. empty string as password). Clearing a TPM can be done through the maintenance screen of the BIOS, which is accessed by booting up the machine with the yellow BIOS jumper removed, and pressing the Esc key when the startup BIOS screen is shown.

Appendix 5 - Tboot

To build tboot you first need to install trousers and libtspi-dev.

Installing TBOOT

1. Download the source tarball from <https://sourceforge.net/projects/tboot/>
2. In the untarred source directory run:
`make`
3. In the same directory as above, run
`sudo make install`
4. Check the `/boot` directory contains the `tboot.gz` file
5. Download the SINIT AC module corresponding to your processor from
<https://software.intel.com/en-us/articles/intel-trusted-execution-technology>
6. Copy the SINIT `*.BIN` file to your `/boot` directory
7. To start using tboot, you need to update your bootloader configuration. If using grub2, run
`sudo update-grub`

And check that `/boot/grub/grub.cfg` now contains an entry for tboot which looks like:

```
menuentry 'Ubuntu GNU/Linux, with tboot 1.9.4 and Linux 4.4.0-24-
generic' --class ubuntu --class gnu-linux --class gnu --class os --class
tboot {
    insmod multiboot2
    insmod part_gpt
    insmod ext2
    set root='hd0,gpt2'
    if [ x$feature_platform_search_hint = xy ]; then
        search --no-floppy --fs-uuid --set=root --hint-bios=hd0,gpt2
--hint-efi=hd0,gpt2 --hint-baremetal=ahci0,gpt2 1df958d9-c01d-43d8-
a1b2-18d022843d3f
    else
        search --no-floppy --fs-uuid --set=root 1df958d9-c01d-43d8-
a1b2-18d022843d3f
    fi
    echo      'Loading tboot 1.9.4 ...'
    multiboot2      /boot/tboot.gz logging=serial,memory
    echo      'Loading Linux 4.4.0-24-generic ...'
    module2 /boot/vmlinuz-4.4.0-24-generic root=UUID=1df958d9-c01d-
43d8-a1b2-18d022843d3f ro quiet splash intel_iommu=on noefi
    echo      'Loading initial ramdisk ...'
    module2 /boot/initrd.img-4.4.0-24-generic
    echo      'Loading sinit 5th_gen_i5_i7_SINIT_79.BIN ...'
    module2 /boot/5th_gen_i5_i7_SINIT_79.BIN
}
```

8. If using tboot with cloudproxy, you need to adjust your bootloader configuration differently (to add the initramfs). Instructions to do so can be found in a subsequent appendix.

More information about Tboot can be found in the README in the Tboot source directory.

Tboot will copy and alter the e820 table provided by GRUB to "reserve" its own memory plus the TXT memory regions. These are marked as E820_UNUSABLE or E820_RESERVED so that the patched Xen code can prevent them from being assigned to dom0. The e820 table is not altered if the measured launch fails for any reason.

Appendi 6 - Initramfs

1. Download and unpack Linux kernel and busybox source

We use Busybox to provide a basic set of Unix utilities to use on the VM instance.

```
mkdir /tmp/tiny_linux
cd /tmp/tiny_linux
curl https://www.kernel.org/pub/linux/kernel/\
v4.x/linux-4.0.3.tar.xz | tar xJf -
curl https://busybox.net/downloads/busybox-1.23.2.tar.bz2 \
| tar xjf -
```

2. Build busybox

```
mkdir -pv ../obj/busybox-x86
make O=../obj/busybox-x86 defconfig
make O=../obj/busybox-x86 menuconfig
```

The last command opens an interactive menu. We will select to build BusyBox statically to ensure it can run on the VM instance.

```
-> Busybox Settings
-> Build Options
[ ] Build BusyBox as a static binary (no shared libs)
```

Select the above option, save your selection and quit the interactive menu. Next we build Busybox as shown below.

```
cd ../obj/busybox-x86
make -j2
make install
```

3. Build initrd image

Next we build the init RAM disk image.

```
mkdir -p /tmp/tiny_linux/initramfs/x86-busybox
cd /tmp/tiny_linux/initramfs/x86-busybox
mkdir -pv {bin,sbin,etc,proc,sys,usr/{bin,sbin}}
cp -av $TOP/obj/busybox-x86/_install/*
```

We will need an init script. The following minimal script works.

```
#!/bin/sh
mount -t proc none /proc
```

```
mount -t sysfs none /sys
mount -t devtmpfs -o mode=0755 udev /dev
```

```
exec /bin/sh
```

Save the above script in a file called init and make it executable.

```
chmod +x init
```

Finally we archive and compress this directory to create our initrd image.

```
find . -print0 \
| cpio --null -ov --format=newc \
| gzip -9 > /tmp/tiny_linux/obj/initramfs-busybox-x86.cpio.gz
```

4. Build linux kernel.

```
cd /tmp/tiny_linux/linux-4.0.3
make O=../obj/linux-x86-basic x86_64_defconfig
make O=../obj/linux-x86-basic kvmconfig
make O=../obj/linux-x86-basic -j2
```

The kernel image can be found at /tmp/tiny_linux/obj/linux-x86-basic/arch/x86_64/boot/bzImage

Appendix 7 - Dmccrypt

Dmccrypt is a standard Linux component, we use it to encrypt page tables so that there is no sensitive information on disks. As a result, while the system paging disk should be in the “cage” during operation, these disks do not contain sensitive information and do not have to be specially handled after operation. Disks do not have to be wiped.

fdisk to make partitions (/dev/sda)

mkfs.ext2 to make file system

change to single user mode

grub-install root-directory=/dev/sda

CRYPTDISKS_ENABLE=Yes in /etc/defaults/cryptdisks

```
/etc/crypttab
swap /dev/hda1 /dev/urandom swap
/etc/fstab
/dev/swapper/swap none
swap sw,pri=1 0 0
```

For swap:

```
swapoff -a
cryptsetup [-c aes -h sha256] -s 128 -d /dev/urandom create swap
/dev/sda1
mkswap /dev/mapper/swap
swapon /dev/mapper/swap
```

Shell script prototype

```
#
#   For debian, first do the following:
#
#   Change to single user mode
#       cat /proc/swaps      (to find out swap device)
#       /sbin/runlevel      (to find runlevel)
#       /sbin/telinit 1     (to go single user)
#       ctrl-alt-f7 puts Linux into console mode.
#
#   /etc/defaults/cryptdisks
#       CRYPTDISKS_ENABLE=Yes
#   /etc/crypttab
#       swap /dev/sda5 /dev/urandom swap
#   /etc/fstab
#       /dev/mapper/swap none
#       swap sw,pri=1 0 0
#   /etc/crypttab
#       swap /dev/sda5 /dev/urandom cipher=aes-cbc-essiv:sha256,size=256,hash=sha256,swap
#   /etc/fstab
#       /dev/mapper/cryptoswap none swap sw 0 0
#
#   /dev/sda5 is swap device
#
# Ref: The whole disk Nov, 2006, Linux-magazine.com, Michael Nerb
#
```

```
swapoff -a
cryptsetup -c aes -h sha256 -s 128 -d /dev/urandom create swap /dev/sda5
mkswap /dev/mapper/swap
swapon /dev/mapper/swap
#
# To remove:
#   swapoff /dev/mapper/swap
#   cryptsetup [luks] remove swap
```


Appendix 8 - Grubby, Grub, Grub (Optional)

This appendix is informational, tboot does all of this automatically in the normal course of events.

The new tboot module must be added as the 'kernel' in the grub.conf file. The existing 'kernel' entry should follow as a 'module'. The SINIT ACM module must be added to the grub.conf boot config as the last module, for example:

```
root (hd0,1)
kernel /tboot.gz logging=serial,vga,memory
module /vmlinuz-2.6.18 root=/dev/VolGroup00/LogVol100 ro
module /initrd-2.6.18.img
module /Q35_SINIT_17.BIN
```

The kernel module is tboot. The next is the Linux (possibly KVM enabled) kernel. Next is the initramfs and finally the ACM module.

GRUB2 does not pass the file name in the command line field of the multiboot entry (module_t::string). Since the tboot code is expecting the file name as the first part of the string, it tries to remove it to determine the command line arguments, which will cause a verification error. The "official" workaround for kernels/etc. that depend on the file name is to duplicate the file name in the grub.config file like below:

```
menuentry 'Xen w/ Intel(R) Trusted Execution Technology' {
    recordfail
    insmod part_msdos
    insmod ext2
    set root='(/dev/sda,msdos5)'
    search --no-floppy --fs-uuid --set=root 4efb64c6-7e11-482e-8bab-
07034a52de39
    multiboot /tboot.gz /tboot.gz logging=vga,memory,serial
    module /xen.gz /xen.gz iommu=required dom0_mem=524288
com1=115200,8n1
    module /vmlinuz-2.6.18-xen /vmlinuz-2.6.18-xen
root=/dev/VolGroup...
    module /initrd-2.6.18-xen.img /initrd-2.6.18-xen.img
    module /Q35_SINIT_17.BIN
}
```

Appendix 9 - Setting up a Cloudproxy enabled KVM host

In this section we will see how to set up a Cloudproxy enabled KVM host, which can launch independent Cloudproxy VMs (i.e. VMs with a stacked Cloudproxy host running)

The root Tao of our CloudProxy KVM host will be a “soft” Tao: i.e. a dummy Tao based on a keyset that is stored on disk, encrypted and integrity protected by a password.

CoreOS Instances

1. Generate soft Tao.

We set up the Tao domain directory and the host directory, where the domain and host configuration information will be stored. We copy in a domain template which we will use in this example and setup the soft Tao which generates a password protected keyset in the `linux_tao_host` directory.

```
mkdir /tmp/domain
cd /tmp/domain
mkdir linux_tao_host
cp $GOPATH/src/github.com/jlmucb/\
    cloudproxy/go/apps/simpleexample/SimpleDomain/\
    domain_template.simpleexample .
tao domain newsoft \
    -config_template domain_template.simpleexample \
    -soft_pass xxx ./linux_tao_host
```

2. Initialize domain

We initialize the domain which will create policy keys, saved in the domain directory and protected by the password. It will also create a `tao.config` file which stores all the domain configuration information and is used by the host to load the domain.

```
tao domain init \
    -tao_domain . \
    -config_template domain_template.simpleexample \
    -pass xxx
```

3. Create linux_host image

We statically build a linux host image (so that it is executable in the VM) and create a tarball of the `linux_host` binary and the host/domain configuration information it needs to run. When launching a VM, the KVM Cloudproxy unpacks this, mounts it on the VM and runs `linux_host` as a stacked Cloudproxy host.

```
CGO_ENABLED=0 go install -a \
    -installsuffix nocgo \
        github.com/jlmucb/cloudproxy/go/apps/linux_host
cp $GOPATH/bin/linux_host .
tar -C . -czf linux_host.img.tgz $(ls .)
```

4. Get CoreOS image

Get the latest stable image of CoreOS.

```
curl -G http://stable.release.core-os.net/amd64-usr/\
    current/coreos_production_qemu_image.img.bz2 > \
    coreos_production_qemu_image.img.bz2
bunzip2 coreos_production_qemu_image.img.bz2
```

5. Setup ssh-agent

We generate a key pair and setup ssh-agent for SSHing into the VM instance.

```
ssh-keygen -t rsa -f ./id_rsa
ssh-add id_rsa
```

6. Initialize and start host

We initialize and start the KVM Cloudproxy host

```
tao host init \
    -root \
    -tao_domain . \
    -pass xxx \
    -hosting kvm_coreos \
    -kvm_coreos_img coreos_production_qemu_image.img \
    -kvm_coreos_ssh_auth_keys id_rsa.pub

tao host start \
    - tao_domain . \
    - host ./linux_tao_host
```

7. Launch hosted program (i.e. VM)

We launch a Cloudproxy VM through the Cloudproxy KVM host. We create another temporary directory where the linux_host tarball is unpacked and mounted on the VM and run from.

```
mkdir another_tmp
tao_launch run \
    -tao_domain . \
    -host ./linux_tao_host \
```

```
-verbose kvm_coreos:linux_host.img.tgz \
-- linux_host.img.tgz another_tmp 2222
```

Launching a VM with custom kernel through KVM Cloudproxy

1. Generate soft Tao.

We set up the Tao domain directory and the host directory, where the domain and host configuration information will be stored. We copy in a domain template which we will use in this example and setup the soft Tao which generates a password protected keyset in the linux_tao_host directory.

```
mkdir /tmp/domain
cd /tmp/domain
mkdir linux_tao_host
cp $GOPATH/src/github.com/jlmucb/\
  cloudproxy/go/apps/simpleexample/SimpleDomain/\
  domain_template.simpleexample .
tao domain newsoft \
  -config_template domain_template.simpleexample \
  -soft_pass xxx ./linux_tao_host
```

2. Initialize domain

We initialize the domain which will create policy keys, saved in the domain directory and protected by the password. It will also create a tao.config file which stores all the domain configuration information and is used by the host to load the domain.

```
tao domain init \
  -tao_domain . \
  -config_template domain_template.simpleexample \
  -pass xxx
```

3. Initialize and start host

We initialize and start the KVM Cloudproxy host

```
tao host init \
  -root \
  -tao_domain . \
  -pass xxx \
  -hosting kvm_custom \

tao host start      -tao_domain .
```

4. Copy Kernel image and initrd image to domain directory

Replace the paths in the following lines with the path to your kernel and initrd image.

```
cp /tmp/tiny_linux/obj/linux-x86-basic/arch/x86_64/boot/bzImage .  
cp /tmp/tiny_linux/obj/initramfs-busybox-x86.cpio.gz .
```

5. Launch hosted program (i.e. VM)

We launch a Cloudproxy VM through the Cloudproxy KVM host.

```
tao_launch run \  
-tao_domain . \  
-verbose kvm_custom: \  
-- bzImage initramfs-busybox-x86.cpio.gz 2222
```

Appendix 10 - SimpleExample with a TPM 2.0 Tao

Here we go through the procedure to run *simpleexample* (as described in Cloudproxy Nuts and Bolts) under a Linux whose root host is a TPM 2.0 rather than a “Soft Tao.” The example in Cloudproxy Nuts and Bolts, employed *SimpleDomainService*, which did not compare the program requesting the Program Certificate against a list of authorized programs as would be required to meet Cloudproxy’s security goals. Here we show how to run *simpleexample* either with the *SimpleDomainServer* or *domain_server*. *domain_server* verifies the attestation before issuing Program Certificates so it is a much more realistic example.

A new aspect of running with TPM 2.0 over the Soft Tao based Tao is that we will produce an endorsement certificate for the TPMs endorsement key, stored in *endorsement_certificate*, corresponding to the certificate for the Soft Tao key, *soft_tao_cert*.

In our example, the endorsement certificate is signed by the policy key but in real deployment one might rely on endorsement certificates provided by the TPM vendor. These certificates can often be found in NvRam in compliant TPMs.

We assume you’ve carried out the *simpleexample* tutorial including the following steps:

As root:

```
mkdir /Domains
chown your-username /Domains
```

As your username

```
cd $CLOUDPROXY/go
go install ...
cd $CLOUDPROXY/go/apps/simpleexample/SimpleDomain
~/bin/copybins (run copybins)
```

We also assume you have removed unneeded files and killed previous instances of Cloudproxy programs as described in the *simpleexample* instructions in “Nuts and Bolts.” Remember that ultimately, in a running version, the provisioned files as well as the Cloudproxy components will be in the *initramfs*.

Having done the foregoing, to run *simpleexample* under a TPM 2.0 Tao, as whatever your user account is, run

```
./initsimpleexampletpm2
```

This will construct the *endorsement certificate*, create the storage hierarchy for *simpleexample* under TPM2 and perform all the same tasks as the corresponding non-tpm *simpleexample* script. Then, as root

```
./initttpm2domain
./runalltpm2simple
```

This will run the entire example with the *SimpleDomainServer*, which doesn't check the Tao Principal Name to ensure the program requesting the Program is one of the programs "trusted" in the domain. Another program, *domain_server* does these checks. To run this full example, first, you must create a file in */Domains/domain.simpleexampletpm2* called *TrustedEntities* naming the authorized programs. Create the file

/Domains/domain.simpleexampletpm2/TrustedEntities and add two lines:

```
trusted_program_tao_names: "Tao-name-of-simpleserver"
trusted_program_tao_names: "Tao-name-of-simpleclient"
```

You can copy the Tao Principal Names of these programs from the output of *runalltpm2simple*, you will need to escape and quotes in the names by preceding the quote with a \. If you want to run the C++ client, you will also need a line for *simpleexample client.exe*.

This will look something like:

```
trusted_program_tao_names:
"tpm2 ([6b6b6304dd984fc9af91569e5a25366eacf66fb8821830665f93e94aaf5943a2
]).PCRs (\ "17, 18", \ "ffffffffffffffffffffffffffffffffffffffff", ... ).TrivialGuard (\ "Liberal \ " ).Program ([a04ae93feaca725df5c48939a70812df4f065a14d4b27e1618976810
aac2e7a7]).PolicyKey ([1bbf737f810e672aee79b82d946f945fe7472dddc982bb17df21f44da1fa50d3]) "
```

Next, as root, run:

```
./inittpm2domain
./runalltpm2full
```

The only difference between `./runalltpm2full` and `./runalltpm2simple` is the invocation of *domain_server* instead of *SimpleDomainServer*.

Appendix 11 - SimpleExample with a TPM 1.2 Tao

Here we go through the procedure to run *simpleexample* (as described in Cloudproxy Nuts and Bolts) under a Linux whose root host is a TPM 1.2 rather than a “Soft Tao.” In addition, while the *simpleexample* employed a *SimpleDomainService*, here we use a domain service which fully verifies the attestation before issuing Program Certificates.

We assume you’ve carried out the *simpleexample* tutorial including the following steps:

As root:

```
mkdir /Domains
chown yourusername /Domains
```

As yourusername

```
cd $CLOUDPROXY/go
go install ...
cd $CLOUDPROXY/go/apps/simpleexample/SimpleDomain
~/bin/copybins (run copybins)
```

The procedure is very similar to running *simpleexample* with TPM 2.0. There is only one substantive difference and that is instead of having the “initdomain” script call *Endorsement* to sign the endorsement certificate, it calls *aiksigner* to sign the TPM1.2 style AIK producing an “*AIK Certificate*.” Ultimately, the *AIK Certificate* plays the same role as the *Quote Certificate* when negotiating the Program Certificates with the domain services.

Remember that ultimately, in a running version, the files as well as the Cloudproxy components will be in the *initramfs*.

Summarizing, to run *simpleexample* under a TPM 1.2 Tao, first (as whatever your user account is),

```
./initsimpleexampletpm1
```

As root

```
./inittpm1domain
./runalltpm1simple
```

Construct the “TrustedEntities” file as described in the previous appendix. Then clean up the previous invocations and files and, as root, run

```
./inittpm1domain
./runalltpm1full
```


Appendix 12 - SimpleExample in VM using stacked host, KVM host under TPM 2.0 Tao

Here we go through the procedure to run *simpleexample* (as described in Cloudproxy Nuts and Bolts) under a Linux in a VM hosted by KVM. The KVM host employs a root host that is a TPM 2.0 host rather than a “Soft Tao.” This is our first “non-trivial” example of a stacked Tao.

We will not labor the details as we did earlier examples since we’ve carried out almost all the steps in previous examples. You must:

1. Build the KVM instance including adding the Cloudproxy components in *initramfs* and incorporating *dmccrypt*.
2. Provision the TPM *endorsement certificate* for the KVM Host.
3. Build the Linux VM including adding the Cloudproxy components in *initramfs* and incorporating *dmccrypt*. Note: The *linux_host* in the VM will be stacked on the KVM Tao. You will have to create a domain in the Linux VM. You can do this as your building the *initramfs* by copying a domain template and calling `tao domain init`, as we did in other examples. For the corresponding call to *init* and *start* host, consult *kvm_coreos_facrotty.go*, it will look like:

```
$BINPATH/tao host init -tao_domain $LINUXDOMAIN -stacked -parent_type
file -parent_spec `tao::RPC+tao::FileMessageChannel(/dev/vport0p1)` -
hosting process
```

4. Run two Linux VM's, one with with the *SimpleClient* and one with *SimpleServer*.

The domain services run as before, but in a third VM.