CMSC 277: Homework #7

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01 December 2005

1. Fix a nonstandard model of analysis * \mathfrak{R} . Let $A \subseteq \mathbb{R}$ and $r \in \mathbb{R}$. Show that r is in the closure of a if and only if there exists $a \in (A^*) \cap \mathcal{F}$ such that $\operatorname{st}(a) = r$.

Let $r \in \overline{A}$. Then the sentence

$$\forall \varepsilon > 0 \exists s \in A(|r-s| < \varepsilon)$$

is in $\operatorname{Th}(\mathfrak{R}) = \operatorname{Th}({}^*\mathfrak{R})$. Let ε be a positive infinitesimal, then there exists some (finite) $a \in {}^*A$ such that $|a-r| < \varepsilon$ so that $a \approx r$. But as $r \in \mathbb{R}$ it follows that $\operatorname{st}(a) = r$.

Similarly, if $a \in {}^*A \cap \mathcal{F}$ such that $\operatorname{st}(a) = r$ then the sentence

$$\forall \varepsilon > 0 \exists s \in \underline{A}(|\underline{r} - s| < \varepsilon)$$

is in $\operatorname{Th}(\mathfrak{R}) = \operatorname{Th}({}^*\mathfrak{R})$ since substituting \underline{a} for s satisfies it. That $r \in \overline{A}$ follows immediately.

2. Fix a nonstandard model of analysis * \mathfrak{R} and let $f: \mathbb{R} \to \mathbb{R}$. Show that f is uniformly continuous on \mathbb{R} if and only if for all $a, b \in {}^*\mathbb{R}$ with $a \approx b$ we have ${}^*f(a) \approx {}^*f(b)$

We proceed using the contrapositive.

Assume that there exist $a, b \in {}^*\mathbb{R}$ such that $a \approx b$ but ${}^*f(a) \not\approx {}^*f(b)$. Then there exists some $\varepsilon > 0$ such that $|{}^*f(a) - {}^*f(b)| > \varepsilon$. The sentence

$$\forall \delta > 0 \exists x, y \in \mathbb{R}(|x - y| < \delta \land |f(x) - f(y)| > \underline{\varepsilon}$$

is in $Th(*\mathfrak{R}) = Th(\mathfrak{R})$ since substituting a for x and b for y satisfies it. But this sentence says precisely that f is not uniformly continuous.

Assume that f is not uniformly continuous and fix $\varepsilon > 0$. Then the sentence

$$\forall \delta > 0 \exists a, b \in \mathbb{R}(|a - b| < \delta \land |f(a) - f(b)| > \underline{\varepsilon}$$

is in $\operatorname{Th}(\mathfrak{R}) = \operatorname{Th}({}^*\mathfrak{R})$. Letting δ be a positive infinitesimal we then have that $a \approx b$ but ${}^*f(a) \not\approx {}^*f(b)$ so that the second property does not hold.

3. Let \mathfrak{L} be the empty language. For each $n \in \mathbb{N}^+$ let $\sigma_n \in \operatorname{Sent}_{\mathfrak{L}}$ be

$$\exists x_1 \cdots \exists x_n \left(\bigwedge_{1 \le i < j \le n} x_i \ne x_j \right)$$

Let $\Sigma = {\sigma_n \mid n \in \mathbb{N}^+}$ and define $T = \operatorname{Cn}(\Sigma)$. Show that T has QE and is complete.

First note that any model of T is necessarily infinite since any model of σ_n has a cardinality of at least n. Hence any model of T cannot be finite.

To show that T has quantifier elimination let $\exists y(\alpha_1 \land \cdots \land \alpha_m)$ be as in Proposition 7.11. Then each α_i is of the form $x_j = y$ or $x_j \neq y$ for some j. If there exists an i such that $\alpha_i = (x_j = y)$ then

$$T \vDash \exists y (\alpha_1 \land \dots \land \alpha_m) \leftrightarrow (\alpha_1 \land \dots \land \alpha_m)_y^{x_j}$$

If $\alpha_i = (x_j \neq y)$ for all i and some j then it is always the case that $T \models \exists y (\alpha_1 \land \dots \land \alpha_m)$ since any model of T is necessarily infinite. Simply pick $\varphi(x_1, \dots, x_k)$ to be any tautology, e.g., $x_1 = x_1$.

Completeness is trivial: the \mathfrak{L} -structure consisting of a single point can be embedded in any model of T. Since T has QE it follows directly from Proposition 7.15 that T is complete.

4. (a) Show that the theory of DLO has QE and is complete.

Recall that any model of DLO is infinite. Let $\exists y (\alpha_1 \wedge \cdots \wedge \alpha_m)$ be as in Proposition 7.11. As above, if there exists an i such that $\alpha_i = (x_j = y)$ then

$$DLO \vDash \exists y(\alpha_1 \land \dots \land \alpha_m) \leftrightarrow (\alpha_1 \land \dots \land \alpha_m)_y^{x_j}$$

Hence we may assume that each α_i is of the form $x_j \neq y$, $x_j < y$, or $x_j \not\leq y$ for some $j \leq k$. But each of these is tautologically satisfied by DLO since the theory contains the axioms that there are no endpoints, i.e., $\forall x \exists y (y < x)$ and $\forall x \exists y (x < y)$.

To see that DLO is complete take the \mathfrak{L} structure consisting of two points in x, y with x < y. Then this structure can be embedded in any model of DLO by mapping x anywhere and mapping y to some element in the model greater than the image of x, which exists by the DLO axioms. It follows from Proposition 7.15 that DLO is complete.

(b) How many definable subsets of \mathbb{R}^2 are there in the model $(\mathbb{R},<)$ of DLO?

 \mathbb{R}^2 and \emptyset are definable as witnessed by x=x and $\neg(x=x)$, respectively. Furthermore, the set of $(x,y) \in \mathbb{R}^2$ satisfying the sentences x < y, y < x, x=y, and $x \neq y$, plus their negations and pairwise conjunction and disjunction (note that some of these are contradictory or redundant, e.g., $(x < y) \land (x = y)$, or $(x < y) \land (x \neq y)$).

Using the fact that DLO has QE and is complete we can show that these are the only such definable sets.