

CONTAINERS: COMPOSITIONALITY FOR TENSORS

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CREATION OF ARRAYS

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- $a : X$

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- $a : \text{Vec } n \ X$
- $a : \text{Ar } s \ X$

ARRAY TYPE

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```
data FV (X : Set) : List ℕ → Set where
  scal : X → FV X []
  nest : Vec (FV X s) n → FV X (n :: s)
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```
S = List ℕ
P = All Fin
```

```
Ar : S → Set → Set
Ar s X = P s → X
```


EXAMPLES

```
transpose : Ar s X → Ar (rev s) X  
transpose a i = a (p-rev i)
```

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```

```
map : (X → Y) → Ar s X → Ar s Y  
map f a i = f (a i)
```

```
zipWith : (X → Y → Z) → Ar s X → Ar s Y → Ar s Z  
zipWith f a b i = f (a i) (b i)
```

```
K : X → Ar s X  
K x y = x
```

EXAMPLES

```
module Matmul
  (σ : ∀ {X s} → (X → X → X) → X → Ar s X → X)
  (_⊗_ : X → Y → Z) (ε : Z) (_⊕_ : Z → Z → Z)
  where

  mm : Ar s (Ar p X)
      → Ar p (Ar q Y)
      → Ar s (Ar q Z)
  mm a b i j = σ _⊕_ ε λ k → a i k ⊗ b k j
```

CONTAINERS

```
record Con : Set1 where
  constructor _◁_
  field
    S : Set
    P : S → Set
```

```
[[_]] : Con → Set → Set
[[ A ◁ B ]] X =  $\Sigma [ a \in A ] (B\ a \rightarrow X)$ 
```

CONTAINER OPERATIONS

$\emptyset : \text{Con}; \quad \mathbb{1} : \text{Con}$
 $\emptyset = \perp \triangleleft \lambda (); \quad \mathbb{1} = \top \triangleleft \lambda _ \rightarrow \top$

$_ \oplus _ : \text{Con} \rightarrow \text{Con} \rightarrow \text{Con}$
 $(A \triangleleft B) \oplus (C \triangleleft D) = (A \cup C)$
 $\triangleleft \lambda \{ (\text{inj}_1 a) \rightarrow B a;$
 $\quad (\text{inj}_2 c) \rightarrow D c \}$

$_ \times^c _ : \text{Con} \rightarrow \text{Con} \rightarrow \text{Con}$
 $(A \triangleleft B) \times^c (C \triangleleft D) = A \times C$
 $\triangleleft \lambda (a, c) \rightarrow B a \cup D c$

$_ \otimes _ : \text{Con} \rightarrow \text{Con} \rightarrow \text{Con}$
 $(A \triangleleft B) \otimes (C \triangleleft D) = A \times C$
 $\triangleleft \lambda (a, c) \rightarrow B a \times D c$

CONTAINER MORPHISMS

```
record Mor (A B : Con) : Set where
  constructor _◁_
  field
    MS : A .S → B .S
    MP : ∀ {a} → B .P (MS a) → A .P a
```

ARRAYS AS CONTAINERS

$\Pi : (A : \text{Set}) \rightarrow (A \rightarrow \text{Set}) \rightarrow \text{Set}$

$\Pi A B = (a : A) \rightarrow B a$

$\otimes : \text{Con} \rightarrow \text{Con} \rightarrow \text{Con}$

$\otimes (A \triangleleft B) (C \triangleleft D) = \llbracket A \triangleleft B \rrbracket C$
 $\triangleleft \lambda (a, f) \rightarrow \Pi (B a) (D \circ f)$

-- (2, f:[5, 6]) < ((i : Fin 2) → Fin (f i))

Ar = $\llbracket \otimes (\mathbb{N} \triangleleft \text{Fin}) (\mathbb{N} \triangleleft \text{Fin}) \rrbracket$

GENERALISATION

Let us use container $(A \triangleleft I)$ as a shape, where:

1. A corresponds to the set of axes;
2. I assigns to a specific axis a indices $I\ a$.

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```
Mat : (m n : ℕ) → Con
```

```
Mat m n = 2 ◁ λ { ff → Fin m; tt → Fin n }
```

ARRAY TYPE

$$\llbracket _ \rrbracket_a : \text{Con} \rightarrow \text{Set} \rightarrow \text{Set}$$
$$\llbracket A \triangleleft B \rrbracket_a X = \prod A B \rightarrow X$$
$$\text{Ar } s = \llbracket s \rrbracket_a$$

ARRAY TYPE

```
[[_]a : Con → Set → Set  
[[ A < B ]]a X = Π A B → X
```

```
Ar s = [[ s ]]a
```

Gives rise to:

```
record Con' : Set where  
  constructor _<'_  
  field  
    S : Cat  
    P : S ⇒ Seto p
```

In the sense that `Ar as Con'` is `Con <' Π`.

PROPERTIES

$\text{nest} : \text{Ar } (s \oplus p) X \rightarrow \text{Ar } s (\text{Ar } p X)$
 $\text{nest } a \ i \ j = a \ \lambda \ \{ \text{inj}_1 \ l \rightarrow i \ l;$
 $\qquad \qquad \qquad (\text{inj}_2 \ r) \rightarrow j \ r \}$

$\text{unnest} : \text{Ar } s (\text{Ar } p X) \rightarrow \text{Ar } (s \oplus p) X$
 $\text{unnest } a \ ij = a \ (ij \circ \text{inj}_1) \ (ij \circ \text{inj}_2)$

$\text{pair} : \text{Ar } s X \rightarrow \text{Ar } p Y \rightarrow \text{Ar } (s \oplus p) (X \times Y)$
 $\text{pair } a \ b = \text{unnest } \lambda \ i \ j \rightarrow a \ i \ , \ b \ j$

PROPERTIES

$\text{map} : (X \rightarrow Y) \rightarrow \text{Ar } s \ X \rightarrow \text{Ar } s \ Y$
 $\text{map } f \ a \ i = f \ (a \ i)$

$\text{reshape} : \text{Mor } s \ p \rightarrow \text{Ar } s \ X \rightarrow \text{Ar } p \ X$
 $\text{reshape } m \ a \ i = a \ (m \cdot \text{MP} \circ i \circ m \cdot \text{MS})$

TENSORS

$$_ \bullet _ : \text{Ar } (s \oplus p) X \rightarrow \text{Ar } (p \oplus q) X \rightarrow \text{Ar } (s \oplus q) X$$

gives rise to a category of tensors: $0b = Con$; morphisms $s \rightarrow p = \text{Ar } (s \oplus p) X$, composition is $_ \bullet _$.

$\text{Ar } (s \oplus p) X$ is a tensor of s covectors and p vectors; 0 and 1 are singletons; $\text{Ar } s X$ can be turned into a “column vector” $\text{Ar } (s \oplus 1) X$ or “row vector” $\text{Ar } (1 \oplus s) X$.

CONCLUSIONS

- New way to understand arrays as containers
- Combinators derived from compositional properties
- Interpretation of tensors
- WIP formalisation

<https://github.com/ashinkarov/2025-types>

THANK YOU!

I am hiring PhD students in Southampton
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Neil is hiring people who are interested in Types/CT/FP
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