

Towards Modular Composition of Inductive Types Using Lean Meta- programming

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Qualgebra

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Motivating Example

```
inductive T where
```

```
  Bool
  N
  Fn ( $\tau_1$   $\tau_2$ : T)
```

```
inductive Term where
```

```
  True
  False
  If (c t1 t2: Term)
  Zero
  Succ (t: Term)
  Pred (t: Term)
  V (x: Var)
  Abs (x: Var) ( $\tau$ : T) (b: Term)
  App (t1 t2: Term)
```

```
inductive Val: Term → Prop
```

```
  T: Val .True
  F: Val .False
  Z: Val .Zero
  S (v: Term): Val v → Val (.Succ v)
  A (x: Var) ( $\tau$ : T) (b: Term):
    Val (.Abs x  $\tau$  b)
```

```
def count: Term → Nat
```

```
  .True      => 1
  .False     => 1
  .If c t1 t2 => 1 + count c + count t1 + count t2
  .Zero      => 1
  .Succ t     => 1 + count t
  .Pred t     => 1 + count t
  .V         => 2
  .Abs b      => 3 + count b
  .App t1 t2 => 1 + count t1 + count t2
```

```
inductive TRel: Term → T → Prop
```

```
  TT: TRel .True .Bool
  FF: TRel .False .Bool
  If: TRel c .Bool → TRel t1  $\tau$  → TRel t2  $\tau$  →
    TRel (.If c t1 t2)  $\tau$ 
  Z: TRel .Zero .N
  S: TRel t .N → TRel (.Succ t) .N
  P: TRel t .N → TRel (.Pred t) .N
  V (x: Var) ( $\tau$ : T):  $\Gamma$  x =  $\tau$  → TRel  $\Gamma$  (.V x)  $\tau$ 
  Abs (x: Var) (b: Term) ( $\tau_1$   $\tau_2$ : T):
    TRel (augment  $\Gamma$  x  $\tau_1$ ) b  $\tau_2$  →
    TRel  $\Gamma$  (.Abs x  $\tau_1$  b) (.Fn  $\tau_1$   $\tau_2$ )
  App (t1 t2: Term) ( $\tau_1$   $\tau_2$ : T):
    TRel  $\Gamma$  t1 (.Fn  $\tau_1$   $\tau_2$ ) → TRel  $\Gamma$  t2  $\tau_1$  →
    TRel  $\Gamma$  (.App t1 t2)  $\tau_2$ 
```

Motivating Example

```
inductive T where
```

```
| Bool
| N
| Fn ( $\tau_1$   $\tau_2$ : T)
```

```
inductive Term where
```

```
| True
| False
| If (c t1 t2: Term)
| Zero
| Succ (t: Term)
| Pred (t: Term)
| V (x: Var)
| Abs (x: Var) ( $\tau$ : T) (b: Term)
| App (t1 t2: Term)
```

```
inductive Val: Term → Prop
```

```
| T: Val .True
| F: Val .False
| Z: Val .Zero
| S (v: Term): Val v → Val (.Succ v)
| A (x: Var) ( $\tau$ : T) (b: Term):
  Val (.Abs x  $\tau$  b)
```

```
def count: Term → Nat
```

```
| .True      => 1
| .False     => 1
| .If c t1 t2 => 1 + count c + count t1 + count t2
| .Zero      => 1
| .Succ t     => 1 + count t
| .Pred t     => 1 + count t
| .V         => 2
| .Abs b      => 3 + count b
| .App t1 t2 => 1 + count t1 + count t2
```

```
inductive TRel: Term → T → Prop
```

```
| TT: TRel .True .Bool
| FF: TRel .False .Bool
| If: TRel c .Bool → TRel t1  $\tau$  → TRel t2  $\tau$  →
  TRel (.If c t1 t2)  $\tau$ 
| Z: TRel .Zero .N
| S: TRel t .N → TRel (.Succ t) .N
| P: TRel t .N → TRel (.Pred t) .N
| V (x: Var) ( $\tau$ : T):  $\Gamma$  x =  $\tau$  → TRel  $\Gamma$  (.V x)  $\tau$ 
| Abs (x: Var) (b: Term) ( $\tau_1$   $\tau_2$ : T):
  TRel (augment  $\Gamma$  x  $\tau_1$ ) b  $\tau_2$  →
  TRel  $\Gamma$  (.Abs x  $\tau_1$  b) (.Fn  $\tau_1$   $\tau_2$ )
| App (t1 t2: Term) ( $\tau_1$   $\tau_2$ : T):
  TRel  $\Gamma$  t1 (.Fn  $\tau_1$   $\tau_2$ ) → TRel  $\Gamma$  t2  $\tau_1$  →
  TRel  $\Gamma$  (.App t1 t2)  $\tau_2$ 
```

Motivating Example

```
inductive T where
```

```
| Bool
| N
| Fn (τ1 τ2: T)
```

```
inductive Term where
```

```
| True
| False
| If (c t1 t2: Term)
| Zero
| Succ (t: Term)
| Pred (t: Term)
| V (x: Var)
| Abs (x: Var) (τ: T) (b: Term)
| App (t1 t2: Term)
```

```
inductive Val: Term → Prop
```

```
| T: Val .True
| F: Val .False
| Z: Val .Zero
| S (v: Term): Val v → Val (.Succ v)
| A (x: Var) (τ: T) (b: Term):
  Val (.Abs x τ b)
```

```
def count: Term → Nat
```

```
| .True      => 1
| .False     => 1
| .If c t1 t2 => 1 + count c + count t1 + count t2
| .Zero      => 1
| .Succ t     => 1 + count t
```

$$\text{scattering} \propto \frac{1}{\text{modularity} \wedge \text{reusability}}$$

```
| t b
| t t1 + count t2
```

```
→ Prop
```

```
t1 τ → TRel t2 τ →
```

```
Z: TRel .Zero .N
S: TRel t .N → TRel (.Succ t) .N
P: TRel t .N → TRel (.Pred t) .N
V (x: Var) (τ: T): Γ x = τ → TRel Γ (.V x) τ
Abs (x: Var) (b: Term) (τ1 τ2: T):
  TRel (augment Γ x τ1) b τ2 →
  TRel Γ (.Abs x τ1 b) (.Fn τ1 τ2)
App (t1 t2: Term) (τ1 τ2: T):
  TRel Γ t1 (.Fn τ1 τ2) → TRel Γ t2 τ1 →
  TRel Γ (.App t1 t2) τ2
```

Boolean Module

```
namespace Boolean

inductive T where
| Bool

inductive Term where
| True
| False
| If (c t1 t2: Term)

def countNodes: Term → Nat
| .True => 1
| .False => 1
| .If c t1 t2 => 1 + countNodes c + countNodes t1 + countNodes t2

inductive Val: Term → Prop
| T: Val .True
| F: Val .False

inductive TRel: Term → T → Prop
| TT: TRel .True .Bool
| FF: TRel .False .Bool
| If: TRel c .Bool → TRel t1 τ → TRel t2 τ → TRel (.If c t1 t2) τ

end Boolean
```


Nat Module

```
namespace Nat

inductive T where
| N

inductive Term where
| Zero
| Succ (t: Term)
| Pred (t: Term)

def countNodes: Term → Nat
| .Zero => 1
| .Succ t => 1 + countNodes t
| .Pred t => 1 + countNodes t

inductive Val: Term → Prop
| Z: Val .Zero
| S (v: Term): Val v → Val (.Succ v)

inductive TRel: Term → T → Prop where
| Z: TRel .Zero .N
| S: TRel t .N → TRel (.Succ t) .N
| P: TRel t .N → TRel (.Pred t) .N

end Nat
```

STLC Module

```
namespace STLC
```

```
inductive T: Type  
| Fn ( $\tau_1$   $\tau_2$ : T)
```

```
abbrev Var := String
```

```
abbrev Context := Var  $\rightarrow$  T
```

```
def augment ( $\Gamma$ : Context) (x: Var) ( $\tau$ : T): Context :=  $\lambda v \mapsto$  if  $v=x$  then  $\tau$  else  $\Gamma$  v
```

```
inductive Term where  
| V (x: Var)  
| Abs (x: Var) ( $\tau$ : T) (b: Term)  
| App ( $t_1$   $t_2$ : Term)
```

```
def countNodes: Term  $\rightarrow$  Nat  
| .V  $\Rightarrow$  2  
| .Abs  $\tau$  b  $\Rightarrow$  3 + countNodes b  
| .App  $\tau_1$   $\tau_2$   $\Rightarrow$  1 + countNodes  $\tau_1$  + countNodes  $\tau_2$ 
```

```
inductive Val: Term  $\rightarrow$  Prop  
| A (x: Var) ( $\tau$ : T) (b: Term): Val (.Abs x  $\tau$  b)
```

```
inductive TRel: Context  $\rightarrow$  Term  $\rightarrow$  T  $\rightarrow$  Prop where  
| V (x: Var) ( $\tau$ : T):  $\Gamma$  x =  $\tau$   $\rightarrow$  TRel  $\Gamma$  (.V x)  $\tau$   
| Abs (x: Var) (b: Term) ( $\tau_1$   $\tau_2$ : T):  
  TRel (augment  $\Gamma$  x  $\tau_1$ ) b  $\tau_2$   $\rightarrow$  TRel  $\Gamma$  (.Abs x  $\tau_1$  b) (.Fn  $\tau_1$   $\tau_2$ )  
| App ( $t_1$   $t_2$ : Term) ( $\tau_1$   $\tau_2$ : T):  
  TRel  $\Gamma$   $t_1$  (.Fn  $\tau_1$   $\tau_2$ )  $\rightarrow$  TRel  $\Gamma$   $t_2$   $\tau_1$   $\rightarrow$  TRel  $\Gamma$  (.App  $t_1$   $t_2$ )  $\tau_2$ 
```

```
end STLC
```

Inductive Type Composition

```
Namespace Boolean
inductive T where
| Bool
```


```
end Boolean
```


```
namespace Nat
inductive T where
| N
```

```
end Nat
```

```
namespace STLC
inductive T: Type
| Fn ( $\tau_1$   $\tau_2$ : T)
```

```
end STLC
```


`inductive T := Boolean.T | Nat.T | STLC.T`


`inductive T where`

Bool
N
Fn (τ_1 τ_2: STLC.T)
Fn (τ_1 τ_2 : T)

Composition and Extension

```
namespace Boolean
```

```
inductive Term where  
| True  
| False  
| If (c t1 t2: Term)
```

```
end Boolean
```

```
namespace Nat
```

```
...  
inductive Term where  
| Zero  
| Succ {t: Term}  
| Pred {t: Term}
```

```
end Nat
```

```
namespace STLC
```

```
inductive Term where  
| V {x: Var}  
| Abs {x: Var} (τ: T) (b: Term)  
| App (t1 t2: Term)
```

```
end STLC
```

Diagram illustrating the composition of namespaces:

```
inductive Term := Boolean.Term | Nat.Term | STLC.Term  
| isZero (t: Term)
```

Arrows indicate that the `Boolean`, `Nat`, and `STLC` namespaces are composed into the `Term` inductive type.

crosscuts Boolean and Nat

```
inductive Term where  
| True  
| False  
| If (c t1 t2: Term)  
| Zero  
| Succ {t: Term}  
| Pred {t: Term}  
| V {x: Var}  
| Abs {x: Var} (τ: T) (b: Term)  
| App (t1 t2: Term)  
| isZero (t: Term)
```

Dependencies

```
Boolean.TRel: Boolean.Term → Boolean.T → Prop
Nat.TRel: Nat.Term → Nat.T → Prop
STLC.TRel: STLC.Term → STLC.T → Prop
```

```
inductive TRel: Context → Term → T → Prop := Boolean.TRel | Nat.TRel | STLC.TRel
| iz: TRel Γ t T.N → TRel Γ (.isZero t) T.Bool
```

```
T = Boolean.T | Nat.T + STLC.T
Term = Boolean.Term | Nat.Term | STLC.Term
      | isZero ...
```

Subtyping and Coercion

```
inductive T := Boolean.T | Nat.T | STLC.T
```

```
Boolean.T <: T  
Nat.T <: T  
STLC.T <: T
```

```
instance: Coe Boolean.T T where  
  coe := λ x ↦ match x with  
    | Boolean.T.Bool => T.Bool
```

```
instance: Coe Nat.T T where  
  coe := λ x ↦ match x with  
    | Nat.T.N => T.N
```

```
instance: Coe STLC.T T where  
  coe := λ x ↦ match x with  
    | STLC.T.Fn τ1 τ2 => T.Fn τ1 τ2
```

```
def τ: T := Boolean.T.Bool
```



recursive coercion

Subtyping and Dependent Coercion

```
inductive T := Boolean.T | Nat.T | STLC.T
```

```
Boolean.T <: T
```

```
Nat.T <: T
```

```
STLC.T <: T
```

```
instance : CoeDep T (T.Bool) Boolean.T where coe := Boolean.T.Bool
```

```
instance : CoeDep T (T.N) Nat.T where coe := Nat.T.N
```

```
instance (a : STLC.T) (b : STLC.T) : CoeDep T (T.Fn a b) STLC.T where coe := STLC.T.Fn a b
```

```
def t: Boolean.T := T.Bool
```



```
def b := T.Bool
```

```
def s: Boolean.T := b
```



Multiplexing Functions

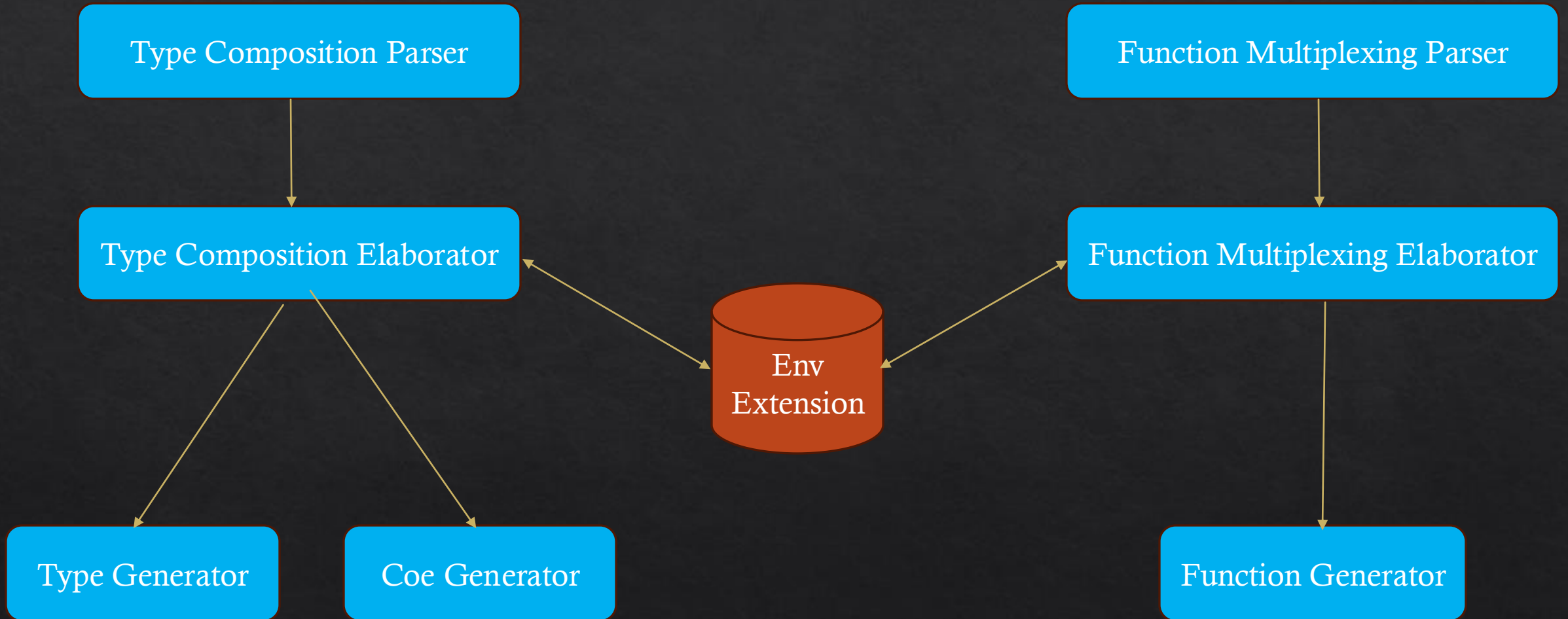
```
fn countNodes := Boolean.countNodes |+ Nat.countNodes |+ STLC.countNodes  
| isZero t => 1 + countNodes t
```

Assumption: same pattern-matching
structure

Adjusting function calls (including
recursive ones)

```
def countNodes: Term → Nat  
| .True => 1  
| .False => 1  
| .If c t1 t2 => 1 + countNodes c + countNodes t1 + countNodes t2  
| .Zero => 1  
| .Succ t => 1 + countNodes t  
| .Pred t => 1 + countNodes t  
| .V _ => 2  
| .Abs b => 3 + countNodes b  
| .App t1 t2 => 1 + countNodes t1 + countNodes t2  
| .isZero t => 1 + countNodes t
```


Architecture



Limitations

- Full support of higher-order types, indexed-types, dependent types
- Assumptions on multiplexed functions
- Mutual recursion
- Composing feature modules instead of individual types/functions
- (Partial?) composition of theorems and proof objects
- Function reuse instead of rewriting
 - Recursion? Modifying fixpoint operators?
 - Cost of function calls? Inlining?

Thank You

Questions

<https://github.com/qualgebra/LeanToolkit/tree/TYPES2025>