# A Ghost Sort for Proof-Relevant yet Erased Data in Rocq and MetaRocq

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## It Sounds Like a Bad Advertisement

## Are you:

- ► Often using Rocq's extraction?
- ► Annoyed by the slowdown of unnecessary arguments?

## It Sounds Like a Bad Advertisement

## Are you:

- Often using Rocq's extraction?
- Annoyed by the slowdown of unnecessary arguments?

If so, then usebuy our new ghost sort!

- ► Takes care of 99.9% of the inconveniences of extraction!
- At the cheap price of a single annotation!

Otherwise, you may still find the talk interesting, I swear (spoiler: there will be drama)!

## **Ghost Sort In Action**

```
Fixpoint lookup {A} (l:list A) (f:fin (length l)): A := match l with
 | nil ⇒ False_rect A f
 | consal \Rightarrow match f with
   | fin0 \Rightarrow a
   | finS f' \Rightarrow lookup l f'
   end
 end.
let rec lookup l f =
 match 1 with
 |Nil \rightarrow assert false
 | Cons (a, 10) \rightarrow (match f with
             | FinO \_ 
ightarrow  a
             | FinS(_, f') \rightarrow lookup l0 f')
```

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                                            Only one annotation needed!
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                                                     | Nil \rightarrow assert fals
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                                                     | Cons (a, 10) \rightarrow (match + with
              |FinO 📈
                                                                  | Fin0 \rightarrow a |
              | \operatorname{FinS}(\sqrt{f'}) \rightarrow \operatorname{lookup} 10 f')
                                                                  | FinS f' \rightarrow lookup l0 f' )
```

## Isn't It Just $\mathbb{N}$ ?

#### Before:

```
Inductive fin: \mathbb{N} \to Set := | fin0 \{n\} : fin (S n) |
| finS \{n\} : fin n \to fin (S n).
```

```
type fin =
| FinO of nat
| FinS of nat * fin
```

## Isn't It Just $\mathbb{N}$ ?

#### Before:

```
\begin{array}{lll} \text{Inductive fin: } \mathbb{N} \to \text{Set:=} & \text{type fin =} \\ | \text{finO} \{n\} \colon \text{fin} (S \, n) & | \text{FinO of } \\ | \text{finS} \{n\} \colon \text{fin n} \to \text{fin} (S \, n). & | \text{FinS of } \\ \end{array}
```

#### After:

```
 \begin{array}{lll} \mbox{Inductive fin}: \mathbb{N}_{\mbox{\scriptsize $0$}} \{\mbox{Ghost}\} \rightarrow \mbox{Set} := & \mbox{type fin} = \\ |\mbox{finO}_{\mbox{\scriptsize $0$}} \{\mbox{n}\}: \mbox{fin}_{\mbox{\scriptsize $(S$ n)$}}. & |\mbox{FinO}_{\mbox{\scriptsize $(S$ n)$}}. & |\mbox{FinS of fin}_{\mbox{\scriptsize $(S$ n)$}}. \\ \end{array}
```

## Isn't It Just $\mathbb{N}$ ?

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\begin{array}{lll} \text{Inductive fin: } \mathbb{N} \to \text{Set:=} & \text{type fin =} \\ | \text{finO} \{ \text{n} \} \colon \text{fin} \, (\text{S n}) & | \text{FinO of } \\ | \text{finS} \{ \text{n} \} \colon \text{fin n} \to \text{fin (S n)}. & | \text{FinS of } \\ \end{array}
```

#### After:

```
Inductive fin: \mathbb{N}_{0}\{Ghost\} \rightarrow Set:= type fin = nat? | fin0 \{n\}: fin (S n) | finS \{n\}: fin n \rightarrow fin (S n).
```

# No Cheating Allowed!

```
\label{eq:constraints} \begin{array}{lll} \text{Inductive vec (A: Type): } \mathbb{N} \to \text{Type:=} & \text{type 'a vec =} \\ | \, \text{VNil: vec A 0} & | \, \text{VNil} \\ | \, \text{VCons: } \forall \, (n:\mathbb{N}), \, \text{A} \to \text{vec A n} \to \text{vec A (S n)}. & | \, \text{VCons of nat *'a *'a vec} \\ \\ \text{Definition length } \{\text{A n}\} \, (\_: \text{vec A n}): \mathbb{N} := \text{n.} & \text{let length n} \, \_ = \text{n} \\ \end{array}
```

# No Cheating Allowed!

```
Inductive vec (A: Type): \mathbb{N} \to \mathsf{Type}:=
                                                                                                                                                                                                                                                                              type 'a vec =
 | VNil : vec A 0
                                                                                                                                                                                                                                                                               l vnil
| VCons : \forall (n : \mathbb{N}), A \rightarrow vec A n
                                                                                                                                                                                                                                                                               Vcons of nat * 'a * 'a vec
                                                                                                                                                                                                                                                                                    Horrible extractions
Definition length \{A \ n\} ( : vec A \ n): \mathbb{N} := n.
                                                                                                                                                                                                                                                                              let length n = n
 Inductive vec (A: Type): \mathbb{N} \cap \{Ghost\} \rightarrow \mathsf{Type} :=
                                                                                                                                                                                                                                                                                 type 'a vec =
 LVNil: vec A 0
                                                                                                                                                                                                                                                                                  I VNil
 | VCons : \forall (n : \mathbb{N}_{0} \{Ghost\}), A \rightarrow vec A n \rightarrow vec A (S n). | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a * 'a vec A (S n) | VCons of 'a vec A (S
 Fixpoint length \{A \ n\} \ (v : vec \ A \ n) : \mathbb{N} :=
                                                                                                                                                                                                                                                                                 let rec length = function
     match v with
                                                                                                                                                                                                                                                                                  |VNil \rightarrow 0|
                                                                                                                                                                                                                                                                                   | VCons(,xs) \rightarrow S(length xs)|
      |VNil \Rightarrow 0|
      | VCons xs \Rightarrow S (length xs)
      end.
```

# No Cheating Allowed!

```
Inductive vec (A: Type): \mathbb{N} \to \mathsf{Type}:=
                                                                                                                                                                                                                                                                                                                                                                                                                                                       type 'a vec =
  | VNil : vec A 0
                                                                                                                                                                                                                                                                                                                                                                                                                                                         l vnil
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                                                                                                                                                                                                                                                                                                                                                                                                                                                          Vcons of nat * 'a * 'a vec
                                                                                                                                                                                                                                                                                                                                                                                                                                                                  Horrible extraction
Definition length \{A \ n\} ( : vec A \ n): \mathbb{N} := n.
                                                                                                                                                                                                                                                                                                                                                                                                                                                       let length n = n
  Inductive vec (A: Type): \mathbb{N} \cap \{Ghost\} \rightarrow Type :=
                                                                                                                                                                                                                                                                                                                                                                                                                                                                  type 'a vec = 'a list?
  LVNil: vec A 0
  | VCons : \forall (n : \mathbb{N} \cap \{Ghost\}), A \rightarrow vec A n \rightarrow vec A (S n).  let length = List.length?
  Fixpoint length \{A \cap A \cap v \in A \cap v \in
         match v with
          |VNil \Rightarrow 0|
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          end.
```

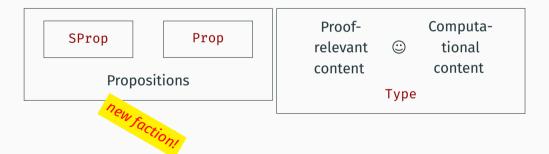
Previously, in La Villa des Sorts...

Strict Acc
propositions eq
Prop

Previously, in La Villa des Sorts...



Previously, in La Villa des Sorts...



Now...



Now...



Now...



# You're Stealing His Job!!

Perks	s = Prop	s = Ghost
Erasure at extraction	✓	✓
Acc accomodation	✓	✓
Consistency	✓	$[Ghost_i] = Type_i \checkmark$
true@{s}≠false@{s}	×	$Ghost_{\mathtt{i}}:Ghost_{\mathtt{i+1}}\checkmark$
Indices erasure	×	✓
Impredicativity	✓	×

## The Ghost Sort

▲ Heavy use of notions seen tuesday ahead!

## Design #1:

- ▶ ⊥ລ{Ghost} elimination,
- redevelopment of logic (or wait for a sort poly. one).

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Design #2:

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## The Ghost Sort

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Design #1:

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Design #3:

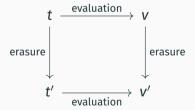
- ► Ghost ~→ Prop,
- ► Prop's logic.

## **Development Plan**

## The plan:

- 1. (More) Examples with designs.
- 2. (Current) MetaRocq Formalization.
- 3. Rocq Implementation.

Main MetaRocq goal — commutation of:



# Ghost **vs.** QTT

	QTT	Ghost
Addition	#(uses of a term)	Ground sort
Erasure	@0 t : A : Set	t:A:Ghost
Type-checking Changes	All rules	Eliminator rules
Erased Accessibility	No	Yes

## Ghost vs. Erasure Modality

```
record Erased (@0 A: Set a): Set a where
  constructor [_]
  field
    @0 erased: A
```

## Key properties:

- ► Monad.
- ► Stable A := Erased A → A.
- ▶ DNE or dec. for A: Stable A.
- ▶ Revival: Erased A  $\simeq$  A.
- ► Few types can be revived.

# Ghost vs. Erasure Modality

```
record Erased (@O A: Set a): Set a where
  constructor[]
  field
   ao erased: A
In Roca:
Inductive Boxa(u) (A: Typea(u)): Ghosta(u):=
  box : A \rightarrow Box A.
Inductive Unbox@{u} (A: Ghost@{u}): Type@{u}:=
 unbox : A \rightarrow Unbox A.
Definition Erased∂{u} (A: Type∂{u}) :=
 Unbox (Box A).
Notation "[x]" := (unbox (box x)).
```

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- ► Monad.
- ▶ Stable A :=  $Erased A \rightarrow A$ .
- ▶ DNE or dec. for A: Stable A.
- ▶ Revival: Erased A  $\simeq$  A.
- ► Few types can be revived.

Ghost satisfy key properties with every design.

(But machinery level differs...)

#### Conclusion

## We have presented:

- ► Development plan.
- Current status.
- ► Comparison with existing systems.

#### Near future:

- ► Settle for a design.
- ► Preferred: Ghost ~> Prop.
- ► Allows: Ghost<sub>i</sub>≈ Type<sub>i</sub> + computational sort.
- ► Focus on the formalization.

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Thanks for your attention!