# Type-safe Bidirectional Channels in Idris 2

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**TYPES** 



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**Typed Communications** 

Stateful Protocols

### **Table of Contents**

**Protocols** 

**Protocol-respecting Programs** 

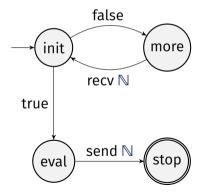
Typed Communications

Stateful Protocols

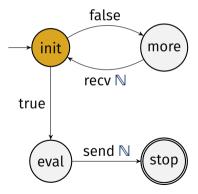
## **Session types**

$$p, q, \ldots := x | \mu x. p | \nu x. p$$
  
| !A.  $p$  | ?A.  $p$  | end  
|  $p \& q | p \oplus q$ 

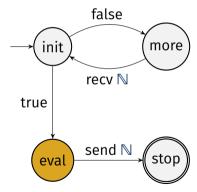
$$\nu i. ((!\mathbb{N}. end) \& (?\mathbb{N}. i))$$



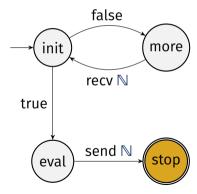




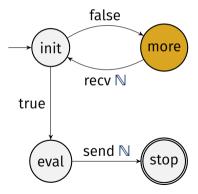
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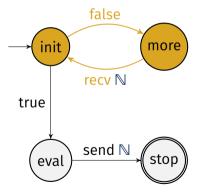
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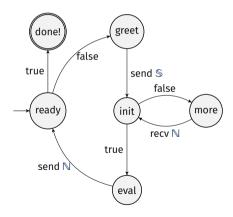
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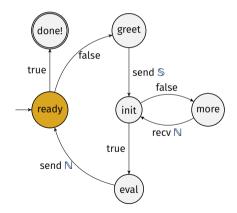
$$\nu$$
i. ((! $\mathbb{N}$ . end)&( $?\mathbb{N}$ .  $i$ ))



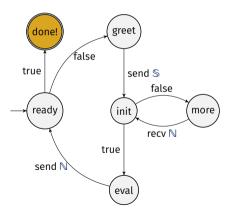
$$\nu r. (\text{end \&}(!S. \nu i. ((!N. r) \& (?N. i))))$$



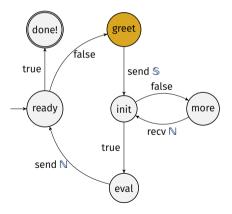
vr. (end & (!S. vi. ((!N. r) & (?N. i))))



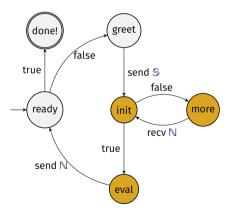
$$\nu r. (\text{end } \& (!S. \nu i. ((!N. r) \& (?N. i))))$$



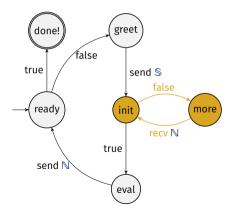
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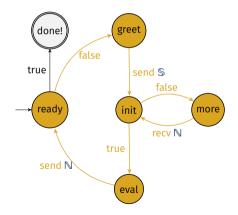
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$$\nu$$
r. (end &(!S.  $\nu$ i. ((!N.  $r$ ) &(?N.  $i$ ))))



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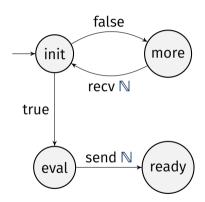
Protocols

**Protocol-respecting Programs** 

Typed Communication:

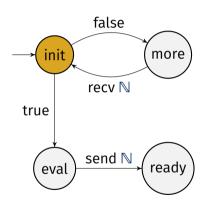
Stateful Protocols

```
adder acc ch = do
  ch <- unroll {nm = Nu "adder"} ch</pre>
  b # ch <- offer ch
  if b then do
      ch <- send ch acc
      rec {nm = Nu "ready"} ch
    else do
      (n # ch) <- recv ch
      ch <- rec {nm = Nu "adder"} ch
      adder (acc + n) ch
```

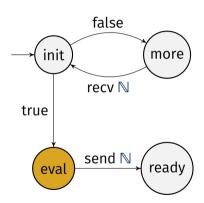


#8

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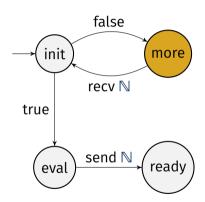


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Protocol-respecting Programs

**Typed Communications** 

Stateful Protocols

## **Idris' Typed Channels**

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data Channel : Type -> Type where [external]
channelGet : HasIO io => Channel a -> io a
channelPut : HasIO io => Channel a -> a -> io ()
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- ▶ Pick one type and stick with it
- No protocol

### **Messages Type**

#### A big (and cheap) union of types

```
data UnionT : (elt : a -> Type) -> (ts : List a) -> Type where
    Element : (k : Nat) -> (0 _ : AtIndex t ts k) -> elt t -> UnionT elt ts
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#### Listing all of the sent and received types

```
SendTypes : Session nms m -> List Type
RecvTypes : Session nms m -> List Type
```

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Communications provoke destructive updates of the channels

- ► (Ab)using linearity to encode uniqueness
- Not unfolding fixpoints by substitution
- Recording where we are in the state machine

## **Syntax of Session Types**

```
data Kind : Type where
  Mu, Nu : Name -> Kind

data Session : List Kind -> Norm -> Type where
  Fix : (kd : Kind) -> Session (kd :: nms) Head -> Session nms Head
  Rec : {kd : Kind} -> Focus kd nms -> Session nms Expr
```

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data Kind : Type where
 Mu. Nu : Name -> Kind
data Session : List Kind -> Norm -> Type where
 Fix : (kd : Kind) -> Session (kd :: nms) Head -> Session nms Head
 Rec : {kd : Kind} -> Focus kd nms -> Session nms Expr
 Send: (ty: Type) -> Session nms n -> Session nms Head
 Recv : (ty : Type) -> Session nms n -> Session nms Head
 End · Session nms Head
 Offer: Session nms m1 -> Session nms m2 -> Session nms Head
 Select: Session nms m1 -> Session nms m2 -> Session nms Head
```

#### **Environment Stack**

To make sense of nested fixpoints, we need a telescopic environment

```
data Env : List Kind -> Type where
Nil : Env []
(::) : Session (nm :: nms) Head -> Env nms -> Env (nm :: nms)
```

#### **Channels**

```
record Channel (s : Session nms k) (e : Env nms) where
  constructor MkChannel
  {sendStep : Nat}
  {recvStep : Nat}
  {0 ogNorm : Norm}
  {0 ogKinds : Kinds}
  {0 ogSession : Session ogKinds ogNorm}
  context : Context recvStep sendStep ogSession s e

  sendChan : Threads.Channel (Union (Bool :: SendTypes ogSession))
  recvChan : Threads.Channel (Union (Bool :: RecvTypes ogSession))
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```
send : Channel (Send ty s) e -@ (ty -> IO1 (Channel s e))
recv : Channel (Recv ty s) e -@ IO1 (Res ty (\ _ => Channel s e))
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recv : Channel (Recv ty s) e -@ IO1 (Res ty (\ _ => Channel s e))

offer : Channel (Offer s1 s2) e -@
   IO1 (Res Bool (\ b => ifThenElse b (Channel s1 e) (Channel s2 e)))
select : Channel (Select s1 s2) e -@
   ((b : Bool) -> IO1 (ifThenElse b (Channel s1 e) (Channel s2 e)))
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unroll : Channel (Fix nm s) e -@ IO1 (Channel s (s :: e))
roll : Channel s (s :: e) -@ IO1 (Channel (Fix nm s) e)
```

```
send : Channel (Send ty s) e -@ (ty -> IO1 (Channel s e))
recv : Channel (Recv tv s) e -@ IO1 (Res tv (\ => Channel s e))
offer: Channel (Offer s1 s2) e -@
  IO1 (Res Bool (\ b => ifThenElse b (Channel s1 e) (Channel s2 e)))
select : Channel (Select s1 s2) e -@
  ((b : Bool) -> IO1 (ifThenElse b (Channel s1 e) (Channel s2 e)))
unroll: Channel (Fix nm s) e -@ IO1 (Channel s (s :: e))
roll : Channel s (s :: e) -@ IO1 (Channel (Fix nm s) e)
rec : {pos : _} -> Channel (Rec pos) e -@
  IO1 (Channel (Fix nm (SessionAt pos e)) (EnvAt pos e))
```

### What's next?

- ► Small runtime overhead
- Ergonomics
- ► Totality Checking