

**CSC 226 SUMMER 2020**  
**ALGORITHMS AND DATA STRUCTURES II**  
**ASSIGNMENT 1**  
**UNIVERSITY OF VICTORIA**

1. An  $n$ -degree *polynomial*  $p(x)$  is an equation of the form

$$p(x) = \sum_{i=0}^n a_i x^i$$

where  $x$  is a real number and each  $a_i$  is a real constant, with  $a_n \neq 0$ .

- a. Describe a simple  $O(n^2)$ -time method for computing  $p(x)$  for a particular value of  $x$ . Justify the runtime.
- b. Consider now the nested form of  $p(x)$ , written

$$p(x) = a_0 + x \left( a_1 + x \left( a_2 + x \left( a_3 + \cdots + x \left( a_{n-1} + x a_n \right) \cdots \right) \right) \right).$$

Using the big-Oh notation, characterize the number of multiplications and additions this method of evaluation uses.

2. Suppose that if the linear selection algorithm used subsequences of 3 instead of 5 or 7, then show that the worst-case runtime would no longer be  $O(n)$ . You may follow the lecture slide examples from lecture 4 when developing your recurrence equation  $T(n)$  for this version of the algorithm. That is, you can use upper bounds to get rid of the ceiling notation and assume an in-place implementation which costs nothing to separate into subsequences. [Note: For a subsequence of 3 elements, it takes at most 3 comparisons to sort them.]
3. Write an in-place version of the linear selection algorithm, using subarrays of size 7 for the median of medians selection. Use the `linearSelect( $S, k$ )` algorithm in slide 4 of lecture 4 and the `inPlaceQuickSort( $S, a, b$ )` algorithm in the Goodrich and Tamassia textbook (page 255) as your guideline for writing this algorithm. The algorithm should have the following heading,

**Algorithm** `inPlaceLinearSelect( $S, a, b, k$ )`

**Input:** An array,  $S$ , of distinct elements; integers  $a$  and  $b$  such  $a \leq b$ ; and integer  $k \in [a + 1, b + 1]$

**Output:** The  $k$ th smallest element of  $S$

4. Consider the experiment of tossing a fair coin four times,
  - a. What is sample space associated with this experiment? Label a heads  $H$  and a tails  $T$ .
  - b. Let  $A$  be the event that heads are flipped first. What is  $\Pr(A)$ ?
  - c. Let  $B$  be the event that exactly two heads and two tails are flipped (in any order). What is  $\Pr(B)$ ?
  - d. Let  $A$  and  $B$  be as defined above, what is  $\Pr(A \cap B)$ ?
  - e. Let  $X$  be the number of heads flipped. What is the expected value of  $X$ ?

5. Suppose you would like to sort  $n$  music files using a comparison-based sorting algorithm (i.e. no bucket sort), but you only have an old, unreliable computer, which you have nicknamed “Rustbucket”. Every time Rustbucket compares two music files,  $x$  and  $y$ , there is an independent 50-50 chance that it has an internal disk fault and returns the value -1, instead of the correct result of 1 for **true** or 0 for **false**, to the question, “ $x \leq y$ ?” Otherwise, Rustbucket correctly performs every other kind of operation (including comparisons not involving music files.) Describe an efficient algorithm that can use Rustbucket to sort  $n$  music files correctly and show that your algorithm has *expected* running time that is  $O(n \log n)$ .