

- 2 longest common subsequences of the 2 strings are: TCAGTTG and TCGATGC. There are more.

X= "TCAAAGATTAAGC"

Y= "TCGATGTCTCGTTG"

		T	C	G	A	T	G	T	C	T	C	G	T	T	G
	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
T	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1
C	0	1	2	2	2	2	2	2	2	2	2	2	2	2	2
A	0	1	2	2	3	3	3	3	3	3	3	3	3	3	3
A	0	1	2	2	3	3	3	3	3	3	3	3	3	3	3
A	0	1	2	2	3	3	3	3	3	3	3	3	3	3	3
G	0	1	2	3	3	3	4	4	4	4	4	4	4	4	4
A	0	1	2	3	4	4	4	4	4	4	4	4	4	4	4
T	0	1	2	3	4	5	5	5	5	5	5	5	5	5	5
T	0	1	2	3	4	5	5	6	6	6	6	6	6	6	6
A	0	1	2	3	4	5	5	6	6	6	6	6	6	6	6
A	0	1	2	3	4	5	5	6	6	6	6	6	6	6	6
G	0	1	2	3	4	5	6	6	6	6	7	7	7	7	7
C	0	1	2	3	4	5	6	6	7	7	7	7	7	7	7

- 2.

```
printLCS(int[][] L, String X, Y)
```

```
String s = ""
```

```
for (int n = X.length, m = Y.length; n > 0 && m > 0;) {
```

```
    while (L[n][m] == L[n-1][m] && n > 0) n--;
```

```
    while (L[n][m] == L[n][m-1] && m > 0) m--;
```

```
    if (n > 0 && m > 0) s += L[n][m];
```

```
    n--;
```

```
    m--;
```

```
}
```

```
return s.reverse()
```

- 3.

Proof by contradiction: Given $X[n] \neq Y[m]$ and $Z[k] \neq X[n]$, assume Z is not an LCS of $X[1]...X[n-1]$ and Y. We know that Z is an LCS of X and Y. For Z not to be an LCS of $X[1]...X[n-1]$ and Y, but to be an LCS of X and Y, the difference between the 2 pairs of strings, or $X[n]$, must be in Z. However, since $X[n]$ is at the end of X, if it were to appear in Z, it must be the last character of Z, or $Z[k]$, for Z to be an LCS of X

and Y. This is a contradiction, therefore Z must be an LCS of $X[1]...X[n-1]$ and Y given the above assumptions.

4.

a. [10, 8, 4, 6, 9]

b.

	0	1	2	3	4	5	6	7	8	9	10
A	1	1	1	4	1	6	1	8	1	1	11
B	0	2	0	0	2	0	2	0	9	0	0
C	0	0	0	0	5	0	0	0	0	0	0
D	0	0	0	0	0	0	7	0	0	0	0
R	0	0	3	0	0	0	0	0	0	10	0

5.

```
consecutiveBlanks(String txt, int M) {  
    for (int i = 0, j = 0; i < txt.length(); i++) { //n loops max  
        if (txt.charAt(i) == " ") { //1 time  
            j++; //1 time  
            if (j >= M) return i-(j-1); //1 time  
        }  
        else j=0; //1 time  
    }  
    return txt.length();  
}
```

This algorithm is $O(3n) = O(n)$, in every case.