

# Elliptic Curve Cryptography

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# Elliptic Curves

Weierstrass equation<sup>1</sup>:

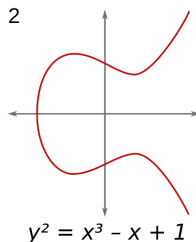
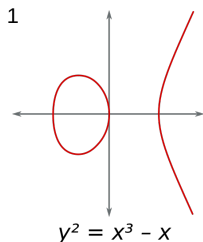
$$y^2 = x^3 + ax + b$$

where

$$\forall a, b \in \mathbb{R}$$

and

$$\Delta = -16(4a^3 + 27b^3) \neq 0$$



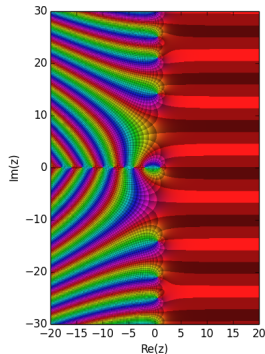
<sup>1</sup>[https://en.wikipedia.org/wiki/Elliptic\\_curve](https://en.wikipedia.org/wiki/Elliptic_curve)

## Utility of Elliptic Curves

Elliptic curves are related to:

- Riemann  $\zeta$  function<sup>2</sup>

$$\begin{aligned}\zeta(s) &= \sum_{n=1}^{\infty} n^{-s} \\ &= \frac{1}{1^s} + \frac{1}{2^s} + \frac{1}{3^s} + \dots \\ &= \frac{1}{\Gamma(s)} \int_0^{\infty} \frac{x^{s-1}}{e^x - 1} dx\end{aligned}$$



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<sup>2</sup>And consequently the Riemann hypothesis: [https://en.wikipedia.org/wiki/Riemann\\_hypothesis#Consequences\\_of\\_the\\_Riemann\\_hypothesis](https://en.wikipedia.org/wiki/Riemann_hypothesis#Consequences_of_the_Riemann_hypothesis)

# Utility of Elliptic Curves

Elliptic curves are related to:

- ▶ Fermat's last theorem

$$\forall x, y, z \in \mathbb{Z}, \quad \exists n \in \mathbb{Z} \quad \text{such that} \quad x^n + y^n = z^n$$

- ▶ Birch and Swinnerton-Dyer conjecture
- ▶ Langlands Program: Galois groups  $\leftrightarrow$  automorphic forms, representation theory, *etc.*
- ▶ Galois Theory
- ▶ Lie Theory
- ▶ *etc.*

# Groups

A group is a 2-tuple consisting of a set  $G$  and an operation  $*$  having a list of properties.

Consider the set of integers and the addition operation,  $(\mathbb{Z}, +)$ .

For any  $a, b, c \in \mathbb{Z}$ :

- ▶ Closure:  $a + b \in \mathbb{Z}$
- ▶ Associativity:  $(a + b) + c = a + (b + c)$
- ▶ Identity:  $\exists 0 \in \mathbb{Z}$  such that  $a + 0 = a$
- ▶ Inverse:  $\exists -a \in \mathbb{Z}$  such that  $a + (-a) = 0$
- ▶ (Commutativity:  $a + b = b + a$ , only for abelian groups)

# Fields

A field is a 3-tuple consisting of a set  $G$  and two operations  $\dagger, *$  having a list of properties.

Consider the set of rational numbers and the addition and multiplication operations,  $(\mathbb{Q}, +, \times)$ . For any  $a, b, c \in \mathbb{Q}$ :

- ▶ the 2-tuple  $(\mathbb{Q}, +)$  is an abelian group
- ▶ the 2-tuple  $(\mathbb{Q} \setminus \{0\}, \times)$  is an abelian group
- ▶ Distributivity:  $a \times (b + c) = a \times b + a \times c$

# Vector Spaces

A vector space is a 3-tuple of a cartesian product of identical sets and two operations.

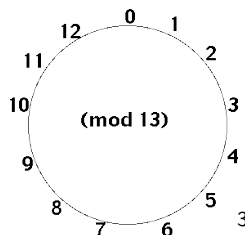


Consider the three copies of the set of integers  $\mathbb{Z} \times \mathbb{Z} \times \mathbb{Z} = \mathbb{Z}^3$ , vector addition and scalar multiplication,  $(\mathbb{Z}^3, +, \cdot)$ . For any  $\mathbf{u}, \mathbf{v}, \mathbf{w} \in \mathbb{Z}^3$  and  $a, b \in \mathbb{Z}$ :

- ▶ the 2-tuple  $(\mathbb{Z}^3, +)$  is a group
- ▶ Compatibility:  $a \cdot (b \cdot \mathbf{v}) = (ab) \cdot \mathbf{v}$
- ▶ Scalar Identity:  $1 \cdot \mathbf{v} = \mathbf{v}$
- ▶ Field Distributivity:  $(a + b) \cdot \mathbf{v} = a \cdot \mathbf{v} + b \cdot \mathbf{v}$
- ▶ Vector Distributivity:  $a \cdot (\mathbf{v} + \mathbf{u}) = a \cdot \mathbf{v} + a \cdot \mathbf{u}$

# Modular Arithmetic and Finite Fields

A finite group or a finite field is a group or a field whose set is finite.



Examples:

$$7 + 8 = 15 \equiv 2 \pmod{13}$$

$$10 \times 10 = 20 \equiv 7 \pmod{13}$$



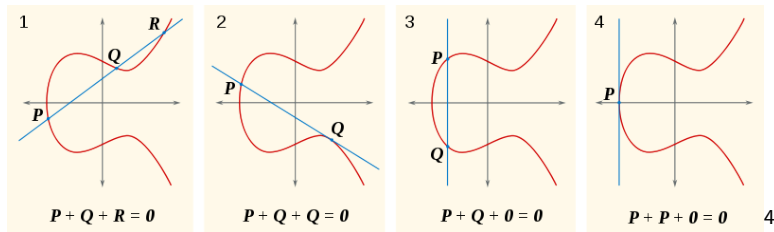
# Projective Plane



The projective plane (the set  $\mathbb{R}^2 + \{P\}$ , where  $P$  is the point at infinity):

- ▶  $\forall$  points  $a, b \in \mathbb{P}$ ,  $\exists$  exactly one line  $L \subset \mathbb{P}$  such that  $a, b \in L$
- ▶  $\forall$  lines  $K, L \in \mathbb{P}$ ,  $\exists$  exactly one point  $a \in \mathbb{P}$  such that  $a \in L$  and  $a \in K$
- ▶  $\forall a, b, c, d \in \mathbb{P}$ ,  $\exists$  no line  $L \subset \mathbb{P}$  such that more than two points  $\in L$

# Projective Plane



The projective plane is needed to do algebra over the points on an elliptic curve. The point at infinity serves as the identity element.

# Discrete Logarithm

- ▶ The discrete logarithm<sup>5</sup>:

$$n^k \bmod p \qquad k, n, p \in \mathbb{Z}$$

- ▶ The discrete logarithm problem:

$$n^k \bmod p = N \equiv m$$

Is there a polynomial time algorithm that can find  $k$  by knowing  $N$ ? Whereas finding  $N$  by knowing  $k$  is easy.

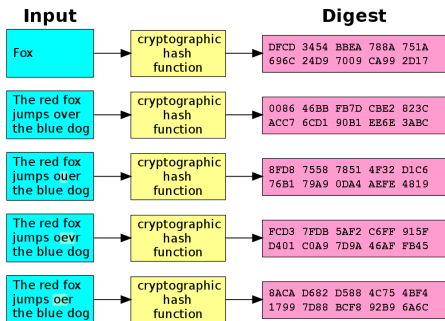
- ▶ ( $n$  is known as a generator of the group  $\mathbb{Z}_p$ )

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<sup>5</sup>Also see:

# Hashing

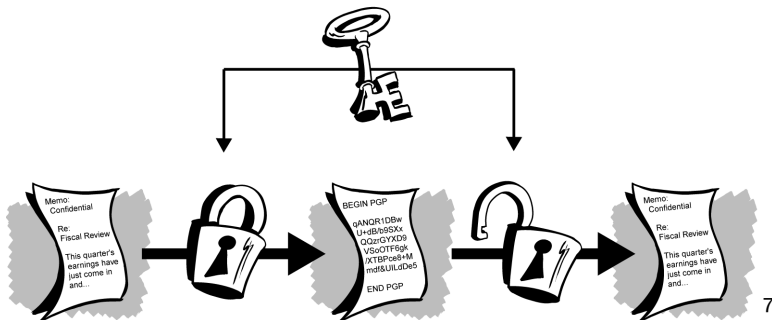
A hash is a function that maps strings of bytes of arbitrary length into a set of strings of bytes that have relatively short and identical lengths.



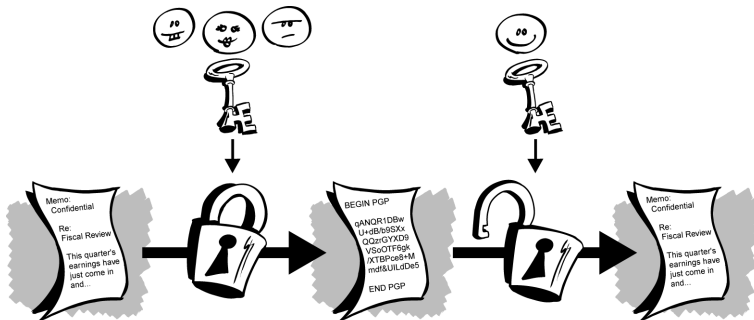
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<sup>6</sup>[https://commons.wikimedia.org/wiki/File:Cryptographic\\_Hash\\_Function.svg](https://commons.wikimedia.org/wiki/File:Cryptographic_Hash_Function.svg)

# Symmetric Keys

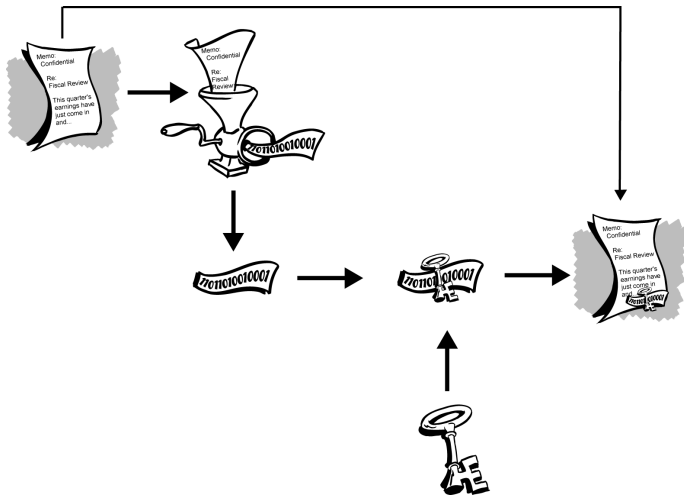


# Asymmetric Keys



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# Signing



# Libraries

- ▶ NaCl
  - ▶ NaCl: <http://nacl.cr.yp.to/>
  - ▶ tweetNaCl: <http://tweetnacl.cr.yp.to/>
  - ▶ libsodium: <https://github.com/jedisct1/libsodium>
- ▶ libnacl
  - ▶ home: <https://github.com/saltstack/libnacl>
  - ▶ docs: <https://libnacl.readthedocs.org/en/latest/>
- ▶ PyNaCl
  - ▶ home: <https://github.com/pyca/pynacl>
  - ▶ docs: <https://pynacl.readthedocs.org/en/latest/>
- ▶ pure\_pynacl: [https://github.com/jfindlay/pure\\_pynacl](https://github.com/jfindlay/pure_pynacl)