

GLIF: A Declarative Framework for Symbolic Natural Language Understanding

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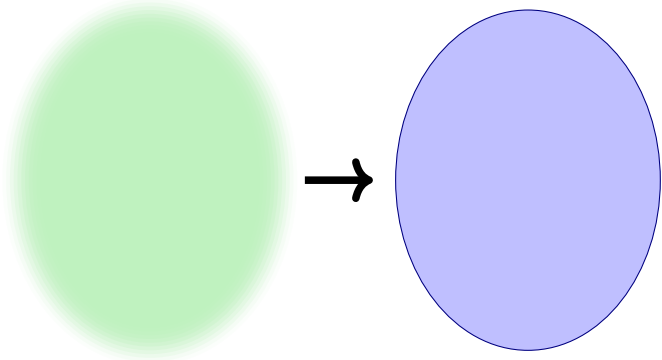
virtual event due to COVID-19

September 22, 2020

Method of Fragments

Natural Language

Logic



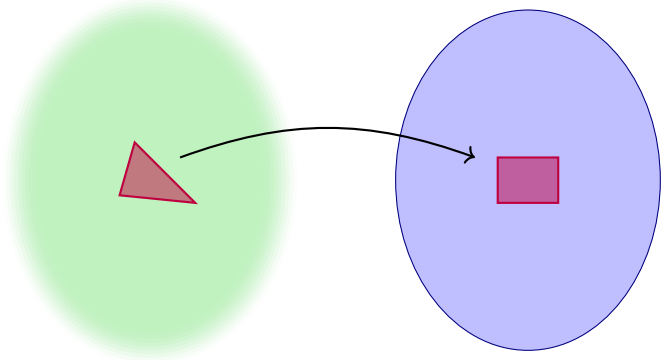
How do we get from messy language to formal logic?

Montague [Mon70]: Look at a “nice” subset and map into logic.

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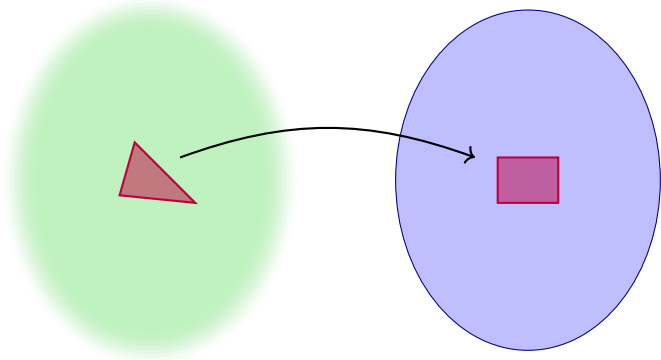
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"Ahmed paints and Berta is quiet."

$p(a) \wedge q(b)$

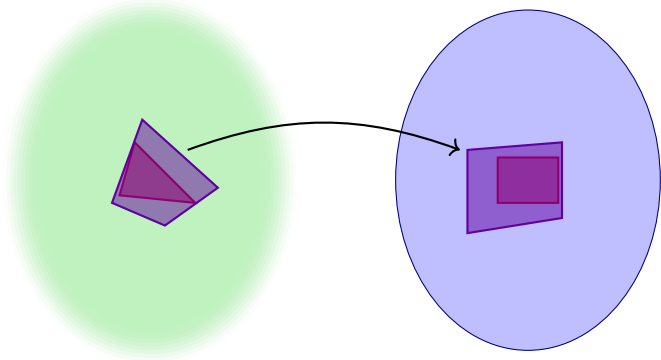
"Ahmed doesn't paint."

$\neg p(a)$

Method of Fragments

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"Every student paints and is quiet."

"Nobody paints."

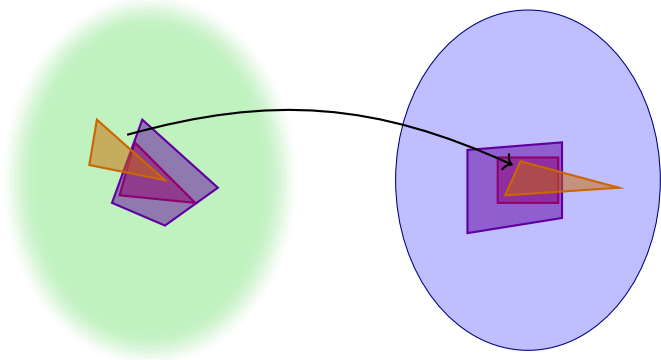
$$\forall x.s(x) \Rightarrow (p(x) \wedge q(x))$$

$$\neg \exists x.p(x)$$

Method of Fragments

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"Ahmed isn't allowed to paint."

$\neg \Diamond p(a)$

"Ahmed and Berta must paint."

$(\Box p(a)) \wedge \Box p(b)$

Method of Fragments

If we only hand-wave, we gloss over problems:

“Ahmed paints. He is quiet.” $\overset{?}{\rightsquigarrow}$ $p(a) \wedge q(a)$

Specify:

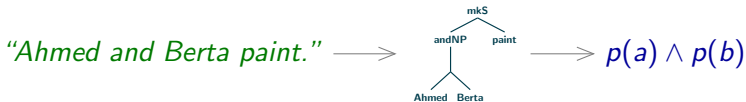
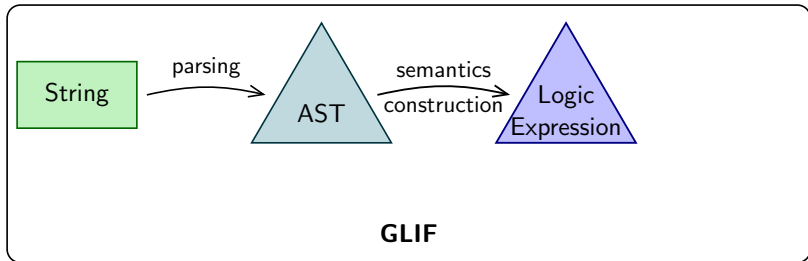
- Grammar *fixes NL subset*
- Target logic
- Semantics construction *maps parse trees to logic*

On paper [Mon74]: *difficult to scale*

- T11. If $\phi, \psi \in \mathbf{P}_t$ and ϕ, ψ translate into ϕ', ψ' respectively, then ϕ **and** ψ translates into $[\phi \wedge \psi]$, ϕ **or** ψ translates into $[\phi \vee \psi]$.
- T12. If $\gamma, \delta \in \mathbf{P}_{IV}$ and γ, δ translate into γ', δ' respectively, then γ **and** δ translates into $\hat{x}[\gamma'(x) \wedge \delta'(x)]$, γ **or** δ translates into $\hat{x}[\gamma'(x) \vee \delta'(x)]$.
- T13. If $\alpha, \beta \in \mathbf{P}_T$ and α, β translate into α', β' respectively, then α **or** β translates into $\hat{P}[\alpha'(P) \vee \beta'(P)]$.

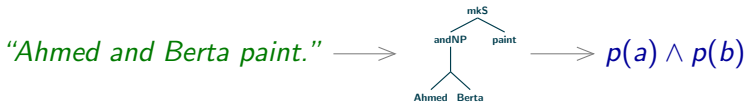
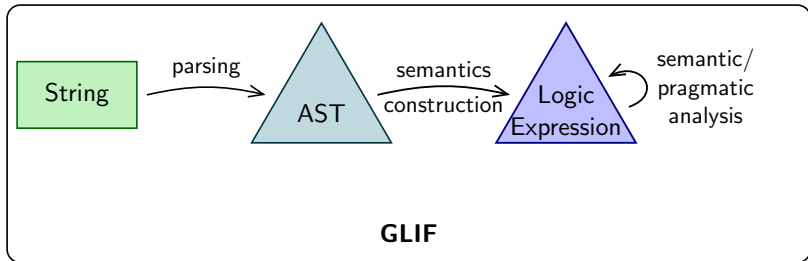
GLIF: Grammatical Logical Inference Framework

We have a tool for this!



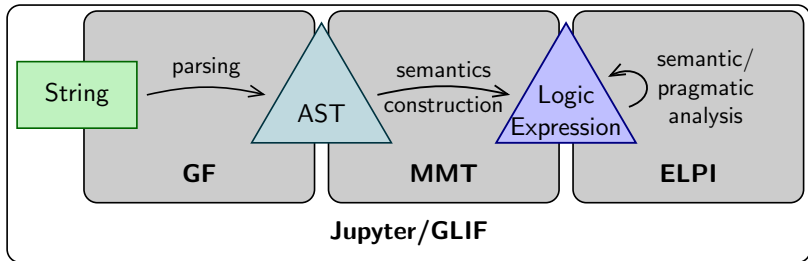
GLIF: Grammatical Logical Inference Framework

We have a tool for this!



GLIF: Grammatical Logical Inference Framework

It combines existing tools.



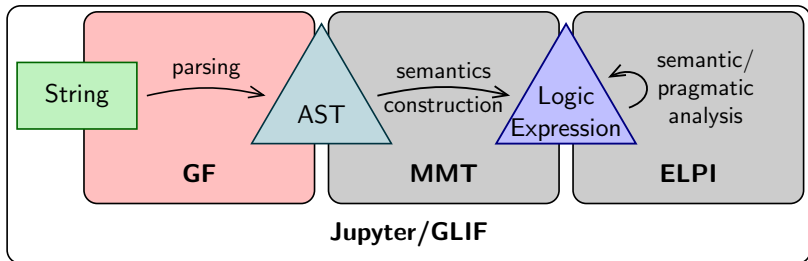
GF (= **grammar** framework)

+ **MMT** (= **logic** framework)

+ **ELPI** (= **inference** framework)

= **GLIF** (= **natural language understanding** framework)

Components of GLIF: GF



Components of GLIF: Grammatical Framework [GF]

- Specialized for developing natural language grammars

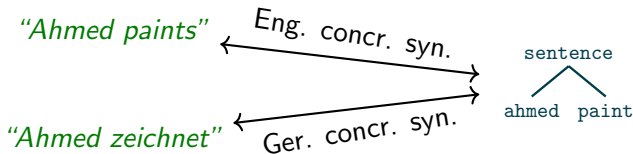
- Separates abstract and concrete syntax

`sentence : NP -> VP -> S;` *--abstract*

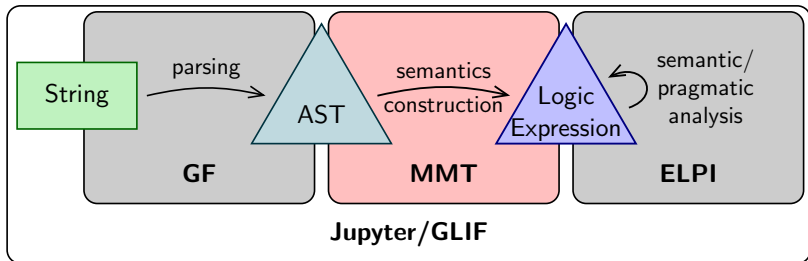
`sentence np vp = np.s ++ vp.s!np.n;` *--concrete*

- Abstract syntax based on type theory

- Comes with large library *≥ 36 languages*



Components of GLIF: MMT



Components of GLIF: MMT

- Modular logic development and knowledge representation
- Not specialized in one logical framework *we use LF*
- We will use MMT to:
 - represent abstract syntax
 - specify target logic and domain theory
 - specify semantics construction

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GF

```
cat
  NP; VP; S;
fun
  sentence :
    NP -> VP -> S;
```



MMT

```
NP : type
VP : type
S : type
sentence :
  NP → VP → S
```

Components of GLIF: MMT

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Logic Syntax

```
o : type //propositions
¬ : o → o
∧ : o → o → o
∨ : o → o → o

ι : type //individuals
∀ : (ι → o) → o
∃ : (ι → o) → o
```

Domain Theory

```
paint : ι → o
quiet : ι → o
ahmed : ι
berta : ι
```

idea: $\forall f$ or $\forall \lambda x.f(x)$
instead of $\forall x.f(x)$

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Semantics Construction

map symbols in abstract syntax to terms in logic/domain theory

Simple setting

S	\mapsto o
NP	\mapsto ι
VP	\mapsto ι \rightarrow o
sentence	\mapsto $\lambda n. \lambda v. v$ n
ahmed	\mapsto ahmed

More advanced

NP	\mapsto (ι \rightarrow o) \rightarrow o
sentence	\mapsto $\lambda n. \lambda v. n$ v
everyone	\mapsto $\lambda p. \forall \lambda x. p$ x
berta	\mapsto $\lambda p. p$ berta

Example: Parsing + Semantics Construction

“Ahmed and Berta paint”

↓ parsing

sentence (andNP ahmed berta) paint

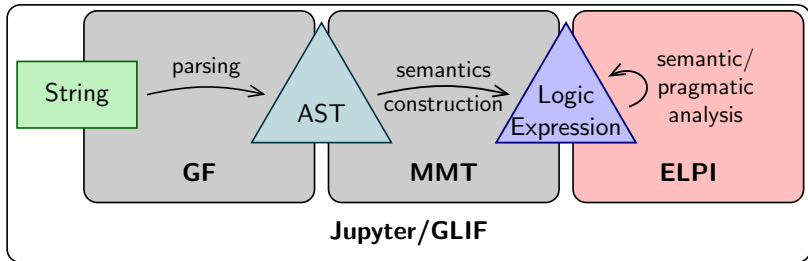
↓ semantics construction

$(\lambda n. \lambda v. n \ v) \ ((\lambda a. \lambda b. \lambda p. a \ p \ \wedge \ b \ p) \ (\lambda p. p \ \text{ahmed}) \ (\lambda p. p \ \text{berta})) \ \text{paint}$

↓ β -reduction

$\text{paint} \ \text{ahmed} \ \wedge \ \text{paint} \ \text{berta}$

Components of GLIF: ELPI

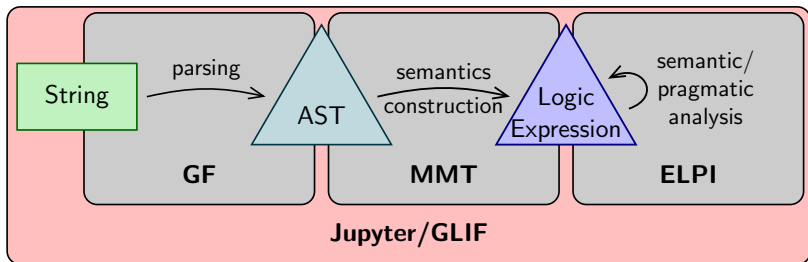


Components of GLIF: ELPI

- Implementation and extension of $\lambda\text{Prolog} \approx \text{Prolog} + \text{HOAS}$
- MMT can generate logic signatures
- First experiments with prover generation
- Generic inference/reasoning step after semantics construction

```
kind o type.  
type not o -> o.  
type and o -> o -> o.  
  
kind i type.  
type forall (i -> o) -> o.
```

Components of GLIF: Jupyter



Components of GLIF: Jupyter

- Unified, notebook-based interface
- Supports implementation and testing
- Useful for prototype, demos, teaching, ...

```
sentence = [n,v] n v |  
andNP = [a,b] [p] (a p)  $\wedge$  (b p) |  
paint = paint |  
ahmed = [p] p ahmed |  
berta = [p] p berta |
```

Created view GrammarSemantics

```
parse "Ahmed and Berta paint" | construct
```

```
(paint ahmed) $\wedge$ (paint berta)
```

Example: Epistemic Q&A

John knows that Mary or Eve knows that Ping has a dog. (S_1)

Mary doesn't know if Ping has a dog. (S_2)

Does Eve know if Ping has a dog? (Q)

$$S_1 = \Box_{john}(\Box_{mary}hd(ping)) \vee \Box_{eve}hd(ping)$$

$$S_2 = \neg((\Box_{mary}hd(ping)) \vee \Box_{mary}\neg hd(ping))$$

$$Q = (\Box_{eve}hd(ping)) \vee \Box_{eve}\neg hd(ping)$$

$$S_1, S_2 \vdash_{S5_n} Q \quad \rightsquigarrow \quad \text{yes}$$

$$S_1, S_2 \vdash_{S5_n} \neg Q \quad \rightsquigarrow \quad \text{no}$$

$$\text{else} \quad \rightsquigarrow \quad \text{maybe}$$

Example: Controlled Natural Languages

- Formal languages
- that are a subset of natural language
- and have fixed semantics

formal verification, ...

"S is a subset of every set iff S is empty"

$\rightsquigarrow (\forall V_{new}. \text{set}(V_{new}) \Rightarrow \text{subset}(V_S, V_{new})) \Leftrightarrow \text{empty}(V_S)$

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Use inference for disambiguation:

"... has a mass of 2m" $\begin{array}{l} \nearrow \text{AST}_1 \longrightarrow \lambda x.mass(x, quant(2, \mathbf{meters})) \\ \searrow \text{AST}_2 \longrightarrow \lambda x.mass(x, mul(2, \mathbf{mVar})) \end{array}$

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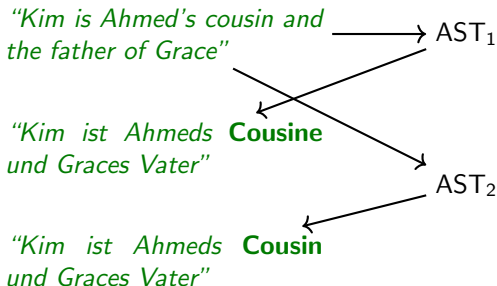
Use inference for disambiguation:

"... has a mass of 2m" \rightarrow AST₁ \rightarrow ~~$\lambda x. \text{mass}(x, \text{quant}(2, \text{meters}))$~~

"... has a mass of 2m" \rightarrow AST₂ \rightarrow $\lambda x. \text{mass}(x, \text{mul}(2, \text{mVar}))$

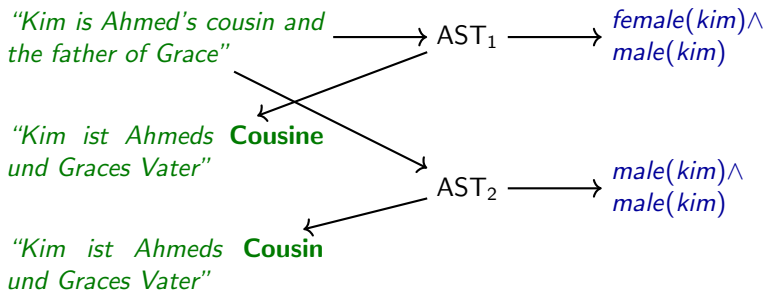
Example: Translation

- Two German words for “*cousin*”, depending on the gender
- Two entries in abstract syntax: `cousin_female` and `cousin_male`
- Use inference to discard ASTs



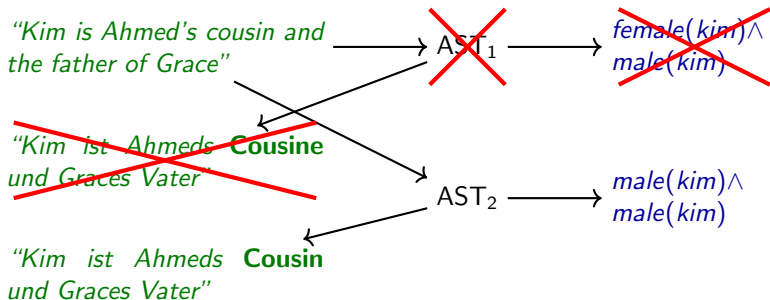
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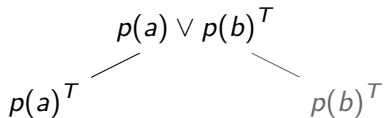
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Example: Tableaux Machine [KK03]

- Can use tableaux for model generation
- Tableau machine: pick “best” branch as model and continue there with next sentence *like a human?*

“Ahmed or Berta paints”

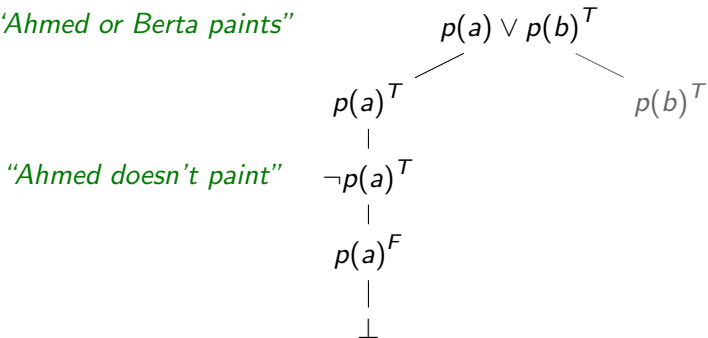


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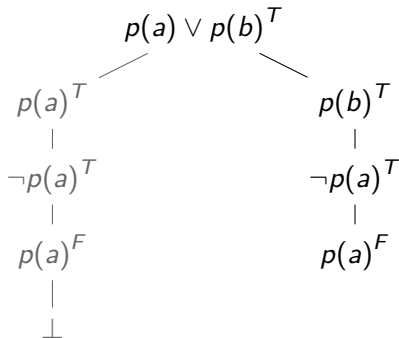
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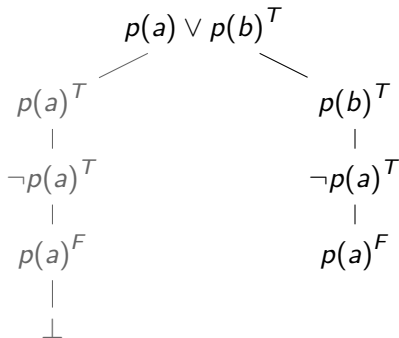


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Example: Tableaux Machine

Background Knowledge

"John talks to Mary."

talkto(j, m)

"Sasha is sad."

sad(s)

$\forall x. fem(x) \Rightarrow \neg male(x)$
male(j)
fem(m)

\downarrow
talkto(j, m)

\downarrow
sad(s)

Example: Tableaux Machine

Background Knowledge

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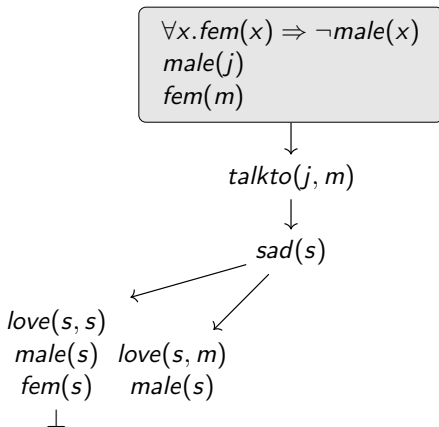
"Sasha is sad."

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"He loves her."

$\exists X. male(X) \wedge$

$\exists Y. fem(Y) \wedge love(X, Y)$



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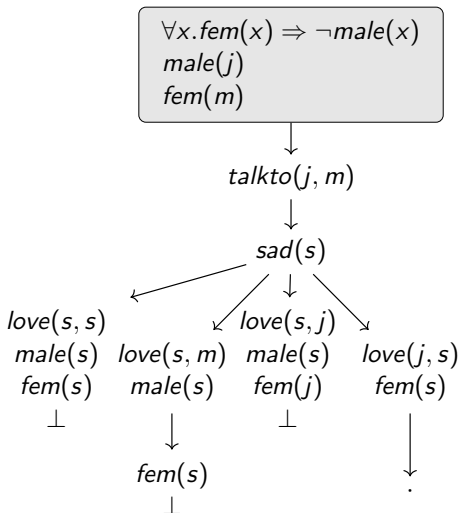
$sad(s)$

"He loves her."

$\exists X.male(X) \wedge$
 $\exists Y.fem(Y) \wedge love(X, Y)$

"Sasha is a woman."

$fem(s)$



Example: Tableaux Machine

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$\forall x.fem(x) \Rightarrow \neg male(x)$
 $male(j)$
 $fem(m)$

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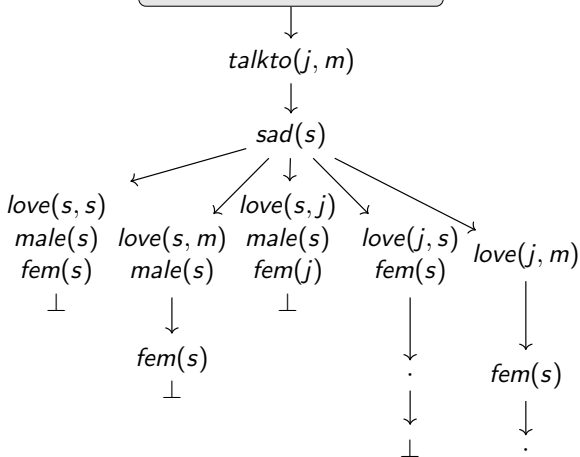
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"Sasha is a woman."

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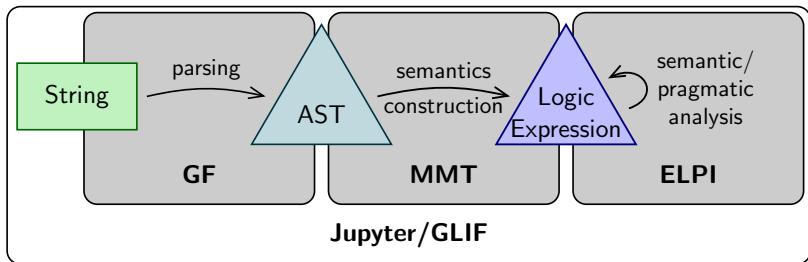
"John doesn't love Sasha."

$\neg love(j, s)$



Conclusion

- GLIF is a tool for prototyping natural language understanding pipelines
- Combines existing, declarative frameworks



References I



GF - Grammatical Framework. URL:

<http://www.grammaticalframework.org> (visited on 09/27/2017).



Michael Kohlhase and Alexander Koller.

“Resource-Adaptive Model Generation as a Performance Model”. In: *Logic Journal of the IGPL* 11.4 (2003), pp. 435–456. URL: <http://jigpal.oxfordjournals.org/cgi/content/abstract/11/4/435>.



R. Montague. “English as a Formal Language”. In: Reprinted in [Tho74], 188–221. Edizioni di Comunità, Milan, 1970, pp. 189–224.

References II



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