

# Decomposing Complete Hypergraphs

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### Problem (Steiner, 1853)

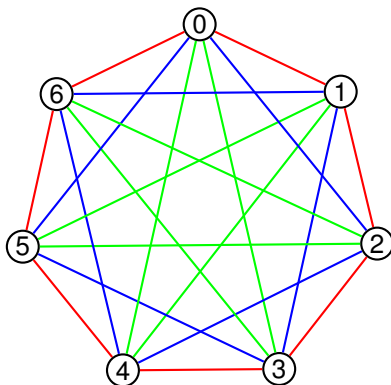
Given a set of  $v$  elements, take subsets of size 3 so that every pair of elements occurs in precisely one of the subsets.

Example: every pair of elements from  $\{0, 1, 2, 3, 4, 5, 6\}$  occurs in exactly one of the following triples:

$$\begin{array}{lll} \{0, 1, 3\} & \{1, 2, 4\} & \{2, 3, 5\} \\ \{3, 4, 6\} & \{0, 4, 5\} & \{1, 5, 6\} \\ & \{0, 2, 6\} & \end{array}$$

### Definition

A *Steiner Triple System* of order  $v$ , denoted  $STS(v)$ , is a set  $V$  of  $v$  *points* with a collection  $\mathcal{B}$  of 3-element subsets of  $V$ , called *blocks*, so that every 2-element subset of  $V$  is contained in exactly one block.



⇒ We can draw the complete graph  $K_7$  as a number of copies of  $K_3$ .

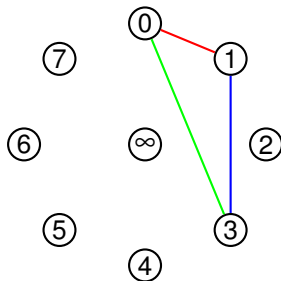
## Definition

Let  $G, H$  be graphs. A  $G$ -decomposition of  $H$  is a collection of graphs isomorphic to  $G$  whose edges partition  $E(H)$ .

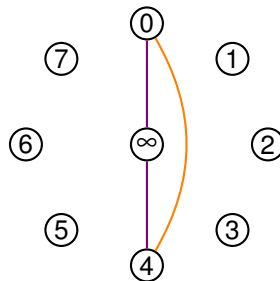
## Lemma

*An STS( $v$ ) is equivalent to a  $K_3$ -decomposition of  $K_v$ .*

An  $STS(9)$ , expressed as a  $K_3$ -decomposition of  $K_9$ :



$\{0, 1, 3\}, \{1, 2, 4\}, \{2, 3, 5\}, \dots$



$\{0, 4, \infty\}, \{1, 5, \infty\}, \dots$

We also have necessary conditions for the existence of a  $K_3$ -decomposition of  $K_v$ :

- 1  $K_v$  has  $\binom{v}{2}$  edges, which are partitioned into parts of size  $|E(K_3)| = 3$ , so 3 must divide  $\binom{v}{2} = \frac{v(v-1)}{2}$ .
- 2 Consider a vertex  $x \in V(K_v)$ . Whenever  $x$  appears in a copy of  $K_3$ , there are 2 edges incident with  $x$ . There are a total of  $\deg_{K_v}(x) = v - 1$  edges incident with  $x$ , so 2 must divide  $v - 1$ .

### Theorem (Kirkman, 1847)

*An STS( $v$ ) exists iff  $v \equiv 1$  or 3 (mod 6).*

### Proof.

( $\Rightarrow$ ) Obvious necessary conditions above.

( $\Leftarrow$ ) Generalise the constructions shown before.



We can generalise “3-element blocks” to “ $k$ -element blocks”, and from “covering every pair of points” to “covering every  $t$ -set of points”:

### Definition

A *Steiner system*, denoted  $S(t, k, v)$ , consists of a set of  $v$  points, and sets of  $k$  points (*blocks*), so that every set of  $t$  points occurs in precisely one block.

### Problem

Given  $k$  and  $t$ , for which  $v$  does there exist an  $S(t, k, v)$ ?

We would like to form an equivalent ‘graph theoretic’ question.

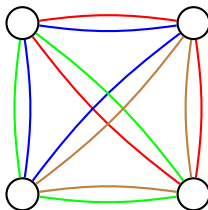
## Definition

A *hypergraph* is a pair  $(V, E)$ , where  $V$  is a set of *vertices* and  $E$  is a set of nonempty subsets of  $V$  called *hyperedges* (or *edges*).

A hypergraph is *t-uniform* if every edge has size  $t$ .

The *complete t-uniform hypergraph of order v*, denoted  $K_v^{(t)}$ , consists of  $v$  vertices and edges given by every  $t$ -set of vertices.

Example:  $K_4^{(3)}$ :





## Lemma

An  $S(t, k, v)$  is a  $K_k^{(t)}$ -decomposition of  $K_v^{(t)}$ .

The necessary divisibility conditions for the existence of an  $S(t, k, v)$  are:

- ①  $|E(K_k^{(t)})| = \binom{k}{t}$  divides  $|E(K_v^{(t)})| = \binom{v}{t}$
- ② For  $x \in V(K_v^{(t)})$ :  $\deg_{K_k^{(t)}}(x) = \binom{k-1}{t-1}$  divides  $\deg_{K_v^{(t)}}(x) = \binom{v-1}{t-1}$
- ③ For  $x, y$ :  $\deg_{K_k^{(t)}}(\{x, y\}) = \binom{k-2}{t-2}$  divides  $\deg_{K_v^{(t)}}(\{x, y\}) = \binom{v-2}{t-2}$
- ④ ...

In summary:

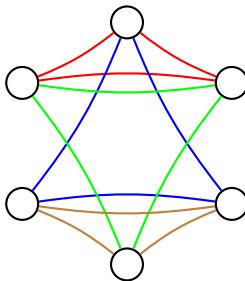
For  $i \geq 0$ ,  $\binom{k-i}{t-i}$  must divide  $\binom{v-i}{t-i}$ .

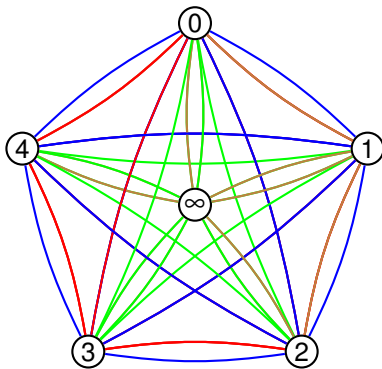
## Problem

Given a  $t$ -uniform hypergraph  $H$ , for which  $v$  does there exist an  $H$ -decomposition of  $K_v^{(t)}$ ?

These can also be referred to as *H*-designs of order  $v$ .

Example: for the hypergraph  $H$  below, does there exist an  $H$ -decomposition of  $K_6^{(3)}$ ?





The necessary divisibility conditions for  $H$ -decompositions of  $K_v^{(t)}$  are:  
For all  $i \geq 0$ ,  $\gcd(\{\deg_H(S) : S \subseteq V(H), |S| = i\})$  must divide  $\binom{v-i}{t-i}$ .

### Example

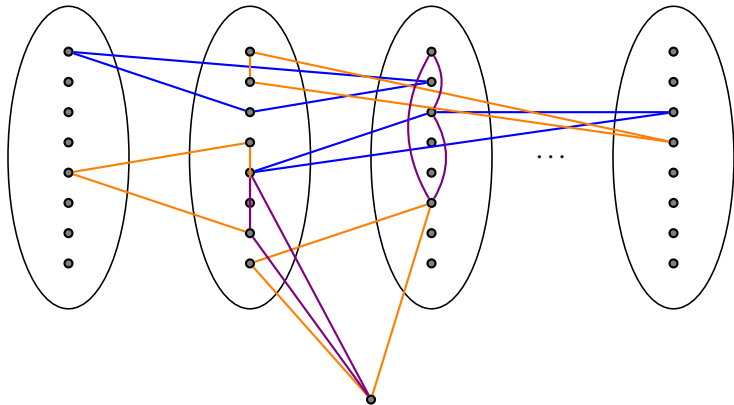
If  $H$  is as before, and there is a  $H$ -design of order  $v$ , then

- ①  $\deg_H(\emptyset) = |E(H)| = 4$  divides  $\binom{v}{3}$
- ② For all  $x \in V(H)$ ,  $\deg_H(x) = 2$ . So 2 divides  $\binom{v-1}{2}$
- ③  $\{\deg_H(\{x, y\}) : x, y \in V(H), x \neq y\} = \{1, 2\}$ . So  $\gcd(\{1, 2\}) = 1$  divides  $\binom{v-2}{1} = v - 2$ .

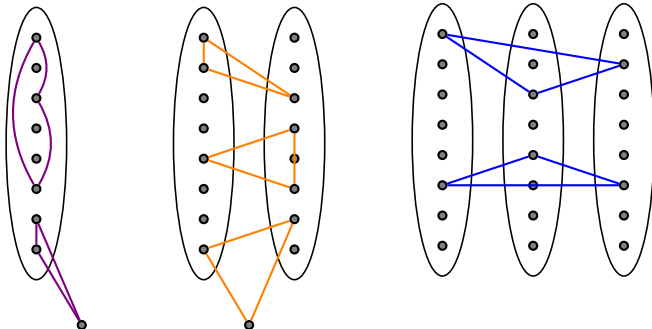
These conditions are equivalent to  $v \equiv 1, 2$  or  $6 \pmod{8}$ .

Does there always exist an  $H$ -design for these values of  $v$ ?

If  $v \equiv 1 \pmod{8}$ , then we need  $n$  sets of 8 points, plus one extra point:



If we can find  $H$ -decompositions of the following hypergraphs, then multiple copies of them will cover all three cases:



It can be shown that all of these exist, so  $H$ -designs of order  $v$  exist for all  $v \equiv 1 \pmod{8}$ . Similarly for  $v \equiv 2$  and  $6 \pmod{8}$ , so

### Theorem

*For the hypergraph  $H$  given before, there is an  $H$ -design of order  $v$  ( $H$ -decomposition of  $K_v^{(3)}$ ) iff  $v \equiv 1, 2$ , or  $6 \pmod{8}$ .*

### Proof.

( $\Rightarrow$ ) Necessary divisibility conditions.

( $\Leftarrow$ ) Divide into three cases:  $1, 2, 6 \pmod{8}$ .

Find  $H$ -decompositions of certain small hypergraphs.

Combine copies of these decompositions as necessary. □