# **Decomposing Complete Hypergraphs**

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### Definition

A *Steiner Triple System* of order v, denoted STS(v), is a set V of v points with a collection  $\mathscr{B}$  of 3-element subsets of V, called *blocks*, so that every 2-element subset of V is contained in exactly one block.





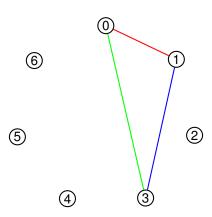


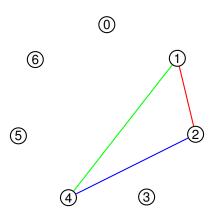








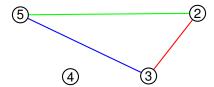


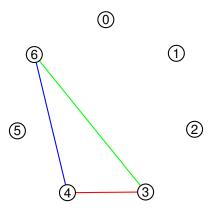


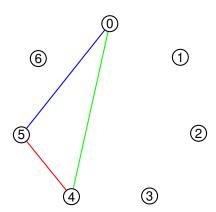


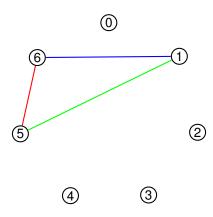


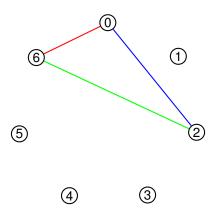


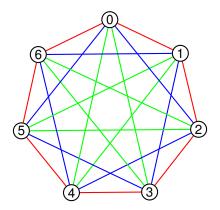












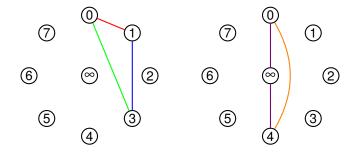
 $\Rightarrow$  We can draw the complete graph  $K_7$  as a number of copies of  $K_3$ .

Let G, H be graphs. A G-decomposition of H is a collection of graphs isomorphic to G whose edges partition E(H).

#### Lemma

An STS(v) is equivalent to a  $K_3$ -decomposition of  $K_v$ .

## An STS(9), expressed as a $K_3$ -decomposition of $K_9$ :



$$\{0,1,3\},\{1,2,4\},\{2,3,5\},\ldots$$

$$\{0,4,\infty\},\{1,5,\infty\},\ldots$$

•  $K_v$  has  $\binom{v}{2}$  edges, which are partitioned into parts of size  $|E(K_3)|=3$ , so 3 must divide  $\binom{v}{2}=\frac{v(v-1)}{2}$ .

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- ② Consider a vertex  $x \in V(K_v)$ . Whenever x appears in a copy of  $K_3$ , there are 2 edges incident with x. There are a total of  $\deg_{K_v}(x) = v 1$  edges incident with x, so 2 must divide v 1.

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## Theorem (Kirkman, 1847)

An STS(v) exists iff  $v \equiv 1$  or 3 (mod 6).

#### Proof.

- (⇒) Obvious necessary conditions above.
- (⇐) Generalise the constructions shown before.

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We would like to form an equivalent 'graph theoretic' question.

A hypergraph is a pair (V, E), where V is a set of vertices and E is a set of nonempty subsets of V called hyperedges (or edges).

A hypergraph is t-uniform if every edge has size t.

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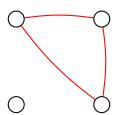
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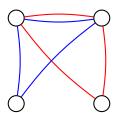
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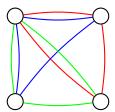
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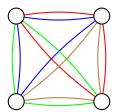
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The necessary divisibility conditions for the existence of an S(t, k, v) are:

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In summary:

For  $i \geq 0$ ,  $\binom{k-i}{t-i}$  must divide  $\binom{v-i}{t-i}$ .

### Problem

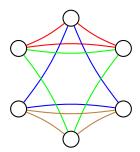
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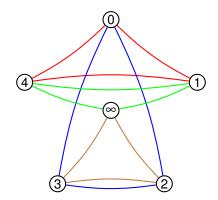
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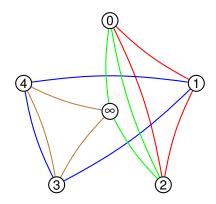
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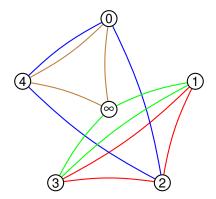
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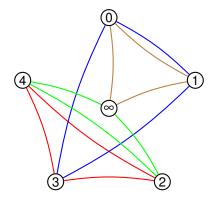
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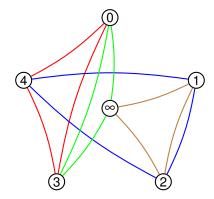












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- ② For all  $x \in V(H)$ ,  $\deg_H(x) = 2$ . So 2 divides  $\binom{v-1}{2}$
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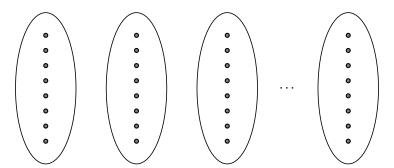
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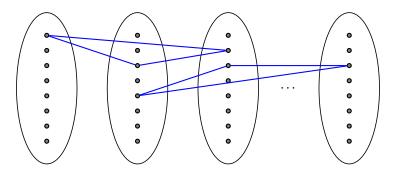
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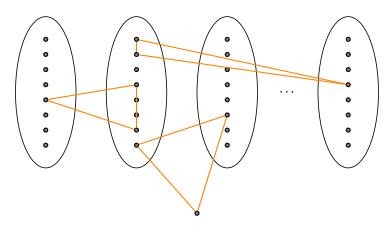
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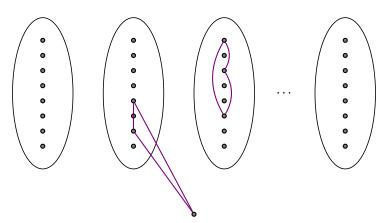
These conditions are equivalent to  $v \equiv 1, 2 \text{ or } 6 \pmod{8}$ .

Does there always exist an H-design for these values of v?





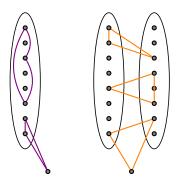




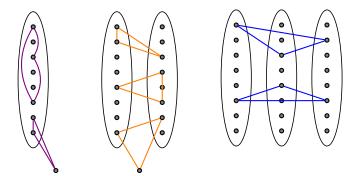
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#### Theorem

For the hypergraph H given before, there is an H-design of order v (H-decomposition of  $K_v^{(3)}$ ) iff  $v \equiv 1, 2, \text{ or } 6 \pmod{8}$ .

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### Proof.

- (⇒) Necessary divisibility conditions.
- $(\Leftarrow)$  Divide into three cases: 1, 2, 6 (mod 8).

Find *H*-decompositions of certain small hypergraphs.

Combine copies of these decompositions as necessary.

