

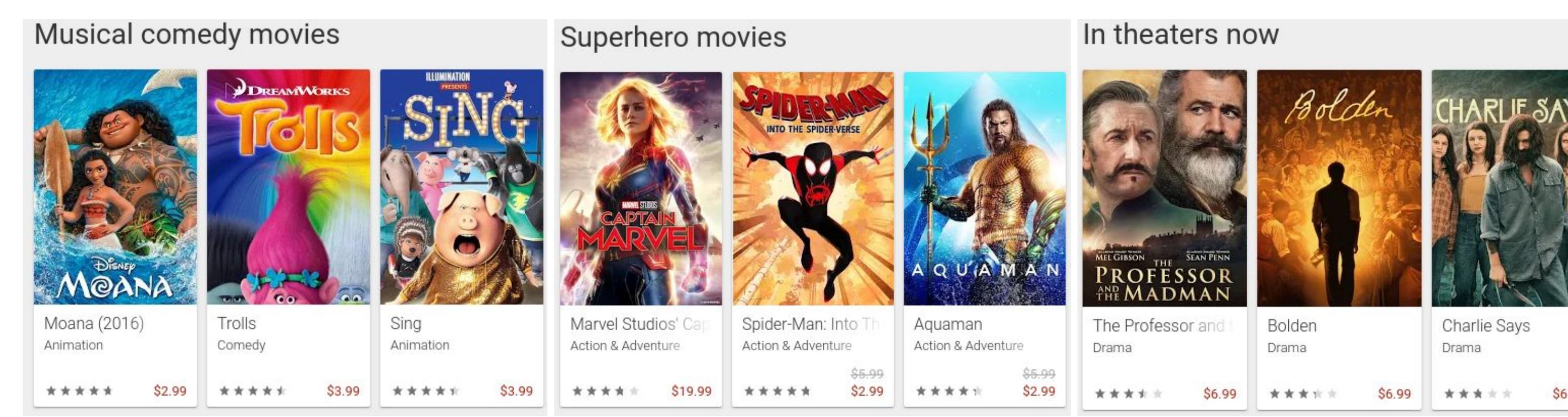
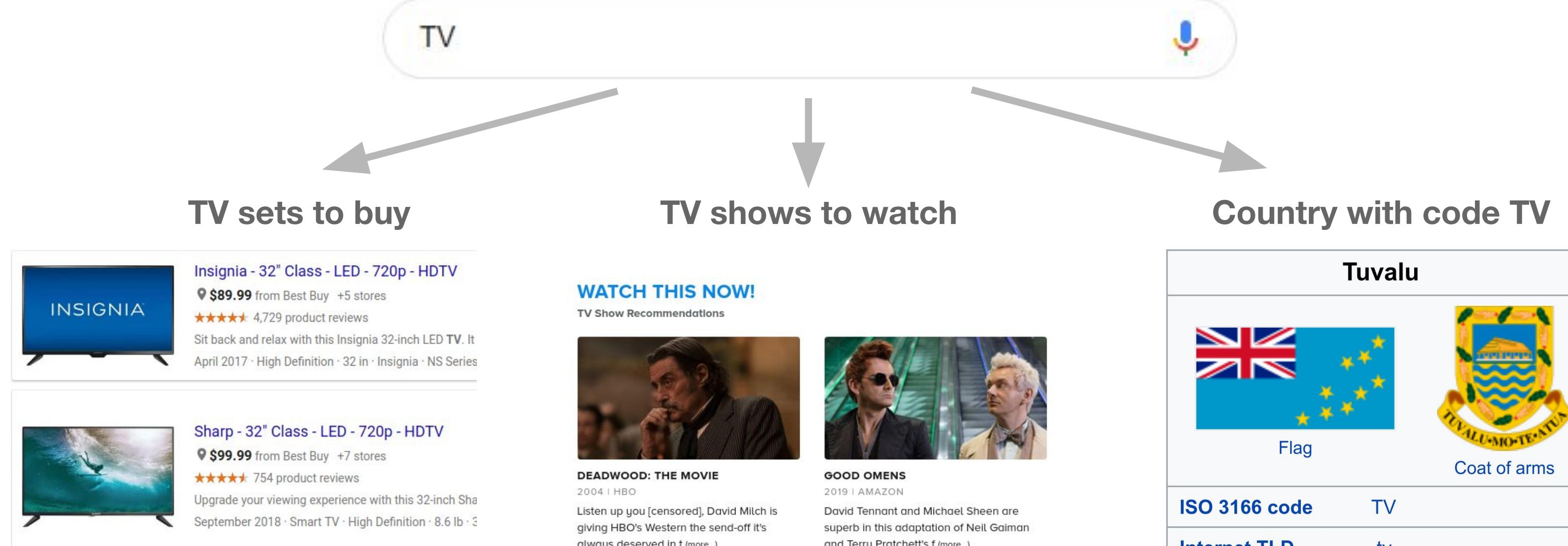
A Tree-Based Method for Fast Repeated Sampling of Determinantal Point Processes

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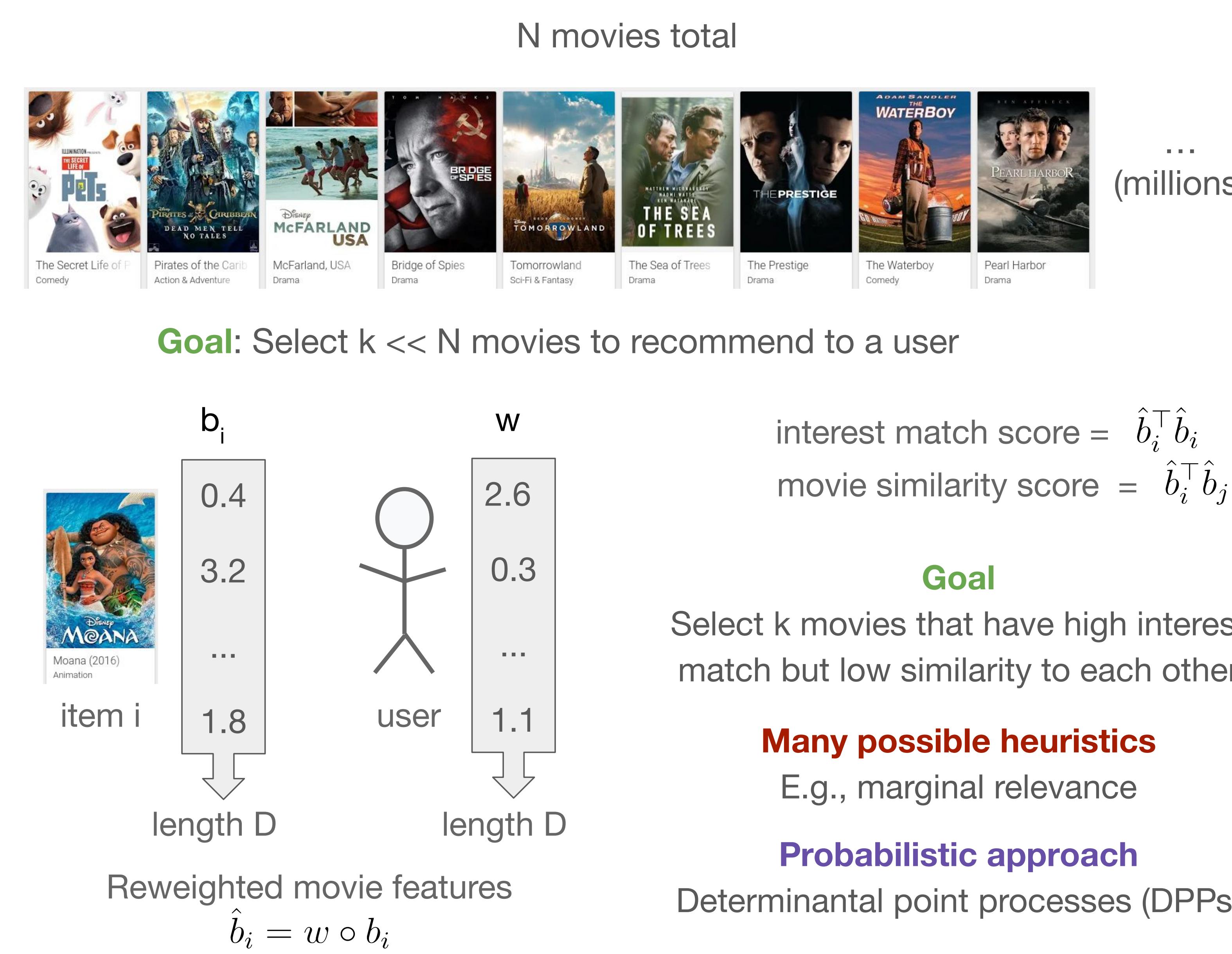
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Motivation

Diverse results are desirable in information retrieval and recommender systems



Determinantal Point Processes



$$W \in \mathbb{R}^{D \times D} \quad \times \quad B \in \mathbb{R}^{D \times N} \quad = \quad \hat{B} = WB$$

Positive semi-definite kernel $\hat{L} = \hat{B}^\top \hat{B}$
 $\hat{L}_{ij} = \hat{b}_i^\top \hat{b}_j$

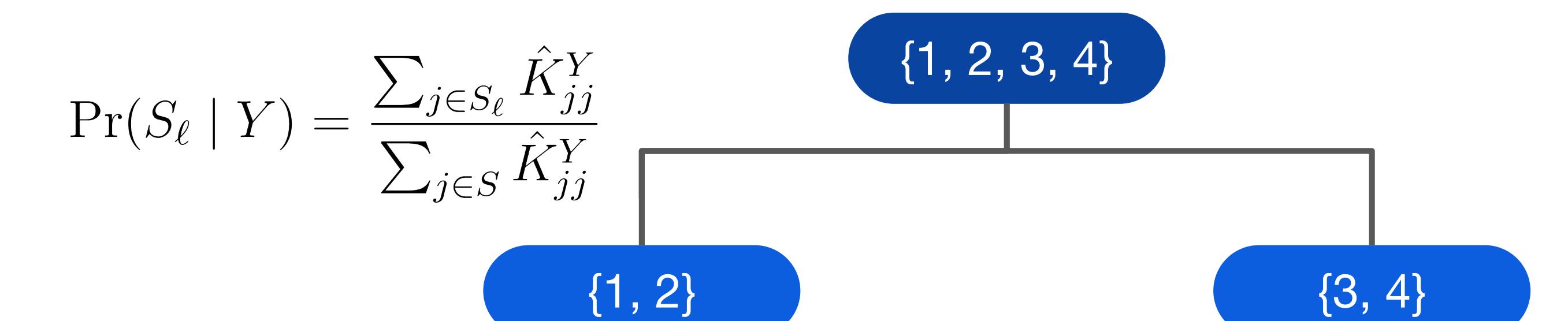
Distribution over all subsets $Y \subseteq [N]$
 $\mathcal{P}_{\hat{L}}(Y) \propto \det(\hat{L}_Y)$

Given $|Y| = 2$: $\det(\hat{L}_{\{i,j\}}) = \hat{b}_i^\top \hat{b}_i \hat{b}_j^\top \hat{b}_j - (\hat{b}_i^\top \hat{b}_j)^2$

Intuition for $|Y| = 2$: $\det(\hat{L}_{\{i,j\}}) = \text{interest in movie } i \cdot \text{interest in movie } j - \text{similarity of movies } i \text{ and } j$

Sampling DPPs

- Goal:** For each user, draw a size-k sample from their DPP
- Problem:** Existing algorithms for k-DPP sampling are too expensive
 - $D << N$ by construction or random projection
 - $O(ND^2)$ preprocessing on $L = B^\top B$
 - $O(Nk^2 + D^3)$ per personalized (W-weighted) sample afterwards
- Our contribution:** Making repeated, personalized k-DPP sampling efficient:
 - $O(ND^2)$ preprocessing on $L = B^\top B$
 - $O(D^2k^2 \log N + D^3)$ per personalized (W-weighted) sample afterwards



$$\sum_{j \in S} \hat{K}_{jj}^Y = \mathbf{1}^\top [R \circ \Sigma^{(S)}] \mathbf{1} - \mathbf{1}^\top [(\hat{K}_Y)^{-1} \circ (F \Sigma^{(S)} F^\top)] \mathbf{1} = f(\Sigma^{(S)}, \hat{\lambda}, \hat{V}, W)$$

where: $\hat{\lambda} = (1/\hat{\lambda})$, $M = \hat{V}_{:,E}^\top W$, $\hat{H} = \hat{V}_E M B_{:,Y}$

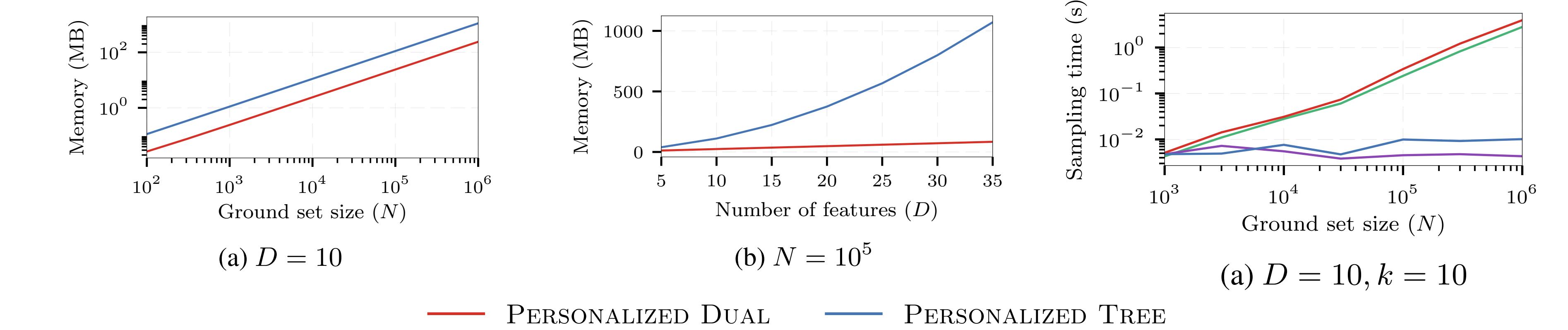
$$R = M^\top \hat{V}_E M, F = \hat{H}^\top M$$

Computable in $O(kD^2)$ time $\Rightarrow O(kD^2 \log N)$ per tree traversal

Overall: $O(k^2D^2 \log N + D^3)$ time to sample

Experiments

Preprocessing:



At $N = 1$ million: standard sampling takes 4 secs; tree-based takes 0.01 secs

Cost: Memory required to store the tree

Approximation

Main idea: If the distribution over items at a tree node is close to uniform, then don't bother moving further down the tree.



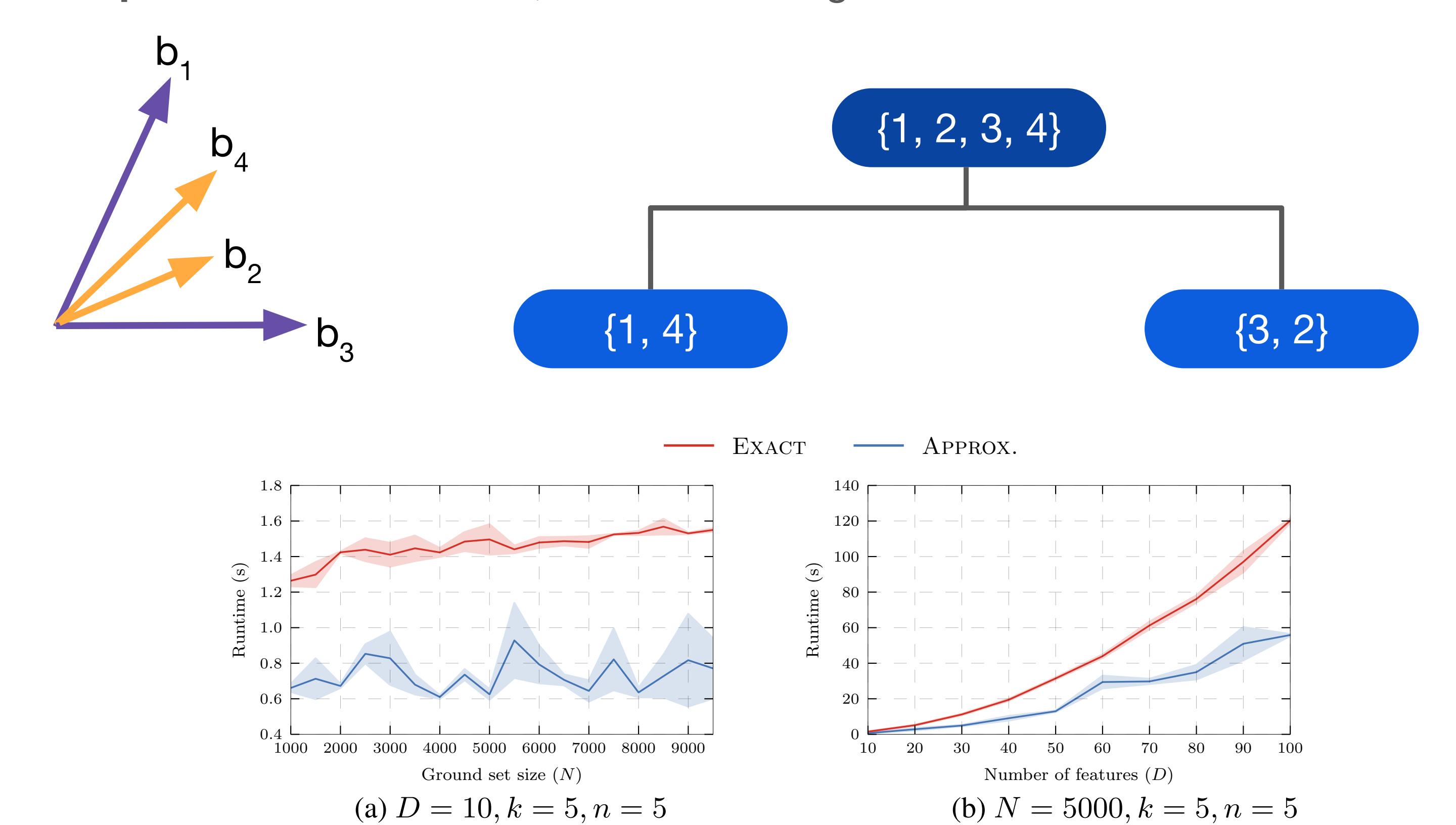
$$\tilde{\Sigma}_{\ell_1 \ell_2}^{(S)} = \max_{j \in S} \left| \sum_{i=1}^{|\ell_1 \ell_2|} \Sigma_{\ell_1 \ell_2}^{(j)} - \frac{1}{|\ell_1 \ell_2|} \sum_{i=1}^{|\ell_1 \ell_2|} \Sigma_{\ell_1 \ell_2}^{(S)} \right|$$

$$\Pr(j | S, Y) - \frac{1}{|\ell_1 \ell_2|} \leq \frac{f(\Sigma^{(S)}, \hat{\lambda}, \hat{V}, W)}{f(\Sigma^{(S)}, \hat{\lambda}, \hat{V}, W)}$$

$$\tilde{\Sigma}_{\ell_1 \ell_2}^{\{1,2,3,4\}} = \max \left(\left| \sum_{i=1}^{|\ell_1 \ell_2|} \Sigma_{\ell_1 \ell_2}^{(\{1\})} - \frac{1}{4} \sum_{i=1}^{|\ell_1 \ell_2|} \Sigma_{\ell_1 \ell_2}^{(\{1,2,3,4\})} \right|, \left| \sum_{i=1}^{|\ell_1 \ell_2|} \Sigma_{\ell_1 \ell_2}^{(\{2\})} - \frac{1}{4} \sum_{i=1}^{|\ell_1 \ell_2|} \Sigma_{\ell_1 \ell_2}^{(\{1,2,3,4\})} \right|, \dots \right)$$

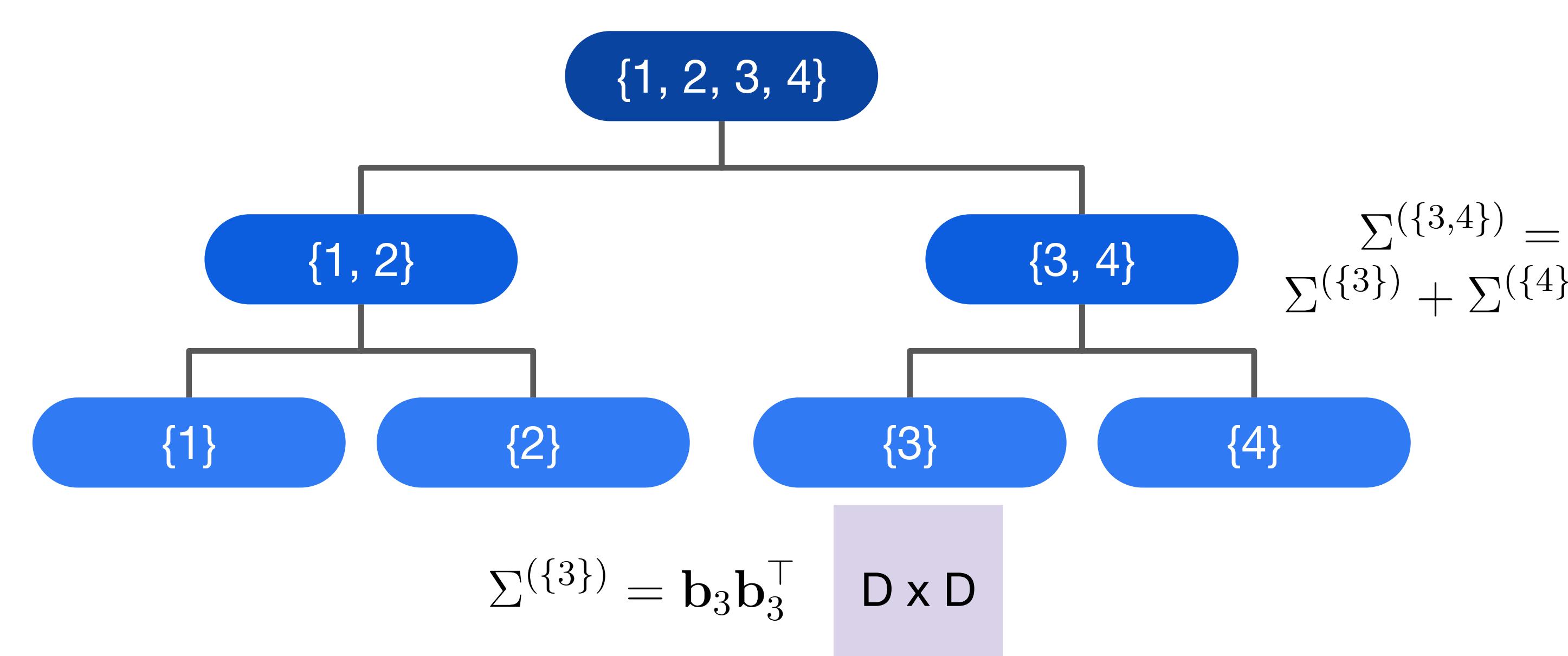
Sampling Y in order y_1, y_2, \dots, y_k :
 $(\text{true probability}) - (\text{probability with early stopping}) \leq (1 + \epsilon)^k - 1$

Idea: Use node-splitting stage of tree construction to increase uniformity
Example: Find distinct items, seed left and right subtrees with these.



Our tree-based algorithm

Key idea: In pre-processing, create a balanced binary tree of depth $\log N$.



Given tree T and $C = BB^\top$, sample from k-DPP with kernel $\hat{L} = (WB)^\top WB$
Add items one at a time, starting from $Y = \{\}$
Traverse tree once for each item addition: