# The Python Language Reference Manual

# **The Python Language Reference Manual**

@ Chris Sheridan 2016

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# INTRODUCTION

This reference manual describes the syntax and "core semantics" of the language. It is terse, but attempts to be exact

and complete. The semantics of non-essential built-in object types and of the built-in functions and modules are described

in library-index. For an informal introduction to the language, please see my other tutorials.

# 1.1 Alternate Implementations

Though there is one Python implementation which is by far the most popular, there are some alternate implementations

which are of particular interest to different audiences.

Known implementations include:

CPython This is the original and most-maintained implementation of Python, written in C. New language features

generally appear here first.

Jython Python implemented in Java. This implementation can be used as a scripting language for Java applications,

or can be used to create applications using the Java class libraries. It is also often used to create tests for Java

libraries. More information can be found at the Jython website.

Python for .NET This implementation actually uses the CPython implementation, but is a managed .NET application

and makes .NET libraries available. It was created by Brian Lloyd. For more information, see the Python for

.NET home page.

IronPython An alternate Python for .NET. Unlike Python.NET, this is a complete Python implementation that generates

IL, and compiles Python code directly to .NET assemblies. It was created by Jim Hugunin, the original

creator of Jython. For more information, see the IronPython website.

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PyPy An implementation of Python written completely in Python. It supports several advanced features not found

in other implementations like stackless support and a Just in Time compiler. One of the goals of the project is

to encourage experimentation with the language itself by making it easier to modify the interpreter (since it is

written in Python). Additional information is available on the PyPy project's home page.

Each of these implementations varies in some way from the language as documented in this manual, or introduces

specific information beyond what's covered in the standard Python documentation. Please refer to the implementation specific

documentation to determine what else you need to know about the specific implementation you're using.

### 1.2 Notation

The descriptions of lexical analysis and syntax use a modified BNF grammar notation. This uses the following style

of definition:

```
name ::= lc_letter (lc_letter | "_")*
lc_letter ::= "a"..."z"
```

The first line says that a name is an lc\_letter followed by a sequence of zero or more lc\_letters and underscores.

An Ic\_letter in turn is any of the single characters 'a' through 'z'. (This rule is actually adhered to for

the names defined in lexical and grammar rules in this document.)

Each rule begins with a name (which is the name defined by the rule) and ::=. A vertical bar (|) is used to separate

alternatives; it is the least binding operator in this notation. A star (\*) means zero or more repetitions of the preceding

item; likewise, a plus (+) means one or more repetitions, and a phrase enclosed in square brackets ([]) means zero

or one occurrences (in other words, the enclosed phrase is optional). The \* and + operators bind as tightly as possible;

parentheses are used for grouping. Literal strings are enclosed in quotes. White space is only meaningful to separate

tokens. Rules are normally contained on a single line; rules with many alternatives may be formatted alternatively

with each line after the first beginning with a vertical bar.

In lexical definitions (as the example above), two more conventions are used: Two literal characters separated by three

dots mean a choice of any single character in the given (inclusive) range of ASCII characters. A phrase between

angular brackets (<...>) gives an informal description of the symbol defined; e.g., this could be used to describe the

notion of 'control character' if needed.

Even though the notation used is almost the same, there is a big difference between the meaning of lexical and

syntactic definitions: a lexical definition operates on the individual characters of

the input source, while a syntax

definition operates on the stream of tokens generated by the lexical analysis. All uses of BNF in the next chapter

("Lexical Analysis") are lexical definitions; uses in subsequent chapters are syntactic definitions.

### **CHAPTER TWO LEXICAL ANALYSIS**

A Python program is read by a parser. Input to the parser is a stream of tokens, generated by the lexical analyzer. This

chapter describes how the lexical analyzer breaks a file into tokens.

Python reads program text as Unicode code points; the encoding of a source file can be given by an encoding declaration

and defaults to UTF-8.

### 2.1 Line structure

A Python program is divided into a number of logical lines.

# 2.1.1 Logical lines

The end of a logical line is represented by the token NEWLINE. Statements cannot cross logical line boundaries

except where NEWLINE is allowed by the syntax (e.g., between statements in compound statements). A logical line

is constructed from one or more physical lines by following the explicit or implicit line joining rules.

# 2.1.2 Physical lines

A physical line is a sequence of characters terminated by an end-of-line sequence. In source files, any of the standard

platform line termination sequences can be used - the Unix form using ASCII LF (linefeed), the Windows form using

the ASCII sequence CR LF (return followed by linefeed), or the old Macintosh form using the ASCII CR (return)

character. All of these forms can be used equally, regardless of platform.

When embedding Python, source code strings should be passed to Python APIs using the standard C conventions for

newline characters (the \n character, representing ASCII LF, is the line terminator).

### 2.1.3 Comments

A comment starts with a hash character (#) that is not part of a string literal, and

ends at the end of the physical line. A

comment signifies the end of the logical line unless the implicit line joining rules are invoked. Comments are ignored

by the syntax; they are not tokens.

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# 2.1.4 Encoding declarations

If a comment in the first or second line of the Python script matches the regular expression

coding[=:]\s\*([-\w.]+), this comment is processed as an encoding declaration; the first group of this

expression names the encoding of the source code file. The recommended forms of this expression are

# -\*- coding: <encoding-name> -\*-

which is recognized also by GNU Emacs, and

# vim:fileencoding=<encoding-name>

which is recognized by Bram Moolenaar's VIM.

If no encoding declaration is found, the default encoding is UTF-8. In addition, if the first bytes of the file are the

UTF-8 byte-order mark (b'\xef\xbb\xbf'), the declared file encoding is UTF-8 (this is supported, among others,

by Microsoft's notepad).

If an encoding is declared, the encoding name must be recognized by Python. The encoding is used for all lexical

analysis, including string literals, comments and identifiers. The encoding declaration must appear on a line of its own.

# 2.1.5 Explicit line joining

Two or more physical lines may be joined into logical lines using backslash characters (\), as follows: when a physical

line ends in a backslash that is not part of a string literal or comment, it is joined with the following forming a single

logical line, deleting the backslash and the following end-of-line character. For

example:

```
if 1900 < year < 2100 and 1 <= month <= 12 \
and 1 <= day <= 31 and 0 <= hour < 24 \
and 0 <= minute < 60 and 0 <= second < 60: # Looks like a valid date
```

return 1

A line ending in a backslash cannot carry a comment. A backslash does not continue a comment. A backslash does

not continue a token except for string literals (i.e., tokens other than string literals cannot be split across physical lines

using a backslash). A backslash is illegal elsewhere on a line outside a string literal.

# 2.1.6 Implicit line joining

Expressions in parentheses, square brackets or curly braces can be split over more than one physical line without using

backslashes. For example:

month\_names = ['Januari', 'Februari', 'Maart', # These are the

'April', 'Mei', 'Juni', # Dutch names

'Juli', 'Augustus', 'September', # for the months

'Oktober', 'November', 'December'] # of the year

Implicitly continued lines can carry comments. The indentation of the continuation lines is not important. Blank

continuation lines are allowed. There is no NEWLINE token between implicit continuation lines. Implicitly continued

lines can also occur within triple-quoted strings (see below); in that case they cannot carry comments.

### 2.1.7 Blank lines

A logical line that contains only spaces, tabs, formfeeds and possibly a comment, is ignored (i.e., no NEWLINE

token is generated). During interactive input of statements, handling of a blank line may differ depending on the implementation of the read-eval-print loop. In the standard interactive interpreter, an entirely blank logical line (i.e.

one containing not even whitespace or a comment) terminates a multi-line statement.

### 2.1.8 Indentation

Leading whitespace (spaces and tabs) at the beginning of a logical line is used to compute the indentation level of the

line, which in turn is used to determine the grouping of statements.

Tabs are replaced (from left to right) by one to eight spaces such that the total number of characters up to and including

the replacement is a multiple of eight (this is intended to be the same rule as used by Unix). The total number of spaces

preceding the first non-blank character then determines the line's indentation. Indentation cannot be split over multiple

physical lines using backslashes; the whitespace up to the first backslash determines the indentation.

Indentation is rejected as inconsistent if a source file mixes tabs and spaces in a way that makes the meaning dependent

on the worth of a tab in spaces; a TabError is raised in that case.

Cross-platform compatibility note: because of the nature of text editors on non-UNIX platforms, it is unwise to use

a mixture of spaces and tabs for the indentation in a single source file. It should also be noted that different platforms

may explicitly limit the maximum indentation level.

A formfeed character may be present at the start of the line; it will be ignored for the indentation calculations above.

Formfeed characters occurring elsewhere in the leading whitespace have an undefined effect (for instance, they may

reset the space count to zero).

The indentation levels of consecutive lines are used to generate INDENT and DEDENT tokens, using a stack, as

follows.

Before the first line of the file is read, a single zero is pushed on the stack; this will never be popped off again. The

numbers pushed on the stack will always be strictly increasing from bottom to top. At the beginning of each logical

line, the line's indentation level is compared to the top of the stack. If it is equal, nothing happens. If it is larger, it is

pushed on the stack, and one INDENT token is generated. If it is smaller, it must

be one of the numbers occurring on

the stack; all numbers on the stack that are larger are popped off, and for each number popped off a DEDENT token is

generated. At the end of the file, a DEDENT token is generated for each number remaining on the stack that is larger than zero.

Here is an example of a correctly (though confusingly) indented piece of Python code:

```
def perm(I):
# Compute the list of all permutations of I
if len(I) \leq 1:
return [l]
r = \Pi
for i in range(len(l)):
s = |[:i] + |[i+1:]
p = perm(s)
for x in p:
r.append(I[i:i+1] + x)
return r
The following example shows various indentation errors:
def perm(I): # error: first line indented
for i in range(len(l)): # error: not indented
s = |[:i] + |[i+1:]
p = perm(I[:i] + I[i+1:]) # error: unexpected indent
for x in p:
r.append(I[i:i+1] + x)
return r # error: inconsistent dedent
```

(Actually, the first three errors are detected by the parser; only the last error is found by the lexical analyzer — the

indentation of return r does not match a level popped off the stack.)

# 2.1.9 Whitespace between tokens

Except at the beginning of a logical line or in string literals, the whitespace characters space, tab and formfeed can be

used interchangeably to separate tokens. Whitespace is needed between two tokens only if their concatenation could

otherwise be interpreted as a different token (e.g., ab is one token, but a b is two tokens).

### 2.2 Other tokens

Besides NEWLINE, INDENT and DEDENT, the following categories of tokens exist: identifiers, keywords, literals,

operators, and delimiters. Whitespace characters (other than line terminators, discussed earlier) are not tokens, but

serve to delimit tokens. Where ambiguity exists, a token comprises the longest possible string that forms a legal token,

when read from left to right.

# 2.3 Identifiers and keywords

Identifiers (also referred to as names) are described by the following lexical definitions.

The syntax of identifiers in Python is based on the Unicode standard annex UAX-31, with elaboration and changes as

defined below;

Within the ASCII range (U+0001..U+007F), the valid characters for identifiers are the same as in Python 2.x: the

uppercase and lowercase letters A through Z, the underscore \_ and, except for the first character, the digits 0 through

9.

Python 3.0 introduces additional characters from outside the ASCII range. For these characters, the classification uses the version of the Unicode Character Database as included in

the unicodedata module.

Identifiers are unlimited in length. Case is significant.

identifier ::= xid\_start xid\_continue\*

id\_start ::= <all characters in general categories Lu, Ll, Lt, Lm, Lo, Nl, the underscore, id\_continue ::= <all characters in id\_start, plus characters in the

categories Mn, Mc, xid\_start ::= <all characters in id\_start whose NFKC normalization is in "id\_start xid\_continue ::= <all characters in id\_continue whose NFKC normalization is in "id\_continue\*">

The Unicode category codes mentioned above stand for:

- Lu uppercase letters
- LI lowercase letters
- Lt titlecase letters
- · Lm modifier letters
- · Lo other letters
- NI letter numbers
- Mn nonspacing marks
- Mc spacing combining marks
- Nd decimal numbers
- Pc connector punctuations
- Other\_ID\_Start explicit list of characters in PropList.txt to support backwards compatibility
- Other\_ID\_Continue likewise

All identifiers are converted into the normal form NFKC while parsing; comparison of identifiers is based on NFKC.

A non-normative HTML file listing all valid identifier characters for Unicode 4.1 can be found at

http://www.dcl.hpi.uni-potsdam.de/home/loewis/table-3131.html.

# 2.3.1 Keywords

The following identifiers are used as reserved words, or keywords of the language, and cannot be used as ordinary

identifiers. They must be spelled exactly as written here:

False class finally is return

None continue for lambda try

True def from nonlocal while

and del global not with

as elif if or yield

assert else import pass

### 2.3.2 Reserved classes of identifiers

Certain classes of identifiers (besides keywords) have special meanings. These classes are identified by the patterns of

leading and trailing underscore characters:

\_\* Not imported by from module import \*. The special identifier \_ is used in the interactive interpreter to

store the result of the last evaluation; it is stored in the builtins module. When not in interactive mode,

has no special meaning and is not defined. See section The import statement.

Note: The name \_ is often used in conjunction with internationalization; refer to the documentation for the

gettext module for more information on this convention.

\_\_\*\_\_ System-defined names. These names are defined by the interpreter and its implementation (including the

standard library). Current system names are discussed in the Special method names section and elsewhere.

More will likely be defined in future versions of Python. Any use of \_\_\_\*\_\_ names, in any context, that does not

follow explicitly documented use, is subject to breakage without warning.

\_\_\* Class-private names. Names in this category, when used within the context of a class definition, are re-written

to use a mangled form to help avoid name clashes between "private" attributes of base and derived classes. See

section Identifiers (Names).

# 2.4 Literals

Literals are notations for constant values of some built-in types.

# 2.4.1 String and Bytes literals

String literals are described by the following lexical definitions:

stringliteral ::= [stringprefix](shortstring | longstring)

```
stringprefix ::= "r" | "u" | "R" | "U"
shortstring ::= "" shortstringitem* "" | "" shortstringitem* ""
longstring ::= "''" longstringitem* "''" longstringitem* """
shortstringitem ::= shortstringchar | stringescapeseg
longstringitem ::= longstringchar | stringescapeseq
shortstringchar ::= <any source character except "\" or newline or the quote>
longstringchar ::= <any source character except "\">
stringescapeseq ::= "\" <any source character>
bytesliteral ::= bytesprefix(shortbytes | longbytes)
bytesprefix ::= "b" | "B" | "br" | "Br" | "bR" | "BR" | "rb" | "rB" | "Rb" | "RB"
shortbytes ::= "" shortbytesitem* "" | "" shortbytesitem* ""
longbytes ::= "''" longbytesitem* "''" | '"" longbytesitem* '""
shortbytesitem ::= shortbyteschar | bytesescapeseq
longbytesitem ::= longbyteschar | bytesescapeseq
shortbyteschar ::= <any ASCII character except "\" or newline or the quote>
longbyteschar ::= <any ASCII character except "\">
```

bytesescapeseg ::= "\" <any ASCII character>

One syntactic restriction not indicated by these productions is that whitespace is not allowed between the

stringprefix or bytesprefix and the rest of the literal. The source character set is defined by the encoding

declaration; it is UTF-8 if no encoding declaration is given in the source file; see section Encoding declarations.

In plain English: Both types of literals can be enclosed in matching single quotes (') or double quotes ("). They can

also be enclosed in matching groups of three single or double quotes (these are generally referred to as triple-quoted

strings). The backslash (\) character is used to escape characters that otherwise have a special meaning, such as

newline, backslash itself, or the quote character.

Bytes literals are always prefixed with 'b' or 'B'; they produce an instance of the bytes type instead of the str

type. They may only contain ASCII characters; bytes with a numeric value of 128 or greater must be expressed with

escapes.

As of Python 3.3 it is possible again to prefix unicode strings with a u prefix to simplify maintenance of dual 2.x and

3.x codebases.

Both string and bytes literals may optionally be prefixed with a letter 'r' or 'R'; such strings are called raw strings

and treat backslashes as literal characters. As a result, in string literals, '\U' and '\u' escapes in raw strings are not

treated specially. Given that Python 2.x's raw unicode literals behave differently than Python 3.x's the 'ur' syntax is

not supported.

New in version 3.3: The 'rb' prefix of raw bytes literals has been added as a synonym of 'br'.New

in version 3.3: Support for the unicode legacy literal (u'value') was reintroduced to simplify the

maintenance of dual Python 2.x and 3.x codebases.

In triple-quoted strings, unescaped newlines and quotes are allowed (and are retained), except that three unescaped

quotes in a row terminate the string. (A "quote" is the character used to open the string, i.e. either ' or ".)

Unless an 'r' or 'R' prefix is present, escape sequences in strings are interpreted according to rules similar to those

used by Standard C. The recognized escape sequences are:

**Escape Sequence Meaning Notes** 

\newline Backslash and newline ignored

\ Backslash (\)

\'Single quote (')

" Double quote (")

\a ASCII Bell (BEL)

\b ASCII Backspace (BS)

\f ASCII Formfeed (FF)

\n ASCII Linefeed (LF)

\r ASCII Carriage Return (CR)

\t ASCII Horizontal Tab (TAB)

\v ASCII Vertical Tab (VT)

\ooo Character with octal value ooo (1,3)

\xhh Character with hex value hh (2,3)

Escape sequences only recognized in string literals are:

**Escape Sequence Meaning Notes** 

\N{name} Character named name in the Unicode database (4)

\uxxxx Character with 16-bit hex value xxxx (5)

\Uxxxxxxxx Character with 32-bit hex value xxxxxxxx (6)

### Notes:

- 1. As in Standard C, up to three octal digits are accepted.
- 2. Unlike in Standard C, exactly two hex digits are required.
- 3. In a bytes literal, hexadecimal and octal escapes denote the byte with the given value. In a string literal, these

escapes denote a Unicode character with the given value.

- 4. Changed in version 3.3: Support for name aliases 1 has been added.
- 5. Individual code units which form parts of a surrogate pair can be encoded using this escape sequence. Exactly

four hex digits are required.

6. Any Unicode character can be encoded this way. Exactly eight hex digits are required.

Unlike Standard C, all unrecognized escape sequences are left in the string unchanged, i.e., the backslash is left in the

string. (This behavior is useful when debugging: if an escape sequence is mistyped, the resulting output is more easily

recognized as broken.) It is also important to note that the escape sequences only recognized in string literals fall into

the category of unrecognized escapes for bytes literals.

Even in a raw string, string quotes can be escaped with a backslash, but the backslash remains in the string; for

example, r"" is a valid string literal consisting of two characters: a backslash and a double quote; r"" is not a

valid string literal (even a raw string cannot end in an odd number of backslashes). Specifically, a raw string cannot

end in a single backslash (since the backslash would escape the following quote

character). Note also that a single

backslash followed by a newline is interpreted as those two characters as part of the string, not as a line continuation.

# 2.4.2 String literal concatenation

Multiple adjacent string or bytes literals (delimited by whitespace), possibly using different quoting conventions,

are allowed, and their meaning is the same as their concatenation. Thus, "hello" 'world' is equivalent to

"helloworld". This feature can be used to reduce the number of backslashes needed, to split long strings conveniently

across long lines, or even to add comments to parts of strings, for example:

1 http://www.unicode.org/Public/6.1.0/ucd/NameAliases.txt

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```
re.compile("[A-Za-z_]" # letter or underscore "[A-Za-z0-9_]*" # letter, digit or underscore
```

Note that this feature is defined at the syntactical level, but implemented at compile time. The '+' operator must be

used to concatenate string expressions at run time. Also note that literal concatenation can use different quoting styles

for each component (even mixing raw strings and triple quoted strings).

# 2.4.3 Numeric literals

There are three types of numeric literals: integers, floating point numbers, and imaginary numbers. There are no

complex literals (complex numbers can be formed by adding a real number and an imaginary number).

Note that numeric literals do not include a sign; a phrase like -1 is actually an expression composed of the unary

operator '-' and the literal 1.

# 2.4.4 Integer literals

Integer literals are described by the following lexical definitions:

```
integer ::= decimalinteger | octinteger | hexinteger | bininteger decimalinteger ::= nonzerodigit digit* | "0"+ nonzerodigit ::= "1"..."9" digit ::= "0"..."9" octinteger ::= "0" ("o" | "O") octdigit+ hexinteger ::= "0" ("x" | "X") hexdigit+ bininteger ::= "0" ("b" | "B") bindigit+ octdigit ::= "0"..."7" hexdigit ::= digit | "a"..."f" | "A"..."F" bindigit ::= "0" | "1"
```

There is no limit for the length of integer literals apart from what can be stored in available memory.

Note that leading zeros in a non-zero decimal number are not allowed. This is for disambiguation with C-style octal

literals, which Python used before version 3.0.

Some examples of integer literals:

7 2147483647 0o177 0b100110111

3 79228162514264337593543950336 0o377 0x100000000

79228162514264337593543950336 Oxdeadbeef

2.4.5 Floating point literals

Floating point literals are described by the following lexical definitions:

```
floatnumber ::= pointfloat | exponentfloat pointfloat ::= [intpart] fraction | intpart "." exponentfloat ::= (intpart | pointfloat) exponent intpart ::= digit+ fraction ::= "." digit+ exponent ::= ("e" | "E") ["+" | "-"] digit+
```

Note that the integer and exponent parts are always interpreted using radix 10. For example, 077e010 is legal, and

denotes the same number as 77e10. The allowed range of floating point literals is implementation-dependent. Some

examples of floating point literals:

3.14 10. .001 1e100 3.14e-10 0e0

Note that numeric literals do not include a sign; a phrase like -1 is actually an expression composed of the unary

operator - and the literal 1.

# 2.4.6 Imaginary literals

Imaginary literals are described by the following lexical definitions:

imagnumber ::= (floatnumber | intpart) ("j" | "J")

An imaginary literal yields a complex number with a real part of 0.0. Complex numbers are represented as a pair of

floating point numbers and have the same restrictions on their range. To create a complex number with a nonzero real

part, add a floating point number to it, e.g., (3+4j). Some examples of imaginary literals:

3.14j 10.j 10j .001j 1e100j 3.14e-10j

2.5 Operators

The following tokens are operators:

## 2.6 Delimiters

The following tokens serve as delimiters in the grammar:

The period can also occur in floating-point and imaginary literals. A sequence of three periods has a special meaning

as an ellipsis literal. The second half of the list, the augmented assignment operators, serve lexically as delimiters, but

also perform an operation.

The following printing ASCII characters have special meaning as part of other tokens or are otherwise significant to

the lexical analyzer:

'"#\

The following printing ASCII characters are not used in Python. Their occurrence outside string literals and comments

is an unconditional error:

\$?'

### 3.0 DATA MODEL

# 3.1 Objects, values and types

Objects are Python's abstraction for data. All data in a Python program is represented by objects or by relations

between objects. (In a sense, and in conformance to Von Neumann's model of a "stored program computer," code is

also represented by objects.)

Every object has an identity, a type and a value. An object's identity never changes once it has been created; you

may think of it as the object's address in memory. The 'is' operator compares the identity of two objects; the id()

function returns an integer representing its identity.

CPython implementation detail: For CPython, id(x) is the memory address where x is stored.

An object's type determines the operations that the object supports (e.g., "does it have a length?") and also defines the

possible values for objects of that type. The type() function returns an object's type (which is an object itself). Like

its identity, an object's type is also unchangeable. 1

The value of some objects can change. Objects whose value can change are said to be mutable; objects whose value is

unchangeable once they are created are called immutable. (The value of an immutable container object that contains a

reference to a mutable object can change when the latter's value is changed; however the container is still considered

immutable, because the collection of objects it contains cannot be changed. So, immutability is not strictly the same

as having an unchangeable value, it is more subtle.) An object's mutability is determined by its type; for instance,

numbers, strings and tuples are immutable, while dictionaries and lists are mutable.

Objects are never explicitly destroyed; however, when they become unreachable they may be garbage-collected. An

implementation is allowed to postpone garbage collection or omit it altogether — it

is a matter of implementation

quality how garbage collection is implemented, as long as no objects are collected that are still reachable.

CPython implementation detail: CPython currently uses a reference-counting scheme with (optional) delayed detection

of cyclically linked garbage, which collects most objects as soon as they become unreachable, but is not

guaranteed to collect garbage containing circular references. See the documentation of the gc module for information

on controlling the collection of cyclic garbage. Other implementations act differently and CPython may change. Do

not depend on immediate finalization of objects when they become unreachable (ex: always close files).

Note that the use of the implementation's tracing or debugging facilities may keep objects alive that would normally

be collectable. Also note that catching an exception with a 'try...except' statement may keep objects alive.

Some objects contain references to "external" resources such as open files or windows. It is understood that these

resources are freed when the object is garbage-collected, but since garbage collection is not guaranteed to happen,

such objects also provide an explicit way to release the external resource, usually a close() method. Programs are

strongly recommended to explicitly close such objects. The 'try...finally' statement and the 'with' statement

provide convenient ways to do this.

Some objects contain references to other objects; these are called containers. Examples of containers are tuples, lists

and dictionaries. The references are part of a container's value. In most cases, when we talk about the value of a

container, we imply the values, not the identities of the contained objects; however, when we talk about the mutability

of a container, only the identities of the immediately contained objects are implied. So, if an immutable container (like

a tuple) contains a reference to a mutable object, its value changes if that mutable object is changed.

Types affect almost all aspects of object behavior. Even the importance of object identity is affected in some sense:

for immutable types, operations that compute new values may actually return a reference to any existing object with

the same type and value, while for mutable objects this is not allowed. E.g., after a = 1; b = 1, a and b may or

may not refer to the same object with the value one, depending on the implementation, but after c = []; d = [],

c and d are guaranteed to refer to two different, unique, newly created empty lists. (Note that c = d = [] assigns

the same object to both c and d.)

3.2 The standard type hierarchy

Below is a list of the types that are built into Python. Extension modules (written in C, Java, or other languages,

depending on the implementation) can define additional types. Future versions of Python may add types to the type

hierarchy (e.g., rational numbers, efficiently stored arrays of integers, etc.), although such additions will often be

provided via the standard library instead.

Some of the type descriptions below contain a paragraph listing 'special attributes.' These are attributes that provide

access to the implementation and are not intended for general use. Their definition may change in the future.

None This type has a single value. There is a single object with this value. This object is accessed through the built-in

name None. It is used to signify the absence of a value in many situations, e.g., it is returned from functions

that don't explicitly return anything. Its truth value is false.

NotImplemented This type has a single value. There is a single object with this value. This object is accessed

through the built-in name NotImplemented. Numeric methods and rich comparison methods may return

this value if they do not implement the operation for the operands provided. (The interpreter will then try the

reflected operation, or some other fallback, depending on the operator.) Its truth value is true.

Ellipsis This type has a single value. There is a single object with this value. This object is accessed through the

literal ... or the built-in name Ellipsis. Its truth value is true.

numbers. Number These are created by numeric literals and returned as results by arithmetic operators and arithmetic

built-in functions. Numeric objects are immutable; once created their value never changes. Python

numbers are of course strongly related to mathematical numbers, but subject to the limitations of numerical

representation in computers.

Python distinguishes between integers, floating point numbers, and complex numbers:

numbers. Integral These represent elements from the mathematical set of integers (positive and negative).

There are two types of integers:

Integers (int)

These represent numbers in an unlimited range, subject to available (virtual) memory only. For

the purpose of shift and mask operations, a binary representation is assumed, and negative numbers

are represented in a variant of 2's complement which gives the illusion of an infinite string

of sign bits extending to the left.

Booleans (bool) These represent the truth values False and True. The two objects representing the

values False and True are the only Boolean objects. The Boolean type is a subtype of the integer type,

and Boolean values behave like the values 0 and 1, respectively, in almost all contexts, the exception

being that when converted to a string, the strings "False" or "True" are returned, respectively.

The rules for integer representation are intended to give the most meaningful interpretation of shift and

mask operations involving negative integers.

numbers.Real (float) These represent machine-level double precision floating point numbers. You are

at the mercy of the underlying machine architecture (and C or Java implementation) for the accepted range

and handling of overflow. Python does not support single-precision floating point numbers; the savings

in processor and memory usage that are usually the reason for using these is dwarfed by the overhead of

using objects in Python, so there is no reason to complicate the language with two kinds of floating point

numbers.

numbers.Complex (complex) These represent complex numbers as a pair of machine-level double precision

floating point numbers. The same caveats apply as for floating point numbers. The real and imaginary

parts of a complex number z can be retrieved through the read-only attributes z.real and z.imag.

Sequences These represent finite ordered sets indexed by non-negative numbers. The built-in function len() returns

the number of items of a sequence. When the length of a sequence is n, the index set contains the numbers 0, 1,

..., n-1. Item i of sequence a is selected by a[i].

Sequences also support slicing: a[i:j] selects all items with index k such that i <= k < j. When used as an

expression, a slice is a sequence of the same type. This implies that the index set is renumbered so that it starts

at 0.

Some sequences also support "extended slicing" with a third "step" parameter: a[i:j:k] selects all items of

a with index x where  $x = i + n^*k$ ,  $n \ge 0$  and  $i \le x \le j$ .

Sequences are distinguished according to their mutability:

Immutable sequences An object of an immutable sequence type cannot change once it is created. (If the object

contains references to other objects, these other objects may be mutable and may be changed; however,

the collection of objects directly referenced by an immutable object cannot change.)

The following types are immutable sequences:

Strings A string is a sequence of values that represent Unicode codepoints. All the codepoints in range

U+0000 - U+10FFFF can be represented in a string. Python doesn't have a chr type, and every

character in the string is represented as a string object with length 1. The built-in function ord() converts

a character to its codepoint (as an integer); chr() converts an integer in range 0 - 10FFFF

to the corresponding character. str.encode() can be used to convert a str to bytes using the

given encoding, and bytes.decode() can be used to achieve the opposite.

Tuples The items of a tuple are arbitrary Python objects. Tuples of two or more items are formed by

comma-separated lists of expressions. A tuple of one item (a 'singleton') can be formed by affixing

a comma to an expression (an expression by itself does not create a tuple, since parentheses must be

usable for grouping of expressions). An empty tuple can be formed by an empty pair of parentheses.

Bytes A bytes object is an immutable array. The items are 8-bit bytes, represented by integers in the

range 0 <= x < 256. Bytes literals (like b'abc') and the built-in function bytes() can be used to

construct bytes objects. Also, bytes objects can be decoded to strings via the decode() method.

Mutable sequences Mutable sequences can be changed after they are created. The subscription and slicing

notations can be used as the target of assignment and del (delete) statements.

There are currently two intrinsic mutable sequence types:

Lists The items of a list are arbitrary Python objects. Lists are formed by placing a comma-separated list

of expressions in square brackets. (Note that there are no special cases needed to form lists of length

0 or 1.)

Byte Arrays A bytearray object is a mutable array. They are created by the built-in bytearray()

constructor. Aside from being mutable (and hence unhashable), byte arrays otherwise provide the

same interface and functionality as immutable bytes objects.

## 3.2. The standard type hierarchy

The extension module array provides an additional example of a mutable sequence type, as does the

collections module.

Set types These represent unordered, finite sets of unique, immutable objects. As such, they cannot be indexed

by any subscript. However, they can be iterated over, and the built-in function len() returns the number of

items in a set. Common uses for sets are fast membership testing, removing duplicates from a sequence, and

computing mathematical operations such as intersection, union, difference, and symmetric difference.

For set elements, the same immutability rules apply as for dictionary keys. Note that numeric types obey the

normal rules for numeric comparison: if two numbers compare equal (e.g., 1 and 1.0), only one of them can

be contained in a set.

There are currently two intrinsic set types:

Sets These represent a mutable set. They are created by the built-in set() constructor and can be modified

afterwards by several methods, such as add().

Frozen sets These represent an immutable set. They are created by the built-in frozenset() constructor.

As a frozenset is immutable and hashable, it can be used again as an element of another set, or as a

dictionary key.

Mappings These represent finite sets of objects indexed by arbitrary index sets. The subscript notation a[k] selects

the item indexed by k from the mapping a; this can be used in expressions and as the target of assignments or

del statements. The built-in function len() returns the number of items in a mapping.

There is currently a single intrinsic mapping type:

Dictionaries These represent finite sets of objects indexed by nearly arbitrary values. The only types of values

not acceptable as keys are values containing lists or dictionaries or other mutable types that are compared

by value rather than by object identity, the reason being that the efficient implementation of dictionaries

requires a key's hash value to remain constant. Numeric types used for keys obey the normal rules for numeric

comparison: if two numbers compare equal (e.g., 1 and 1.0) then they can be used interchangeably

to index the same dictionary entry.

Dictionaries are mutable; they can be created by the {...} notation (see section Dictionary displays).

The extension modules dbm.ndbm and dbm.gnu provide additional examples of mapping types, as does

the collections module.

Callable types These are the types to which the function call operation (see section Calls) can be applied:

User-defined functions A user-defined function object is created by a function definition (see section Function

definitions). It should be called with an argument list containing the same number of items as the function's

formal parameter list.

Special attributes:

Attribute Meaning
doc The function's documentation string, or None if unavailable Writable
name The function's name Writable
qualname The function's qualified name New in version 3.3. Writable
module. The name of the module the function was defined in or None

if unavailable.
Writable
defaults A tuple containing default argument values for those
arguments that have defaults, or None if no arguments have a
default value
Writable
code The code object representing the compiled function body. Writable
globals A reference to the dictionary that holds the function's global
variables— the global namespace of the module in which the
function was defined.
Read-only
dict The namespace supporting arbitrary function attributes. Writable
closure None or a tuple of cells that contain bindings for the function's
free variables.
Read-only
annotations A dict containing annotations of parameters. The keys of the
dict are the parameter names, or 'return' for the return
annotation, if provided.
Writable
kwdefaults A dict containing defaults for keyword-only parameters. Writable
Most of the attributes labelled "Writable" check the type of the assigned value.
Function objects also support getting and setting arbitrary attributes, which can be used, for example, to
attach metadata to functions. Regular attribute dot-notation is used to get and set such attributes. Note that
the current implementation only supports function attributes on user-defined functions. Function attributes
on built-in functions may be supported in the future.
Additional information about a function's definition can be retrieved from its code object; see the description
of internal types below.

Instance methods An instance method object combines a class, a class instance and any callable object (normally

a user-defined function).
Special read-only attributes:self is the class instance object,func is the function object;
doc is the method's documentation (same asfuncdoc);name is the method
name (same asfuncname);module is the name of the module the method was
defined in, or None if unavailable.
Methods also support accessing (but not setting) the arbitrary function attributes on the underlying function
object.
User-defined method objects may be created when getting an attribute of a class (perhaps via an instance
of that class), if that attribute is a user-defined function object or a class method object.
When an instance method object is created by retrieving a user-defined function object from a class via
one of its instances, itsself attribute is the instance, and the method object is said to be bound.
The new method'sfunc attribute is the original function object.
When a user-defined method object is created by retrieving another method object from a class or instance,
the behaviour is the same as for a function object, except that thefunc attribute of the new instance
is not the original method object but itsfunc attribute.
When an instance method object is created by retrieving a class method object from a class or instance, its
self attribute is the class itself, and itsfunc attribute is the function object underlying the
class method.
When an instance method object is called, the underlying function (func) is called, inserting the
class instance (self) in front of the argument list. For instance, when C is a class which contains a

definition for a function $f()$ , and $x$ is an instance of $C$ , calling $x.f(1)$ is equivalent to calling $C.f(x)$ ,
1).
When an instance method object is derived from a class method object, the "class instance" stored in
self will actually be the class itself, so that calling either x.f(1) or C.f(1) is equivalent to
calling f(C,1) where f is the underlying function.
Note that the transformation from function object to instance method object happens each time the attribute
is retrieved from the instance. In some cases, a fruitful optimization is to assign the attribute to a local
variable and call that local variable. Also notice that this transformation only happens for user-defined
functions; other callable objects (and all non-callable objects) are retrieved without transformation. It is
also important to note that user-defined functions which are attributes of a class instance are not converted
to bound methods; this only happens when the function is an attribute of the class.
Generator functions A function or method which uses the yield statement (see section The yield statement)
is called a generator function. Such a function, when called, always returns an iterator object which can
be used to execute the body of the function: calling the iterator's iteratornext() method will
cause the function to execute until it provides a value using the yield statement. When the function
executes a return statement or falls off the end, a StopIteration exception is raised and the iterator
will have reached the end of the set of values to be returned.
Built-in functions A built-in function object is a wrapper around a C function. Examples of built-in functions
are len() and math.sin() (math is a standard built-in module). The number and type of the
arguments are determined by the C function. Special read-only attributes: doc is the function's documentation

string, or None if unavailable; \_\_name\_\_ is the function's name; \_\_self\_\_ is set to None (but see the next item); \_\_module\_\_ is the name of the module the function was defined in or None if unavailable. Built-in methods This is really a different disguise of a built-in function, this time containing an object passed to the C function as an implicit extra argument. An example of a built-in method is alist.append(), assuming alist is a list object. In this case, the special read-only attribute self is set to the object denoted by alist. Classes Classes are callable. These objects normally act as factories for new instances of themselves, but variations are possible for class types that override new (). The arguments of the call are passed to \_new\_\_() and, in the typical case, to \_\_init\_\_() to initialize the new instance. Class Instances Instances of arbitrary classes can be made callable by defining a call () method in their class. Modules Modules are a basic organizational unit of Python code, and are created by the import system as invoked either by the import statement (see import), or by calling functions such as importlib.import module() and built-in import (). A module object has a namespace implemented by a dictionary object (this is the dictionary referenced by the globals attribute of functions defined in the module). Attribute references are translated to lookups in this dictionary, e.g., m.x is equivalent to m. dict ["x"]. A module object does not contain the code object used to initialize the module (since it isn't needed once the initialization is done). Attribute assignment updates the module's namespace dictionary, e.g., m.x = 1 is equivalent to m.\_\_dict\_\_["x"] = 1.

Special read-only attribute: \_\_dict\_\_ is the module's namespace as a dictionary object. CPython implementation detail: Because of the way CPython clears module dictionaries, the module dictionary will be cleared when the module falls out of scope even if the dictionary still has live references. To avoid this, copy the dictionary or keep the module around while using its dictionary directly. Predefined (writable) attributes: name is the module's name; doc is the module's documentation string, or None if unavailable; file is the pathname of the file from which the module was loaded, if it was loaded from a file. The \_\_file\_\_ attribute may be missing for certain types of modules, such as C modules that are statically linked into the interpreter; for extension modules loaded dynamically from a shared library, it is the pathname of the shared library file. Custom classes Custom class types are typically created by class definitions (see section Class definitions). A class has a namespace implemented by a dictionary object. Class attribute references are translated to lookups in this dictionary, e.g., C.x is translated to C. dict ["x"] (although there are a number of hooks which allow for other means of locating attributes). When the attribute name is not found there, the attribute search continues in the base classes. This search of the base classes uses the C3 method resolution order which behaves correctly even in the presence of 'diamond' inheritance structures where there are multiple inheritance paths leading back to a common ancestor. Additional details on the C3 MRO used by Python can be found in the documentation accompanying the 2.3 release at http://www.python.org/download/releases/2.3/mro/. When a class attribute reference (for class C, say) would yield a class method object, it is transformed into an instance method object whose \_\_self\_\_ attributes is C. When it would yield a static method object, it

is transformed into the object wrapped by the static method object. See section Implementing Descriptors for
another way in which attributes retrieved from a class may differ from those actually contained in itsdict
Class attribute assignments update the class's dictionary, never the dictionary of a base class.
A class object can be called (see above) to yield a class instance (see below).
Special attributes:name is the class name;module is the module name in which the class was
defined;dict is the dictionary containing the class's namespace;bases is a tuple (possibly empty
or a singleton) containing the base classes, in the order of their occurrence in the base class list;doc is
the class's documentation string, or None if undefined.
Class instances A class instance is created by calling a class object (see above). A class instance has a namespace
implemented as a dictionary which is the first place in which attribute references are searched. When an
attribute is not found there, and the instance's class has an attribute by that name, the search continues with
the class attributes. If a class attribute is found that is a user-defined function object, it is transformed into an
instance method object whoseself attribute is the instance. Static method and class method objects
are also transformed; see above under "Classes". See section Implementing Descriptors for another way in
which attributes of a class retrieved via its instances may differ from the objects actually stored in the class's
dict If no class attribute is found, and the object's class has agetattr() method, that is called
to satisfy the lookup.
Attribute assignments and deletions update the instance's dictionary, never a class's dictionary. If the class has
asetattr() ordelattr() method, this is called instead of updating the instance dictionary
directly.

Class instances can pretend to be numbers, sequences, or mappings if they have methods with certain special

names. See section Special method names.

Special attributes: \_\_dict\_\_ is the attribute dictionary; \_\_class\_\_ is the instance's class.

I/O objects (also known as file objects) A file object represents an open file. Various shortcuts are available to create

file objects: the open() built-in function, and also os.popen(), os.fdopen(), and the makefile()

method of socket objects (and perhaps by other functions or methods provided by extension modules).

The objects sys.stdin, sys.stdout and sys.stderr are initialized to file objects corresponding to the

interpreter's standard input, output and error streams; they are all open in text mode and therefore follow the

interface defined by the io. TextIOBase abstract class.

Internal types A few types used internally by the interpreter are exposed to the user. Their definitions may change

with future versions of the interpreter, but they are mentioned here for completeness.

Code objects Code objects represent byte-compiled executable Python code, or bytecode. The difference

between a code object and a function object is that the function object contains an explicit reference to

the function's globals (the module in which it was defined), while a code object contains no context:

also the default argument values are stored in the function object, not in the code object (because they

represent values calculated at run-time). Unlike function objects, code objects are immutable and contain

no references (directly or indirectly) to mutable objects.

Special read-only attributes: co\_name gives the function name; co\_argcount is the number of positional

arguments (including arguments with default values); co\_nlocals is the number of local variables

used by the function (including arguments); co\_varnames is a tuple containing the

names of the local

variables (starting with the argument names); co\_cellvars is a tuple containing the names of local variables that are referenced by nested functions; co\_freevars is a tuple containing the names of free

variables; co\_code is a string representing the sequence of bytecode instructions; co\_consts is a tuple

containing the literals used by the bytecode; co\_names is a tuple containing the names used by the

bytecode; co\_filename is the filename from which the code was compiled; co\_firstlineno is the

first line number of the function; co\_Inotab is a string encoding the mapping from bytecode offsets to

line numbers (for details see the source code of the interpreter); co\_stacksize is the required stack

size (including local variables); co\_flags is an integer encoding a number of flags for the interpreter.

The following flag bits are defined for co\_flags: bit 0x04 is set if the function uses the \*arguments

syntax to accept an arbitrary number of positional arguments; bit 0x08 is set if the function uses the

\*\*keywords syntax to accept arbitrary keyword arguments; bit 0x20 is set if the function is a generator.

Future feature declarations (from \_\_future\_\_ import division) also use bits in co\_flags to

indicate whether a code object was compiled with a particular feature enabled: bit 0x2000 is set if the

function was compiled with future division enabled; bits 0x10 and 0x1000 were used in earlier versions

of Python.

Other bits in co\_flags are reserved for internal use.

If a code object represents a function, the first item in co\_consts is the documentation string of the

function, or None if undefined.

Frame objects Frame objects represent execution frames. They may occur in traceback objects (see below).

Special read-only attributes: f\_back is to the previous stack frame (towards the

caller), or None if

this is the bottom stack frame; f\_code is the code object being executed in this frame; f\_locals is the

dictionary used to look up local variables; f\_globals is used for global variables; f\_builtins is used

for built-in (intrinsic) names; f\_lasti gives the precise instruction (this is an index into the bytecode

string of the code object).

Special writable attributes: f\_trace, if not None, is a function called at the start of each source code

line (this is used by the debugger); f\_lineno is the current line number of the frame — writing to this

from within a trace function jumps to the given line (only for the bottom-most frame). A debugger can

implement a Jump command (aka Set Next Statement) by writing to f\_lineno.

Traceback objects Traceback objects represent a stack trace of an exception. A traceback object is created

when an exception occurs. When the search for an exception handler unwinds the execution stack, at

each unwound level a traceback object is inserted in front of the current traceback. When an exception

handler is entered, the stack trace is made available to the program. (See section The try statement.) It is

accessible as the third item of the tuple returned by sys.exc\_info(). When the program contains no

suitable handler, the stack trace is written (nicely formatted) to the standard error stream; if the interpreter

is interactive, it is also made available to the user as sys.last\_traceback.

Special read-only attributes: tb\_next is the next level in the stack trace (towards the frame where the

exception occurred), or None if there is no next level; tb\_frame points to the execution frame of the

current level; tb\_lineno gives the line number where the exception occurred; tb\_lasti indicates the

precise instruction. The line number and last instruction in the traceback may differ from the line number

of its frame object if the exception occurred in a try statement with no matching except clause or with a

finally clause.

Slice objects Slice objects are used to represent slices for \_\_getitem\_\_() methods. They are also created

by the built-in slice() function. Special read-only attributes: start is the lower bound; stop is the upper bound; step is the step value;

each is None if omitted. These attributes can have any type.

Slice objects support one method:

slice.indices(self, length)

This method takes a single integer argument length and computes information about the slice that the

slice object would describe if applied to a sequence of length items. It returns a tuple of three integers;

respectively these are the start and stop indices and the step or stride length of the slice. Missing or

out-of-bounds indices are handled in a manner consistent with regular slices.

Static method objects Static method objects provide a way of defeating the transformation of function objects

to method objects described above. A static method object is a wrapper around any other object, usually a

user-defined method object. When a static method object is retrieved from a class or a class instance, the

object actually returned is the wrapped object, which is not subject to any further transformation. Static

method objects are not themselves callable, although the objects they wrap usually are. Static method

objects are created by the built-in staticmethod() constructor.

Class method objects A class method object, like a static method object, is a wrapper around another object

that alters the way in which that object is retrieved from classes and class instances. The behaviour of

class method objects upon such retrieval is described above, under "User-defined methods". Class method

objects are created by the built-in classmethod() constructor.

#### 3.3 Special method names

A class can implement certain operations that are invoked by special syntax (such as arithmetic operations or subscripting

and slicing) by defining methods with special names. This is Python's approach to operator overloading,

allowing classes to define their own behavior with respect to language operators. For instance, if a class defines

a method named  $\_$  getitem $\_$ (), and x is an instance of this class, then x[i] is roughly equivalent to

type(x).\_\_getitem\_\_(x, i). Except where mentioned, attempts to execute an operation raise an exception

when no appropriate method is defined (typically AttributeError or TypeError).

When implementing a class that emulates any built-in type, it is important that the emulation only be implemented

to the degree that it makes sense for the object being modelled. For example, some sequences may work well with

retrieval of individual elements, but extracting a slice may not make sense. (One example of this is the NodeList

interface in the W3C's Document Object Model.)

#### 3.3.1 Basic customization

object	_new_	_(cls[,		])
--------	-------	---------	--	----

Called to create a new instance of class cls. \_\_new\_\_() is a static method (special-cased so you need not

declare it as such) that takes the class of which an instance was requested as its first argument. The remaining

arguments are those passed to the object constructor expression (the call to the class). The return value of

\_\_new\_\_() should be the new object instance (usually an instance of cls).

Typical implementations create a new instance of the class by invoking the superclass's \_\_new\_\_() method

using super(currentclass, cls).\_\_new\_\_(cls[, ...]) with appropriate arguments and then

modifying the newly-created instance as necessary before returning it.

Ifnew() returns an instance of cls, then the new instance'sinit() method will be invoked like
init(self[,]), where self is the new instance and the remaining arguments are the same as
were passed tonew().
Ifnew() does not return an instance of cls, then the new instance'sinit() method will not be invoked.
3.3. Special method names
new() is intended mainly to allow subclasses of immutable types (like int, str, or tuple) to customize
instance creation. It is also commonly overridden in custom metaclasses in order to customize class creation.
objectinit(self [, ])
Called when the instance is created. The arguments are those passed to the class constructor expression.
If a base class has aninit() method, the derived class'sinit() method, if any,
must explicitly call it to ensure proper initialization of the base class part of the instance; for example:
BaseClassinit(self, [args]). As a special constraint on constructors, no value may be
returned; doing so will cause a TypeError to be raised at runtime.
objectdel(self)
Called when the instance is about to be destroyed. This is also called a destructor. If a base class has a
del() method, the derived class'sdel() method, if any, must explicitly call it to ensure proper
deletion of the base class part of the instance. Note that it is possible (though not recommended!) for the
del() method to postpone destruction of the instance by creating a new reference to it. It may then be
called at a later time when this new reference is deleted. It is not guaranteed thatdel() methods are
called for objects that still exist when the interpreter exits.

Note: del x doesn't directly call x. \_\_del\_\_() — the former decrements the reference count for x by one, and the latter is only called when x's reference count reaches zero. Some common situations that may prevent the reference count of an object from going to zero include: circular references between objects (e.g., a doubly-linked list or a tree data structure with parent and child pointers); a reference to the object on the stack frame of a function that caught an exception (the traceback stored in sys.exc info()[2] keeps the stack frame alive); or a reference to the object on the stack frame that raised an unhandled exception in interactive mode (the traceback stored in sys.last traceback keeps the stack frame alive). The first situation can only be remedied by explicitly breaking the cycles; the latter two situations can be resolved by storing None in sys.last traceback. Circular references which are garbage are detected when the option cycle detector is enabled (it's on by default), but can only be cleaned up if there are no Pythonlevel del () methods involved. Refer to the documentation for the gc module for more information about how del () methods are handled by the cycle detector, particularly the description of the garbage value. Warning: Due to the precarious circumstances under which del () methods are invoked, exceptions that occur during their execution are ignored, and a warning is printed to sys.stderr instead. Also, when del () is invoked in response to a module being deleted (e.g., when execution of the program is done), other globals referenced by the \_\_\_del\_\_\_() method may already have been deleted or in the process of being torn down (e.g. the import machinery shutting down). For this reason, del () methods should do the absolute minimum needed to maintain external invariants. Starting with version 1.5, Python

guarantees that globals whose name begins with a single underscore are deleted

from their module before other globals are deleted; if no other references to such globals exist, this may help in assuring that imported modules are still available at the time when the del () method is called. object. repr\_(self) Called by the repr() built-in function to compute the "official" string representation of an object. If at all possible, this should look like a valid Python expression that could be used to recreate an object with the same value (given an appropriate environment). If this is not possible, a string of the form <...some useful description...> should be returned. The return value must be a string object. If a class defines repr () but not str (), then repr () is also used when an "informal" string representation of instances of that class is required. This is typically used for debugging, so it is important that the representation is information-rich and unambiguous. object. str (self) Called by str(object) and the built-in functions format() and print() to compute the "informal" or nicely printable string representation of an object. The return value must be a string object. This method differs from object. repr () in that there is no expectation that str () return a valid Python expression: a more convenient or concise representation can be used. The default implementation defined by the built-in type object calls object. repr (). object.\_\_bytes\_\_(self) Called by bytes() to compute a byte-string representation of an object. This should return a bytes object. object. format (self, format spec)

Called by the format() built-in function (and by extension, the str.format() method

of class str) to
produce a "formatted" string representation of an object. The format_spec argument is a string that contains
a description of the formatting options desired. The interpretation of the format_spec argument is up to the
type implementingformat(), however most classes will either delegate formatting to one of the built-in
types, or use a similar formatting option syntax.
See formatspec for a description of the standard formatting syntax.
The return value must be a string object.
objectlt(self, other)
objectle(self, other)
objecteq(self, other)
objectne(self, other)
objectgt(self, other)
objectge(self, other)
These are the so-called "rich comparison" methods. The correspondence between operator symbols and method
names is as follows: x <y calls="" td="" x="=y" x<="y" xeq(y),<="" xle(y),="" xlt(y),=""></y>
x!=y calls xne(y), x>y calls xgt(y), and x>=y calls xge(y).
A rich comparison method may return the singleton NotImplemented if it does not implement the operation
for a given pair of arguments. By convention, False and True are returned for a successful comparison.
However, these methods can return any value, so if the comparison operator is used in a Boolean context (e.g.,
in the condition of an if statement), Python will call bool() on the value to determine if the result is true or
false.
There are no implied relationships among the comparison operators. The truth of x==y does not imply that
x!=y is false. Accordingly, when definingeq(), one should also definene() so that the operators
will behave as expected. See the paragraph onhash() for some important

notes on creating hashable
objects which support custom comparison operations and are usable as dictionary keys.
There are no swapped-argument versions of these methods (to be used when the left argument does not support
the operation but the right argument does); rather,lt() andgt() are each other's reflection,
le() andge() are each other's reflection, andeq() andne() are their own reflection.
Arguments to rich comparison methods are never coerced.
To automatically generate ordering operations from a single root operation, see
functools.total_ordering().
objecthash(self)
Called by built-in function hash() and for operations on members of hashed collections including set,
frozenset, and dicthash() should return an integer. The only required property is that objects
which compare equal have the same hash value; it is advised to somehow mix together (e.g. using exclusive or)
the hash values for the components of the object that also play a part in comparison of objects.
If a class does not define aneq() method it should not define ahash() operation either; if it
defineseq() but nothash(), its instances will not be usable as items in hashable collections. If a
class defines mutable objects and implements aneq() method, it should not implementhash(),
3.3. Special method names
since the implementation of hashable collections requires that a key's hash value is immutable (if the object's
hash value changes, it will be in the wrong hash bucket).
User-defined classes haveeq() andhash() methods by default; with them, all objects compare
unequal (except with themselves) and xhash() returns an appropriate value

such that x == y implies
both that x is y and $hash(x) == hash(y)$ .
A class that overrideseq() and does not definehash() will have itshash() implicitly set
to None. When thehash() method of a class is None, instances of the class will raise an appropriate
TypeError when a program attempts to retrieve their hash value, and will also be correctly identified as
unhashable when checking isinstance(obj, collections.Hashable).
If a class that overrideseq() needs to retain the implementation ofhash() from a parent class,
the interpreter must be told this explicitly by settinghash = <parentclass>hash</parentclass>
If a class that does not overrideeq() wishes to suppress hash support, it should includehash =
None in the class definition. A class which defines its ownhash() that explicitly raises a TypeError
would be incorrectly identified as hashable by an isinstance(obj, collections.Hashable) call.
Note: By default, thehash() values of str, bytes and datetime objects are "salted" with an unpredictable
random value. Although they remain constant within an individual Python process, they are not predictable
between repeated invocations of Python.
This is intended to provide protection against a denial-of-service caused by carefully-chosen inputs that exploit
the worst case performance of a dict insertion, O(n^2) complexity. See http://www.ocert.org/advisories/ocert-2011-003.html for details.
Changing hash values affects the iteration order of dicts, sets and other mappings. Python has never made
guarantees about this ordering (and it typically varies between 32-bit and 64-bit builds).
Hash randomization is enabled by default.
objectbool(self)
Called to implement truth value testing and the built-in operation bool(); should return False or True.

When this method is not defined,len() is called, if it is defined, and the object is considered true if its
result is nonzero. If a class defines neitherlen() norbool(), all its instances are considered true.
3.3.2 Customizing attribute access
The following methods can be defined to customize the meaning of attribute access (use of, assignment to, or deletion
of x.name) for class instances.
objectgetattr(self, name)
Called when an attribute lookup has not found the attribute in the usual places (i.e. it is not an instance attribute
nor is it found in the class tree for self). name is the attribute name. This method should return the (computed)
attribute value or raise an AttributeError exception.
Note that if the attribute is found through the normal mechanism,getattr() is not called. (This is an
intentional asymmetry betweengetattr() andsetattr().) This is done both for efficiency
reasons and because otherwisegetattr() would have no way to access other attributes of the instance.
Note that at least for instance variables, you can fake total control by not inserting any values in the instance
attribute dictionary (but instead inserting them in another object). See thegetattribute() method
below for a way to actually get total control over attribute access.
objectgetattribute(self, name)
Called unconditionally to implement attribute accesses for instances of the class.
If the class also defines
getattr(), the latter will not be called unlessgetattribute() either calls it explicitly
or raises an AttributeError. This method should return the (computed) attribute value or raise
an AttributeError exception. In order to avoid infinite recursion in this method, its

implementation

should always call the base class method with the same name to access any attributes it needs, for example,

object. getattribute (self, name).

Note: This method may still be bypassed when looking up special methods as the result of implicit invocation

via language syntax or built-in functions. See Special method lookup.

object. setattr (self, name, value)

Called when an attribute assignment is attempted. This is called instead of the normal mechanism (i.e. store the

value in the instance dictionary). name is the attribute name, value is the value to be assigned to it.

If \_\_setattr\_\_() wants to assign to an instance attribute, it should call the base class method with the same

name, for example, object. \_\_setattr\_\_(self, name, value).

object.\_\_delattr\_\_(self, name)

Like \_\_setattr\_\_() but for attribute deletion instead of assignment. This should only be implemented if

del obj.name is meaningful for the object.

object.\_\_dir\_\_(self)

Called when dir() is called on the object. A sequence must be returned. dir() converts the returned

sequence to a list and sorts it.

Implementing Descriptors

The following methods only apply when an instance of the class containing the method (a so-called descriptor class)

appears in an owner class (the descriptor must be in either the owner's class dictionary or in the class dictionary for

one of its parents). In the examples below, "the attribute" refers to the attribute whose name is the key of the property

in the owner class' \_\_dict\_\_.

object. get (self, instance, owner)

Called to get the attribute of the owner class (class attribute access) or of an instance of that class (instance

attribute access). owner is always the owner class, while instance is the instance

that the attribute was accessed

through, or None when the attribute is accessed through the owner. This method should return the (computed)

attribute value or raise an AttributeError exception.

object.\_\_set\_\_(self, instance, value)

Called to set the attribute on an instance instance of the owner class to a new value, value.

object.\_\_delete\_\_(self, instance)

Called to delete the attribute on an instance instance of the owner class.

**Invoking Descriptors** 

In general, a descriptor is an object attribute with "binding behavior", one whose attribute access has been overridden

by methods in the descriptor protocol: \_\_get\_\_(), \_\_set\_\_(), and \_\_delete\_\_(). If any of those methods

are defined for an object, it is said to be a descriptor.

The default behavior for attribute access is to get, set, or delete the attribute from an object's dictionary. For instance,

a.x has a lookup chain starting with a. \_\_dict\_\_['x'], then type(a). \_\_dict\_\_['x'], and continuing

through the base classes of type(a) excluding metaclasses.

However, if the looked-up value is an object defining one of the descriptor methods, then Python may override the

default behavior and invoke the descriptor method instead. Where this occurs in the precedence chain depends on

which descriptor methods were defined and how they were called.

The starting point for descriptor invocation is a binding, a.x. How the arguments are assembled depends on a:

Direct Call The simplest and least common call is when user code directly invokes a descriptor method:

x.\_\_get\_\_(a).

Instance Binding If binding to an object instance, a.x is transformed into the call:

type(a).\_\_dict\_\_['x'].\_\_get\_\_(a, type(a)).

Class Binding If binding to a class, A.x is transformed into the call:

A.\_\_dict\_\_['x'].\_\_get\_\_(None,

A).
Super Binding If a is an instance of super, then the binding super(B, obj).m() searches
objclassmro for the base class A immediately preceding B and then invokes the descriptor
with the call: Adict['m']get(obj, objclass).
For instance bindings, the precedence of descriptor invocation depends on the which descriptor methods are defined.
A descriptor can define any combination ofget(),set() anddelete(). If it does not define
get(), then accessing the attribute will return the descriptor object itself unless there is a value in the object's
instance dictionary. If the descriptor definesset() and/ordelete(), it is a data descriptor; if
it defines neither, it is a non-data descriptor. Normally, data descriptors define bothget() andset(),
while non-data descriptors have just theget() method. Data descriptors withset() andget()
defined always override a redefinition in an instance dictionary. In contrast, non- data descriptors can be overridden by
instances.
Python methods (including staticmethod() and classmethod()) are implemented as non-data descriptors.
Accordingly, instances can redefine and override methods. This allows individual instances to acquire behaviors that
differ from other instances of the same class.
The property() function is implemented as a data descriptor. Accordingly, instances cannot override the behavior
of a property.
slots
By default, instances of classes have a dictionary for attribute storage. This wastes space for objects having very few
instance variables. The space consumption can become acute when creating large numbers of instances.
The default can be overridden by definingslots in a class definition. The slotsdeclaration takes a sequence

of instance variables and reserves just enough space in each instance to hold a value for each variable. Space is saved
becausedict is not created for each instance.
objectslots
This class variable can be assigned a string, iterable, or sequence of strings with variable names used by instances.
If defined in a class,slots reserves space for the declared variables and prevents the automatic
creation ofdict andweakref for each instance.
Notes on usingslots
<ul> <li>When inheriting from a class withoutslots, thedict attribute of that class will always be accessible, so</li> </ul>
aslots definition in the subclass is meaningless.
<ul> <li>Without adict variable, instances cannot be assigned new variables not listed in theslots definition.</li> </ul>
Attempts to assign to an unlisted variable name raises AttributeError. If dynamic assignment of new
variables is desired, then add 'dict' to the sequence of strings in theslots declaration.
<ul> <li>Without aweakref variable for each instance, classes definingslots do</li> </ul>
not support weak references to
its instances. If weak reference support is needed, then add 'weakref' to the sequence of strings in the
slots declaration.
<ul> <li>slots are implemented at the class level by creating descriptors (Implementing Descriptors) for each variable</li> </ul>
name. As a result, class attributes cannot be used to set default values for instance variables defined byslots;
otherwise, the class attribute would overwrite the descriptor assignment.
• The action of aslots declaration is limited to the class where it is defined. As a result, subclasses will have
adict unless they also defineslots (which must only contain names of any additional slots).

• If a class defines a slot also defined in a base class, the instance variable defined by the base class slot is

inaccessible (except by retrieving its descriptor directly from the base class). This renders the meaning of the

program undefined. In the future, a check may be added to prevent this.

 Nonempty \_\_slots\_\_ does not work for classes derived from "variable-length" built-in types such as int, str

and tuple.

• Any non-string iterable may be assigned to \_\_slots\_\_. Mappings may also be used; however, in the future,

special meaning may be assigned to the values corresponding to each key.

\_\_class\_\_ assignment works only if both classes have the same \_\_slots\_\_.

### 3.3.3 Customizing class creation

By default, classes are constructed using type(). The class body is executed in a new namespace and the class name

is bound locally to the result of type(name, bases, namespace).

The class creation process can be customised by passing the metaclass keyword argument in the class definition

line, or by inheriting from an existing class that included such an argument. In the following example, both MyClass

and MySubclass are instances of Meta:

class Meta(type):

pass

class MyClass(metaclass=Meta):

pass

class MySubclass(MyClass):

pass

Any other keyword arguments that are specified in the class definition are passed through to all metaclass operations

described below.

When a class definition is executed, the following steps occur:

- the appropriate metaclass is determined
- the class namespace is prepared

- the class body is executed
- the class object is created

Determining the appropriate metaclass

The appropriate metaclass for a class definition is determined as follows:

- if no bases and no explicit metaclass are given, then type() is used
- if an explicit metaclass is given and it is not an instance of type(), then it is used directly as the metaclass
- if an instance of type() is given as the explicit metaclass, or bases are defined, then the most derived metaclass

is used

The most derived metaclass is selected from the explicitly specified metaclass (if any) and the metaclasses (i.e.

type(cls)) of all specified base classes. The most derived metaclass is one which is a subtype of all of these

candidate metaclasses. If none of the candidate metaclasses meets that criterion, then the class definition will fail with

TypeError.

Preparing the class namespace

Once the appropriate metaclass has been identified, then the class namespace is prepared. If the metaclass

has a \_\_prepare\_\_ attribute, it is called as namespace = metaclass.\_\_prepare\_\_(name, bases,

\*\*kwds) (where the additional keyword arguments, if any, come from the class definition).

If the metaclass has no \_\_prepare\_\_ attribute, then the class namespace is initialised as an empty dict() instance.

Metaclasses in Python 3000 Introduced the \_\_prepare\_\_ namespace hook

Executing the class body

The class body is executed (approximately) as exec(body, globals(), namespace). The key difference

from a normal call to exec() is that lexical scoping allows the class body (including any methods) to reference

names from the current and outer scopes when the class definition occurs inside a function.

However, even when the class definition occurs inside the function, methods

defined inside the class still cannot see

names defined at the class scope. Class variables must be accessed through the first parameter of instance or class

methods, and cannot be accessed at all from static methods.

Creating the class object

Once the class namespace has been populated by executing the class body, the class object is created by calling

metaclass(name, bases, namespace, \*\*kwds) (the additional keywords passed here are the same as

those passed to \_\_prepare\_\_).

This class object is the one that will be referenced by the zero-argument form of super(). \_\_class\_\_ is an implicit

closure reference created by the compiler if any methods in a class body refer to either \_\_class\_\_ or super. This

allows the zero argument form of super() to correctly identify the class being defined based on lexical scoping,

while the class or instance that was used to make the current call is identified based on the first argument passed to the

method.

After the class object is created, it is passed to the class decorators included in the class definition (if any) and the

resulting object is bound in the local namespace as the defined class.

New super Describes the implicit \_\_class\_\_ closure reference

Metaclass example

The potential uses for metaclasses are boundless. Some ideas that have been explored include logging, interface

checking, automatic delegation, automatic property creation, proxies, frameworks, and automatic resource locking/

synchronization.

Here is an example of a metaclass that uses an collections. Ordered Dict to remember the order that class members were defined: class OrderedClass(type): @classmethod def prepare (metacls, name, bases, \*\*kwds): return collections.OrderedDict() def \_\_new\_\_(cls, name, bases, namespace, \*\*kwds): result = type. new (cls, name, bases, dict(namespace)) result.members = tuple(namespace) return result class A(metaclass=OrderedClass): def one(self): pass def two(self): pass def three(self): pass def four(self): pass >>> A.members (' module ', 'one', 'two', 'three', 'four') When the class definition for A gets executed, the process begins with calling the metaclass's \_\_prepare\_\_() method which returns an empty collections. Ordered Dict. That mapping records the methods and attributes of A as they are defined within the body of the class statement. Once those definitions are executed, the ordered dictionary is fully populated and the metaclass's \_\_new\_\_() method gets invoked. That method builds the new type

## 3.3.4 Customizing instance and subclass checks

and it saves the ordered dictionary keys in an attribute called members.

The following methods are used to override the default behavior of the isinstance() and issubclass() builtin

functions.

In particular, the metaclass abc.ABCMeta implements these methods in order to allow the addition of Abstract Base

Classes (ABCs) as "virtual base classes" to any class or type (including built-in types), including other ABCs.

class.\_\_instancecheck\_\_(self, instance)

Return true if instance should be considered a (direct or indirect) instance of class. If defined, called to implement

isinstance(instance, class).

class.\_\_subclasscheck\_\_(self, subclass)

Return true if subclass should be considered a (direct or indirect) subclass of class. If defined, called to implement

issubclass(subclass, class).

Note that these methods are looked up on the type (metaclass) of a class. They cannot be defined as class methods in

the actual class. This is consistent with the lookup of special methods that are called on instances, only in this case the

instance is itself a class.

Introducing Abstract Base Classes Includes the specification for customizing isinstance() and

issubclass() behavior through \_\_instancecheck\_\_() and \_\_subclasscheck\_\_(), with motivation

for this functionality in the context of adding Abstract Base Classes (see the abc module) to the language.

# 3.3.5 Emulating callable objects

```
object.__call__(self [, args...])
```

Called when the instance is "called" as a function; if this method is defined, x(arg1, arg2, ...) is a

shorthand for x.\_\_call\_\_(arg1, arg2, ...).

3.3.6 Emulating container types

The following methods can be defined to implement container objects. Containers usually are sequences (such as

lists or tuples) or mappings (like dictionaries), but can represent other containers

as well. The first set of methods is used either to emulate a sequence or to emulate a mapping; the difference is that for a sequence, the allowable keys should be the integers k for which  $0 \le k \le N$  where N is the length of the sequence, or slice objects, which define a range of items. It is also recommended that mappings provide the methods keys(), values(), items(), get(), clear(), setdefault(), pop(), popitem(), copy(), and update() behaving similar to those for Python's standard dictionary objects. The collections module provides a MutableMapping abstract base class to help create those methods from a base set of \_\_getitem\_\_(), setitem (), delitem (), and keys(). Mutable sequences should provide methods append(), count(), index(), extend(), insert(), pop(), remove(), reverse() and sort(), like Python standard list objects. Finally, sequence types should implement addition (meaning concatenation) and multiplication (meaning repetition) by defining the methods add (), radd (), iadd (), mul (), rmul () and imul () described below: they should not define other numerical operators. It is recommended that both mappings and sequences implement the contains () method to allow efficient use of the in operator; for mappings, in should search the mapping's keys; for sequences, it should search through the values. It is further recommended that both mappings and sequences implement the \_\_iter\_\_() method to allow efficient iteration through the container; for mappings, \_\_iter\_\_() should be the same as keys(); for sequences, it should iterate through the values. object. len (self) Called to implement the built-in function len(). Should return the length of the

Also, an object that doesn't define a bool () method and whose len ()

considered to be false in a Boolean context.

object, an integer >= 0.

method returns zero is

Note: Slicing is done exclusively with the following three methods. A call like
a[1:2] = b
is translated to
a[slice(1, 2, None)] = b
and so forth. Missing slice items are always filled in with None.
objectgetitem(self, key)
Called to implement evaluation of self[key]. For sequence types, the accepted keys should be integers and
slice objects. Note that the special interpretation of negative indexes (if the class wishes to emulate a sequence
type) is up to thegetitem() method. If key is of an inappropriate type, TypeError may be raised;
if of a value outside the set of indexes for the sequence (after any special interpretation of negative values),
IndexError should be raised. For mapping types, if key is missing (not in the container), KeyError should
be raised.
Note: for loops expect that an IndexError will be raised for illegal indexes to allow proper detection of
the end of the sequence.
objectsetitem(self, key, value)
Called to implement assignment to self[key]. Same note as forgetitem(). This should only be implemented for mappings if the objects support changes to the values for keys, or if new keys can be added, or
for sequences if elements can be replaced. The same exceptions should be raised for improper key values as for
thegetitem() method.
objectdelitem(self, key)
Called to implement deletion of self[key]. Same note as forgetitem(). This should only be
implemented for mappings if the objects support removal of keys, or for sequences if elements can be removed
from the sequence. The same exceptions should be raised for improper key values as for thegetitem()
method.

objectiter(self)
This method is called when an iterator is required for a container. This method should return a new iterator
object that can iterate over all the objects in the container. For mappings, it should iterate over the keys of the
container, and should also be made available as the method keys().
Iterator objects also need to implement this method; they are required to return themselves. For more information
on iterator objects, see typeiter.
objectreversed(self)
Called (if present) by the reversed() built-in to implement reverse iteration. It should return a new iterator
object that iterates over all the objects in the container in reverse order.
If thereversed() method is not provided, the reversed() built-in will fall back to using the sequence
protocol (len() andgetitem()). Objects that support the sequence protocol should only
providereversed() if they can provide an implementation that is more efficient than the one provided
by reversed().
The membership test operators (in and not in) are normally implemented as an iteration through a sequence.
However, container objects can supply the following special method with a more efficient implementation, which also
does not require the object be a sequence.
objectcontains(self, item)
Called to implement membership test operators. Should return true if item is in self, false otherwise. For
mapping objects, this should consider the keys of the mapping rather than the values or the key-item pairs.
For objects that don't definecontains(), the membership test first tries iteration viaiter(),
then the old sequence iteration protocol viagetitem(), see this section in the language reference.

# 3.3.7 Emulating numeric types

The following methods can be defined to emulate numeric objects. Methods corresponding to operations that are not

supported by the particular kind of number implemented (e.g., bitwise operations for non-integral numbers) should be

left undefined.
objectadd(self, other)
objectsub(self, other)
objectmul(self, other)
objecttruediv(self, other)
objectfloordiv(self, other)
objectmod(self, other)
objectdivmod(self, other)
objectpow(self, other[, modulo ])
objectlshift(self, other)
objectrshift(self, other)
objectand(self, other)
objectxor(self, other)
objector(self, other)
These methods are called to implement the binary arithmetic operations (+, -, *, /, $//$ , %, divmod(),
pow(), **, <<, >>, &, ^,  ). For instance, to evaluate the expression x + y, where x is an instance of a class
that has anadd() method, xadd(y) is called. Thedivmod() method should be the
equivalent to usingfloordiv() andmod(); it should not be related totruediv(). Note
thatpow() should be defined to accept an optional third argument if the ternary version of the built-in
pow() function is to be supported.
If one of those methods does not support the operation with the supplied

NotImplemented.

arguments, it should return

```
object. radd (self, other)
object.__rsub__(self, other)
object. rmul (self, other)
object. rtruediv (self, other)
object.__rfloordiv__(self, other)
object.__rmod__(self, other)
object. rdivmod (self, other)
object. rpow (self, other)
object. rlshift (self, other)
object.__rrshift__(self, other)
object. rand (self, other)
object.__rxor__(self, other)
object. ror (self, other)
These methods are called to implement the binary arithmetic operations (+, -, *, /,
//, %, divmod(), pow(),
**, <<, >>, &, ^, |) with reflected (swapped) operands. These functions are only
called if the left operand does
not support the corresponding operation and the operands are of different types. 2
For instance, to evaluate the
expression x - y, where y is an instance of a class that has an rsub () method,
y.__rsub__(x) is
called if x. sub (y) returns NotImplemented.
Note that ternary pow() will not try calling __rpow__() (the coercion rules would
become too complicated).
Note: If the right operand's type is a subclass of the left operand's type and that
subclass provides the reflected
method for the operation, this method will be called before the left operand's non-
reflected method. This
behavior allows subclasses to override their ancestors' operations.
object. iadd (self, other)
object. isub (self, other)
object.__imul__(self, other)
object.__itruediv__(self, other)
```

```
object.__ifloordiv__(self, other)
object. imod (self, other)
object.__ipow__(self, other[, modulo ])
object.__ilshift__(self, other)
object.__irshift__(self, other)
object.__iand__(self, other)
object. ixor (self, other)
object. ior (self, other)
These methods are called to implement the augmented arithmetic assignments
(+=, -=, *=, /=, //=, %=, **=,
<=, >>=, &=, ^=, |=). These methods should attempt to do the operation in-place
(modifying self ) and return
the result (which could be, but does not have to be, self). If a specific method is
not defined, the augmented
assignment falls back to the normal methods. For instance, to execute the
statement x += y, where x is an
instance of a class that has an <u>iadd</u>() method, x.<u>iadd</u>(y) is called. If x is
an instance of a class
that does not define a iadd () method, x. add (y) and y. radd (x) are
considered, as with
the evaluation of x + y.
2 For operands of the same type, it is assumed that if the non-reflected method
(such as __add__()) fails the operation is not supported, which
is why the reflected method is not called.
object. neg (self)
object. pos (self)
object. abs (self)
object. invert (self)
Called to implement the unary arithmetic operations (-, +, abs() and ~).
object. complex (self)
object.__int__(self)
object. float (self)
```

```
object.__round__(self [, n ])
```

Called to implement the built-in functions complex(), int(), float() and round(). Should return a

value of the appropriate type.

```
object.__index__(self)
```

Called to implement operator.index(). Also called whenever Python needs an integer object (such as in

slicing, or in the built-in bin(), hex() and oct() functions). Must return an integer.

# 3.3.8 With Statement Context Managers

A context manager is an object that defines the runtime context to be established when executing a with statement.

The context manager handles the entry into, and the exit from, the desired runtime context for the execution of the

block of code. Context managers are normally invoked using the with statement (described in section The with

statement), but can also be used by directly invoking their methods.

Typical uses of context managers include saving and restoring various kinds of global state, locking and unlocking

resources, closing opened files, etc.

For more information on context managers, see typecontextmanager.

```
object. enter (self)
```

Enter the runtime context related to this object. The with statement will bind this method's return value to the

target(s) specified in the as clause of the statement, if any.

```
object.__exit__(self, exc_type, exc_value, traceback)
```

Exit the runtime context related to this object. The parameters describe the exception that caused the context to

be exited. If the context was exited without an exception, all three arguments will be None.

If an exception is supplied, and the method wishes to suppress the exception (i.e., prevent it from being propagated),

it should return a true value. Otherwise, the exception will be processed normally upon exit from this

method.

Note that \_\_exit\_\_() methods should not reraise the passed-in exception; this is the caller's responsibility.

### 3.3.9 Special method lookup

For custom classes, implicit invocations of special methods are only guaranteed to work correctly if defined on an

object's type, not in the object's instance dictionary. That behaviour is the reason why the following code raises an

```
exception:
>>> class C:
... pass
>>> c = C()
>>> c. len = lambda: 5
>>> len(c)
Traceback (most recent call last):
File "<stdin>", line 1, in <module>
TypeError: object of type 'C' has no len()
The rationale behind this behaviour lies with a number of special methods such as
hash () and repr ()
that are implemented by all objects, including type objects. If the implicit lookup of
these methods used the conventional
lookup process, they would fail when invoked on the type object itself:
>>> 1 .__hash__() == hash(1)
True
>>> int.__hash__() == hash(int)
Traceback (most recent call last):
File "<stdin>", line 1, in <module>
TypeError: descriptor '__hash__' of 'int' object needs an argument
```

Incorrectly attempting to invoke an unbound method of a class in this way is

sometimes referred to as 'metaclass

```
methods:
>>> type(1).__hash__(1) == hash(1)
True
>>> type(int). hash (int) == hash(int)
True
In addition to bypassing any instance attributes in the interest of correctness,
implicit special method lookup generally
also bypasses the __getattribute__() method even of the object's metaclass:
>>> class Meta(type):
... def __getattribute__(*args):
... print("Metaclass getattribute invoked")
... return type. getattribute (*args)
>>> class C(object, metaclass=Meta):
... def len (self):
... return 10
... def __getattribute__(*args):
... print("Class getattribute invoked")
... return object. getattribute (*args)
>>> c = C()
>>> c.__len__() # Explicit lookup via instance
Class getattribute invoked
10
>>> type(c).__len__(c) # Explicit lookup via type
Metaclass getattribute invoked
10
>>> len(c) # Implicit lookup
10
Bypassing the getattribute () machinery in this fashion provides significant
scope for speed optimisations
```

confusion', and is avoided by bypassing the instance when looking up special

within the interpreter, at the cost of some flexibility in the handling of special methods (the special method must be set

on the class object itself in order to be consistently invoked by the interpreter).

### **CHAPTER FOUR EXECUTION MODEL**

### 4.1 Naming and binding

Names refer to objects. Names are introduced by name binding operations. Each occurrence of a name in the program

text refers to the binding of that name established in the innermost function block containing the use.

A block is a piece of Python program text that is executed as a unit. The following are blocks: a module, a function

body, and a class definition. Each command typed interactively is a block. A script file (a file given as standard input

to the interpreter or specified on the interpreter command line the first argument) is a code block. A script command (a

command specified on the interpreter command line with the '-c' option) is a code block. The string argument passed

to the built-in functions eval() and exec() is a code block.

A code block is executed in an execution frame. A frame contains some administrative information (used for debugging)

and determines where and how execution continues after the code block's execution has completed.

A scope defines the visibility of a name within a block. If a local variable is defined in a block, its scope includes

that block. If the definition occurs in a function block, the scope extends to any blocks contained within the defining

one, unless a contained block introduces a different binding for the name. The scope of names defined in a class block

is limited to the class block; it does not extend to the code blocks of methods – this includes comprehensions and

generator expressions since they are implemented using a function scope. This means that the following will fail:

class A:

a = 42

b = list(a + i for i in range(10))

When a name is used in a code block, it is resolved using the nearest enclosing scope. The set of all such scopes visible

to a code block is called the block's environment.

If a name is bound in a block, it is a local variable of that block, unless declared as nonlocal. If a name is bound at

the module level, it is a global variable. (The variables of the module code block are local and global.) If a variable is

used in a code block but not defined there, it is a free variable.

When a name is not found at all, a NameError exception is raised. If the name refers to a local variable that has not

been bound, a UnboundLocalError exception is raised. UnboundLocalError is a subclass of NameError.

The following constructs bind names: formal parameters to functions, import statements, class and function definitions

(these bind the class or function name in the defining block), and targets that are identifiers if occurring in

an assignment, for loop header, or after as in a with statement or except clause. The import statement of

the form from ... import \* binds all names defined in the imported module, except those beginning with an

underscore. This form may only be used at the module level.

A target occurring in a del statement is also considered bound for this purpose (though the actual semantics are to

unbind the name).

Each assignment or import statement occurs within a block defined by a class or function definition or at the module

level (the top-level code block).

If a name binding operation occurs anywhere within a code block, all uses of the name within the block are treated as

references to the current block. This can lead to errors when a name is used within a block before it is bound. This rule

is subtle. Python lacks declarations and allows name binding operations to occur anywhere within a code block. The

local variables of a code block can be determined by scanning the entire text of the block for name binding operations. If the global statement occurs within a block, all uses of the name specified in the statement refer to the binding

of that name in the top-level namespace. Names are resolved in the top-level namespace by searching the global

namespace, i.e. the namespace of the module containing the code block, and the builtins namespace, the namespace of

the module builtins. The global namespace is searched first. If the name is not found there, the builtins namespace

is searched. The global statement must precede all uses of the name.

The builtins namespace associated with the execution of a code block is actually found by looking up the name

builtins	in its global namespace; this should be a dictionary or a module (in	n
the latter ca	se the module's	

dictionary is used). By default, when in the \_\_main\_\_ module, \_\_builtins\_\_ is the built-in module

builtins; when in any other module, \_\_builtins\_\_ is an alias for the dictionary of the builtins module

itself. \_\_builtins\_\_ can be set to a user-created dictionary to create a weak form of restricted execution.

CPython implementation detail: Users should not touch \_\_builtins\_\_; it is strictly an implementation detail.

Users wanting to override values in the builtins namespace should import the builtins module and modify its

attributes appropriately.

The namespace for a module is automatically created the first time a module is imported. The main module for a script

is always called \_\_main\_\_.

The global statement has the same scope as a name binding operation in the same block. If the nearest enclosing

scope for a free variable contains a global statement, the free variable is treated as a global.

A class definition is an executable statement that may use and define names. These references follow the normal

rules for name resolution. The namespace of the class definition becomes the attribute dictionary of the class. Names

defined at the class scope are not visible in methods.

### 4.1.1 Interaction with dynamic features

There are several cases where Python statements are illegal when used in conjunction with nested scopes that contain

free variables.

If a variable is referenced in an enclosing scope, it is illegal to delete the name. An error will be reported at compile

time.

If the wild card form of import — import \* — is used in a function and the function contains or is a nested block

with free variables, the compiler will raise a SyntaxError.

The eval() and exec() functions do not have access to the full environment for resolving names. Names may

be resolved in the local and global namespaces of the caller. Free variables are not resolved in the nearest enclosing

namespace, but in the global namespace. 1 The exec() and eval() functions have optional arguments to override

the global and local namespace. If only one namespace is specified, it is used for both.

# 4.2 Exceptions

Exceptions are a means of breaking out of the normal flow of control of a code block in order to handle errors or

other exceptional conditions. An exception is raised at the point where the error is detected; it may be handled by the surrounding code block or by any code block that directly or indirectly invoked the code block where the error

occurred.

The Python interpreter raises an exception when it detects a run-time error (such as division by zero). A Python

program can also explicitly raise an exception with the raise statement. Exception handlers are specified with the

try ... except statement. The finally clause of such a statement can be used to specify cleanup code which does

not handle the exception, but is executed whether an exception occurred or not in the preceding code.

Python uses the "termination" model of error handling: an exception handler can find out what happened and continue

execution at an outer level, but it cannot repair the cause of the error and retry the failing operation (except by reentering

the offending piece of code from the top).

When an exception is not handled at all, the interpreter terminates execution of the program, or returns to its interactive

main loop. In either case, it prints a stack backtrace, except when the exception is SystemExit.

Exceptions are identified by class instances. The except clause is selected depending on the class of the instance: it

must reference the class of the instance or a base class thereof. The instance can be received by the handler and can

carry additional information about the exceptional condition.

Note: Exception messages are not part of the Python API. Their contents may change from one version of Python to

the next without warning and should not be relied on by code which will run under multiple versions of the interpreter.

See also the description of the try statement in section The try statement and raise statement in section The raise

statement.

### **CHAPTER FIVE THE IMPORT SYSTEM**

Python code in one module gains access to the code in another module by the process of importing it. The import
statement is the most common way of invoking the import machinery, but it is not the only way. Functions such as
importlib.import_module() and built-inimport() can also be used to invoke the import machinery.
The import statement combines two operations; it searches for the named module, then it binds the results of
that search to a name in the local scope. The search operation of the import statement is defined as a call to the
import() function, with the appropriate arguments. The return value ofimport() is used to perform
the name binding operation of the import statement. See the import statement for the exact details of that name
binding operation.
A direct call toimport() performs only the module search and, if found, the module creation operation. While
certain side-effects may occur, such as the importing of parent packages, and the updating of various caches (including
sys.modules), only the import statement performs a name binding operation.
When callingimport() as part of an import statement, the import system first checks the module global
namespace for a function by that name. If it is not found, then the standard builtinimport() is called. Other
mechanisms for invoking the import system (such as importlib.import_module()) do not perform this check
and will always use the standard import system.
When a module is first imported, Python searches for the module and if found, it creates a module object 1, initializing

for the named module when the import machinery is invoked. These strategies can be modified and extended by using

it. If the named module cannot be found, an ImportError is raised. Python

implements various strategies to search

various hooks described in the sections below. Changed in version 3.3: There is no longer any implicit import machinery - the full import system is

exposed through sys.meta\_path. In addition, native namespace package support has been implemented

#### 5.1 importlib

The importlib module provides a rich API for interacting with the import system. For example

importlib.import\_module() provides a recommended, simpler API than built-in \_\_import\_\_() for invoking

the import machinery. Refer to the importlib library documentation for additional detail.

#### 5.2 Packages

Python has only one type of module object, and all modules are of this type, regardless of whether the module is

implemented in Python, C, or something else. To help organize modules and provide a naming hierarchy, Python has a concept of packages.

You can think of packages as the directories on a file system and modules as files within directories, but don't take

this analogy too literally since packages and modules need not originate from the file system. For the purposes of this

documentation, we'll use this convenient analogy of directories and files. Like file system directories, packages are

organized hierarchically, and packages may themselves contain subpackages, as well as regular modules.

It's important to keep in mind that all packages are modules, but not all modules are packages. Or put another way,

packages are just a special kind of module. Specifically, any module that contains a \_\_path\_\_ attribute is considered

a package.

All modules have a name. Subpackage names are separated from their parent package name by dots, akin to Python's

standard attribute access syntax. Thus you might have a module called sys and a package called email, which in

turn has a subpackage called email.mime and a module within that subpackage called email.mime.text.

### 5.2.1 Regular packages

Python defines two types of packages, regular packages and namespace packages. Regular packages are traditional packages as they existed in Python 3.2 and earlier. A regular package is typically implemented as a directory containing an init .py file. When a regular package is imported, this init .py file is implicitly executed, and the objects it defines are bound to names in the package's namespace. The init .py file can contain the same Python code that any other module can contain, and Python will add some additional attributes to the module when it is imported. For example, the following file system layout defines a top level parent package with three subpackages: parent/ \_\_init\_\_.py one/ \_\_init\_\_.py two/ \_\_init\_\_.py three/ init .py Importing parent.one will implicitly execute parent/ init .py and parent/one/\_\_init\_\_.py. Subsequent imports of parent.two or parent.three will execute

### 5.2.2 Namespace packages

parent/three/\_\_init\_\_.py respectively.

parent/two/\_\_init\_\_.py and

A namespace package is a composite of various portions, where each portion contributes a subpackage to the parent

package. Portions may reside in different locations on the file system. Portions may also be found in zip files, on

the network, or anywhere else that Python searches during import. Namespace packages may or may not correspond

directly to objects on the file system; they may be virtual modules that have no concrete representation.

Namespace packages do not use an ordinary list for their \_\_path\_\_ attribute. They instead use a custom iterable type

which will automatically perform a new search for package portions on the next import attempt within that package if

the path of their parent package (or sys.path for a top level package) changes.

With namespace packages, there is no parent/\_\_init\_\_.py file. In fact, there may be multiple parent directories

found during import search, where each one is provided by a different portion. Thus parent/one may not

be physically located next to parent/two. In this case, Python will create a namespace package for the top-level

parent package whenever it or one of its subpackages is imported.

### 5.3 Searching

To begin the search, Python needs the fully qualified name of the module (or package, but for the purposes of this

discussion, the difference is immaterial) being imported. This name may come from various arguments to the import

statement, or from the parameters to the importlib.import\_module() or \_\_import\_\_() functions.

This name will be used in various phases of the import search, and it may be the dotted path to a submodule, e.g.

foo.bar.baz. In this case, Python first tries to import foo, then foo.bar, and finally foo.bar.baz. If any of

the intermediate imports fail, an ImportError is raised.

#### 5.3.1 The module cache

The first place checked during import search is sys.modules. This mapping serves as a cache of all modules that

have been previously imported, including the intermediate paths. So if foo.bar.baz

was previously imported,

sys.modules will contain entries for foo, foo.bar, and foo.bar.baz. Each key will have as its value the

corresponding module object.

During import, the module name is looked up in sys.modules and if present, the associated value is the module

satisfying the import, and the process completes. However, if the value is None, then an ImportError is raised. If

the module name is missing, Python will continue searching for the module.

sys.modules is writable. Deleting a key may not destroy the associated module (as other modules may hold

references to it), but it will invalidate the cache entry for the named module, causing Python to search anew for the

named module upon its next import. The key can also be assigned to None, forcing the next import of the module to

result in an ImportError.

Beware though, as if you keep a reference to the module object, invalidate its cache entry in sys.modules, and then

re-import the named module, the two module objects will not be the same. By contrast, imp.reload() will reuse

the same module object, and simply reinitialise the module contents by rerunning the module's code.

### 5.3.2 Finders and loaders

If the named module is not found in sys.modules, then Python's import protocol is invoked to find and load the

module. This protocol consists of two conceptual objects, finders and loaders. A finder's job is to determine whether

it can find the named module using whatever strategy it knows about. Objects that implement both of these interfaces

are referred to as importers - they return themselves when they find that they can load the requested module.

Python includes a number of default finders and importers. The first one knows how to locate built-in modules, and

the second knows how to locate frozen modules. A third default finder searches an import path for modules. The

import path is a list of locations that may name file system paths or zip files. It can also be extended to search for any

locatable resource, such as those identified by URLs.

The import machinery is extensible, so new finders can be added to extend the range and scope of module searching.

Finders do not actually load modules. If they can find the named module, they return a loader, which the import

machinery then invokes to load the module and create the corresponding module object.

The following sections describe the protocol for finders and loaders in more detail, including how you can create and

register new ones to extend the import machinery.

### 5.3.3 Import hooks

The import machinery is designed to be extensible; the primary mechanism for this are the import hooks. There are

two types of import hooks: meta hooks and import path hooks.

Meta hooks are called at the start of import processing, before any other import processing has occurred, other than

sys.modules cache look up. This allows meta hooks to override sys.path processing, frozen modules, or even

built-in modules. Meta hooks are registered by adding new finder objects to sys.meta\_path, as described below.

Import path hooks are called as part of sys.path (or package.\_\_path\_\_) processing, at the point where their

associated path item is encountered. Import path hooks are registered by adding new callables to sys.path\_hooks

as described below.

### 5.3.4 The meta path

When the named module is not found in sys.modules, Python next searches sys.meta path, which contains

a list of meta path finder objects. These finders are queried in order to see if they know how to handle the named

module. Meta path finders must implement a method called find\_module() which takes two arguments, a name

and an import path. The meta path finder can use any strategy it wants to determine whether it can handle the named

module or not.

If the meta path finder knows how to handle the named module, it returns a loader object. If it cannot handle the named

module, it returns None. If sys.meta\_path processing reaches the end of its list without returning a loader, then

an ImportError is raised. Any other exceptions raised are simply propagated up, aborting the import process.

The find\_module() method of meta path finders is called with two arguments. The first is the fully qualified name

of the module being imported, for example foo.bar.baz. The second argument is the path entries to use for the

module search. For top-level modules, the second argument is None, but for submodules or subpackages, the second

argument is the value of the parent package's \_\_path\_\_ attribute. If the appropriate \_\_path\_\_ attribute cannot be

accessed, an ImportError is raised.

The meta path may be traversed multiple times for a single import request. For example, assuming

none of the modules involved has already been cached, importing foo.bar.baz will first perform a

top level import, calling mpf.find\_module("foo", None) on each meta path finder (mpf). After

foo has been imported, foo.bar will be imported by traversing the meta path a second time, calling

mpf.find\_module("foo.bar", foo.\_\_path\_\_\_). Once foo.bar has been imported, the final traversal

will call mpf.find\_module("foo.bar.baz", foo.bar.\_\_path\_\_\_).

Some meta path finders only support top level imports. These importers will always return None when anything other

than None is passed as the second argument.

Python's default sys.meta\_path has three meta path finders, one that knows how to import built-in modules, one

that knows how to import frozen modules, and one that knows how to import modules from an import path (i.e. the

path based finder).

### 5.4 Loaders

If and when a module loader is found its load\_module() method is called, with a single argument, the fully

qualified name of the module being imported. This method has several responsibilities, and should return the module

object it has loaded 2. If it cannot load the module, it should raise an ImportError, although any other exception

raised during load\_module() will be propagated.

In many cases, the finder and loader can be the same object; in such cases the finder.find\_module() would

just return self.

Loaders must satisfy the following requirements:

• If there is an existing module object with the given name in sys.modules, the loader must use that existing

module. (Otherwise, imp.reload() will not work correctly.) If the named module does not exist in

sys.modules, the loader must create a new module object and add it to sys.modules.

Note that the module must exist in sys.modules before the loader executes the module code. This is crucial

because the module code may (directly or indirectly) import itself; adding it to sys.modules beforehand

prevents unbounded recursion in the worst case and multiple loading in the best.

If loading fails, the loader must remove any modules it has inserted into sys.modules, but it must remove

only the failing module, and only if the loader itself has loaded it explicitly. Any module already in

the sys.modules cache, and any module that was successfully loaded as a sideeffect, must remain in the

cache.

• The loader may set the \_\_file\_\_ attribute of the module. If set, this attribute's

value must be a string. The
loader may opt to leavefile unset if it has no semantic meaning (e.g. a module loaded from a database).
Iffile is set, it may also be appropriate to set thecached attribute which is the path to any
compiled version of the code (e.g. byte-compiled file). The file does not need to exist to set this attribute; the
path can simply point to whether the compiled file would exist .
• The loader may set thename attribute of the module. While not required, setting this attribute is highly
recommended so that the repr() of the module is more informative.
<ul> <li>If the module is a package (either regular or namespace), the loader must set the module object'spath</li> </ul>
attribute. The value must be iterable, but may be empty ifpath has no further significance to the loader.
Ifpath is not empty, it must produce strings when iterated over. More details on the semantics of
path are given below.
• Theloader attribute must be set to the loader object that loaded the module. This is mostly for introspection
and reloading, but can be used for additional loader-specific functionality, for example getting data
associated with a loader.
• The module'spackage attribute should be set. Its value must be a string, but it can be the same value
as itsname If the attribute is set to None or is missing, the import system will fill it in with a more
appropriate value. When the module is a package, itspackage value should be set to itsname
When the module is not a package,package should be set to the empty string for top-level modules, or
for submodules, to the parent package's name.
This attribute is used instead ofname to calculate explicit relative imports for main modules.
If the module is a Python module (as opposed to a built-in module or a dynamically loaded extension), the loader

should execute the module's code in the module's global name space (moduledict).
5.4.1 Module reprs
By default, all modules have a usable repr, however depending on the attributes set above, and hooks in the loader,
you can more explicitly control the repr of module objects.
Loaders may implement a module_repr() method which takes a single argument, the module object. When
repr(module) is called for a module with a loader supporting this protocol, whatever is returned from
moduleloadermodule_repr(module) is returned as the module's repr without further processing.
This return value must be a string.
If the module has noloader attribute, or the loader has no module_repr() method, then the module
object implementation itself will craft a default repr using whatever information is available. It will try to use the
modulename, modulefile, and moduleloader as input into the repr, with defaults for
whatever information is missing.
Here are the exact rules used:
<ul> <li>If the module has aloader and that loader has a module_repr() method, call it with a single argument,</li> </ul>
which is the module object. The value returned is used as the module's repr.
• If an exception occurs in module_repr(), the exception is caught and discarded, and the calculation of the
module's repr continues as if module_repr() did not exist.
• If the module has afile attribute, this is used as part of the module's repr.
• If the module has nofile but does have aloader, then the loader's repr is used as part of the

• Otherwise, just use the module's \_\_name\_\_ in the repr.

module's repr.

This example, shows how a loader can craft its own module repr:

class NamespaceLoader:
@classmethod
def module_repr(cls, module):
return " <module '{}'="" (namespace)="">".format(modulename)</module>
5.4.2 modulepath
By definition, if a module has anpath attribute, it is a package, regardless of its value.
A package'spath attribute is used during imports of its subpackages. Within the import machinery, it functions
much the same as sys.path, i.e. providing a list of locations to search for modules during import. However,
path is typically much more constrained than sys.path.
path must be an iterable of strings, but it may be empty. The same rules used for sys.path also apply
to a package'spath, and sys.path_hooks (described below) are consulted when traversing a package's
path
automatically setspath correctly for the namespace package.

#### 5.5 The Path Based Finder

As mentioned previously, Python comes with several default meta path finders. One of these, called the path based

finder, searches an import path, which contains a list of path entries. Each path entry names a location to search for

modules.

The path based finder itself doesn't know how to import anything. Instead, it traverses the individual path entries,

associating each of them with a path entry finder that knows how to handle that particular kind of path.

The default set of path entry finders implement all the semantics for finding modules on the file system, handling

special file types such as Python source code (.py files), Python byte code (.pyc and .pyo files) and shared libraries

(e.g. .so files). When supported by the zipimport module in the standard library, the default path entry finders

also handle loading all of these file types (other than shared libraries) from zipfiles.

Path entries need not be limited to file system locations. They can refer to URLs, database queries, or any other

location that can be specified as a string.

The path based finder provides additional hooks and protocols so that you can extend and customize the types of

searchable path entries. For example, if you wanted to support path entries as network URLs, you could write a hook

that implements HTTP semantics to find modules on the web. This hook (a callable) would return a path entry finder

supporting the protocol described below, which was then used to get a loader for the module from the web.

A word of warning: this section and the previous both use the term finder, distinguishing between them by using the

terms meta path finder and path entry finder. These two types of finders are very similar, support similar protocols,

and function in similar ways during the import process, but it's important to keep in mind that they are subtly different.

In particular, meta path finders operate at the beginning of the import process, as keyed off the sys.meta\_path

traversal.

By contrast, path entry finders are in a sense an implementation detail of the path based finder, and in fact, if the path

based finder were to be removed from sys.meta\_path, none of the path entry finder semantics would be invoked.

### 5.5.1 Path entry finders

The path based finder is responsible for finding and loading Python modules and packages whose location is specified

with a string path entry. Most path entries name locations in the file system, but they need not be limited to this.

As a meta path finder, the path based finder implements the find\_module() protocol previously described, however

it exposes additional hooks that can be used to customize how modules are found and loaded from the import

path.

Three variables are used by the path based finder, sys.path, sys.path\_hooks and sys.path\_importer\_cache. The \_\_path\_\_ attributes on package objects are also used. These provide

additional ways that the import machinery can be customized.

sys.path contains a list of strings providing search locations for modules and packages. It is initialized from the

PYTHONPATH environment variable and various other installation- and implementation-specific defaults. Entries in

sys.path can name directories on the file system, zip files, and potentially other "locations" (see the site module)

that should be searched for modules, such as URLs, or database queries. Only strings and bytes should be present on

sys.path; all other data types are ignored. The encoding of bytes entries is determined by the individual path entry

finders.

The path based finder is a meta path finder, so the import machinery begins the import path search by calling the path

based finder's find\_module() method as described previously. When the path argument to find\_module()

is given, it will be a list of string paths to traverse - typically a package's \_\_path\_\_ attribute for an import within

that package. If the path argument is None, this indicates a top level import and sys.path is used.

The path based finder iterates over every entry in the search path, and for each of these, looks for an appropriate path

entry finder for the path entry. Because this can be an expensive operation (e.g. there may be stat() call overheads

for this search), the path based finder maintains a cache mapping path entries to path entry finders. This cache is

maintained in sys.path\_importer\_cache (despite the name, this cache actually stores finder objects rather

than being limited to importer objects). In this way, the expensive search for a particular path entry location's path

entry finder need only be done once. User code is free to remove cache entries from sys.path\_importer\_cache

forcing the path based finder to perform the path entry search again 3.

If the path entry is not present in the cache, the path based finder iterates over every callable in sys.path\_hooks.

Each of the path entry hooks in this list is called with a single argument, the path entry to be searched. This callable may

either return a path entry finder that can handle the path entry, or it may raise ImportError. An ImportError is

used by the path based finder to signal that the hook cannot find a path entry finder for that path entry. The exception

is ignored and import path iteration continues. The hook should expect either a string or bytes object; the encoding

of bytes objects is up to the hook (e.g. it may be a file system encoding, UTF-8, or something else), and if the hook

cannot decode the argument, it should raise ImportError.

If sys.path\_hooks iteration ends with no path entry finder being returned, then the path based finder's

find\_module() method will store None in sys.path\_importer\_cache (to indicate that there is no finder

for this path entry) and return None, indicating that this meta path finder could not find the module.

3 In legacy code, it is possible to find instances of imp.NullImporter in the sys.path\_importer\_cache. It is recommended that

code be changed to use None instead. See portingpythoncode for more details.

If a path entry finder is returned by one of the path entry hook callables on sys.path\_hooks, then the following

protocol is used to ask the finder for a module loader, which is then used to load the module.

### 5.5.2 Path entry finder protocol

In order to support imports of modules and initialized packages and also to contribute portions to namespace packages,

path entry finders must implement the find\_loader() method.

find\_loader() takes one argument, the fully qualified name of the module being imported. find\_loader()

returns a 2-tuple where the first item is the loader and the second item is a namespace portion. When the first item

(i.e. the loader) is None, this means that while the path entry finder does not have a loader for the named module, it

knows that the path entry contributes to a namespace portion for the named module. This will almost always be the

case where Python is asked to import a namespace package that has no physical presence on the file system. When a

path entry finder returns None for the loader, the second item of the 2-tuple return value must be a sequence, although

it can be empty.

If find\_loader() returns a non-None loader value, the portion is ignored and the loader is returned from the path

based finder, terminating the search through the path entries.

For backwards compatibility with other implementations of the import protocol, many path entry finders also support

the same, traditional find\_module() method that meta path finders support. However path entry finder

find\_module() methods are never called with a path argument (they are expected to record the appropriate

path information from the initial call to the path hook).

The find\_module() method on path entry finders is deprecated, as it does not allow the path entry finder to contribute

portions to namespace packages. Instead path entry finders should implement the find\_loader() method

as described above. If it exists on the path entry finder, the import system will always call find\_loader() in

preference to find\_module().

### 5.6 Replacing the standard import system

The most reliable mechanism for replacing the entire import system is to delete the default contents of

sys.meta\_path, replacing them entirely with a custom meta path hook.

If it is acceptable to only alter the behaviour of import statements without affecting

other APIs that access the import

system, then replacing the builtin \_\_import\_\_() function may be sufficient. This technique may also be employed

at the module level to only alter the behaviour of import statements within that module.

To selectively prevent import of some modules from a hook early on the meta path (rather than disabling the standard

import system entirely), it is sufficient to raise ImportError directly from find\_module() instead of returning

None. The latter indicates that the meta path search should continue. while raising an exception terminates it

immediately.

#### 5.7 References

The import machinery has evolved considerably since Python's early days. The original specification for packages is

still available to read, although some details have changed since the writing of that document.

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### **CHAPTER SIX EXPRESSIONS**

This chapter explains the meaning of the elements of expressions in Python.

Syntax Notes: In this and the following chapters, extended BNF notation will be used to describe syntax, not lexical

analysis. When (one alternative of) a syntax rule has the form

name ::= othername

and no semantics are given, the semantics of this form of name are the same as for othername.

#### 6.1 Arithmetic conversions

When a description of an arithmetic operator below uses the phrase "the numeric arguments are converted to a common

type," this means that the operator implementation for built-in types works that way:

- If either argument is a complex number, the other is converted to complex;
- otherwise, if either argument is a floating point number, the other is converted to floating point;
- otherwise, both must be integers and no conversion is necessary.

Some additional rules apply for certain operators (e.g., a string left argument to the '%' operator). Extensions must

define their own conversion behavior.

#### 6.2 Atoms

Atoms are the most basic elements of expressions. The simplest atoms are identifiers or literals. Forms enclosed in

parentheses, brackets or braces are also categorized syntactically as atoms. The syntax for atoms is:

atom ::= identifier | literal | enclosure

enclosure ::= parenth\_form | list\_display | dict\_display | set\_display

| generator\_expression | yield\_atom

### 6.2.1 Identifiers (Names)

An identifier occurring as an atom is a name. See section Identifiers and keywords for lexical definition and section

Naming and binding for documentation of naming and binding.

When the name is bound to an object, evaluation of the atom yields that object. When a name is not bound, an attempt

to evaluate it raises a NameError exception.

Private name mangling: When an identifier that textually occurs in a class definition begins with two or more

underscore characters and does not end in two or more underscores, it is considered a private name of that class.

Private names are transformed to a longer form before code is generated for them. The transformation inserts the

class name, with leading underscores removed and a single underscore inserted, in front of the name. For example,

the identifier \_\_spam occurring in a class named Ham will be transformed to \_Ham\_spam. This transformation is

independent of the syntactical context in which the identifier is used. If the transformed name is extremely long (longer

than 255 characters), implementation defined truncation may happen. If the class name consists only of underscores,

no transformation is done.

### 6.2.2 Literals

Python supports string and bytes literals and various numeric literals:

literal ::= stringliteral | bytesliteral

| integer | floatnumber | imagnumber

Evaluation of a literal yields an object of the given type (string, bytes, integer, floating point number, complex number)

with the given value. The value may be approximated in the case of floating point and imaginary (complex) literals.

See section Literals for details.

All literals correspond to immutable data types, and hence the object's identity is less important than its value. Multiple

evaluations of literals with the same value (either the same occurrence in the program text or a different occurrence)

may obtain the same object or a different object with the same value.

### 6.2.3 Parenthesized forms

A parenthesized form is an optional expression list enclosed in parentheses:

```
parenth_form ::= "(" [expression_list] ")"
```

A parenthesized expression list yields whatever that expression list yields: if the list contains at least one comma, it

yields a tuple; otherwise, it yields the single expression that makes up the expression list.

An empty pair of parentheses yields an empty tuple object. Since tuples are immutable, the rules for literals apply

(i.e., two occurrences of the empty tuple may or may not yield the same object).

Note that tuples are not formed by the parentheses, but rather by use of the comma operator. The exception is the

empty tuple, for which parentheses are required — allowing unparenthesized "nothing" in expressions would cause

ambiguities and allow common typos to pass uncaught.

## 6.2.4 Displays for lists, sets and dictionaries

For constructing a list, a set or a dictionary Python provides special syntax called "displays", each of them in two

flavors:

- either the container contents are listed explicitly, or
- they are computed via a set of looping and filtering instructions, called a comprehension.

Common syntax elements for comprehensions are:

```
comprehension ::= expression comp_for
comp_for ::= "for" target_list "in" or_test [comp_iter]
comp_iter ::= comp_for | comp_if
comp_if ::= "if" expression_nocond [comp_iter]
```

The comprehension consists of a single expression followed by at least one for clause and zero or more for or if

clauses. In this case, the elements of the new container are those that would be produced by considering each of the

for or if clauses a block, nesting from left to right, and evaluating the expression to produce an element each time

the innermost block is reached.

Note that the comprehension is executed in a separate scope, so names assigned to in the target list don't "leak" in the

enclosing scope.

# 6.2.5 List displays

A list display is a possibly empty series of expressions enclosed in square brackets:

list\_display ::= "[" [expression\_list | comprehension] "]"

A list display yields a new list object, the contents being specified by either a list of expressions or a comprehension.

When a comma-separated list of expressions is supplied, its elements are evaluated from left to right and placed into

the list object in that order. When a comprehension is supplied, the list is constructed from the elements resulting from

the comprehension.

## 6.2.6 Set displays

A set display is denoted by curly braces and distinguishable from dictionary displays by the lack of colons separating

keys and values:

set\_display ::= "{" (expression\_list | comprehension) "}"

A set display yields a new mutable set object, the contents being specified by either a sequence of expressions or a

comprehension. When a comma-separated list of expressions is supplied, its elements are evaluated from left to right

and added to the set object. When a comprehension is supplied, the set is constructed from the elements resulting from

the comprehension.

An empty set cannot be constructed with {}; this literal constructs an empty dictionary.

6.2.7 Dictionary displays

A dictionary display is a possibly empty series of key/datum pairs enclosed in curly braces:

dict\_display ::= "{" [key\_datum\_list | dict\_comprehension] "}"

key\_datum\_list ::= key\_datum ("," key\_datum)\* [","]

key\_datum ::= expression ":" expression

dict\_comprehension ::= expression ":" expression comp\_for

A dictionary display yields a new dictionary object.

If a comma-separated sequence of key/datum pairs is given, they are evaluated from left to right to define the entries

of the dictionary: each key object is used as a key into the dictionary to store the corresponding datum. This means

that you can specify the same key multiple times in the key/datum list, and the final dictionary's value for that key will

be the last one given.

A dict comprehension, in contrast to list and set comprehensions, needs two expressions separated with a colon followed

by the usual "for" and "if" clauses. When the comprehension is run, the resulting key and value elements are

inserted in the new dictionary in the order they are produced.

Restrictions on the types of the key values are listed earlier in section The standard type hierarchy. (To summarize, the

key type should be hashable, which excludes all mutable objects.) Clashes between duplicate keys are not detected;

the last datum (textually rightmost in the display) stored for a given key value prevails.

# 6.2.8 Generator expressions

A generator expression is a compact generator notation in parentheses:

generator\_expression ::= "(" expression comp\_for ")"

A generator expression yields a new generator object. Its syntax is the same as for comprehensions, except that it is

enclosed in parentheses instead of brackets or curly braces.

Variables used in the generator expression are evaluated lazily when the

\_\_next\_\_() method is called for generator

object (in the same fashion as normal generators). However, the leftmost for clause is immediately evaluated, so that

an error produced by it can be seen before any other possible error in the code that handles the generator expression.

Subsequent for clauses cannot be evaluated immediately since they may depend on the previous for loop. For

example:  $(x^*y \text{ for } x \text{ in range}(10) \text{ for } y \text{ in bar}(x)).$ 

The parentheses can be omitted on calls with only one argument. See section Calls for the detail.

# 6.2.9 Yield expressions

yield\_atom ::= "(" yield\_expression ")"

yield expression ::= "yield" [expression list | "from" expression]

The yield expression is only used when defining a generator function, and can only be used in the body of a function

definition. Using a yield expression in a function definition is sufficient to cause that definition to create a generator

function instead of a normal function.

When a generator function is called, it returns an iterator known as a generator. That generator then controls the execution

of a generator function. The execution starts when one of the generator's methods is called. At that time, the execution

proceeds to the first yield expression, where it is suspended again, returning the value of expression\_list

to generator's caller. By suspended we mean that all local state is retained, including the current bindings of local variables,

the instruction pointer, and the internal evaluation stack. When the execution is resumed by calling one of the

generator's methods, the function can proceed exactly as if the yield expression was just another external call. The

value of the yield expression after resuming depends on the method which resumed the execution. If next ()

is used (typically via either a for or the next() builtin) then the result is None, otherwise, if send() is used, then

the result will be the value passed in to that method.

All of this makes generator functions quite similar to coroutines; they yield multiple times, they have more than one

entry point and their execution can be suspended. The only difference is that a generator function cannot control where

should the execution continue after it yields; the control is always transferred to the generator's caller.

yield expressions are allowed in the try clause of a try ... finally construct. If the generator is not resumed

before it is finalized (by reaching a zero reference count or by being garbage collected), the generator-iterator's

close() method will be called, allowing any pending finally clauses to execute.

When yield from <expr> is used, it treats the supplied expression as a subiterator. All values produced by that

subiterator are passed directly to the caller of the current generator's methods. Any values passed in with send()

and any exceptions passed in with throw() are passed to the underlying iterator if it has the appropriate methods. If this is not the case, then send() will raise AttributeError or TypeError, while throw() will just raise the

passed in exception immediately.

When the underlying iterator is complete, the value attribute of the raised StopIteration instance becomes the

value of the yield expression. It can be either set explicitly when raising StopIteration, or automatically when

the sub-iterator is a generator (by returning a value from the sub-generator).

Changed in version 3.3: Added yield from <expr> to delegate control flow to a subiterator

The parentheses can be omitted when the yield expression is the sole expression on the right hand side of an

assignment statement.

Generator-iterator methods

This subsection describes the methods of a generator iterator. They can be used to control the execution of a generator

function.

Note that calling any of the generator methods below when the generator is already executing raises a ValueError

exception.

generator.\_\_next\_\_()

Starts the execution of a generator function or resumes it at the last executed yield expression. When a

generator function is resumed with a \_\_next\_\_() method, the current yield expression always evaluates to

None. The execution then continues to the next yield expression, where the generator is suspended again, and

the value of the expression\_list is returned to next()'s caller. If the generator exits without yielding

another value, a Stoplteration exception is raised.

This method is normally called implicitly, e.g. by a for loop, or by the built-in next() function.

generator.send(value)

Resumes the execution and "sends" a value into the generator function. The value argument becomes the

result of the current yield expression. The send() method returns the next value yielded by the generator, or

raises StopIteration if the generator exits without yielding another value. When send() is called to start

the generator, it must be called with None as the argument, because there is no yield expression that could

receive the value.

generator.throw(type[, value[, traceback ]])

Raises an exception of type type at the point where generator was paused, and returns the next value yielded

by the generator function. If the generator exits without yielding another value, a StopIteration exception

is raised. If the generator function does not catch the passed-in exception, or raises a different exception, then

that exception propagates to the caller.

generator.close()

Raises a GeneratorExit at the point where the generator function was paused. If the generator function then

raises StopIteration (by exiting normally, or due to already being closed) or GeneratorExit (by not

catching the exception), close returns to its caller. If the generator yields a value, a RuntimeError is raised.

If the generator raises any other exception, it is propagated to the caller. close() does nothing if the generator

has already exited due to an exception or normal exit.

#### Examples

Here is a simple example that demonstrates the behavior of generators and generator functions:

```
>>> def echo(value=None):
... print("Execution starts when 'next()' is called for the first time.")
... try:
... while True:
... try:
... value = (yield value)
... except Exception as e:
... value = e
... finally:
... print("Don't forget to clean up when 'close()' is called.")
>>> generator = echo(1)
>>> print(next(generator))
Execution starts when 'next()' is called for the first time.
1
>>> print(next(generator))
None
>>> print(generator.send(2))
2
>>> generator.throw(TypeError, "spam")
TypeError('spam',)
>>> generator.close()
```

Don't forget to clean up when 'close()' is called.

#### 6.3 Primaries

Primaries represent the most tightly bound operations of the language. Their syntax is:

primary ::= atom | attributeref | subscription | slicing | call

### 6.3.1 Attribute references

An attribute reference is a primary followed by a period and a name:

attributeref ::= primary "." identifier

The primary must evaluate to an object of a type that supports attribute references, which most objects do. This

object is then asked to produce the attribute whose name is the identifier (which can be customized by overriding the

\_\_getattr\_\_() method). If this attribute is not available, the exception AttributeError is raised. Otherwise,

the type and value of the object produced is determined by the object. Multiple evaluations of the same attribute

reference may yield different objects.

# 6.3.2 Subscriptions

A subscription selects an item of a sequence (string, tuple or list) or mapping (dictionary) object:

subscription ::= primary "[" expression list "]"

The primary must evaluate to an object that supports subscription, e.g. a list or dictionary. User-defined objects can

support subscription by defining a \_\_getitem\_\_() method.

For built-in objects, there are two types of objects that support subscription:

If the primary is a mapping, the expression list must evaluate to an object whose value is one of the keys of the

mapping, and the subscription selects the value in the mapping that corresponds to that key. (The expression list is a

tuple except if it has exactly one item.)

If the primary is a sequence, the expression (list) must evaluate to an integer or a slice (as discussed in the following

section).

The formal syntax makes no special provision for negative indices in sequences; however, built-in sequences all

provide a \_\_getitem\_\_() method that interprets negative indices by adding the length of the sequence to the index

(so that x[-1] selects the last item of x). The resulting value must be a nonnegative integer less than the number of

items in the sequence, and the subscription selects the item whose index is that value (counting from zero). Since the

support for negative indices and slicing occurs in the object's \_\_getitem\_\_() method, subclasses overriding this

method will need to explicitly add that support.

A string's items are characters. A character is not a separate data type but a string of exactly one character.

## 6.3.3 Slicings

A slicing selects a range of items in a sequence object (e.g., a string, tuple or list). Slicings may be used as expressions

or as targets in assignment or del statements. The syntax for a slicing:

```
slicing ::= primary "[" slice_list "]"

slice_list ::= slice_item ("," slice_item)* [","]

slice_item ::= expression | proper_slice

proper_slice ::= [lower_bound] ":" [upper_bound] [ ":" [stride] ]

lower_bound ::= expression

upper_bound ::= expression

stride ::= expression
```

There is ambiguity in the formal syntax here: anything that looks like an expression list also looks like a slice list, so

any subscription can be interpreted as a slicing. Rather than further complicating the syntax, this is disambiguated by

defining that in this case the interpretation as a subscription takes priority over the interpretation as a slicing (this is

the case if the slice list contains no proper slice).

The semantics for a slicing are as follows. The primary must evaluate to a mapping object, and it is indexed (using the

same \_\_getitem\_\_() method as normal subscription) with a key that is constructed from the slice list, as follows.

If the slice list contains at least one comma, the key is a tuple containing the conversion of the slice items; otherwise,

the conversion of the lone slice item is the key. The conversion of a slice item that is an expression is that expression.

The conversion of a proper slice is a slice object (see section The standard type hierarchy) whose start, stop

and step attributes are the values of the expressions given as lower bound, upper bound and stride, respectively,

substituting None for missing expressions.

### 6.3.4 Calls

A call calls a callable object (e.g., a function) with a possibly empty series of arguments:

```
call ::= primary "(" [argument_list [","] | comprehension] ")"
argument_list ::= positional_arguments ["," keyword_arguments]
["," "**" expression] ["," keyword_arguments]
["," "**" expression]
[keyword_arguments ["," "**" expression]
["," keyword_arguments] ["," "**" expression]
["**" expression ["," keyword_arguments] ["," "**" expression]
["**" expression
positional_arguments ::= expression ("," expression)*
keyword_arguments ::= keyword_item ("," keyword_item)*
keyword_item ::= identifier "=" expression
```

A trailing comma may be present after the positional and keyword arguments but does not affect the semantics.

The primary must evaluate to a callable object (user-defined functions, built-in functions, methods of built-in objects,

class objects, methods of class instances, and all objects having a \_\_call\_\_() method are callable). All argument

expressions are evaluated before the call is attempted. Please refer to section Function definitions for the syntax of

formal parameter lists.

If keyword arguments are present, they are first converted to positional arguments, as follows. First, a list of unfilled

slots is created for the formal parameters. If there are N positional arguments, they are placed in the first N slots. Next,

for each keyword argument, the identifier is used to determine the corresponding slot (if the identifier is the same as

the first formal parameter name, the first slot is used, and so on). If the slot is already filled, a TypeError exception

is raised. Otherwise, the value of the argument is placed in the slot, filling it (even if the expression is None, it fills the

slot). When all arguments have been processed, the slots that are still unfilled are filled with the corresponding default

value from the function definition. (Default values are calculated, once, when the function is defined; thus, a mutable

object such as a list or dictionary used as default value will be shared by all calls that don't specify an argument value

for the corresponding slot; this should usually be avoided.) If there are any unfilled slots for which no default value is

specified, a TypeError exception is raised. Otherwise, the list of filled slots is used as the argument list for the call.

CPython implementation detail: An implementation may provide built-in functions whose positional parameters do

not have names, even if they are 'named' for the purpose of documentation, and which therefore cannot be supplied

by keyword. In CPython, this is the case for functions implemented in C that use PyArg ParseTuple() to parse

their arguments.

If there are more positional arguments than there are formal parameter slots, a TypeError exception is raised, unless

a formal parameter using the syntax \*identifier is present; in this case, that formal parameter receives a tuple

containing the excess positional arguments (or an empty tuple if there were no excess positional arguments).

If any keyword argument does not correspond to a formal parameter name, a TypeError exception is raised, unless a

formal parameter using the syntax \*\*identifier is present; in this case, that formal parameter receives a dictionary

containing the excess keyword arguments (using the keywords as keys and the argument values as corresponding

values), or a (new) empty dictionary if there were no excess keyword arguments.

If the syntax \*expression appears in the function call, expression must evaluate to an iterable. Elements from

this iterable are treated as if they were additional positional arguments; if there are positional arguments x1, ..., xN,

and expression evaluates to a sequence y1, ..., yM, this is equivalent to a call with M+N positional arguments x1,

```
..., xN, y1, ..., yM.
```

A consequence of this is that although the \*expression syntax may appear after some keyword arguments, it is

processed before the keyword arguments (and the \*\*expression argument, if any – see below). So:

```
>>> def f(a, b):
... print(a, b)
...
>>> f(b=1, *(2,))
2 1
```

```
>>> f(a=1, *(2,))
```

Traceback (most recent call last):

File "<stdin>", line 1, in?

TypeError: f() got multiple values for keyword argument 'a'

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It is unusual for both keyword arguments and the \*expression syntax to be used in

the same call, so in practice

this confusion does not arise.

If the syntax \*\*expression appears in the function call, expression must evaluate to a mapping, the contents

of which are treated as additional keyword arguments. In the case of a keyword appearing in both expression and

as an explicit keyword argument, a TypeError exception is raised.

Formal parameters using the syntax \*identifier or \*\*identifier cannot be used as positional argument slots

or as keyword argument names.

A call always returns some value, possibly None, unless it raises an exception. How this value is computed depends

on the type of the callable object.

If it is—

a user-defined function: The code block for the function is executed, passing it the argument list. The first thing

the code block will do is bind the formal parameters to the arguments; this is described in section Function

definitions. When the code block executes a return statement, this specifies the return value of the function

call.

a built-in function or method: The result is up to the interpreter; see built-in-funcs for the descriptions of built-in

functions and methods.

a class object: A new instance of that class is returned.

a class instance method: The corresponding user-defined function is called, with an argument list that is one longer

than the argument list of the call: the instance becomes the first argument.

a class instance: The class must define a \_\_call\_\_() method; the effect is then the same as if that method was

called.

# **6.4** The power operator

The power operator binds more tightly than unary operators on its left; it binds less tightly than unary operators on its

right. The syntax is:

power ::= primary ["\*\*" u\_expr]

Thus, in an unparenthesized sequence of power and unary operators, the operators are evaluated from right to left (this

does not constrain the evaluation order for the operands): -1\*\*2 results in -1.

The power operator has the same semantics as the built-in pow() function, when called with two arguments: it yields

its left argument raised to the power of its right argument. The numeric arguments are first converted to a common

type, and the result is of that type.

For int operands, the result has the same type as the operands unless the second argument is negative; in that case, all

arguments are converted to float and a float result is delivered. For example, 10\*\*2 returns 100, but 10\*\*-2 returns

0.01.

Raising 0.0 to a negative power results in a ZeroDivisionError. Raising a negative number to a fractional

power results in a complex number. (In earlier versions it raised a ValueError.)

# 6.5 Unary arithmetic and bitwise operations

All unary arithmetic and bitwise operations have the same priority:

u expr::= power | "-" u expr | "+" u expr | "~" u expr

The unary - (minus) operator yields the negation of its numeric argument.

The unary + (plus) operator yields its numeric argument unchanged.

The unary ~ (invert) operator yields the bitwise inversion of its integer argument. The bitwise inversion of x is defined

as -(x+1). It only applies to integral numbers.

In all three cases, if the argument does not have the proper type, a TypeError exception is raised.

# 6.6 Binary arithmetic operations

The binary arithmetic operations have the conventional priority levels. Note that some of these operations also apply to

certain non-numeric types. Apart from the power operator, there are only two levels, one for multiplicative operators

and one for additive operators:

m\_expr ::= u\_expr | m\_expr "\*" u\_expr | m\_expr "//" u\_expr | m\_expr "/" u\_expr | m\_expr "%" u\_expr

a\_expr ::= m\_expr | a\_expr "+" m\_expr | a\_expr "-" m\_expr

The \* (multiplication) operator yields the product of its arguments. The arguments must either both be numbers, or

one argument must be an integer and the other must be a sequence. In the former case, the numbers are converted to a

common type and then multiplied together. In the latter case, sequence repetition is performed; a negative repetition

factor yields an empty sequence.

The / (division) and // (floor division) operators yield the quotient of their arguments. The numeric arguments are

first converted to a common type. Integer division yields a float, while floor division of integers results in an integer;

the result is that of mathematical division with the 'floor' function applied to the result. Division by zero raises the

ZeroDivisionError exception.

The % (modulo) operator yields the remainder from the division of the first argument by the second. The numeric

arguments are first converted to a common type. A zero right argument raises the ZeroDivisionError exception.

The arguments may be floating point numbers, e.g., 3.14%0.7 equals 0.34 (since 3.14 equals 4\*0.7 + 0.34.)

The modulo operator always yields a result with the same sign as its second operand (or zero); the absolute value of

the result is strictly smaller than the absolute value of the second operand 1.

The floor division and modulo operators are connected by the following identity:  $x = (x//y)^*y + (x\%y)$ .

Floor division and modulo are also connected with the built-in function divmod(): divmod(x, y) == (x//y,

x%y). 2.

In addition to performing the modulo operation on numbers, the % operator is also

overloaded by string objects to

perform old-style string formatting (also known as interpolation). The syntax for string formatting is described in the

Python Library Reference, section old-string-formatting.

1 While abs(x%y) < abs(y) is true mathematically, for floats it may not be true numerically due to roundoff. For example, and assuming

a platform on which a Python float is an IEEE 754 double-precision number, in order that -1e-100 % 1e100 have the same sign as 1e100,

the computed result is -1e-100 + 1e100, which is numerically exactly equal to 1e100. The function math.fmod() returns a result whose

sign matches the sign of the first argument instead, and so returns -1e-100 in this case. Which approach is more appropriate depends on the application.

2 If x is very close to an exact integer multiple of y, it's possible for x//y to be one larger than (x-x%y)//y due to rounding. In such cases,

Python returns the latter result, in order to preserve that div mod(x,y)[0] \* y + x % y be very close to x.

The floor division operator, the modulo operator, and the divmod() function are not defined for complex numbers.

Instead, convert to a floating point number using the abs() function if appropriate.

The + (addition) operator yields the sum of its arguments. The arguments must either both be numbers or both

sequences of the same type. In the former case, the numbers are converted to a common type and then added together.

In the latter case, the sequences are concatenated.

The - (subtraction) operator yields the difference of its arguments. The numeric arguments are first converted to a common type.

## 6.7 Shifting operations

The shifting operations have lower priority than the arithmetic operations:

shift\_expr ::= a\_expr | shift\_expr ( "<<" | ">>" ) a\_expr

These operators accept integers as arguments. They shift the first argument to

the left or right by the number of bits

given by the second argument.

A right shift by n bits is defined as division by pow(2,n). A left shift by n bits is defined as multiplication with

pow(2,n).

Note: In the current implementation, the right-hand operand is required to be at most sys.maxsize. If the righthand

operand is larger than sys.maxsize an OverflowError exception is raised.

## 6.8 Binary bitwise operations

Each of the three bitwise operations has a different priority level:

```
and_expr ::= shift_expr | and_expr "&" shift_expr
xor_expr ::= and_expr | xor_expr "^" and_expr
or_expr ::= xor_expr | or_expr "|" xor_expr
```

The & operator yields the bitwise AND of its arguments, which must be integers.

The ^ operator yields the bitwise XOR (exclusive OR) of its arguments, which must be integers.

The | operator yields the bitwise (inclusive) OR of its arguments, which must be integers.

# 6.9 Comparisons

Unlike C, all comparison operations in Python have the same priority, which is lower than that of any arithmetic,

shifting or bitwise operation. Also unlike C, expressions like a < b < c have the interpretation that is conventional

in mathematics:

```
comparison ::= or_expr ( comp_operator or_expr )*
comp_operator ::= "<" | ">" | "==" | ">=" | "<=" | "!="
| "is" ["not"] | ["not"] "in"
```

Comparisons yield boolean values: True or False.

Comparisons can be chained arbitrarily, e.g.,  $x < y \le z$  is equivalent to x < y and  $y \le z$ , except that y is

evaluated only once (but in both cases z is not evaluated at all when x < y is found to be false).

Formally, if a, b, c, ..., y, z are expressions and op1, op2, ..., opN are comparison operators, then a op1 b op2 c

... y opN z is equivalent to a op1 b and b op2 c and ... y opN z, except that each expression

is evaluated at most once.

Note that a op1 b op2 c doesn't imply any kind of comparison between a and c, so that, e.g., x < y > z is

perfectly legal (though perhaps not pretty).

The operators <, >, ==, >=, <=, and != compare the values of two objects. The objects need not have the same

type. If both are numbers, they are converted to a common type. Otherwise, the == and != operators always consider

objects of different types to be unequal, while the <, >, >= and <= operators raise a TypeError when comparing

objects of different types that do not implement these operators for the given pair of types. You can control comparison

behavior of objects of non-built-in types by defining rich comparison methods like \_\_gt\_\_(), described in section

Basic customization.

Comparison of objects of the same type depends on the type:

- Numbers are compared arithmetically.
- The values float('NaN') and Decimal('NaN') are special. The are identical to themselves, x is x

but are not equal to themselves, x != x. Additionally, comparing any value to a nota-number value will return

False. For example, both 3 < float('NaN') and float('NaN') < 3 will return False.

- Bytes objects are compared lexicographically using the numeric values of their elements.
- Strings are compared lexicographically using the numeric equivalents (the result of the built-in function ord())

of their characters. 3 String and bytes object can't be compared!

 Tuples and lists are compared lexicographically using comparison of corresponding elements. This means that to compare equal, each element must compare equal and the two sequences must be of the same type and have

the same length.

If not equal, the sequences are ordered the same as their first differing elements. For example,  $[1,2,x] \le$ 

[1,2,y] has the same value as  $x \le y$ . If the corresponding element does not exist, the shorter sequence is

ordered first (for example, [1,2] < [1,2,3]).

• Mappings (dictionaries) compare equal if and only if they have the same (key, value) pairs. Order comparisons

```
('<', '<=', '>=', '>') raise TypeError.
```

• Sets and frozensets define comparison operators to mean subset and superset tests. Those relations do not define

total orderings (the two sets {1,2} and {2,3} are not equal, nor subsets of one another, nor supersets of one

another). Accordingly, sets are not appropriate arguments for functions which depend on total ordering. For

example, min(), max(), and sorted() produce undefined results given a list of sets as inputs.

 Most other objects of built-in types compare unequal unless they are the same object; the choice whether one

object is considered smaller or larger than another one is made arbitrarily but consistently within one execution

of a program.

Comparison of objects of the differing types depends on whether either of the types provide explicit support for the

comparison. Most numeric types can be compared with one another. When cross-type comparison is not supported,

the comparison method returns NotImplemented. The operators in and not in test for membership. x in

s evaluates to true if x is a member of s, and false otherwise. x not in s returns the negation of x in s. All

built-in sequences and set types support this as well as dictionary, for which in tests whether a the dictionary has a

given key. For container types such as list, tuple, set, frozenset, dict, or collections.deque, the expression x in y is

equivalent to any(x is e or x == e for e in y).

3 While comparisons between strings make sense at the byte level, they may be counter-intuitive to users. For example, the strings "\u00C7"

and "\u0327\u0043" compare differently, even though they both represent the same unicode character (LATIN CAPITAL LETTER C WITH

CEDILLA). To compare strings in a human recognizable way, compare using unicodedata.normalize().

For the string and bytes types, x in y is true if and only if x is a substring of y. An equivalent test is y.find(x)

!= -1. Empty strings are always considered to be a substring of any other string, so "" in "abc" will return

True.

For user-defined classes which define the \_\_contains\_\_() method, x in y is true if and only if

y.\_\_contains\_\_(x) is true.

For user-defined classes which do not define \_\_contains\_\_() but do define \_\_iter\_\_(), x in y is true if

some value z with x == z is produced while iterating over y. If an exception is raised during the iteration, it is as if

in raised that exception.

Lastly, the old-style iteration protocol is tried: if a class defines \_\_getitem\_\_(), x in y is true if and only if

there is a non-negative integer index i such that x == y[i], and all lower integer indices do not raise IndexError

exception. (If any other exception is raised, it is as if in raised that exception).

The operator not in is defined to have the inverse true value of in.

The operators is and is not test for object identity: x is y is true if and only if x and y are the same object. x

is not y yields the inverse truth value. 4

## 6.10 Boolean operations

or\_test ::= and\_test | or\_test "or" and\_test and\_test ::= not\_test | and\_test "and" not\_test not\_test ::= comparison | "not" not\_test

In the context of Boolean operations, and also when expressions are used by control flow statements, the following

values are interpreted as false: False, None, numeric zero of all types, and empty strings and containers (including

strings, tuples, lists, dictionaries, sets and frozensets). All other values are interpreted as true. User-defined objects

can customize their truth value by providing a \_\_bool\_\_() method.

The operator not yields True if its argument is false, False otherwise.

The expression x and y first evaluates x; if x is false, its value is returned; otherwise, y is evaluated and the resulting

value is returned.

The expression x or y first evaluates x; if x is true, its value is returned; otherwise, y is evaluated and the resulting

value is returned.

(Note that neither and nor or restrict the value and type they return to False and True, but rather return the last

evaluated argument. This is sometimes useful, e.g., if s is a string that should be replaced by a default value if it is

empty, the expression s or 'foo' yields the desired value. Because not has to invent a value anyway, it does not

bother to return a value of the same type as its argument, so e.g., not 'foo' yields False, not ".)

# 6.11 Conditional expressions

conditional\_expression ::= or\_test ["if" or\_test "else" expression]

expression ::= conditional\_expression | lambda\_form

expression\_nocond ::= or\_test | lambda\_form\_nocond

Conditional expressions (sometimes called a "ternary operator") have the lowest priority of all Python operations.

4 Due to automatic garbage-collection, free lists, and the dynamic nature of descriptors, you may notice seemingly unusual behaviour in certain

uses of the is operator, like those involving comparisons between instance methods, or constants. Check their documentation for more info.

The expression x if C else y first evaluates the condition, C (not x); if C is true, x is evaluated and its value is

returned; otherwise, y is evaluated and its value is returned.

#### 6.12 Lambdas

lambda\_form ::= "lambda" [parameter\_list]: expression

lambda\_form\_nocond ::= "lambda" [parameter\_list]: expression\_nocond

Lambda forms (lambda expressions) have the same syntactic position as expressions. They are a shorthand to create

anonymous functions; the expression lambda arguments: expression yields a function object. The

unnamed object behaves like a function object defined with

def <lambda>(arguments):

return expression

See section Function definitions for the syntax of parameter lists. Note that functions created with lambda forms

cannot contain statements or annotations.

# 6.13 Expression lists

expression\_list ::= expression ( "," expression )\* [","]

An expression list containing at least one comma yields a tuple. The length of the tuple is the number of expressions

in the list. The expressions are evaluated from left to right.

The trailing comma is required only to create a single tuple (a.k.a. a singleton); it is optional in all other cases. A

single expression without a trailing comma doesn't create a tuple, but rather yields the value of that expression. (To

create an empty tuple, use an empty pair of parentheses: ().)

#### 6.14 Evaluation order

Python evaluates expressions from left to right. Notice that while evaluating an assignment, the right-hand side is

evaluated before the left-hand side.

In the following lines, expressions will be evaluated in the arithmetic order of their

```
suffixes:
expr1, expr2, expr3, expr4
(expr1, expr2, expr3, expr4)
{expr1: expr2, expr3: expr4}
expr1 + expr2 * (expr3 - expr4)
expr1(expr2, expr3, *expr4, **expr5)
expr3, expr4 = expr1, expr2
```

## 6.15 Operator precedence

The following table summarizes the operator precedences in Python, from lowest precedence (least binding) to highest

precedence (most binding). Operators in the same box have the same precedence. Unless the syntax is explicitly given,

operators are binary. Operators in the same box group left to right (except for comparisons, including tests, which all

have the same precedence and chain from left to right — see section Comparisons — and exponentiation, which

groups from right to left).

Operator Description

lambda Lambda expression

if - else Conditional expression

or Boolean OR

and Boolean AND

not x Boolean NOT

in, not in, is, is not, <, <=, >, >=, !=,

==

Comparisons, including membership tests and identity tests

| Bitwise OR

^ Bitwise XOR

& Bitwise AND

<<, >> Shifts

```
+, - Addition and subtraction

*, /, //, % Multiplication, division, remainder 5
+x, -x, ~x Positive, negative, bitwise NOT

** Exponentiation 6
x[index], x[index:index],
x(arguments...), x.attribute
Subscription, slicing, call, attribute reference
(expressions...),
```

[expressions...], {key:

value...}, {expressions...}

Binding or tuple display, list display, dictionary display, set

display

5The % operator is also used for string formatting; the same precedence applies.

6The power operator \*\* binds less tightly than an arithmetic or bitwise unary operator on its right, that is, 2\*\*-1 is 0.5.

#### **CHAPTER SEVEN SIMPLE STATEMENTS**

Simple statements are comprised within a single logical line. Several simple statements may occur on a single line

```
separated by semicolons. The syntax for simple statements is:
```

```
simple_stmt ::= expression_stmt
| assert_stmt
| assignment_stmt
| augmented_assignment_stmt
| pass_stmt
| del_stmt
| return_stmt
| yield_stmt
| raise_stmt
| break_stmt
| continue_stmt
| import_stmt
| global_stmt
| nonlocal_stmt
```

# 7.1 Expression statements

Expression statements are used (mostly interactively) to compute and write a value, or (usually) to call a procedure (a

function that returns no meaningful result; in Python, procedures return the value None). Other uses of expression

statements are allowed and occasionally useful. The syntax for an expression statement is:

```
expression_stmt ::= expression_list
```

An expression statement evaluates the expression list (which may be a single expression).

In interactive mode, if the value is not None, it is converted to a string using the built-in repr() function and the

resulting string is written to standard output on a line by itself (except if the result is None, so that procedure calls do

## 7.2 Assignment statements

Assignment statements are used to (re)bind names to values and to modify attributes or items of mutable objects:

```
assignment_stmt ::= (target_list "=")+ (expression_list | yield_expression)
target_list ::= target ("," target)* [","]
target ::= identifier
| "(" target_list ")"
| "[" target_list "]"
| attributeref
| subscription
| slicing
| "*" target
```

(See section Primaries for the syntax definitions for the last three symbols.)

An assignment statement evaluates the expression list (remember that this can be a single expression or a commaseparated

list, the latter yielding a tuple) and assigns the single resulting object to each of the target lists, from left to

right.

Assignment is defined recursively depending on the form of the target (list). When a target is part of a mutable object

(an attribute reference, subscription or slicing), the mutable object must ultimately perform the assignment and decide

about its validity, and may raise an exception if the assignment is unacceptable. The rules observed by various types

and the exceptions raised are given with the definition of the object types (see section The standard type hierarchy).

Assignment of an object to a target list, optionally enclosed in parentheses or square brackets, is recursively defined

as follows.

• If the target list is a single target: The object is assigned to that target.

 If the target list is a comma-separated list of targets: The object must be an iterable with the same number of

items as there are targets in the target list, and the items are assigned, from left to right, to the corresponding

targets.

 If the target list contains one target prefixed with an asterisk, called a "starred" target: The object must be

a sequence with at least as many items as there are targets in the target list, minus one. The first items of

the sequence are assigned, from left to right, to the targets before the starred target. The final items of the

sequence are assigned to the targets after the starred target. A list of the remaining items in the sequence

is then assigned to the starred target (the list can be empty).

 Else: The object must be a sequence with the same number of items as there are targets in the target list,

and the items are assigned, from left to right, to the corresponding targets.

Assignment of an object to a single target is recursively defined as follows.

- If the target is an identifier (name):
- If the name does not occur in a global or nonlocal statement in the current code block: the name is

bound to the object in the current local namespace.

 Otherwise: the name is bound to the object in the global namespace or the outer namespace determined by

nonlocal, respectively.

The name is rebound if it was already bound. This may cause the reference count for the object previously

bound to the name to reach zero, causing the object to be deallocated and its destructor (if it has one) to be

called.

• If the target is a target list enclosed in parentheses or in square brackets: The object must be an iterable with the

same number of items as there are targets in the target list, and its items are assigned, from left to right, to the

corresponding targets.

• If the target is an attribute reference: The primary expression in the reference is evaluated. It should yield an

object with assignable attributes; if this is not the case, TypeError is raised. That object is then asked to assign

the assigned object to the given attribute; if it cannot perform the assignment, it raises an exception (usually but

not necessarily AttributeError). Note: If the object is a class instance and the attribute reference occurs

on both sides of the assignment operator, the RHS expression, a.x can access either an instance attribute or (if

no instance attribute exists) a class attribute. The LHS target a.x is always set as an instance attribute, creating

it if necessary. Thus, the two occurrences of a.x do not necessarily refer to the same attribute: if the RHS

expression refers to a class attribute, the LHS creates a new instance attribute as the target of the assignment:

class Cls:

x = 3 # class variable

inst = Cls()

inst.x = inst.x + 1 # writes inst.x as 4 leaving Cls.x as 3

This description does not necessarily apply to descriptor attributes, such as properties created with

property().

• If the target is a subscription: The primary expression in the reference is evaluated. It should yield either

a mutable sequence object (such as a list) or a mapping object (such as a dictionary). Next, the subscript

expression is evaluated.

If the primary is a mutable sequence object (such as a list), the subscript must yield an integer. If it is negative,

the sequence's length is added to it. The resulting value must be a nonnegative integer less than the sequence's

length, and the sequence is asked to assign the assigned object to its item with

that index. If the index is out of

range, IndexError is raised (assignment to a subscripted sequence cannot add new items to a list).

If the primary is a mapping object (such as a dictionary), the subscript must have a type compatible with the

mapping's key type, and the mapping is then asked to create a key/datum pair which maps the subscript to the

assigned object. This can either replace an existing key/value pair with the same key value, or insert a new

key/value pair (if no key with the same value existed).

For user-defined objects, the \_\_setitem\_\_() method is called with appropriate arguments.

• If the target is a slicing: The primary expression in the reference is evaluated. It should yield a mutable sequence

object (such as a list). The assigned object should be a sequence object of the same type. Next, the lower and

upper bound expressions are evaluated, insofar they are present; defaults are zero and the sequence's length. The

bounds should evaluate to integers. If either bound is negative, the sequence's length is added to it. The resulting

bounds are clipped to lie between zero and the sequence's length, inclusive. Finally, the sequence object is asked

to replace the slice with the items of the assigned sequence. The length of the slice may be different from the

length of the assigned sequence, thus changing the length of the target sequence, if the object allows it.

CPython implementation detail: In the current implementation, the syntax for targets is taken to be the same as for

expressions, and invalid syntax is rejected during the code generation phase, causing less detailed error messages.

WARNING: Although the definition of assignment implies that overlaps between the left-hand side and the righthand

side are 'safe' (for example a, b = b, a swaps two variables), overlaps within the collection of assigned-to

variables are not safe! For instance, the following program prints [0, 2]:

$$x = [0, 1]$$

```
i = 0
i, x[i] = 1, 2
print(x)
```

# 7.2.1 Augmented assignment statements

Augmented assignment is the combination, in a single statement, of a binary operation and an assignment statement:

augmented\_assignment\_stmt ::= augtarget augop (expression\_list |
yield\_expression)

augtarget ::= identifier | attributeref | subscription | slicing

(See section Primaries for the syntax definitions for the last three symbols.)

An augmented assignment evaluates the target (which, unlike normal assignment statements, cannot be an unpacking)

and the expression list, performs the binary operation specific to the type of assignment on the two operands, and

assigns the result to the original target. The target is only evaluated once.

An augmented assignment expression like x += 1 can be rewritten as x = x + 1 to achieve a similar, but not

exactly equal effect. In the augmented version, x is only evaluated once. Also, when possible, the actual operation is

performed in-place, meaning that rather than creating a new object and assigning that to the target, the old object is

modified instead.

With the exception of assigning to tuples and multiple targets in a single statement, the assignment done by augmented

assignment statements is handled the same way as normal assignments. Similarly, with the exception of the possible inplace

behavior, the binary operation performed by augmented assignment is the same as the normal binary operations.

For targets which are attribute references, the same caveat about class and instance attributes applies as for regular

assignments.

#### 7.3 The assert statement

Assert statements are a convenient way to insert debugging assertions into a program:

assert\_stmt ::= "assert" expression ["," expression]

The simple form, assert expression, is equivalent to

if \_\_debug\_\_:

if not expression: raise AssertionError

The extended form, assert expression1, expression2, is equivalent to

if \_\_debug\_\_:

if not expression1: raise AssertionError(expression2)

These equivalences assume that \_\_debug\_\_ and AssertionError refer to the built-in variables with those names.

In the current implementation, the built-in variable \_\_debug\_\_ is True under normal circumstances, False when

optimization is requested (command line option -O). The current code generator emits no code for an assert statement

when optimization is requested at compile time. Note that it is unnecessary to include the source code for the

expression that failed in the error message; it will be displayed as part of the stack trace.

Assignments to \_\_debug\_\_ are illegal. The value for the built-in variable is determined when the interpreter starts.

### 7.4 The pass statement

pass stmt ::= "pass"

pass is a null operation — when it is executed, nothing happens. It is useful as a placeholder when a statement is

required syntactically, but no code needs to be executed, for example:

def f(arg): pass # a function that does nothing (yet)

class C: pass # a class with no methods (yet)

#### 7.5 The del statement

del\_stmt ::= "del" target\_list

Deletion is recursively defined very similar to the way assignment is defined. Rather than spelling it out in full details,

here are some hints.

Deletion of a target list recursively deletes each target, from left to right.

Deletion of a name removes the binding of that name from the local or global namespace, depending on whether the

name occurs in a global statement in the same code block. If the name is unbound, a NameError exception will

be raised.

Deletion of attribute references, subscriptions and slicings is passed to the primary object involved; deletion of a slicing

is in general equivalent to assignment of an empty slice of the right type (but even this is determined by the sliced

object). Changed in version 3.2: Previously it was illegal to delete a name from the local namespace if it occurs as a

free variable in a nested block.

### 7.6 The return statement

return\_stmt ::= "return" [expression\_list]

return may only occur syntactically nested in a function definition, not within a nested class definition.

If an expression list is present, it is evaluated, else None is substituted.

return leaves the current function call with the expression list (or None) as return value.

When return passes control out of a try statement with a finally clause, that finally clause is executed

before really leaving the function.

In a generator function, the return statement indicates that the generator is done and will cause StopIteration

to be raised. The returned value (if any) is used as an argument to construct StopIteration and becomes the

StopIteration.value attribute.

# 7.7 The yield statement

yield\_stmt ::= yield\_expression

The yield statement is only used when defining a generator function, and is only used in the body of the generator

function. Using a yield statement in a function definition is sufficient to cause that definition to create a generator

function instead of a normal function.

When a generator function is called, it returns an iterator known as a generator iterator, or more commonly, a generator.

The body of the generator function is executed by calling the next() function on the generator repeatedly until it

raises an exception.

When a yield statement is executed, the state of the generator is frozen and the value of expression\_list is

returned to next()'s caller. By "frozen" we mean that all local state is retained, including the current bindings of

local variables, the instruction pointer, and the internal evaluation stack: enough information is saved so that the next

time next() is invoked, the function can proceed exactly as if the yield statement were just another external call.

The yield statement is allowed in the try clause of a try ... finally construct. If the generator is not resumed

before it is finalized (by reaching a zero reference count or by being garbage collected), the generator-iterator's

close() method will be called, allowing any pending finally clauses to execute.

When yield from <expr> is used, it treats the supplied expression as a subiterator, producing values from it

until the underlying iterator is exhausted.

#### 7.8 The raise statement

raise\_stmt ::= "raise" [expression ["from" expression]]

If no expressions are present, raise re-raises the last exception that was active in the current scope. If no exception

is active in the current scope, a RuntimeError exception is raised indicating that this is an error.

Otherwise, raise evaluates the first expression as the exception object. It must be either a subclass or an instance of

BaseException. If it is a class, the exception instance will be obtained when needed by instantiating the class with

no arguments.

The type of the exception is the exception instance's class, the value is the instance itself.

A traceback object is normally created automatically when an exception is raised and attached to it as the

\_\_traceback\_\_ attribute, which is writable. You can create an exception and set your own traceback in one step

using the with\_traceback() exception method (which returns the same exception instance, with its traceback set

to its argument), like so:

raise Exception("foo occurred").with\_traceback(tracebackobj)

The from clause is used for exception chaining: if given, the second expression must be another exception class or

instance, which will then be attached to the raised exception as the \_\_cause\_\_ attribute (which is writable). If the

raised exception is not handled, both exceptions will be printed:

```
>>> try:
```

```
... print(1 / 0)
```

... except Exception as exc:

... raise RuntimeError("Something bad happened") from exc

. . .

Traceback (most recent call last):

File "<stdin>", line 2, in <module>

ZeroDivisionError: int division or modulo by zero

The above exception was the direct cause of the following exception:

Traceback (most recent call last):

File "<stdin>", line 4, in <module>

RuntimeError: Something bad happened

A similar mechanism works implicitly if an exception is raised inside an exception handler: the previous exception is

then attached as the new exception's context attribute:

>>> try:

... print(1 / 0)

... except:

... raise RuntimeError("Something bad happened")

. . .

Traceback (most recent call last):

File "<stdin>", line 2, in <module>

ZeroDivisionError: int division or modulo by zero

During handling of the above exception, another exception occurred:

Traceback (most recent call last):

File "<stdin>", line 4, in <module>

RuntimeError: Something bad happened

Additional information on exceptions can be found in section Exceptions, and information about handling exceptions

is in section The try statement.

### 7.9 The break statement

break\_stmt ::= "break"

break may only occur syntactically nested in a for or while loop, but not nested in a function or class definition

within that loop.

It terminates the nearest enclosing loop, skipping the optional else clause if the loop has one.

If a for loop is terminated by break, the loop control target keeps its current value.

When break passes control out of a try statement with a finally clause, that finally clause is executed before

really leaving the loop.

### 7.10 The continue statement

continue\_stmt ::= "continue"

continue may only occur syntactically nested in a for or while loop, but not nested

in a function or class

definition or finally clause within that loop. It continues with the next cycle of the nearest enclosing loop.

When continue passes control out of a try statement with a finally clause, that finally clause is executed

before really starting the next loop cycle.

# 7.11 The import statement

```
import_stmt ::= "import" module ["as" name] ( "," module ["as" name] )*
| "from" relative_module "import" identifier ["as" name]
( "," identifier ["as" name] )*
| "from" relative_module "import" "(" identifier ["as" name]
( "," identifier ["as" name] )* [","] ")"
| "from" module "import" "*"
module ::= (identifier ".")* identifier
relative_module ::= "."* module | "."+
name ::= identifier
```

The basic import statement (no from clause) is executed in two steps:

- 1. find a module, loading and initializing it if necessary
- 2. define a name or names in the local namespace for the scope where the import statement occurs.

When the statement contains multiple clauses (separated by commas) the two steps are carried out separately for each

clause, just as though the clauses had been separated out into individual import statements.

The details of the first step, finding and loading modules is described in greater detail in the section on the import

system, which also describes the various types of packages and modules that can be imported, as well as all the hooks

that can be used to customize the import system. Note that failures in this step may indicate either that the module

could not be located, or that an error occurred while initializing the module, which

includes execution of the module's code.

If the requested module is retrieved successfully, it will be made available in the local namespace in one of three ways:

- If the module name is followed by as, then the name following as is bound directly to the imported module.
- If no other name is specified, and the module being imported is a top level module, the module's name is bound

in the local namespace as a reference to the imported module

• If the module being imported is not a top level module, then the name of the top level package that contains the

module is bound in the local namespace as a reference to the top level package. The imported module must be

accessed using its full qualified name rather than directly

The from form uses a slightly more complex process:

- 1. find the module specified in the from clause loading and initializing it if necessary;
- 2. for each of the identifiers specified in the import clauses:
- (a) check if the imported module has an attribute by that name
- (b) if not, attempt to import a submodule with that name and then check the imported module again for that

attribute

- (c) if the attribute is not found, ImportError is raised.
- (d) otherwise, a reference to that value is bound in the local namespace, using the name in the as clause if it

is present, otherwise using the attribute name

### Examples:

import foo # foo imported and bound locally

import foo.bar.baz # foo.bar.baz imported, foo bound locally

import foo.bar.baz as fbb # foo.bar.baz imported and bound as fbb

from foo.bar import baz # foo.bar.baz imported and bound as baz

from foo import attr # foo imported and foo.attr bound as attr

If the list of identifiers is replaced by a star ('\*'), all public names defined in the module are bound in the local

namespace for the scope where the import statement occurs.

The public names defined by a module are determined by checking the module's namespace for a variable named

\_\_all\_\_\_; if defined, it must be a sequence of strings which are names defined or imported by that module. The

names given in \_\_all\_\_ are all considered public and are required to exist. If \_\_all\_\_ is not defined, the set of

public names includes all names found in the module's namespace which do not begin with an underscore character

('\_'). \_\_all\_\_ should contain the entire public API. It is intended to avoid accidentally exporting items that are

not part of the API (such as library modules which were imported and used within the module).

The from form with \* may only occur in a module scope. The wild card form of import — import \* — is only

allowed at the module level. Attempting to use it in class or function definitions will raise a SyntaxError.

When specifying what module to import you do not have to specify the absolute name of the module. When a module

or package is contained within another package it is possible to make a relative import within the same top package

without having to mention the package name. By using leading dots in the specified module or package after from

you can specify how high to traverse up the current package hierarchy without specifying exact names. One leading dot

means the current package where the module making the import exists. Two dots means up one package level. Three

dots is up two levels, etc. So if you execute from . import mod from a module in the pkg package then you

will end up importing pkg.mod. If you execute from ..subpkg2 import mod from within pkg.subpkg1

you will import pkg.subpkg2.mod. importlib.import\_module() is provided to support applications that determine which modules need to be

loaded dynamically.

#### 7.11.1 Future statements

A future statement is a directive to the compiler that a particular module should be compiled using syntax or semantics

that will be available in a specified future release of Python. The future statement is intended to ease migration to

future versions of Python that introduce incompatible changes to the language. It allows use of the new features on a

per-module basis before the release in which the feature becomes standard.

```
future_statement ::= "from" "__future__" "import" feature ["as" name]
("," feature ["as" name])*
| "from" "__future__" "import" "(" feature ["as" name]
("," feature ["as" name])* [","] ")"
feature ::= identifier
name ::= identifier
```

A future statement must appear near the top of the module. The only lines that can appear before a future statement

#### are:

- the module docstring (if any),
- comments,
- · blank lines, and
- other future statements.

The features recognized by Python 3.0 are absolute\_import, division, generators, unicode\_literals, print\_function, nested\_scopes and with\_statement. They are all redundant

because they are always enabled, and only kept for backwards compatibility.

A future statement is recognized and treated specially at compile time: Changes to the semantics of core constructs

are often implemented by generating different code. It may even be the case that a new feature introduces new incompatible

syntax (such as a new reserved word), in which case the compiler may need to parse the module differently.

Such decisions cannot be pushed off until runtime.

For any given release, the compiler knows which feature names have been defined, and raises a compile-time error if

a future statement contains a feature not known to it.

The direct runtime semantics are the same as for any import statement: there is a standard module \_\_future\_\_,

described later, and it will be imported in the usual way at the time the future statement is executed.

The interesting runtime semantics depend on the specific feature enabled by the future statement.

Note that there is nothing special about the statement:

import \_\_future\_\_ [as name]

That is not a future statement; it's an ordinary import statement with no special semantics or syntax restrictions.

Code compiled by calls to the built-in functions exec() and compile() that occur in a module M containing a

future statement will, by default, use the new syntax or semantics associated with the future statement. This can be

controlled by optional arguments to compile()— see the documentation of that function for details.

A future statement typed at an interactive interpreter prompt will take effect for the rest of the interpreter session. If an

interpreter is started with the -i option, is passed a script name to execute, and the script includes a future statement,

it will be in effect in the interactive session started after the script is executed.

### 7.12 The global statement

global\_stmt ::= "global" identifier ("," identifier)\*

The global statement is a declaration which holds for the entire current code block. It means that the listed identifiers

are to be interpreted as globals. It would be impossible to assign to a global variable without global, although free

variables may refer to globals without being declared global.

Names listed in a global statement must not be used in the same code block textually preceding that global

statement.

Names listed in a global statement must not be defined as formal parameters or in a for loop control target, class

definition, function definition, or import statement.

CPython implementation detail: The current implementation does not enforce the latter two restrictions, but programs

should not abuse this freedom, as future implementations may enforce them or silently change the meaning of

the program.

Programmer's note: the global is a directive to the parser. It applies only to code parsed at the same time as the

global statement. In particular, a global statement contained in a string or code object supplied to the built-in

exec() function does not affect the code block containing the function call, and code contained in such a string is

unaffected by global statements in the code containing the function call. The same applies to the eval() and

compile() functions.

#### 7.13 The nonlocal statement

nonlocal\_stmt ::= "nonlocal" identifier ("," identifier)\*

The nonlocal statement causes the listed identifiers to refer to previously bound variables in the nearest enclosing

scope. This is important because the default behavior for binding is to search the local namespace first. The statement

allows encapsulated code to rebind variables outside of the local scope besides the global (module) scope.

Names listed in a nonlocal statement, unlike to those listed in a global statement, must refer to pre-existing bindings

in an enclosing scope (the scope in which a new binding should be created cannot be determined unambiguously).

Names listed in a nonlocal statement must not collide with pre-existing bindings in the local scope.

#### CHAPTER EIGHT COMPOUND STATEMENTS

Compound statements contain (groups of) other statements; they affect or control the execution of those other statements

in some way. In general, compound statements span multiple lines, although in simple incarnations a whole

compound statement may be contained in one line.

The if, while and for statements implement traditional control flow constructs. try specifies exception handlers

and/or cleanup code for a group of statements, while the with statement allows the execution of initialization and

finalization code around a block of code. Function and class definitions are also syntactically compound statements.

Compound statements consist of one or more 'clauses.' A clause consists of a header and a 'suite.' The clause headers

of a particular compound statement are all at the same indentation level. Each clause header begins with a uniquely

identifying keyword and ends with a colon. A suite is a group of statements controlled by a clause. A suite can be one

or more semicolon-separated simple statements on the same line as the header, following the header's colon, or it can

be one or more indented statements on subsequent lines. Only the latter form of suite can contain nested compound

statements; the following is illegal, mostly because it wouldn't be clear to which if clause a following else clause

would belong:

if test1: if test2: print(x)

Also note that the semicolon binds tighter than the colon in this context, so that in the following example, either all or

none of the print() calls are executed:

if x < y < z: print(x); print(y); print(z)

Summarizing:

compound\_stmt ::= if\_stmt

| while\_stmt

| for\_stmt

```
| try_stmt
| with_stmt
| funcdef
| classdef
suite ::= stmt_list NEWLINE | NEWLINE INDENT statement+ DEDENT
statement ::= stmt_list NEWLINE | compound_stmt
stmt_list ::= simple_stmt (";" simple_stmt)* [";"]
```

Note that statements always end in a NEWLINE possibly followed by a DEDENT. Also note that optional continuation

clauses always begin with a keyword that cannot start a statement, thus there are no ambiguities (the 'dangling else'

problem is solved in Python by requiring nested if statements to be indented).

The formatting of the grammar rules in the following sections places each clause on a separate line for clarity.

#### 8.1 The if statement

The if statement is used for conditional execution:

```
if_stmt ::= "if" expression ":" suite
( "elif" expression ":" suite )*
["else" ":" suite]
```

It selects exactly one of the suites by evaluating the expressions one by one until one is found to be true (see section

Boolean operations for the definition of true and false); then that suite is executed (and no other part of the if

statement is executed or evaluated). If all expressions are false, the suite of the else clause, if present, is executed.

#### 8.2 The while statement

The while statement is used for repeated execution as long as an expression is true:

```
while_stmt ::= "while" expression ":" suite
["else" ":" suite]
```

This repeatedly tests the expression and, if it is true, executes the first suite; if the expression is false (which may be

the first time it is tested) the suite of the else clause, if present, is executed and the loop terminates.

A break statement executed in the first suite terminates the loop without executing the else clause's suite. A

continue statement executed in the first suite skips the rest of the suite and goes back to testing the expression.

#### 8.3 The for statement

The for statement is used to iterate over the elements of a sequence (such as a string, tuple or list) or other iterable

object:

for\_stmt ::= "for" target\_list "in" expression\_list ":" suite

["else" ":" suite]

The expression list is evaluated once; it should yield an iterable object. An iterator is created for the result of the

expression\_list. The suite is then executed once for each item provided by the iterator, in the order of ascending

indices. Each item in turn is assigned to the target list using the standard rules for assignments (see Assignment

statements), and then the suite is executed. When the items are exhausted (which is immediately when the sequence is

empty or an iterator raises a StopIteration exception), the suite in the else clause, if present, is executed, and

the loop terminates.

A break statement executed in the first suite terminates the loop without executing the else clause's suite. A

continue statement executed in the first suite skips the rest of the suite and continues with the next item, or with the

else clause if there was no next item.

The suite may assign to the variable(s) in the target list; this does not affect the next item assigned to it.

Names in the target list are not deleted when the loop is finished, but if the sequence is empty, it will not have been

assigned to at all by the loop. Hint: the built-in function range() returns an iterator

of integers suitable to emulate

the effect of Pascal's for i := a to b do; e.g., list(range(3)) returns the list [0, 1, 2].

Note: There is a subtlety when the sequence is being modified by the loop (this can only occur for mutable sequences,

i.e. lists). An internal counter is used to keep track of which item is used next, and this is incremented on each iteration.

When this counter has reached the length of the sequence the loop terminates. This means that if the suite deletes the

current (or a previous) item from the sequence, the next item will be skipped (since it gets the index of the current

item which has already been treated). Likewise, if the suite inserts an item in the sequence before the current item, the

current item will be treated again the next time through the loop. This can lead to nasty bugs that can be avoided by

making a temporary copy using a slice of the whole sequence, e.g.,

for x in a[:]:

if x < 0: a.remove(x)

## 8.4 The try statement

The try statement specifies exception handlers and/or cleanup code for a group of statements:

```
try_stmt ::= try1_stmt | try2_stmt
try1_stmt ::= "try" ":" suite
("except" [expression ["as" target]] ":" suite)+
["else" ":" suite]
["finally" ":" suite]
try2_stmt ::= "try" ":" suite
"finally" ":" suite
```

The except clause(s) specify one or more exception handlers. When no exception occurs in the try clause, no

exception handler is executed. When an exception occurs in the try suite, a search for an exception handler is started.

This search inspects the except clauses in turn until one is found that matches the exception. An expression-less except

clause, if present, must be last; it matches any exception. For an except clause with an expression, that expression is

evaluated, and the clause matches the exception if the resulting object is "compatible" with the exception. An object

is compatible with an exception if it is the class or a base class of the exception object or a tuple containing an item

compatible with the exception.

If no except clause matches the exception, the search for an exception handler continues in the surrounding code and

on the invocation stack. 1

If the evaluation of an expression in the header of an except clause raises an exception, the original search for a handler

is canceled and a search starts for the new exception in the surrounding code and on the call stack (it is treated as if

the entire try statement raised the exception).

When a matching except clause is found, the exception is assigned to the target specified after the as keyword in that

except clause, if present, and the except clause's suite is executed. All except clauses must have an executable block.

When the end of this block is reached, execution continues normally after the entire try statement. (This means that if

two nested handlers exist for the same exception, and the exception occurs in the try clause of the inner handler, the

outer handler will not handle the exception.)

When an exception has been assigned using as target, it is cleared at the end of the except clause. This is as if

except E as N:

foo

was translated to

except E as N:

try:

foo

1 The exception is propagated to the invocation stack unless there is a finally

clause which happens to raise another exception. That new exception causes the old one to be lost.

finally:

del N

This means the exception must be assigned to a different name to be able to refer to it after the except clause. Exceptions

are cleared because with the traceback attached to them, they form a reference cycle with the stack frame,

keeping all locals in that frame alive until the next garbage collection occurs.

Before an except clause's suite is executed, details about the exception are stored in the sys module and can be

access via sys.exc\_info(). sys.exc\_info() returns a 3-tuple consisting of the exception class, the exception

instance and a traceback object (see section The standard type hierarchy) identifying the point in the program where the

exception occurred. sys.exc\_info() values are restored to their previous values (before the call) when returning

from a function that handled an exception.

The optional else clause is executed if and when control flows off the end of the try clause. 2 Exceptions in the

else clause are not handled by the preceding except clauses.

If finally is present, it specifies a 'cleanup' handler. The try clause is executed, including any except and

else clauses. If an exception occurs in any of the clauses and is not handled, the exception is temporarily saved.

The finally clause is executed. If there is a saved exception it is re-raised at the end of the finally clause. If

the finally clause raises another exception, the saved exception is set as the context of the new exception. If the

finally clause executes a return or break statement, the saved exception is discarded:

def f():

try:

```
1/0
```

finally:

return 42

>>> f()

42

The exception information is not available to the program during execution of the finally clause.

When a return, break or continue statement is executed in the try suite of a try...finally statement, the

finally clause is also executed 'on the way out.' A continue statement is illegal in the finally clause. (The

reason is a problem with the current implementation— this restriction may be lifted in the future).

Additional information on exceptions can be found in section Exceptions, and information on using the raise statement

to generate exceptions may be found in section The raise statement.

#### 8.5 The with statement

The with statement is used to wrap the execution of a block with methods defined by a context manager (see section

With Statement Context Managers). This allows common try...except...finally usage patterns to be encapsulated

for convenient reuse.

```
with_stmt ::= "with" with_item ("," with_item)* ":" suite with_item ::= expression ["as" target]
```

The execution of the with statement with one "item" proceeds as follows:

- 1. The context expression (the expression given in the with\_item) is evaluated to obtain a context manager.
- 2. The context manager's \_\_exit\_\_() is loaded for later use.
- 3. The context manager's \_\_enter\_\_() method is invoked.
- 2 Currently, control "flows off the end" except in the case of an exception or the execution of a return, continue, or break statement.

4. If a target was included in the with statement, the return value from enter () is assigned to it. Note: The with statement guarantees that if the enter () method returns without an error, then exit () will always be called. Thus, if an error occurs during the assignment to the target list, it will be treated the same as an error occurring within the suite would be. See step 6 below. 5. The suite is executed. 6. The context manager's \_\_exit\_\_() method is invoked. If an exception caused the suite to be exited, its type, value, and traceback are passed as arguments to \_\_exit\_\_(). Otherwise, three None arguments are supplied. If the suite was exited due to an exception, and the return value from the exit () method was false, the exception is reraised. If the return value was true, the exception is suppressed, and execution continues with the statement following the with statement. If the suite was exited for any reason other than an exception, the return value from \_\_exit\_\_() is ignored, and execution proceeds at the normal location for the kind of exit that was taken. With more than one item, the context managers are processed as if multiple with statements were nested: with A() as a, B() as b: suite is equivalent to with A() as a: with B() as b: suite 8.6 Function definitions

A function definition defines a user-defined function object (see section The standard type hierarchy):

funcdef ::= [decorators] "def" funcname "(" [parameter\_list] ")" ["->" expression] decorators ::= decorator+

```
decorator ::= "@" dotted_name ["(" [parameter_list [","]] ")"] NEWLINE
dotted_name ::= identifier ("." identifier)*
parameter_list ::= (defparameter ",")*
( "*" [parameter] ("," defparameter)* ["," "**" parameter]
| "**" parameter
| defparameter [","] )
parameter ::= identifier [":" expression]
defparameter ::= parameter ["=" expression]
funcname ::= identifier
```

A function definition is an executable statement. Its execution binds the function name in the current local namespace

to a function object (a wrapper around the executable code for the function). This function object contains a reference

to the current global namespace as the global namespace to be used when the function is called.

The function definition does not execute the function body; this gets executed only when the function is called. 3

A function definition may be wrapped by one or more decorator expressions. Decorator expressions are evaluated

when the function is defined, in the scope that contains the function definition. The result must be a callable, which

is invoked with the function object as the only argument. The returned value is bound to the function name instead of

the function object. Multiple decorators are applied in nested fashion. For example, the following code

```
@f1(arg)
@f2
def func(): pass
is equivalent to
def func(): pass
func = f1(arg)(f2(func))
```

When one or more parameters have the form parameter = expression, the

function is said to have "default parameter

values." For a parameter with a default value, the corresponding argument may be omitted from a call, in which case

the parameter's default value is substituted. If a parameter has a default value, all following parameters up until the

"\*" must also have a default value— this is a syntactic restriction that is not expressed by the grammar.

Default parameter values are evaluated when the function definition is executed. This means that the expression

is evaluated once, when the function is defined, and that the same "precomputed" value is used for each call. This

is especially important to understand when a default parameter is a mutable object, such as a list or a dictionary: if

the function modifies the object (e.g. by appending an item to a list), the default value is in effect modified. This is

generally not what was intended. A way around this is to use None as the default, and explicitly test for it in the body

of the function, e.g.:

def whats\_on\_the\_telly(penguin=None):

if penguin is None:

penguin = []

penguin.append("property of the zoo")

return penguin

Function call semantics are described in more detail in section Calls. A function call always assigns values to all

parameters mentioned in the parameter list, either from position arguments, from keyword arguments, or from default

values. If the form "\*identifier" is present, it is initialized to a tuple receiving any excess positional parameters,

defaulting to the empty tuple. If the form "\*\*identifier" is present, it is initialized to a new dictionary receiving

any excess keyword arguments, defaulting to a new empty dictionary. Parameters after "\*" or "\*identifier" are

keyword-only parameters and may only be passed used keyword arguments.

Parameters may have annotations of the form ": expression" following the

parameter name. Any parameter

may have an annotation even those of the form \*identifier or \*\*identifier. Functions may have "return"

annotation of the form "-> expression" after the parameter list. These annotations can be any valid Python

expression and are evaluated when the function definition is executed.

Annotations may be evaluated in a different

order than they appear in the source code. The presence of annotations does not change the semantics of a function. The

annotation values are available as values of a dictionary keyed by the parameters' names in the \_\_annotations\_\_

attribute of the function object.

It is also possible to create anonymous functions (functions not bound to a name), for immediate use in expressions.

This uses lambda forms, described in section Lambdas. Note that the lambda form is merely a shorthand for a

simplified function definition; a function defined in a "def" statement can be passed around or assigned to another

name just like a function defined by a lambda form. The "def" form is actually more powerful since it allows the

execution of multiple statements and annotations.

Programmer's note: Functions are first-class objects. A "def" form executed inside a function definition defines a

local function that can be returned or passed around. Free variables used in the nested function can access the local

3 A string literal appearing as the first statement in the function body is transformed into the function's \_\_doc\_\_ attribute and therefore the function's docstring.

variables of the function containing the def. See section Naming and binding for details.

#### 8.7 Class definitions

A class definition defines a class object (see section The standard type hierarchy): classdef ::= [decorators] "class" classname [inheritance] ":" suite

inheritance ::= "(" [parameter\_list] ")"

classname ::= identifier

A class definition is an executable statement. The inheritance list usually gives a list of base classes (see Customizing

class creation for more advanced uses), so each item in the list should evaluate to a class object which allows

subclassing. Classes without an inheritance list inherit, by default, from the base class object; hence,

class Foo:

pass

is equivalent to

class Foo(object):

pass

The class's suite is then executed in a new execution frame (see Naming and binding), using a newly created local

namespace and the original global namespace. (Usually, the suite contains mostly function definitions.) When the

class's suite finishes execution, its execution frame is discarded but its local namespace is saved. 4 A class object is

then created using the inheritance list for the base classes and the saved local namespace for the attribute dictionary.

The class name is bound to this class object in the original local namespace.

Class creation can be customized heavily using metaclasses.

Classes can also be decorated: just like when decorating functions,

@f1(arg)

@f2

class Foo: pass

is equivalent to

class Foo: pass

Foo = f1(arg)(f2(Foo))

The evaluation rules for the decorator expressions are the same as for function decorators. The result must be a class

object, which is then bound to the class name.

Programmer's note: Variables defined in the class definition are class attributes;

they are shared by instances. Instance

attributes can be set in a method with self.name = value. Both class and instance attributes are accessible

through the notation "self.name", and an instance attribute hides a class attribute with the same name when accessed

in this way. Class attributes can be used as defaults for instance attributes, but using mutable values there can

lead to unexpected results. Descriptors can be used to create instance variables with different implementation details.

#### CHAPTER NINE TOP-LEVEL COMPONENTS

The Python interpreter can get its input from a number of sources: from a script passed to it as standard input or as

program argument, typed in interactively, from a module source file, etc. This chapter gives the syntax used in these cases.

## 9.1 Complete Python programs

While a language specification need not prescribe how the language interpreter is invoked, it is useful to have a notion

of a complete Python program. A complete Python program is executed in a minimally initialized environment: all

built-in and standard modules are available, but none have been initialized, except for sys (various system services),

builtins (built-in functions, exceptions and None) and \_\_main\_\_. The latter is used to provide the local and

global namespace for execution of the complete program.

The syntax for a complete Python program is that for file input, described in the next section.

The interpreter may also be invoked in interactive mode; in this case, it does not read and execute a complete program

but reads and executes one statement (possibly compound) at a time. The initial environment is identical to that of a

complete program; each statement is executed in the namespace of \_\_main\_\_.

Under Unix, a complete program can be passed to the interpreter in three forms: with the -c string command line

option, as a file passed as the first command line argument, or as standard input. If the file or standard input is a tty

device, the interpreter enters interactive mode; otherwise, it executes the file as a complete program.

### 9.2 File input

All input read from non-interactive files has the same form:

file input ::= (NEWLINE | statement)\*

This syntax is used in the following situations:

- when parsing a complete Python program (from a file or from a string);
- · when parsing a module;
- when parsing a string passed to the exec() function;

# 9.3 Interactive input

Input in interactive mode is parsed using the following grammar:

interactive\_input ::= [stmt\_list] NEWLINE | compound\_stmt NEWLINE

Note that a (top-level) compound statement must be followed by a blank line in interactive mode; this is needed to

help the parser detect the end of the input.

# 9.4 Expression input

There are two forms of expression input. Both ignore leading whitespace. The string argument to eval() must have

the following form:

eval input ::= expression list NEWLINE\*

Note: to read 'raw' input line without interpretation, you can use the readline() method of file objects, including

sys.stdin.

### **CHAPTER TEN FULL GRAMMAR SPECIFICATION**

This is the full Python grammar, as it is read by the parser generator and used to parse Python source files:

```
# Grammar for Python
# Note: Changing the grammar specified in this file will most likely
# require corresponding changes in the parser module
# (../Modules/parsermodule.c). If you can't make the changes to
# that module yourself, please co-ordinate the required changes
# with someone who can; ask around on python-dev for help. Fred
# Drake <fdrake@acm.org> will probably be listening there.
# NOTE WELL:
# "How to Change Python's Grammar"
# Start symbols for the grammar:
# single input is a single interactive statement;
# file_input is a module or sequence of commands read from an input file;
# eval input is the input for the eval() functions.
# NB: compound stmt in single input is followed by extra NEWLINE!
single_input: NEWLINE | simple_stmt | compound_stmt NEWLINE
file input: (NEWLINE | stmt)* ENDMARKER
eval input: testlist NEWLINE* ENDMARKER
decorator: '@' dotted_name [ '(' [arglist] ')' ] NEWLINE
decorators: decorator+
decorated: decorators (classdef | funcdef)
funcdef: 'def' NAME parameters ['->' test] ':' suite
parameters: '(' [typedargslist] ')'
typedargslist: (tfpdef ['=' test] (',' tfpdef ['=' test])* [','
['*' [tfpdef] (',' tfpdef ['=' test])* [',' '**' tfpdef] | '**' tfpdef]]
| '*' [tfpdef] (',' tfpdef ['=' test])* [',' '**' tfpdef] | '**' tfpdef)
tfpdef: NAME [':' test]
varargslist: (vfpdef ['=' test] (',' vfpdef ['=' test])* [','
['*' [vfpdef] (',' vfpdef ['=' test])* [',' '**' vfpdef] | '**' vfpdef]]
```

```
| '*' [vfpdef] (',' vfpdef ['=' test])* [',' '**' vfpdef] | '**' vfpdef)
vfpdef: NAME
stmt: simple stmt | compound stmt
simple stmt: small stmt (';' small stmt)* [';'] NEWLINE
small stmt: (expr stmt | del stmt | pass stmt | flow stmt |
import stmt | global stmt | nonlocal stmt | assert stmt)
expr_stmt: testlist_star_expr (augassign (yield_expr|testlist) |
('=' (yield_expr|testlist_star_expr))*)
testlist star expr: (test|star expr) (',' (test|star expr))* [',']
augassign: ('+=' | '-=' | '*=' | '/=' | '%=' | '&=' | '|=' | '^=' |
'<<=' | '>>=' | '**=' | '//=')
# For normal assignments, additional restrictions enforced by the interpreter
del_stmt: 'del' exprlist
pass stmt: 'pass'
flow_stmt: break_stmt | continue_stmt | return_stmt | raise_stmt | yield_stmt
break_stmt: 'break'
continue_stmt: 'continue'
return stmt: 'return' [testlist]
yield stmt: yield expr
raise stmt: 'raise' [test ['from' test]]
import stmt: import name | import from
import name: 'import' dotted as names
# note below: the ('.' | '...') is necessary because '...' is tokenized as ELLIPSIS
import from: ('from' (('.' | '...')* dotted name | ('.' | '...')+)
'import' ('*' | '(' import_as_names ')' | import_as_names))
import as name: NAME ['as' NAME]
dotted_as_name: dotted_name ['as' NAME]
import_as_names: import_as_name (',' import_as_name)* [',']
dotted as names: dotted as name (',' dotted as name)*
```

```
dotted name: NAME ('.' NAME)*
global stmt: 'global' NAME (',' NAME)*
nonlocal stmt: 'nonlocal' NAME (',' NAME)*
assert stmt: 'assert' test [',' test]
compound stmt: if stmt | while stmt | for stmt | try stmt | with stmt | funcdef |
classdef if stmt: 'if' test ':' suite ('elif' test ':' suite)* ['else' ':' suite]
while stmt: 'while' test ':' suite ['else' ':' suite]
for stmt: 'for' exprlist 'in' testlist ':' suite ['else' ':' suite]
try stmt: ('try' ':' suite
((except clause ':' suite)+
['else' ':' suite]
['finally' ':' suite] |
'finally' ':' suite))
with stmt: 'with' with item (',' with item)* ':' suite
with item: test ['as' expr]
# NB compile.c makes sure that the default except clause is last
except clause: 'except' [test ['as' NAME]]
suite: simple stmt | NEWLINE INDENT stmt+ DEDENT
test: or test ['if' or test 'else' test] | lambdef
test_nocond: or_test | lambdef_nocond
lambdef: 'lambda' [varargslist] ':' test
lambdef nocond: 'lambda' [varargslist] ':' test nocond
or test: and test ('or' and test)*
and test: not test ('and' not test)*
not_test: 'not' not_test | comparison
comparison: expr (comp_op expr)*
# <> isn't actually a valid comparison operator in Python.
comp op: '<'|'>'|'=='|'>='|'<>'|'!='|'in'|'not' 'in'|'is'|'is' 'not'
star_expr: '*' expr
expr: xor_expr ('|' xor_expr)*
xor expr: and expr ('A' and expr)*
and expr: shift expr ('&' shift expr)*
```

```
shift_expr: arith_expr (('<<'|'>>') arith_expr)*
arith expr: term (('+'|'-') term)*
term: factor (('*'|'/'|'%'|'//') factor)*
factor: ('+'|'-'|'~') factor | power
power: atom trailer* ['**' factor]
atom: ('(' [yield_expr|testlist_comp] ')' |
'[' [testlist comp] ']' |
'{' [dictorsetmaker] '}' |
NAME | NUMBER | STRING+ | '...' | 'None' | 'True' | 'False')
testlist_comp: (test|star_expr) ( comp_for | (',' (test|star_expr))* [','] )
trailer: '(' [arglist] ')' | '[' subscriptlist ']' | '.' NAME
subscriptlist: subscript (',' subscript)* [',']
subscript: test | [test] ':' [test] [sliceop]
sliceop: ':' [test]
exprlist: (expr|star_expr) (',' (expr|star_expr))* [',']
testlist: test (',' test)* [',']
dictorsetmaker: ( (test ':' test (comp_for | (',' test ':' test)* [','])) |
(test (comp_for | (',' test)* [','])) )
classdef: 'class' NAME ['(' [arglist] ')'] ':' suite
arglist: (argument ',')* (argument [',']
|'*' test (',' argument)* [',' '**' test]
|'**' test)
# The reason that keywords are test nodes instead of NAME is that using NAME
# results in an ambiguity. ast.c makes sure it's a NAME.
argument: test [comp for] | test '=' test # Really [keyword '='] test
comp iter: comp for | comp if
comp_for: 'for' exprlist 'in' or_test [comp_iter]
comp if: 'if' test nocond [comp iter]
# not used in grammar, but may appear in "node" passed from Parser to Compiler
encoding_decl: NAME
yield expr: 'yield' [yield arg]
```

yield\_arg: 'from' test | testlist

#### APPENDIX A GLOSSARY

>>> The default Python prompt of the interactive shell. Often seen for code examples which can be executed

interactively in the interpreter.

... The default Python prompt of the interactive shell when entering code for an indented code block or within a

pair of matching left and right delimiters (parentheses, square brackets or curly braces).

2to3 A tool that tries to convert Python 2.x code to Python 3.x code by handling most of the incompatibilities which

can be detected by parsing the source and traversing the parse tree.

2to3 is available in the standard library as lib2to3; a standalone entry point is provided as

Tools/scripts/2to3. See 2to3-reference.

abstract base class Abstract base classes complement duck-typing by providing a way to define interfaces when

other techniques like hasattr() would be clumsy or subtly wrong (for example with magic methods).

ABCs introduce virtual subclasses, which are classes that don't inherit from a class but are still recognized

by isinstance() and issubclass(); see the abc module documentation. Python comes with many

built-in ABCs for data structures (in the collections.abc module), numbers (in the numbers module),

streams (in the io module), import finders and loaders (in the importlib.abc module). You can create your

own ABCs with the abc module.

argument A value passed to a function (or method) when calling the function. There are two types of arguments:

• keyword argument: an argument preceded by an identifier (e.g. name=) in a function call or passed as a

value in a dictionary preceded by \*\*. For example, 3 and 5 are both keyword arguments in the following

calls to complex():

complex(real=3, imag=5)
complex(\*\*{'real': 3, 'imag': 5})

• positional argument: an argument that is not a keyword argument. Positional arguments can appear at the

beginning of an argument list and/or be passed as elements of an iterable preceded by \*. For example, 3

and 5 are both positional arguments in the following calls:

complex(3, 5)

complex(\*(3, 5))

Arguments are assigned to the named local variables in a function body. See the Calls section for the rules

governing this assignment. Syntactically, any expression can be used to represent an argument; the evaluated

value is assigned to the local variable.

See also the parameter glossary entry, the FAQ question on the difference between arguments and parameters,

attribute A value associated with an object which is referenced by name using dotted expressions. For example, if

an object o has an attribute a it would be referenced as o.a.

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BDFL Benevolent Dictator For Life, a.k.a. Guido van Rossum, Python's creator.

bytes-like object An object that supports the bufferobjects, like bytes, bytearray or memoryview. Bytes-like

objects can be used for various operations that expect binary data, such as compression, saving to a binary file

or sending over a socket. Some operations need the binary data to be mutable, in which case not all bytes-like

objects can apply.

bytecode Python source code is compiled into bytecode, the internal representation of a Python program in the

CPython interpreter. The bytecode is also cached in .pyc and .pyo files so that executing the same file is

faster the second time (recompilation from source to bytecode can be avoided). This "intermediate language" is

said to run on a virtual machine that executes the machine code corresponding to each bytecode. Do note that

bytecodes are not expected to work between different Python virtual machines, nor to be stable between Python

releases.

A list of bytecode instructions can be found in the documentation for the dis module.

class A template for creating user-defined objects. Class definitions normally contain method definitions which

operate on instances of the class.

coercion The implicit conversion of an instance of one type to another during an operation which involves two

arguments of the same type. For example, int(3.15) converts the floating point number to the integer 3, but

in 3+4.5, each argument is of a different type (one int, one float), and both must be converted to the same type

before they can be added or it will raise a TypeError. Without coercion, all arguments of even compatible

types would have to be normalized to the same value by the programmer, e.g., float(3)+4.5 rather than just

3+4.5.

complex number An extension of the familiar real number system in which all numbers are expressed as a sum of

a real part and an imaginary part. Imaginary numbers are real multiples of the imaginary unit (the square root

of -1), often written i in mathematics or j in engineering. Python has built-in support for complex numbers,

which are written with this latter notation; the imaginary part is written with a j suffix, e.g., 3+1j. To get

access to complex equivalents of the math module, use cmath. Use of complex numbers is a fairly advanced

mathematical feature. If you're not aware of a need for them, it's almost certain you can safely ignore them.

context manager An object which controls the environment seen in a with statement by defining \_\_enter\_\_()

CPython The canonical implementation of the Python programming language, as

distributed on python.org. The term

"CPython" is used when necessary to distinguish this implementation from others such as Jython or IronPython.

decorator A function returning another function, usually applied as a function transformation using the @wrapper

syntax. Common examples for decorators are classmethod() and staticmethod().

The decorator syntax is merely syntactic sugar, the following two function definitions are semantically equivalent:

```
def f(...):
...
f = staticmethod(f)
@staticmethod
def f(...):
```

The same concept exists for classes, but is less commonly used there. See the documentation for function

definitions and class definitions for more about decorators.

descriptor Any object which defines the methods \_\_get\_\_(), \_\_set\_\_(), or \_\_delete\_\_(). When a class

attribute is a descriptor, its special binding behavior is triggered upon attribute lookup. Normally, using a.b to

get, set or delete an attribute looks up the object named b in the class dictionary for a, but if b is a descriptor,

the respective descriptor method gets called. Understanding descriptors is a key to a deep understanding of Python because they are the basis for many features including functions, methods, properties, class methods,

static methods, and reference to super classes.

For more information about descriptors' methods, see Implementing Descriptors.

dictionary An associative array, where arbitrary keys are mapped to values. The keys can be any object with

hash() andeq() me	hods. Called a hash in Perl.
-------------------	------------------------------

docstring A string literal which appears as the first expression in a class, function or module. While ignored when

the suite is executed, it is recognized by the compiler and put into the \_\_doc\_\_ attribute of the enclosing class,

function or module. Since it is available via introspection, it is the canonical place for documentation of the

object.

duck-typing A programming style which does not look at an object's type to determine if it has the right interface;

instead, the method or attribute is simply called or used ("If it looks like a duck and quacks like a duck, it must

be a duck.") By emphasizing interfaces rather than specific types, well-designed code improves its flexibility

by allowing polymorphic substitution. Duck-typing avoids tests using type() or isinstance(). (Note,

however, that duck-typing can be complemented with abstract base classes.) Instead, it typically employs

hasattr() tests or EAFP programming.

EAFP Easier to ask for forgiveness than permission. This common Python coding style assumes the existence

of valid keys or attributes and catches exceptions if the assumption proves false. This clean and fast style is

characterized by the presence of many try and except statements. The technique contrasts with the LBYL

style common to many other languages such as C.

expression A piece of syntax which can be evaluated to some value. In other words, an expression is an accumulation

of expression elements like literals, names, attribute access, operators or function calls which all return a value.

In contrast to many other languages, not all language constructs are expressions. There are also statements

which cannot be used as expressions, such as if. Assignments are also statements, not expressions.

extension module A module written in C or C++, using Python's C API to interact with the core and with user code.

file object An object exposing a file-oriented API (with methods such as read() or write()) to an underlying

resource. Depending on the way it was created, a file object can mediate access to a real on-disk file or to another

type of storage or communication device (for example standard input/output, in-

memory buffers, sockets, pipes,

etc.). File objects are also called file-like objects or streams.

There are actually three categories of file objects: raw binary files, buffered binary files and text files. Their

interfaces are defined in the io module. The canonical way to create a file object is by using the open()

function.

file-like object A synonym for file object.

finder An object that tries to find the loader for a module. It must implement either a method named

find\_loader() or a method named find\_module().

importlib.abc.Finder for an abstract base class.

floor division Mathematical division that rounds down to nearest integer. The floor division operator is //. For

example, the expression 11 // 4 evaluates to 2 in contrast to the 2.75 returned by float true division. Note

that (-11) // 4 is -3 because that is -2.75 rounded downward.

function A series of statements which returns some value to a caller. It can also be passed zero or more arguments

which may be used in the execution of the body. See also parameter, method, and the Function definitions

section.

function annotation An arbitrary metadata value associated with a function parameter or return value. Its syntax is

explained in section Function definitions. Annotations may be accessed via the \_\_annotations\_\_ special

attribute of a function object.

Python itself does not assign any particular meaning to function annotations. They are intended to be interpreted

by third-party libraries or tools.

\_\_future\_\_ A pseudo-module which programmers can use to enable new language features which are not compatible

with the current interpreter.

By importing the \_\_future\_\_ module and evaluating its variables, you can see when a new feature was first

added to the language and when it becomes the default:

```
>>> import __future__
>>> future .division
```

\_Feature((2, 2, 0, 'alpha', 2), (3, 0, 0, 'alpha', 0), 8192)

garbage collection The process of freeing memory when it is not used anymore. Python performs garbage collection

via reference counting and a cyclic garbage collector that is able to detect and break reference cycles.

generator A function which returns an iterator. It looks like a normal function except that it contains yield

statements for producing a series a values usable in a for-loop or that can be retrieved one at a time with the

next() function. Each yield temporarily suspends processing, remembering the location execution state

(including local variables and pending try-statements). When the generator resumes, it picks-up where it left-off

(in contrast to functions which start fresh on every invocation).

generator expression An expression that returns an iterator. It looks like a normal expression followed by a for

expression defining a loop variable, range, and an optional if expression. The combined expression generates

values for an enclosing function:

>>> sum(i\*i for i in range(10)) # sum of squares 0, 1, 4, ... 81

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GIL See global interpreter lock.

global interpreter lock The mechanism used by the CPython interpreter to assure that only one thread executes

Python bytecode at a time. This simplifies the CPython implementation by making the object model (including

critical built-in types such as dict) implicitly safe against concurrent access. Locking the entire interpreter

makes it easier for the interpreter to be multi-threaded, at the expense of much of the parallelism afforded by

multi-processor machines.

However, some extension modules, either standard or third-party, are designed

so as to release the GIL when

doing computationally-intensive tasks such as compression or hashing. Also, the GIL is always released when

doing I/O.

Past efforts to create a "free-threaded" interpreter (one which locks shared data at a much finer granularity)

have not been successful because performance suffered in the common singleprocessor case. It is believed

that overcoming this performance issue would make the implementation much more complicated and therefore

costlier to maintain.

hashable An object is hashable if it has a hash value which never changes during its lifetime (it needs a

\_\_hash\_\_() method), and can be compared to other objects (it needs an \_\_eq\_\_() method). Hashable

objects which compare equal must have the same hash value.

Hashability makes an object usable as a dictionary key and a set member, because these data structures use the

hash value internally.

All of Python's immutable built-in objects are hashable, while no mutable containers (such as lists or dictionaries)

are. Objects which are instances of user-defined classes are hashable by default; they all compare unequal

(except with themselves), and their hash value is their id().

IDLE An Integrated Development Environment for Python. IDLE is a basic editor and interpreter environment which

ships with the standard distribution of Python.

immutable An object with a fixed value. Immutable objects include numbers, strings and tuples. Such an object

cannot be altered. A new object has to be created if a different value has to be stored. They play an important

role in places where a constant hash value is needed, for example as a key in a dictionary.

import path A list of locations (or path entries) that are searched by the path based finder for modules to import.

During import, this list of locations usually comes from sys.path, but for subpackages it may also come from

the parent package's \_\_path\_\_ attribute.

importing The process by which Python code in one module is made available to Python code in another module.

importer An object that both finds and loads a module; both a finder and loader object.

interactive Python has an interactive interpreter which means you can enter statements and expressions at the interpreter

prompt, immediately execute them and see their results. Just launch python with no arguments

(possibly by selecting it from your computer's main menu). It is a very powerful way to test out new ideas or

inspect modules and packages (remember help(x)).

interpreted Python is an interpreted language, as opposed to a compiled one, though the distinction can be blurry

because of the presence of the bytecode compiler. This means that source files can be run directly without explicitly

creating an executable which is then run. Interpreted languages typically have a shorter development/debug

cycle than compiled ones, though their programs generally also run more slowly. See also interactive.

iterable An object capable of returning its members one at a time. Examples of iterables include all sequence types

(such as list, str, and tuple) and some non-sequence types like dict, file objects, and objects of any

classes you define with an \_\_iter\_\_() or \_\_getitem\_\_() method. Iterables can be used in a for loop

and in many other places where a sequence is needed (zip(), map(), ...). When an iterable object is passed

as an argument to the built-in function iter(), it returns an iterator for the object. This iterator is good for one

pass over the set of values. When using iterables, it is usually not necessary to call iter() or deal with iterator

objects yourself. The for statement does that automatically for you, creating a temporary unnamed variable to

hold the iterator for the duration of the loop. See also iterator, sequence, and generator.

iterator An object representing a stream of data. Repeated calls to the iterator's \_\_next\_\_() method (or passing

it to the built-in function next()) return successive items in the stream. When no more data are available

a StopIteration exception is raised instead. At this point, the iterator object is exhausted and any further

calls to its \_\_next\_\_() method just raise StopIteration again. Iterators are required to have an

\_\_iter\_\_() method that returns the iterator object itself so every iterator is also iterable and may be used in

most places where other iterables are accepted. One notable exception is code which attempts multiple iteration

passes. A container object (such as a list) produces a fresh new iterator each time you pass it to the iter()

function or use it in a for loop. Attempting this with an iterator will just return the same exhausted iterator

object used in the previous iteration pass, making it appear like an empty container.

More information can be found in typeiter.

key function A key function or collation function is a callable that returns a value used for sorting or ordering. For

example, locale.strxfrm() is used to produce a sort key that is aware of locale specific sort conventions.

A number of tools in Python accept key functions to control how elements are ordered or grouped. They include

min(), max(), sorted(), list.sort(), heapq.nsmallest(), heapq.nlargest(), and itertools.groupby().

There are several ways to create a key function. For example. the str.lower() method can serve as a key

function for case insensitive sorts. Alternatively, an ad-hoc key function can be built from a lambda expression

such as lambda r: (r[0], r[2]). Also, the operator module provides three key

function constructors:

attrgetter(), itemgetter(), and methodcaller(). See the Sorting HOW TO for examples

of how to create and use key functions.

keyword argument See argument.

lambda An anonymous inline function consisting of a single expression which is evaluated when the function is

called. The syntax to create a lambda function is lambda [arguments]: expression

LBYL Look before you leap. This coding style explicitly tests for pre-conditions before making calls or lookups.

This style contrasts with the EAFP approach and is characterized by the presence of many if statements.

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In a multi-threaded environment, the LBYL approach can risk introducing a race condition between "the looking"

and "the leaping". For example, the code, if key in mapping: return mapping[key] can

fail if another thread removes key from mapping after the test, but before the lookup. This issue can be solved

with locks or by using the EAFP approach.

list A built-in Python sequence. Despite its name it is more akin to an array in other languages than to a linked list

since access to elements are O(1).

list comprehension A compact way to process all or part of the elements in a sequence and return a list with the

results. result = [' $\{:\#04x\}$ '.format(x) for x in range(256) if x % 2 == 0] generates a list of strings containing even hex numbers (0x..) in the range from 0 to 255. The if clause is optional.

If omitted, all elements in range(256) are processed.

loader An object that loads a module. It must define a method named load\_module(). A loader is typically

mapping A container object that supports arbitrary key lookups and implements the methods specified

in the Mapping or MutableMapping abstract base classes. Examples include dict,

collections.defaultdict, collections.OrderedDict and collections.Counter.

meta path finder A finder returned by a search of sys.meta\_path. Meta path finders are related to, but different

from path entry finders.

metaclass The class of a class. Class definitions create a class name, a class dictionary, and a list of base classes.

The metaclass is responsible for taking those three arguments and creating the class. Most object oriented

programming languages provide a default implementation. What makes Python special is that it is possible to

create custom metaclasses. Most users never need this tool, but when the need arises, metaclasses can provide

powerful, elegant solutions. They have been used for logging attribute access, adding thread-safety, tracking

object creation, implementing singletons, and many other tasks.

More information can be found in Customizing class creation.

method A function which is defined inside a class body. If called as an attribute of an instance of that class, the

method will get the instance object as its first argument (which is usually called self). See function and nested

scope.

method resolution order Method Resolution Order is the order in which base classes are searched for a member

during lookup. See The Python 2.3 Method Resolution Order.

module An object that serves as an organizational unit of Python code. Modules have a namespace containing

arbitrary Python objects. Modules are loaded into Python by the process of importing.

MRO See method resolution order.

mutable Mutable objects can change their value but keep their id(). See also immutable.

named tuple Any tuple-like class whose indexable elements are also accessible using named attributes (for example,

time.localtime() returns a tuple-like object where the year is accessible either with an index such as

t[0] or with a named attribute like t.tm\_year).

A named tuple can be a built-in type such as time.struct\_time, or it can be created with a

regular class definition. A full featured named tuple can also be created with the factory function

collections.namedtuple(). The latter approach automatically provides extra features such as a selfdocumenting

representation like Employee(name='jones', title='programmer').

namespace The place where a variable is stored. Namespaces are implemented as dictionaries. There are the local,

global and built-in namespaces as well as nested namespaces in objects (in methods). Namespaces support

modularity by preventing naming conflicts. For instance, the functions builtins.open() and os.open()

are distinguished by their namespaces. Namespaces also aid readability and maintainability by making it clear

which module implements a function. For instance, writing random.seed() or itertools.islice()

makes it clear that those functions are implemented by the random and itertools modules, respectively.

nested scope The ability to refer to a variable in an enclosing definition. For instance, a function defined inside

another function can refer to variables in the outer function. Note that nested scopes by default work only for

reference and not for assignment. Local variables both read and write in the innermost scope. Likewise, global

variables read and write to the global namespace. The nonlocal allows writing to outer scopes.

new-style class Old name for the flavor of classes now used for all class objects. In earlier Python versions,

only new-style classes could use Python's newer, versatile features like
slots, descriptors, properties,
getattribute(), class methods, and static methods.

object Any data with state (attributes or value) and defined behavior (methods). Also the ultimate base class of any

new-style class.

package A Python module which can contain submodules or recursively, subpackages. Technically, a package is a

Python module with an \_\_path\_\_ attribute.

parameter A named entity in a function (or method) definition that specifies an argument (or in some cases, arguments)

that the function can accept. There are five types of parameters:

• positional-or-keyword: specifies an argument that can be passed either positionally or as a keyword argument.

This is the default kind of parameter, for example foo and bar in the following: def func(foo, bar=None): ...

• positional-only: specifies an argument that can be supplied only by position. Python has no syntax for

defining positional-only parameters. However, some built-in functions have positional-only parameters

(e.g. abs()).

• keyword-only: specifies an argument that can be supplied only by keyword. Keyword-only parameters can

be defined by including a single var-positional parameter or bare \* in the parameter list of the function

definition before them, for example kw\_only1 and kw\_only2 in the following: def func(arg, \*, kw\_only1, kw\_only2): ...

• var-positional: specifies that an arbitrary sequence of positional arguments can be provided (in addition

to any positional arguments already accepted by other parameters). Such a parameter can be defined by

prepending the parameter name with \*, for example args in the following: def func(\*args, \*\*kwargs): ...

 var-keyword: specifies that arbitrarily many keyword arguments can be provided (in addition to any keyword

arguments already accepted by other parameters). Such a parameter can be defined by prepending

the parameter name with \*\*, for example kwargs in the example above.

Parameters can specify both optional and required arguments, as well as default values for some optional arguments.

See also the argument glossary entry, the FAQ question on the difference

between arguments and parameters,

the inspect. Parameter class, the Function definitions section.

path entry A single location on the import path which the path based finder consults to find modules for importing.

path entry finder A finder returned by a callable on sys.path\_hooks (i.e. a path entry hook) which knows how

to locate modules given a path entry.

path entry hook A callable on the sys.path\_hook list which returns a path entry finder if it knows how to find

modules on a specific path entry.

path based finder One of the default meta path finders which searches an import path for modules.

portion A set of files in a single directory (possibly stored in a zip file) that contribute to a namespace package.

provisional package A provisional package is one which has been deliberately excluded from the standard library's

backwards compatibility guarantees. While major changes to such packages are not expected, as long as they

are marked provisional, backwards incompatible changes (up to and including removal of the package) may

occur if deemed necessary by core developers. Such changes will not be made gratuitously – they will occur

only if serious flaws are uncovered that were missed prior to the inclusion of the package.

This process allows the standard library to continue to evolve over time, without locking in problematic design

errors for extended periods of time.

Python 3000 Nickname for the Python 3.x release line (coined long ago when the release of version 3 was something

in the distant future.) This is also abbreviated "Py3k".

Pythonic An idea or piece of code which closely follows the most common idioms of the Python language, rather

than implementing code using concepts common to other languages. For

example, a common idiom in Python

is to loop over all elements of an iterable using a for statement. Many other languages don't have this type of

construct, so people unfamiliar with Python sometimes use a numerical counter instead:

for i in range(len(food)): print(food[i]) As opposed to the cleaner, Pythonic method: for piece in food: print(piece) qualified name A dotted name showing the "path" from a module's global scope to a class, function or method defined in that module. >>> class C: ... class D: ... def meth(self): ... pass >>> C. qualname\_\_\_ 'C' >>> C.D. qualname 'C.D' >>> C.D.meth. qualname 'C.D.meth' When used to refer to modules, the fully qualified name means the entire dotted path to the module, including any parent packages, e.g. email.mime.text:

'email.mime.text'

>>> import email.mime.text

>>> email.mime.text.\_\_name\_\_

reference count The number of references to an object. When the reference count of an object drops to zero, it is

deallocated. Reference counting is generally not visible to Python code, but it is a

key element of the CPython

implementation. The sys module defines a getrefcount() function that programmers can call to return

the reference count for a particular object.

regular package A traditional package, such as a directory containing an \_\_init\_\_.py file.

\_\_slots\_\_ A declaration inside a class that saves memory by pre-declaring space for instance attributes and eliminating

instance dictionaries. Though popular, the technique is somewhat tricky to get right and is best reserved for

rare cases where there are large numbers of instances in a memory-critical application.

sequence An iterable which supports efficient element access using integer indices via the \_\_getitem\_\_() special

method and defines a \_\_len\_\_() method that returns the length of the sequence. Some built-in sequence types

are list, str, tuple, and bytes. Note that dict also supports \_\_getitem\_\_() and \_\_len\_\_(), but

is considered a mapping rather than a sequence because the lookups use arbitrary immutable keys rather than

integers.

slice An object usually containing a portion of a sequence. A slice is created using the subscript notation, []

with colons between numbers when several are given, such as in variable\_name[1:3:5]. The bracket

(subscript) notation uses slice objects internally.

special method A method that is called implicitly by Python to execute a certain operation on a type, such as addition.

Such methods have names starting and ending with double underscores. Special methods are documented in

Special method names.

statement A statement is part of a suite (a "block" of code). A statement is either an expression or a one of several

constructs with a keyword, such as if, while or for.

struct sequence A tuple with named elements. Struct sequences expose an interface similar to named tuple in that

elements can either be accessed either by index or as an attribute. However, they do not have any of the named

tuple methods like \_make() or \_asdict(). Examples of struct sequences include sys.float info and

the return value of os.stat().

triple-quoted string A string which is bound by three instances of either a quotation mark (") or an apostrophe

('). While they don't provide any functionality not available with single-quoted strings, they are useful for a

number of reasons. They allow you to include unescaped single and double quotes within a string and they can

span multiple lines without the use of the continuation character, making them especially useful when writing

docstrings.

type The type of a Python object determines what kind of object it is; every object has a type. An object's type is

accessible as its \_\_class\_\_ attribute or can be retrieved with type(obj).

universal newlines A manner of interpreting text streams in which all of the following are recognized as ending a

line: the Unix end-of-line convention '\n', the Windows convention '\r\n', and the old Macintosh convention

'\r'. See str.splitlines() for an additional use.

view The objects returned from dict.keys(), dict.values(), and dict.items() are called dictionary

views. They are lazy sequences that will see changes in the underlying dictionary. To force the dictionary view

to become a full list use list(dictview). See dict-views.

virtual machine A computer defined entirely in software. Python's virtual machine executes the bytecode emitted

by the bytecode compiler.

Zen of Python Listing of Python design principles and philosophies that are helpful in understanding and using the

language. The listing can be found by typing "import this" at the interactive prompt.