

Procedures 201:

Higher Level Calling Conventions

Slides revised 3/25/2014 by Patrick Kelley

Procedures

- Higher Level languages have adopted a standard
 - Referred to as C-style calling
 - Uses the stack to pass parameters and returns
 - Keeps local variables on the stack
 - Allows for recursive calling
- Every call to a procedure maintains a unique stack frame
- Registers are always preserved by the callee

Procedure Design

- Begins with the calling parameters and returns
- Avoid complexity
 - Use 32-bit data and pointers
 - Avoid bytes, half-words, strings, and arrays
- Add the data required; this is the stack-frame size
- Caller puts arguments onto the stack
 - In reverse order
 - Without moving stack pointer
 - Must ensure that there is no stack overflow

Ex. Factorial function

- Imagine an hll (say, C++) function for factorial
- The prototype is: `int factorial(int input)`
- Let's use 32-bit unsigned ints
- So we need 2 words, one for input and one for the result
- Without knowing anything more, we can define the calling code for the fact function.

The Calling Stack

Special	Address	Data
	0x00377238	
	0x0037723C	
Input (sp-8)	0x00377240	0x0
return (sp-4)	0x00377244	0x0
Stack Pointer	0x00377248	0x000045F7
	0x0037724C	

- Remember, the top of the stack grows toward smaller addresses.
- Caller does not change the stack pointer
- Must check for stack overflow (grow past top of stack) if possible (not easily on our MIPS)

Example Call

```
.text
```

```
... # some code before this, then set up for the call
    li    $t0, 0xB          # get the parameter
    sw    $t0, -8(sp)       # put it on the stack frame
                                # return value is at -4(sp)
    jal   Factorial         # make the call
    lw    $t0, -4(sp)       # get the return value into $t0
```

...

- Note that we didn't care about preserving \$t0
- If we had some way to know the top of the stack, we could compare it plus our framesize to \$sp
- Caller did not change \$sp or initialize the return value
- Multiple parameters are put on stack in reverse order

Push and Pop

```
.text
```

```
...
```

```
# push a value (register) onto the stack
```

```
sw      $t0, -4(sp)          # any register will do
```

```
addiu   $sp, $sp, -4         # adjust stack pointer
```

```
...
```

```
# pop a value (register) from the stack
```

```
lw      $t0, 0(sp)          # any register will do
```

```
addiu   $sp, $sp, 4         # adjust stack pointer
```

```
...
```

- Pop values in the reverse order they were pushed
- Should check before pushing if there is room:
 - Platform dependent
 - On QTSPIM, put a label at end of data section:

```
.data
```

```
...    # data declarations
```

```
enddata:
```

```
.text
```

```
...
```

```
la      $t0, enddata        # any register will do
subu    $t0, $sp, $t0        # the register now holds the
                              # available stack space
```

Pushes and Pops

.text

...

```
# push multiple values (registers) onto the stack
sw      $t0, -4(sp)           # any registers
addiu   $sp, $sp, -4          # adjust stack pointer
sw      $t1, -4(sp)
addiu   $sp, $sp, -4          # adjust stack pointer
sw      $s5, -4(sp)
addiu   $sp, $sp, -4          # adjust stack pointer
```

...

```
# pop multiple values (registers) from the stack
# in reverse order
lw      $s5, 0(sp)           # any registers
addiu   $sp, $sp, 4          # adjust stack pointer
lw      $t1, 0(sp)
addiu   $sp, $sp, 4          # adjust stack pointer
lw      $t0, 0(sp)
addiu   $sp, $sp, 4          # adjust stack pointer
```

...

Pushes and Pops(alternate)

.text

...

push multiple values (registers) onto the stack

sw \$t0, -4(sp) # any registers

sw \$t1, -8(sp)

sw \$s5, -12(sp)

addiu \$sp, \$sp, -12 # adjust stack pointer

...

pop multiple values (registers) from the stack

lw \$t0, 8(sp) # any registers

lw \$t1, 4(sp)

lw \$s5, 0(sp)

addiu \$sp, \$sp, 12 # adjust stack pointer

...

- OR -

pop multiple values (registers) from the stack

addiu \$sp, \$sp, 12 # adjust stack pointer

lw \$t0, -4(sp) # any registers

lw \$t1, -8(sp)

lw \$s5, -12(sp)

...

Notice that pop order does not matter now

Procedure Design 2

- The called procedure:
 - Saves frame pointer
 - Copies the stack pointer to the frame pointer (so we can get it back later)
 - Moves the stack pointer to the top of the passed params
 - If it uses registers, it pushes them onto the stack
- Local variables are put on the stack as well
- Before return:
 - local variables are deallocated
 - Registers are popped off the stack
 - Copy frame pointer to stack pointer
 - Restore original frame pointer

Ex. Factorial function

- Let's implement this:

```
int Factorial (int input)
{
    int dummy; // to make it a little interesting
    dummy = 5;
    if (input == 0)
        return 1;
    else return input * Factorial(input - 1);
}
```
- And we'll use registers \$s0, \$s1, \$s2, and \$s3

The Factorial Stack Frame (at call)

	Special	Address	Data
		0x00377220	?
		0x00377224	?
		0x00377228	?
		0x0037722C	?
		0x00377230	?
		0x00377234	?
		0x00377238	?
		0x0037723C	?
	Input (sp-8)	0x00377240	0xB
	return (sp-4)	0x00377244	?
SP		0x00377248	0x000045F7
		0x0037724C	

Setting the Frame Pointer

```
.text
```

```
...
```

```
# Factorial expects an input at $sp - 8 and computes the factorial  
# of that input. It returns the value at $sp - 4. Since the  
# factorial algorithm is recursive, a stack frame is used...
```

```
Factorial:
```

```
# push current frame pointer onto stack ($fp or $30)
```

```
sw      $fp, -12(sp)          # remember -8 and -4 currently in use
```

```
move    $fp, $sp             # copy stack pointer to frame pointer
```

```
addiu   $sp, $sp, -12        # adjust stack pointer
```

```
...
```

The Factorial Stack Frame (\$fp set)

	Special	Address	Data
		0x00377220	?
		0x00377224	?
		0x00377228	?
		0x0037722C	?
		0x00377230	?
		0x00377234	?
		0x00377238	?
SP	Old Frame Pointer	0x0037723C	\$fp (old)
	Input (fp-8)	0x00377240	0xB
	return (fp-4)	0x00377244	?
FP		0x00377248	0x000045F7
		0x0037724C	

Saving Registers

`.text`

...

```
# Factorial expects an input at $sp - 8 and computes the factorial  
# of that input. It returns the value at $sp - 4. Since the  
# factorial algorithm is recursive, a stack frame is used...
```

Factorial:

```
# push current frame pointer onto stack ($fp or $30)
```

```
sw      $fp, -12(sp)          # remember -8 and -4 currently in use
```

```
move    $fp, $sp             # copy stack pointer to frame pointer
```

```
addiu   $sp, $sp, -12        # adjust stack pointer
```

```
# save $ra and any other registers we need
```

```
sw      $ra, -4(sp)          # return address and our other regs
```

```
sw      $s0, -8(sp)
```

```
sw      $s1, -12(sp)
```

```
sw      $s2, -16(sp)
```

```
sw      $s3, -20(sp)
```

```
addiu   $sp, $sp, -20        # adjust stack pointer
```

...

The Factorial Stack Frame (save regs)

	Special	Address	Data
		0x00377220	?
		0x00377224	?
SP	\$s3	0x00377228	\$s3 (old)
	\$s2	0x0037722C	\$s2 (old)
	\$s1	0x00377230	\$s1 (old)
	\$s0	0x00377234	\$s0 (old)
	\$ra	0x00377238	\$ra (old)
	Old Frame Pointer	0x0037723C	\$fp (old)
	Input (fp-8)	0x00377240	0xB
	return (fp-4)	0x00377244	?
FP		0x00377248	0x000045F7
		0x0037724C	

Space for Local Variables

`.text`

...

```
# Factorial expects an input at $sp - 8 and computes the factorial
# of that input. It returns the value at $sp - 4. Since the
# factorial algorithm is recursive, a stack frame is used...
```

Factorial:

```
# push current frame pointer onto stack ($fp or $30)
sw      $fp, -12(sp)          # remember -8 and -4 currently in use
move    $fp, $sp             # copy stack pointer to frame pointer
addiu   $sp, $sp, -12        # adjust stack pointer
```

```
# save $ra and any other registers we need
sw      $ra, -4(sp)           # any registers
sw      $s0, -8(sp)
sw      $s1, -12(sp)
sw      $s2, -16(sp)
sw      $s3, -20(sp)
addiu   $sp, $sp, -20        # adjust stack pointer
```

```
# reserve space for the local variable 'dummy' at 0($sp)
addiu   $sp, $sp, -4
li      $s0, 5                # for storing into 'dummy'
sw      $s0, 0($sp)          # store the local value
```

...

The Factorial Stack Frame (locals)

	Special	Address	Data
		0x00377220	?
SP	'dummy'	0x00377224	5
	\$s3	0x00377228	\$s3 (old)
	\$s2	0x0037722C	\$s2 (old)
	\$s1	0x00377230	\$s1 (old)
	\$s0	0x00377234	\$s0 (old)
	\$ra	0x00377238	\$ra (old)
	Old Frame Pointer	0x0037723C	\$fp (old)
	Input (fp-8)	0x00377240	0xB
	return (fp-4)	0x00377244	?
FP		0x00377248	0x000045F7
		0x0037724C	

The Factorial Algorithm

`.text`

`...`

`# reserve space for the local variable 'dummy' at 0($sp)`

`addiu $sp, $sp, -4`

`li $s0, 5 # for storing into 'dummy'`

`sw $s0, 0($sp) # store the local value`

`# load the input parameter into a register`

`lw $s1, -8($fp) # remember, $fp points where $sp was`
`# on the procedure call`

`# see if the input is 0 or not`

`bnez $s1, callFact # if not 0, do recursive call`

`li $s2, 1 # otherwise set the return ($s2) to 1`

`j doneFact # jump to return code`

`callFact:`

`addiu $s3, $s1, -1 # $s3 is parameter for recursive call`

`sw $s3, -8(sp) # put it on the stack frame`

`jal Factorial # make the call`

`lw $s3, -4(sp) # get the return value into $s3`

`multu $s3, $s1 # multiply the return * input`

`mflo $s2 # assume not bigger than LO and put`

`# in $s2 for return`

`...`

Preparing to Return

`.text`

`...`

```
# see if the input is 0 or not
bnez    $s1, callFact
li      $s2, 1
j       doneFact
```

```
# if not 0, do recursive call
# otherwise set the return ($s2) to 1
# jump to return code
```

`callFact:`

```
addiu   $s3, $s1, -1
sw      $s3, -8(sp)
jal     Factorial
lw      $s3, -4(sp)
multu   $s3, $s1
mflo    $s2
```

```
# $s3 is parameter for recursive call
# put it on the stack frame
# make the call
# get the return value into $s3
# multiply the return * input
# assume not bigger than LO and put
# in $s2 for return
```

`doneFact:`

```
sw      $s2, -4(fp)
```

```
# put our return value relative to $fp
```

```
# now we can begin cleanup prior to return
```

```
addiu   $sp, $sp, 4
```

```
# done with local variables so adjust $sp
```

`...`

The Factorial Stack Frame (free vars)

	Special	Address	Data
		0x00377220	?
		0x00377224	5
SP	\$s3	0x00377228	\$s3 (old)
	\$s2	0x0037722C	\$s2 (old)
	\$s1	0x00377230	\$s1 (old)
	\$s0	0x00377234	\$s0 (old)
	\$ra	0x00377238	\$ra (old)
	Old Frame Pointer	0x0037723C	\$fp (old)
	Input (fp-8)	0x00377240	0xB
	return (fp-4)	0x00377244	0x2611500
FP		0x00377248	0x000045F7
		0x0037724C	

Restoring Registers

```
.text
```

```
...
```

```
doneFact:
```

```
    sw    $s2, -4(fp)           # put our return value relative to $fp
```

```
    # now we can begin cleanup prior to return
```

```
    addiu $sp, $sp, 4           # done with local variables so adjust $sp
```

```
    # pop saved registers
```

```
    addiu $sp, $sp, 20
```

```
    # adjust stack pointer
```

```
    sw    $ra, -4(sp)
```

```
    # restore return address
```

```
    sw    $s0, -8(sp)
```

```
    # restore the other registers we used
```

```
    sw    $s1, -12(sp)
```

```
    sw    $s2, -16(sp)
```

```
    sw    $s3, -20(sp)
```

```
...
```

The Factorial Stack Frame (pop regs)

	Special	Address	Data
		0x00377220	?
		0x00377224	5
		0x00377228	\$s3 (old)
		0x0037722C	\$s2 (old)
		0x00377230	\$s1 (old)
		0x00377234	\$s0 (old)
		0x00377238	\$ra (old)
SP	Old Frame Pointer	0x0037723C	\$fp (old)
	Input (fp-8)	0x00377240	0xD
	return (fp-4)	0x00377244	0x2611500
FP		0x00377248	0x000045F7
		0x0037724C	

Collapse Stack Frame

```
.text
```

```
...
```

```
doneFact:
```

```
    sw      $s2, -4(fp)           # put our return value relative to $fp
```

```
    # now we can begin cleanup prior to return
```

```
    addiu   $sp, $sp, 4           # done with local variables so adjust $sp
```

```
    # pop saved registers
```

```
    addiu   $sp, $sp, -20         # adjust stack pointer
```

```
    sw      $ra, -4(sp)          # restore return address
```

```
    sw      $s0, -8(sp)          # restore the other registers we used
```

```
    sw      $s1, -12(sp)
```

```
    sw      $s2, -16(sp)
```

```
    sw      $s3, -20(sp)
```

```
    # restore stack pointer and frame pointer to collapse stack frame
```

```
    move    $sp, $fp            # stack pointer is back where it was
```

```
    lw      $fp, -12(sp)         # get old $fp from where we stored it
```

```
    jr      $ra                 # everything back like it was, so return
```

```
# END OF Factorial Procedure
```

```
...
```


The Factorial Stack Frame (collapsed)

	Special	Address	Data
		0x00377220	?
		0x00377224	5
		0x00377228	\$s3 (old)
		0x0037722C	\$s2 (old)
		0x00377230	\$s1 (old)
		0x00377234	\$s0 (old)
		0x00377238	\$ra (old)
		0x0037723C	\$fp (old)
	Input (sp-8)	0x00377240	0xD
	return (sp-4)	0x00377244	0x2611500
SP		0x00377248	0x000045F7
		0x0037724C	

Stack Frame Summary

Caller:

- Pushes \$ra if not already on stack before anything else
- Puts parameters on stack above stack pointer location
- Leaves room for returns below params
- Does NOT adjust the stack pointer after storing params
- Calls the procedure with the 'jal' instruction

Stack Frame Summary (cont.)

Callee:

- Saves \$fp on stack, copies \$sp to \$fp, and then adjusts \$sp to point to saved \$fp
- Pushes \$ra and other registers on stack
- Adjusts stack pointer to allow for local variables

When ready to return:

- Adjusts stack pointer back to before local variables
- Pops \$ra and other registers from stack
- Copies \$fp to \$sp and then restores \$fp from stack
- Returns by 'jr \$ra' to the Caller