

First work example: the accelerated scalar Himeno code

Having a more detailed view

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Profiling the first steps



Comparing versions 01a and 02

- \$> module load craype-accel-nvidia60
- \$> module load perftools-base
- \$> module load perftools-lite-gpu
- \$> make VERSION=XX OMP=yes
- \$> sbatch submit.wlm # if necessary change to himeno_vXX.x

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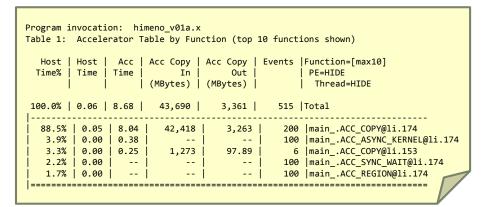
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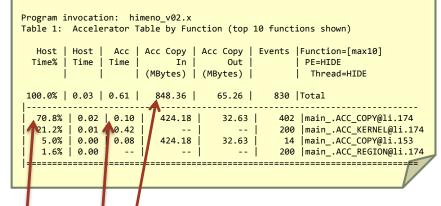
Why is version 2 so much faster?



v01a: Single kernel

v02: data region in jacobi()





reduced data copy

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First stages to accelerating an application



Understand and characterise the application

 Profiling tools, code inspection, speaking to developers if you can

1. Introduce first OpenMP kernels

Already done

2. Introduce data regions in subprograms

- reduce unnecessary data movements
- will probably require more OpenMP kernels

Next stages to accelerating an application

- Move up the calltree, adding higher-level data regions
 - ideally, port entire application so data arrays live entirely on the GPU
 - otherwise, minimise traffic between CPU and GPU
 - This will give the single biggest performance gain

Only now think about performance tuning for kernels

- First correct any obviously inefficient scheduling on the GPU
 - This will give some good performance improvements
- Optionally, experiment with OpenMP tuning clauses
 - You may gain some final additional performance from this
- And remember Amdahl's law...



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Step 3: Further optimising data movements

- The code times the calls to jacobi() routine
- Each call contains a data region
 - So we are still timing some data movements
- Solution: move up the call tree to parent routine
 - Add a second, outer data region
 - Spans calls to iteration routines
 - Specified arrays then only move on boundaries of outer data region
 - moves the data copies outside of the timed region

first call jacobi(

second call jacobi()

```
$> grep "End transfer" hiveno v02.log
ACC: End transfer (to acc 444780648 bytes, to host 0 bytes)
ACC: End transfer (to acc 48 bytes, to host 0 bytes)
ACC: End transfer (to acc 0 bytes, to host 24 bytes)
ACC: End transfer (to acc 12 bytes, to host 0 bytes)
ACC: End transfer (to acc 0 bytes, to host 0 bytes)
. . . # nn1 times
ACC: End transfer (to acc 0 bytes, to hos 34213896 bytes)
ACC: End transfer (to acc 444780648 bytes, to host 0 bytes)
ACC: End transfer (to acc 44 bytes, to host 0 bytes)
ACC: End transfer (to acc 0 bytes, to host 24 bytes)
ACC: End transfer (to acc 12 bytes, to host 0 bytes)
ACC: End transfer (to acc 0 bytes, to host 0 bytes)
. . . # nn2 times
ACC: End transfer (to acc 0 bytes, to host 34213896 bytes)
```

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Step3: Adding a data region

- Spans both calls to jacobi
- Arrays just copyin now
 - and transfers not timed
 - p can be to
 - wrk2 can be alloc
- data region in jacobi()
 - not necessary
 - but you keep it (nest data regions)
 - if present data is not touched

- Drawback: arrays have to be in scope for this to work
 - may need to unpick clever use of module data
 - or use OpenMP4.5 unstructured data regions

```
PROGRAM himeno
    CALL initmt
 !$omp target data &
 !$omp& map(to:a,b,c,bnd,wrk1)&
 !$omp& map(to:p) map(alloc:wrk2)
    cpu0 = gettime()
                                         outer data region
    CALL jacobi(3,gosa)
    cpu1 = gettime()
                                           p is defined as
    cpu0 = gettime()
                                                 copyin
    CALL jacobi(nn,gosa)
    cpu1 = gettime()
 !$omp end target data
 END PROGRAM himeno
   data clauses
                            SUBROUTINE jacobi(nn,gosa)
   here optional
                            !$omp target data map(tofrom:p)&
                            !$omp& map(to:a,b,c,bnd,wrk1)&
  if jacobi() is called
                            !$omp& map(alloc:wrk2)
outside the outer data
                              iter lp: DO loop = 1,nn
region, data movement
      is defined
                                 <...>
                               ENDDO iter lp
                            !$omp end target data
                            END SUBROUTINE jacobi
```

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Step 3: Performance with the first kernel



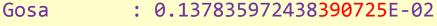
Original performance:

Gosa : 0.137835972438443027E-02

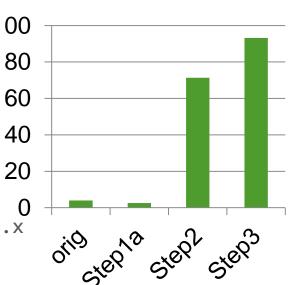
MFLOPS : 3750.2014219750636

With outer OpenMP data region:

- \$> make VERSION=03 OMP=yes
- \$> sbatch submit.wlm # change to himeno_v03.x



MFLOPS : 93190.630878365104



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Step 3a: Going further

- Best solution is to port entire application to GPU
 - data regions span entire use of arrays
 - all enclosed loopnests accelerated with OpenMP device constructs
 - no significant data transfers
- Expand outer data region
 - to include call to initialisation routine as well
 - arrays can now all be declared as scratch space with alloc clause
 - need to accelerated loopnests in initmt(), no data scoping risk over-cautious
- N.B. No easy way to ONLY allocate arrays in GPU memory
 - CPU version is now dead space, but
 - GPU memory is usually the limiting factor, so usually not a problem

```
PROGRAM himeno
!$omp target data &
!$omp& map(alloc:a,b,c,p)&
!$omp& map(alloc:bnd,wrk1,wrk2)
    CALL initmt
                            SUBROUTINE initmt
    cpu0 = gettime()
    CALL jacobi(3,gosa)
                            !$omp target teams distribute
    cpu1 = gettime()
                               DO k = 1,mkmax
                                  a(i,j,k,1) = 0d0
    cpu0 = gettime()
                                  . . .
    CALL jacobi(nn,gcsa)
                               ENDDO
    cpu1 = gettime()
                            !$omp end target teams distribute
!$omp end target dat
END PROGRAM himeno
                            !$omp target teams distribute
                               DO k = 1.kmax
                               ENDDO
                            !$omp end target teams distribute
 SUBROUTINE jacobi(nn,
                            END SUBROUTINE initmt
    iter_lp: DO loop = '
 !$omp target teams dist
       DO k = 2, kmax-1
       <...>
      ENDDO
 !$omp end target teams distribute
    ENDDO iter lp
 END SUBROUTINE jacobi
```

Step 3a: Performance with the first kernel



Original performance:

Gosa : 0.137835972438443027E-02

MFLOPS : 3750.2014219750636

whole main in OpenMP data region:

\$> make VERSION=03a OMP=yes

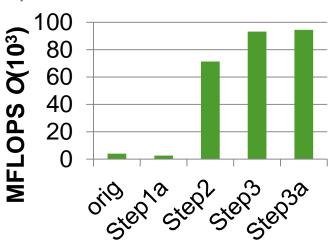
\$> sbatch submit.wlm # change to himeno_v03a.x

Gosa : 0.137835972438396428E-02

MFLOPS : 94522.234347635153

 No significant data transfers now

> doesn't improve measured performance in this case



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Step 4: Performance with larger problem size



Original performance:

```
$> make VERSION=00 OMP=yes \
    PROBLEM_SIZE=4
```

\$> sbatch submit.wlm # change to himeno_v00.x

Gosa : 0.775850492248915911E-03

MFLOPS : 3368.2488685152798

whole main in OpenMP data region:

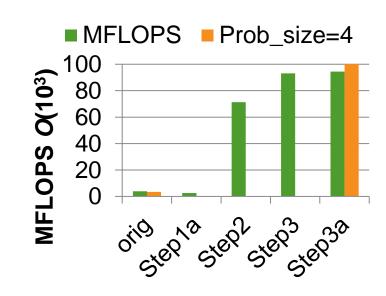
```
$> make VERSION=03a OMP=yes \
    PROBLEM SIZE=4
```

\$> sbatch submit.wlm # change to himeno_v03a.x

Gosa : 0.775850492248780278E-03

MFLOPS : 100140.70427375865

 Why does this bigger problem perform better?



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Now it's your turn



- vary problem size
- advanced loop scheduling
 - tuning clauses
 - vary threads per block
 - loop collapsing



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In summary

- We ported the entire Himeno code to the GPU
 - chiefly to avoid data transfers
 - 4 OpenMP kernels (only 1 significant for compute performance)
 - 1 outer data region
 - 1 inner data regions (nested within this)
 - 6 directive pairs for 200 lines of Fortran/C
 - Profiling frequently showed the bottlenecks
 - Correctness was also frequently checked
- First step was optimising data transfers
- Next steps
 - Checking kernels are scheduling sensibly
 - Look at kernel optimisation



In summary... continued

Further performance tuning

- data region gave a 22x speedup; kernel tuning is secondary
- Low-level languages like CUDA
 - offer more direct control of the hardware
 - but OpenMP is much easier to use
 - should get close to CUDA performance
- Remember Amdahl's Law:
 - speed up the compute of a parallel application,
 - and soon become network bound
 - Don't waste time trying to get an extra 10% in the compute
 - You are better concentrating your efforts on tuning the comms or I/O

Bottom line:

- 3-4x speedup from 7 directive pairs in 200 lines of Fortran/C
 - performance comparing GPU to the complete CPU

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