
MODULE *HKFM*

EXTENDS *Integers, Sequences*
 CONSTANTS *Client, Song*
 VARIABLES *inbox, state*

Type definitions (kinda sorta) and other useful stuff...

There are various places where we want to refer to all variables at once, so it's useful to define a *vars* tuple.

$$vars \triangleq \langle inbox, state \rangle$$

The constant *Client* is the set of all clients, represented however we want. It's defined externally, in the model.

We define *Node* to be the set of all nodes in the system, including the server. We don't care how the server is represented, only that it doesn't clash with any of the clients. We use the TLA+ CHOOSE operator to express this.

$$\begin{aligned} Server &\triangleq \text{CHOOSE } x : x \notin Client \\ Node &\triangleq Client \cup \{Server\} \end{aligned}$$

These terms all relate to the *playhead*. A *playhead* has two fields:

i: the current track in the *playlist*
t: the number of seconds into that track

When *i* = -1 it means we're not playing anything.

Every node has a *State* consisting of two fields:

playlist: a sequence of songs from the constant set *Song*
playhead: as described above

$$\begin{aligned} Idx &\triangleq Nat \cup \{-1\} \\ Playlist &\triangleq Seq(Song) \\ Playhead &\triangleq [i : Idx, t : Nat] \\ Stopped &\triangleq [i \mapsto -1, t \mapsto 0] \\ State &\triangleq [playlist : Playlist, playhead : Playhead] \\ InitState &\triangleq [playlist \mapsto \langle \rangle, playhead \mapsto Stopped] \end{aligned}$$

Clients send "add", "seek", and "skip" messages to the server and the server sends "sync" messages to all clients whenever its state changes. The term *Message* is the set of all possible messages that can occur.

$$\begin{aligned} Message &\triangleq [action : \{\text{"sync"}\}, data : State] \cup \\ &\quad [action : \{\text{"add"}\}, data : Song, sender : Client] \cup \\ &\quad [action : \{\text{"seek"}, \text{"skip"}\}, data : Playhead, sender : Client] \end{aligned}$$

The *TypeOK* formula states that *inbox* must be a function from nodes to sequences of messages and state must be a function from nodes to states. We can ask *TLC* to check that *TypeOK* is an invariant of every behaviour, meaning it will find circumstances where *inbox* and state end up looking wonky. It's also useful to have as a high level type definition for these variables.

$$\begin{aligned} TypeOK &\triangleq \wedge inbox \in [Node \rightarrow Seq(Message)] \\ &\quad \wedge state \in [Node \rightarrow State] \end{aligned}$$

Message Constructors

These operators are just for convenience when creating messages in actions below.

$$\begin{aligned} SyncMsg &\triangleq [action \mapsto \text{"sync"}, data \mapsto state'[Server]] \\ AddMsg(client, song) &\triangleq [action \mapsto \text{"add"}, data \mapsto song, sender \mapsto client] \\ SeekMsg(client, playhead) &\triangleq [action \mapsto \text{"seek"}, data \mapsto playhead, sender \mapsto client] \\ SkipMsg(client, playhead) &\triangleq [action \mapsto \text{"skip"}, data \mapsto playhead, sender \mapsto client] \end{aligned}$$

Client Actions

$$\begin{aligned} SendAdd(self, song) &\triangleq \\ &\text{LET} \\ &\quad msg \triangleq AddMsg(self, song) \\ &\text{IN} \\ &\quad \wedge inbox' = [inbox \text{ EXCEPT } ![Server] = Append(inbox[Server], msg)] \\ &\quad \wedge \text{UNCHANGED } state \\ RecvSync(self) &\triangleq \\ &\quad \wedge inbox[self] \neq \langle \rangle \\ &\quad \wedge \text{LET} \\ &\quad \quad msg \triangleq Head(inbox[self]) \\ &\quad \quad tail \triangleq Tail(inbox[self]) \\ &\quad \text{IN} \\ &\quad \quad \wedge msg.action = \text{"sync"} \\ &\quad \quad \wedge inbox' = [inbox \text{ EXCEPT } ![self] = tail] \\ &\quad \quad \wedge state' = [state \text{ EXCEPT } ![self] = msg.data] \\ SendSeek(self) &\triangleq \\ &\quad \text{LET} \\ &\quad \quad playhead \triangleq state[self].playhead \\ &\quad \quad msg \triangleq SeekMsg(self, [playhead \text{ EXCEPT } !.t = playhead.t + 1]) \end{aligned}$$

IN
 $\wedge \text{playhead} \neq \text{Stopped}$
 $\wedge \text{inbox}' = [\text{inbox} \text{ EXCEPT } ![Server] = \text{Append}(\text{inbox}[Server], \text{msg})]$
 $\wedge \text{UNCHANGED state}$

 $\text{SendSkip}(\text{self}) \triangleq$
 LET
 $\text{playhead} \triangleq \text{state}[\text{self}].\text{playhead}$
 $\text{msg} \triangleq \text{SkipMsg}(\text{self}, \text{playhead})$
 IN
 $\wedge \text{playhead} \neq \text{Stopped}$
 $\wedge \text{inbox}' = [\text{inbox} \text{ EXCEPT } ![Server] = \text{Append}(\text{inbox}[Server], \text{msg})]$
 $\wedge \text{UNCHANGED state}$

Server Actions

$\text{BroadcastSync} \triangleq$
 $\wedge \text{inbox}' = [n \in \text{Node} \mapsto \text{IF } n = \text{Server}$
 $\quad \text{THEN } \text{Tail}(\text{inbox}[n])$
 $\quad \text{ELSE } \text{Append}(\text{inbox}[n], \text{SyncMsg})]$

 $\text{RecvAdd} \triangleq$
 $\wedge \text{inbox}[Server] \neq \langle \rangle$
 $\wedge \text{LET}$
 $\quad \text{server} \triangleq \text{state}[Server]$
 $\quad \text{msg} \triangleq \text{Head}(\text{inbox}[Server])$
 IN
 $\wedge \text{msg.action} = \text{"add"}$
 $\wedge \text{LET}$
 $\quad \text{newPlaylist} \triangleq \text{Append}(\text{server.playlist}, \text{msg.data})$
 $\quad \text{newPlayhead} \triangleq \text{IF } \text{server.playhead} = \text{Stopped}$
 $\quad \quad \text{THEN } [i \mapsto \text{Len}(\text{server.playlist}), t \mapsto 0]$
 $\quad \quad \text{ELSE } \text{server.playhead}$
 IN
 $\wedge \text{state}' = [\text{state} \text{ EXCEPT } ![Server] = [\text{playlist} \mapsto \text{newPlaylist},$
 $\quad \text{playhead} \mapsto \text{newPlayhead}]]$
 $\wedge \text{BroadcastSync}$

 $\text{RecvSeek} \triangleq$
 $\wedge \text{inbox}[Server] \neq \langle \rangle$
 $\wedge \text{LET}$
 $\quad \text{server} \triangleq \text{state}[Server]$
 $\quad \text{msg} \triangleq \text{Head}(\text{inbox}[Server])$
 IN
 $\wedge \text{msg.action} = \text{"seek"}$

$$\wedge state' = [state \text{ EXCEPT } ![Server].playhead.t = msg.data.t]$$

$$\wedge BroadcastSync$$

$$RecvSkip \triangleq$$

$$\wedge inbox[Server] \neq \langle \rangle$$

$$\wedge \text{LET}$$

$$server \triangleq state[Server]$$

$$msg \triangleq Head(inbox[Server])$$

$$\text{IN}$$

$$\wedge msg.action = \text{"skip"}$$

$$\wedge \text{LET}$$

$$newIndex \triangleq server.playhead.i + 1$$

$$newPlayhead \triangleq \text{IF } newIndex < Len(server.playlist)$$

$$\text{THEN } [i \mapsto newIndex, t \mapsto 0]$$

$$\text{ELSE } Stopped$$

$$\text{IN}$$

$$\wedge state' = [state \text{ EXCEPT } ![Server].playhead = newPlayhead]$$

$$\wedge BroadcastSync$$

Randomly lose a message from an *inbox*

$$Remove(i, seq) \triangleq$$

$$[j \in 1 \dots (Len(seq) - 1) \mapsto \text{IF } j < i \text{ THEN } seq[j] \text{ ELSE } seq[j + 1]]$$

$$LoseMsg \triangleq$$

$$\exists n \in \text{DOMAIN } inbox :$$

$$\exists i \in \text{DOMAIN } inbox[n] :$$

$$\wedge inbox' = [inbox \text{ EXCEPT } ![n] = Remove(i, inbox[n])]$$

$$\wedge \text{UNCHANGED } state$$

Spec

$$Init \triangleq$$

$$\wedge inbox = [n \in Node \mapsto \langle \rangle]$$

$$\wedge state = [n \in Node \mapsto InitState]$$

$$Next \triangleq$$

$$\vee \exists self \in Client, song \in Song : SendAdd(self, song)$$

$$\vee \exists self \in Client : RecvSync(self)$$

$$\vee \exists self \in Client : SendSeek(self)$$

$$\vee \exists self \in Client : SendSkip(self)$$

$$\vee RecvAdd$$

$$\vee RecvSeek$$

$$\vee RecvSkip$$

$$\vee LoseMsg$$

$$Spec \triangleq \\ Init \wedge \Box[Next]_{vars}$$

THEOREM $Spec \Rightarrow \Box TypeOK$
