Concurrency in Go

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Concurrency... in Go

- Concurrency is about dealing with lots of things at once.
 - The opposite of sequential, or dealing with one thing at a time.
 - A web server deals with concurrent incoming and outgoing work whose lifetimes overlap.
- The Go programming language was open sourced in 2009.
 - o In the Algol family (C, C++, Python, Java, JavaScript, C#). Example code:

```
func printTheSquare(x int) {
   logger.Println(x, "squared is", x*x)
}
```

- Emphasis on software engineering in practice.
- One practical concern is concurrency, and this talk is about that aspect of Go.

Concurrency

- Concurrency is the composition of independently executing things.
 - Concurrency is about dealing with lots of things at once.
 - Think of a web server handling overlapping requests.
- Parallelism is the simultaneous execution of (possibly related) things.
 - Parallelism is about doing lots of things at once.
 - Think of vector math or GPUs, doing the same operation on multiple elements.
- Related, but separate ideas.
 - https://blog.golang.org/concurrency-is-not-parallelism
 - This talk is about concurrency.
- For example, Unix pipes connect processes.
 - find /usr/local/go/src/image | grep _test.go\$ | xargs wc -l
 - o On a single-core CPU, this will be concurrent without being parallel.

Concurrency

- Two popular approaches to concurrency.
 - Equally expressive; duals of one another.
 - o Dan Kegel, "The C10k problem", 1999, surveys handling 10,000 network connections.
- Threads are procedure oriented.
 - Synchronous.
 - 'Obvious' extension of sequential, imperative programming: multiple flows of control.
 - Think of Unix processes... that share memory.
- Events are message oriented.
 - Asynchronous.
 - Think of typical GUIs (Graphical User Interfaces).

Threads

- Threads are like sub-processes that share memory (the address space).
 - Context switches are cheaper than between processes.
 - Can share more state, e.g. caches, intermediate work.
- One problem: race conditions on that memory.
 - "numRequests++" is not atomic.
 - One solution: mutexes enforce mutual exclusion.
 - Another solution: only share immutable things.
- Another problem: fault isolation.
 - The Google Chrome web browser deliberately uses multiple (sandboxed) processes.
 - Not covered in this talk.
- Larger design question: put concurrency in the language or the libraries?
 - Not covered in this talk: functional languages, immutability.

Mutexes

- Protects shared, mutable state.
 - A necessary evil. (Or is it?)
- One problem: holding the lock for too short a time.
 - TOCTTOU (Time Of Check To Time Of Use).
 - Or, more commonly, simply forgetting to take the lock when needed.
- Another problem: holding the lock for too long a time.

```
o mutex.Lock()
    x++
    logger.Println("x is", x)  // I/O can take arbitrarily long.
    mutex.Unlock()
```

- More problems: deadlock (e.g. dining philosophers), livelock, starvation.
 - O How coarse or fine should the locks be?

Events

- Instead of actively doing this and that, loop:
 - Select what you're interested in, and then
 - React to things that happen.
- Commonly seen in GUI programming and in JavaScript.
 - Mouse, keyboard, paint events are all interleaved in the one event loop.
 - In web browser JavaScript, XMLHttpRequest events are also interleaved.
 - For web servers (e.g. Node.JS), incoming requests and back-end responses are interleaved.
- Asynchronous APIs.
 - When reading from a file or socket, don't wait for the result, tell me later.
- Often only one OS-level thread.
 - No need to create, schedule or mutually exclude multiple threads.

Events

- Events have their own problems.
 - Callbacks, callbacks everywhere.
 - Context lost as a single conceptual operation broken into multiple handlers:
 - I drag the mouse from A to B, but "A's location" isn't part of:
 - B's mouse event,
 - a local variable in the mouse event handler or
 - the call stack.
 - Somewhat easier with closures (e.g. JavaScript), but not ideal (e.g. nesting).
- Asynchronous APIs are viral.
 - http://journal.stuffwithstuff.com/2015/02/01/what-color-is-your-function/

Go

- Concurrency in Go is built on three things:
 - Goroutines, a concurrent computation model,
 - Channels, a concurrent communication model,
 - Select, a concurrent control structure.
- These build on old concepts that were half-forgotten.
 - Tony Hoare, "Communicating Sequential Processes" (CSP), 1978.
- Threads or Events? Both!
 - The Go programmer uses 'threads', which is simpler conceptually.
 - The Go runtime is implemented as 'events', which is more efficient.
- No silver bullet. You can still have race conditions and deadlocks in Go.
 - But the programming model is higher level and it's easier to avoid bugs a priori.
 - Primary benefit is clarity, not efficiency per se, but clarity can lead to efficiency.

Goroutines

- Unit of execution: PC (program counter) and the stack.
 - Many goroutines are multiplexed onto few OS-level threads.
 - Cheaper than threads to create and switch between.
 - Goroutine stacks start small (2 kilobytes) and grow on demand. On Linux, by default, threads take 2 megabytes up front.
 - Scheduling doesn't require going through the kernel and back.
 - Language makes it easy and idiomatic to create new goroutines.
 - Feasible for a web server process to have a million goroutines.
- This requires compiler and runtime support.
 - We already require a runtime for garbage collection.
 - Other, leaner languages choose no mandatory runtime. That's OK too.
 - Programming language design is about trade-offs. There's more than one 'right' answer.

Goroutines

Do foo and bar sequentially (similar to ; in a Unix shell):

```
o foo(42)
bar()
```

 Do foo and bar concurrently (similar to & in a Unix shell), where foo runs in a new goroutine while bar runs in this one:

```
o go foo(42)
bar()
```

• The order matters. This does not do foo concurrently with bar, as bar runs in a new goroutine *after* foo finishes:

```
o foo(42)
go bar()
```

Goroutines

- I/O APIs are synchronous. The runtime makes it efficient.
 - o nBytesRead, err := conn.Read(buffer)
- Synchronous is simpler.
 - o Go decoders (e.g. decompressors, image codecs) can just Read and block.
 - o C/C++ decoders have to spill their intermediate state (local variables) and resume later.
- In Go, the basic web server model is simple, concurrent and efficient.

```
o for {
    conn, err := socket.Accept()
    if err != nil {
       return err
    }
    go serve(conn)
}
```

The "go" keyword starts a new, concurrent goroutine.

```
o func main() {
    go expensiveComputation(x, y, z)
    anotherExpensiveComputation(a, b, c)
}
```

- This story is incomplete.
 - O How do we know when the two computations are done?
 - O What are their values?

Goroutines communicate with other goroutines via channels:

```
func computeAndSend(ch chan int, x, y, z int) {
  ch <- expensiveComputation(x, y, z)</pre>
func main() {
  ch := make(chan int)
  go computeAndSend(ch, x, y, z)
  v2 := anotherExpensiveComputation(a, b, c)
  v1 := <-ch
  fmt.Println(v1, v2)
```

- Communication (the <- operator) is sharing and synchronization.
 - Sharing is often giving.
 - Receive blocks until there is a sender.
 - Send blocks until there is a receiver.
 - Multiple readers and multiple writers are OK.
- Channels can also be buffered, so that sends don't always block.
 - Not covered in this talk.

- Do not communicate by sharing memory; instead, share memory by communicating.
 - https://blog.golang.org/share-memory-by-communicating
- Threads and locks are concurrency primitives; CSP is a concurrency model:
 - Analogy: Edgar Dijkstra, "Go To Statement Considered Harmful", 1968.
 - goto is a control flow primitive; structured programming (if statements, for loops, function calls)
 is a control flow model.

- Let's distribute some work over a pool of workers.
- Compare communicating by sharing memory...

```
type Work struct {
  x, y, z int
  assigned bool
type WorkSet struct {
       sync.Mutex
  mu
  work []*Work
```

- ...with sharing memory by communicating.
- Each worker receives and sends pieces of work:

```
type Work struct { x, y, z int }
 func worker(in, out chan *Work) {
   for {
    work := <-in
    work.z = work.x + work.y
     out <- work
```

The manager connects the workers:

```
func main() {
  in := make(chan *Work)
  out := make(chan *Work)
  for i := 0; i < nWorkers; i++ {
    go worker(in, out)
  go workProducer(in)
 workConsumer(out)
```

Select

Select chooses 1 of N communications. It blocks until one can proceed:

```
for {
  select {
  case work := <-in:</pre>
     doSome(work)
     out <- work
  case <-done:</pre>
     return
```

- "Done" could be a timeout, an explicit cancellation, etc.
 - https://blog.golang.org/context

Select

- Let's make a "chat roulette" server.
 - https://blog.golang.org/two-recent-go-talks
- Recall the basic web server model:

```
o for {
    conn, err := socket.Accept()
    if err != nil {
       return err
    }
    go serve(conn)
}
```

Select

Matching two connections can select on the same channel:

```
var matcher = make(chan net.Conn)
 func serve(conn net.Conn) {
   select {
   case matcher <- conn: // Sending myself means please match with me.
     // Nothing else to do. We're now handled by the other goroutine.
   case partnerConn := <-matcher: // Receiving means I met my partner.</pre>
     chat(partnerConn, conn)
```

Demo!

Summary

- Threads+locks and events can work, but Go offers a different approach:
 - o Goroutines, a concurrent computation model.
 - Channels, a concurrent communication model.
 - Select, a concurrent control structure.
- Do not communicate by sharing memory; instead, share memory by communicating.
- Not a law, just a different way to think about concurrency.
 - Decompose problems into small, simple, self-contained things, working in concert.
 - Concurrency patterns: http://talks.golang.org/2012/concurrency.slide
 - Don't overdo it. Sometimes all you need is a mutex (and Go provides them).
- http://golang.org/ has tutorials, downloads, and more!