

##

Array

→ It is used to store same type of data elements in the continuous manner. "Name of the array is the Base Address."

A	B	C	D	E
0	1	2	3	4

** Advantages

① Random Access \Rightarrow Get i^{th} element in $O(1)$ Time

because all have address of the first element.

② Cache Friendliness \Rightarrow Cache is fastest and very close

to the CPU. The cache friendliness refers to that to

access the nearby set of element if you access element of array

that means, if you access element A, then B, C, D also get fetched.

##

Array Types

① Static/Fixed Sized \Rightarrow Size of Array is Fixed

Eg: \rightarrow ① int arr[100] ② int arr[n] ③ int arr[7] = {10, 20, 30, 40}

Stack Allocation

④ int *arr = new int[n]

Heap Allocation.

② Dynamic Sized Array \Rightarrow Array size is not fixed, it can be changed

\rightarrow Vector is used to implement the Dynamic Sized Array

* Resize automatically

Operation in Array

* Searching in Array

* For Sorted Array search is done by Binary Search

* For Unsorted Array search is done by Linear Search

Linear Search

```
int Search (int arr[], int n, int k)
{
    for (int i = 0; i < n; i++)
        if (arr[i] == k)
            return i;
}
```

return -1;

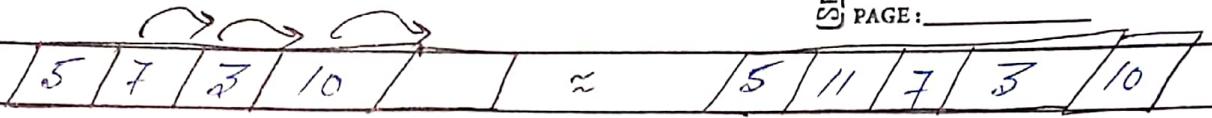
Time Complexity $\Rightarrow O(n)$

} Insertion in Array

\rightarrow We can't insert if array size is fixed and it is already full.

* For that we need to check whether array is full or not.

Good Write



pos = 2, x = 11, idx = 1

```
int insert (int arr[], int n, int x, int cap, int pos)
{
    if (n == cap)
        return n;

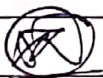
    int idx = pos - 1;

    for (int i = n - 1; i >= idx; i--)
        arr[i + 1] = arr[i];

    arr[idx] = x;

    return (n + 1);
}
```

Time Complexity $\Rightarrow O(N)$, At the End $\Rightarrow O(1)$.



Deletion in array

```
int delete (int arr[], int n, int x)
{
    int i;

    for (i = 0; i < n; i++)
        if (arr[i] == x) // Searching for 1st occurrence
            break;
}
```

Good Write

if (s == n) // Element Not found

return n;

for (int j = s; j < n-1; j++) // If element present loop will run

arr[j] = arr[j+1]; // Shifting of elements

return s-1; // Return deleted array

}

Summary of operations

Insert $O(n)$

Search $O(n)$ [Unsorted] $O(\log n)$ [Sorted]

Delete $O(n)$

Get i^{th} element $O(1)$

Update i^{th} element $O(1)$.

Insert at end & Delete at end done in $O(1)$ Time

Largest Element in Array

int largest (int arr[], int n)

{
int max = arr[0];

for (s=1; s < n; s++)

Good Write

if (arr[i] > max)

max = arr[i]

return max;

(*) Method for finding size of any array

int n = sizeof(arr) / sizeof(arr[0]).

(*) Problem: 1 \Rightarrow Reverse the Array / String

int reversearray(int arr[], int n)

{
 //id
 int start = 0; int end = n-1; int temp;

while (arr[start] != arr[end])

{
 temp = arr[start];

arr[start] = arr[end];

arr[end] = temp;

start ++;

end --;

}

return

(*) Swapping with the help of temp variable

Good Write