



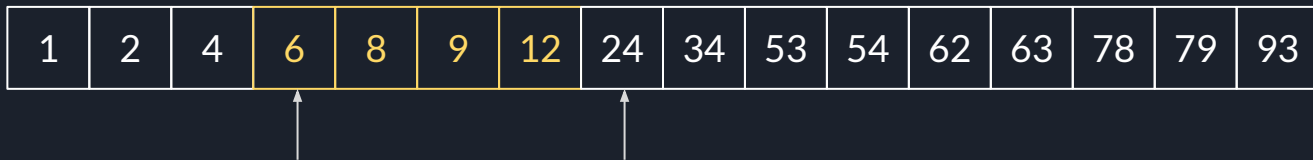
## 20. 滑動窗口



# 什麼是滑動窗口

滑動窗口(sliding window)為雙指針(two pointers)的延伸，透過窗口的身長或縮短維持窗口內資料的特性。

滑動窗口常用於子序列、子字串或連續性資料之分析。





# 滑動窗口的架構

寫滑動窗口一定要先定義：

窗口的意義- 窗口內的資料代表什麼？要記錄什麼？

窗口的型態- 固定大小？動態變化？

窗口的資料變化- 何時擴大窗口？何時縮小窗口？



# Leetcode 1456 - Maximum Number of Vowels in a Substring of Given Length

給予字串 $s$ 及正整數 $k$ , 回傳任意長度為 $k$ 之子字串中, 具有最多英文母音字母的數量。

## Example 1:

**Input:**  $s = \text{"abciidef"}, k = 3$

**Output:** 3

**Explanation:** The substring "iii" contains 3 vowel letters.

## Example 2:

**Input:**  $s = \text{"aeiou"}, k = 2$

**Output:** 2

**Explanation:** Any substring of length 2 contains 2 vowels.

## Example 3:

**Input:**  $s = \text{"leetcode"}, k = 3$

**Output:** 2

**Explanation:** "lee", "eet" and "ode" contain 2 vowels.



# Leetcode 1456 - Maximum Number of Vowels in a Substring of Given Length

窗口的意義 - 窗口為各個子字串，並記錄含有之母音字母數量

窗口的型態 - 固定大小為 $k$

窗口的資料變化 - 每次均將窗口右移左側邊界縮小、右側邊界擴大



# Leetcode 1456 - Maximum Number of Vowels in a Substring of Given Length

```
class Solution {
private:
    bool isVowel(char &c) {
        return c == 'a' || c == 'e' || c == 'i' || c == 'o' || c == 'u';
    }
public:
    int maxVowels(string s, int k) {
        int max_c = 0, c = 0;
        int l = 0, r = 0;

        while (r < s.size() && r - l < k) {
            r++;
            c += (int) isVowel(s[r - 1]);
            max_c = max(max_c, c);
        }

        while (r < s.size()) {
            r++;
            c += (int) isVowel(s[r - 1]);

            while (r - l > k) {
                c -= (int) isVowel(s[l]);
                l++;
            }

            max_c = max(max_c, c);
        }

        return max_c;
    }
};
```



# Leetcode 187 - Repeated DNA Sequences

給予由 'A', 'C', 'G', 'T' 組成的字串, 請找出所有重複、10字元長的序列, 順序不限。

## Example 1:

**Input:** s = "AAAAACCCCCAAAAACCCCCCAAAAGGGTTT"

**Output:** ["AAAAACCCCC", "CCCCCAAAAA"]

## Example 2:

**Input:** s = "AAAAAAAAAAAA"

**Output:** ["AAAAAAAAAA"]



# Leetcode 187 - Repeated DNA Sequences

窗口的意義 - 窗口為各DNA子序列

窗口的型態 - 固定大小為10

窗口的資料變化 - 每次均將窗口右移(左側邊界縮小、右側邊界擴大)



# Leetcode 187 - Repeated DNA Sequences

作法一 unordered\_map + string

把所有的子序列截下來放進map中計數

```
class Solution {
public:
    vector<string> findRepeatedDnaSequences(string s) {
        if (s.size() < 10) return vector<string> ();
        unordered_map<string, int> seq;
        vector<string> rtn;

        for (int i = 0; i <= s.size() - 10; i++) {
            seq[string(s, i, 10)]++;
        }

        for (const pair<string, int> &p: seq) {
            if (p.second > 1) {
                rtn.push_back(p.first);
            }
        }

        return rtn;
    }
};
```



# Leetcode 187 - Repeated DNA Sequences

作法二 unordered\_map + 位元運算

位元運算是壓縮空間常用的技巧

設 'A' 為0, 'C' 為1, 'G' 為2, 'T' 為3

每次將數值乘以4 (左移2位元)

並去掉第20位元以後的值 ( $\% 4^{10}$ )

再加上新加入窗口內的值

這樣可以省去建構新字串的時間、空間, 以及減少 unordered\_map 雜湊運算的時間

# Leetcode 187 - Repeated DNA Sequences

```
class Solution {
public:
    vector<string> findRepeatedDnaSequences(string s) {
        if (s.size() < 10) return vector<string> ();
        unordered_map<unsigned int, int> seq;
        vector<string> rtn;

        unsigned int val = 0;

        for (int i = 0; i < 10; i++) {
            val <= 2;
            val &= 1048575; //  $4^{10} - 1$ 


            switch (s[i]) {
                case 'A': {
                    val += 0;
                    break;
                }
                case 'C': {
                    val += 1;
                    break;
                }
                case 'G': {
                    val += 2;
                    break;
                }
                case 'T': {
                    val += 3;
                    break;
                }
            }
            seq[val] = 1;
        }

        int l = 0, r = 10; // window [l, r)
        while (r < s.size()) {
            l++; r++;
            val <= 2;
            val &= 1048575;

            switch (s[r - 1]) {
                case 'A': {
                    val += 0;
                    break;
                }
                case 'C': {
                    val += 1;
                    break;
                }
                case 'G': {
                    val += 2;
                    break;
                }
                case 'T': {
                    val += 3;
                    break;
                }
            }

            seq[val]++;
            if (seq[val] == 2) {
                rtn.push_back(string(s, l, 10));
            }
        }

        return rtn;
    }
};
```



### 3. Longest Substring Without Repeating Characters

給予一個字串，回傳不含重複字元的最長子字串長度

字串可能含有字母、數字或符號

#### Example 1:

**Input:** `s = "abcabcbb"`

**Output:** `3`

**Explanation:** The answer is "abc", with the length of 3. Note that "bca" and "cab" are also correct answers.

#### Example 2:

**Input:** `s = "bbbbbb"`

**Output:** `1`


**Explanation:** The answer is "b", with the length of 1.

#### Example 3:

**Input:** `s = "pwwkew"`

**Output:** `3`

**Explanation:** The answer is "wke", with the length of 3. Notice that the answer must be a substring, "pwke" is a subsequence and not a substring.




### 3. Longest Substring Without Repeating Characters

窗口的意義 - 窗口為不含重複字元的子字串，紀錄最大長度

窗口的型態 - 動態

窗口的資料變化 - 持續擴大右側邊界，一但發現有重複便縮小左側邊界



### 3. Longest Substring Without Repeating Characters

a : 0個

b : 0個

c : 0個

len : 0

max : 0



### 3. Longest Substring Without Repeating Characters

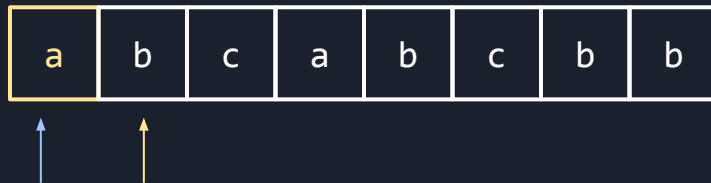
a : 1個

b : 0個

c : 0個

len : 1

max : 1



### 3. Longest Substring Without Repeating Characters

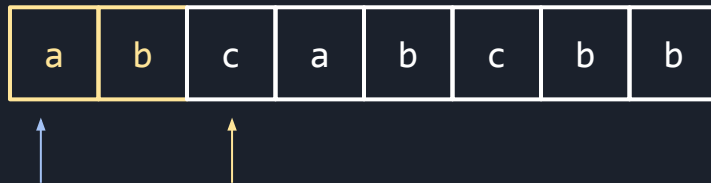
a : 1個

b : 1個

c : 0個

len : 2

max : 2





### 3. Longest Substring Without Repeating Characters

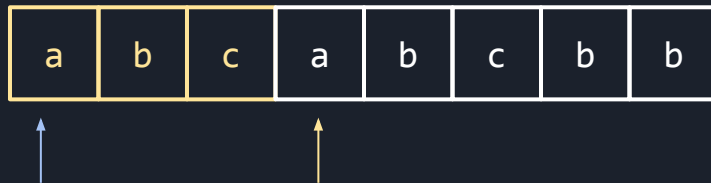
a : 1個

b : 1個

c : 1個

len : 3

max : 3



### 3. Longest Substring Without Repeating Characters

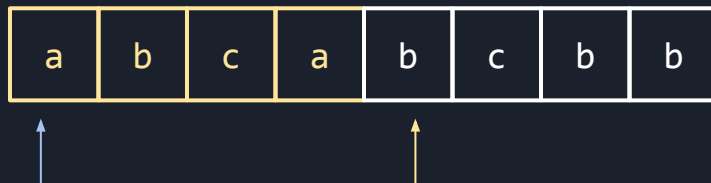
a : 2個

b : 1個

c : 1個

~~len : 4~~

max : 3



### 3. Longest Substring Without Repeating Characters

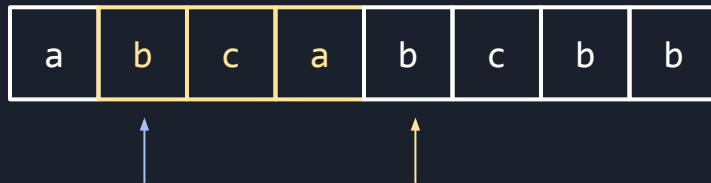
a : 2 -> 1個

b : 1個

c : 1個

len : 3

max : 3



### 3. Longest Substring Without Repeating Characters

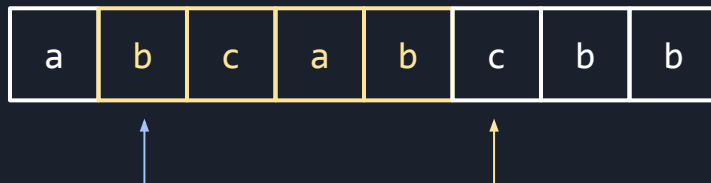
a : 1個

b : 2個

c : 1個

~~len~~ : 4

max : 3

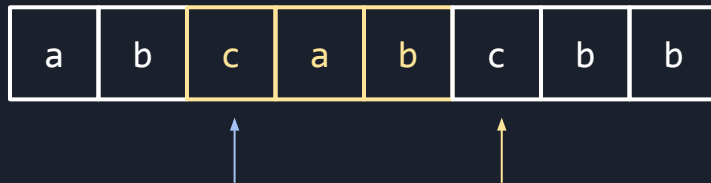


### 3. Longest Substring Without Repeating Characters

a : 1個  
b : 2 -> 1個  
c : 1個

len : 3

max : 3



### 3. Longest Substring Without Repeating Characters

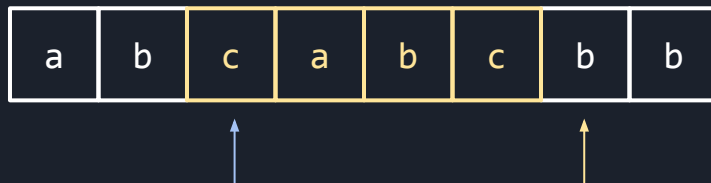
a : 1個

b : 1個

c : 2個

~~len~~ : 4

max : 3

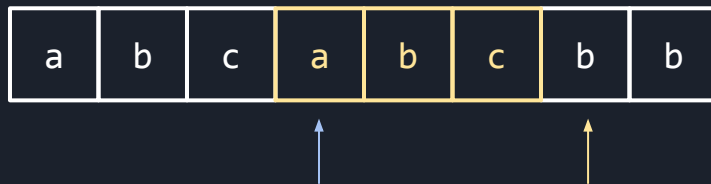


### 3. Longest Substring Without Repeating Characters

a : 1個  
b : 1個  
c : 2 -> 1個

len : 3

max : 3



### 3. Longest Substring Without Repeating Characters

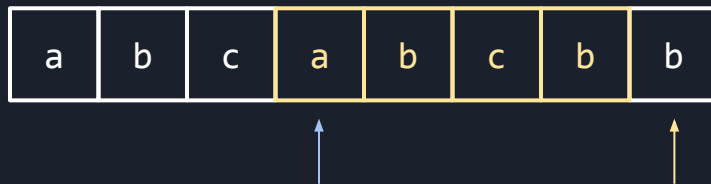
a : 1個

b : 2個

c : 1個

~~len~~ : 4

max : 3





### 3. Longest Substring Without Repeating Characters

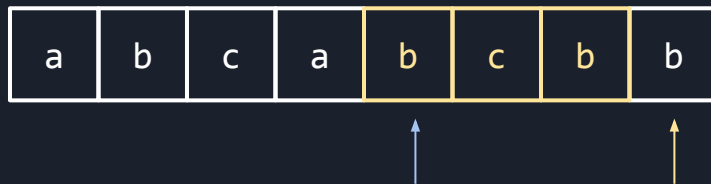
a : 1 -> 0個

b : 2個

c : 1個

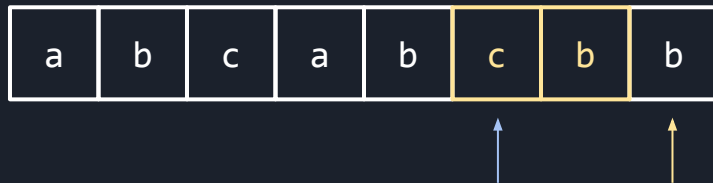
~~len : 3~~

max : 3



### 3. Longest Substring Without Repeating Characters

a : 0個  
b : 2 -> 1個  
c : 1個  
  
len : 2  
max : 3



### 3. Longest Substring Without Repeating Characters

a : 0個

b : 2個

c : 1個

~~len : 3~~

max : 3



### 3. Longest Substring Without Repeating Characters

a : 0個  
b : 2個  
c : 1 -> 0個

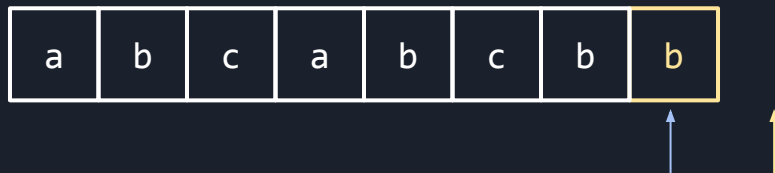
~~len : 2~~


max : 3



### 3. Longest Substring Without Repeating Characters

a : 0個  
b : 2 -> 1個  
c : 0個  
  
len : 1  
max : 3





### 3. Longest Substring Without Repeating Characters

```
class Solution {
public:
    int lengthOfLongestSubstring(string s) {
        if (s.size() == 0 || s.size() == 1) {
            return s.size();
        }
        vector<bool> seen(128, false);
        int l = 0, r = 0, max_len = 0;

        while (r < s.size()) {
            while (seen[s[r]]) {
                seen[s[l]] = false;
                l++;
                max_len = max(max_len, r - l + 1);
            }
            max_len = max(max_len, r - l + 1);
            seen[s[r]] = true;
            r++;
        }

        return max_len;
    }
};
```