

# Software Model Checking

Exercises

## Assignment 2

**Georg Weissenbacher**  
**(credit: Henning Goes)**  
**184.747**



## What you have to do

1. Translate a small C program into a transition relation
2. Write simple a bounded model checker
  - ▶ Input: A transition relation and an unwinding depth  $N$
  - ▶ Output: Is there a bug within  $N$  steps of the transition relation?

## Encoding the transition relation

The encoding is based on *SMTLib* syntax:

```
(program
  (state (x Int))
  (init (= x 0))
  (transition (= x_ (+ x 1)))
  (property (< x 10))
)
```

Encodes:  $I(x) = (x = 0)$ ,  $T(x, x') = (x' = x + 1)$ ,  $P(x) = x < 10$

See <http://smtlib.org> for the syntax & semantics of expressions & types.

## How to talk to an SMT solver

### User/Application

“I want to talk about the following variables:  
 $x$  which is an integer and  $c$  which is boolean.”

“The formula  $x > 5 \wedge c$  holds.”

“Can you find a solution?”

“What is the value of  $x$ ?”

### Solver

“Ok.”

“Ok.”

“Yes, I can!”

“It’s 20!”

## How to talk to an SMT solver

### User/Application

(declare-fun x () Int)

(declare-fun c () Bool)

“The formula  $x > 5 \wedge c$  holds.”

“Can you find a solution?”

“What is the value of  $x$ ?”

### Solver

“Ok.”

“Yes, I can!”

“It’s 20!”

### User/Application

(declare-fun x () Int)

(declare-fun c () Bool)

(assert (and (> x 5) c))

“Can you find a solution?”

“What is the value of  $x$ ?”

### Solver

“Yes, I can!”

“It’s 20!”

## How to talk to an SMT solver

### User/Application

(declare-fun x () Int)

(declare-fun c () Bool)

(assert (and (> x 5) c))

(check-sat)

“What is the value of  $x$ ?”

### Solver

sat

“It’s 20!”

## How to talk to an SMT solver

### User/Application

(declare-fun x () Int)

(declare-fun c () Bool)

(assert (and (> x 5) c))

(check-sat)

(get-value (x))

### Solver

sat

((x 20))



The SMTLib format is based on *s-expressions* (made famous by *lisp*). An expression can be:

- ▶ A constant: 10, true, etc.
- ▶ A function application:  $(f\ arg_1\ \dots\ arg_n)$

For example, the expression  $x = y \cdot z$  would be represented as:  
`(= x (* y z))`

SMTLib supports various types:

- ▶ Booleans: `Bool`
- ▶ Natural numbers: `Int`
- ▶ Bitvectors of length  $n$ : `(_ BitVec n)`
- ▶ Arrays with index type  $i$  and element type  $e$ : `(Array i e)`

Two ways to introduce new variables:

- ▶ Introduce a fresh unconstrained variable:  
`(declare-fun f (Int Bool) Int)`  
declares  $f$  to be an uninterpreted function with two arguments.
- ▶ Define an alias for a function:  
`(define-fun f ((x Int) (y Int)) Int (* x y))`  
defines  $f$  to be an alias for the multiplication function.

SMTLib supports multiple theories which provide you with function symbols you can use:

- ▶ Core: Boolean constants, `and`, `or`, `not`, etc.
- ▶ Bitvectors: Constants: `(_ bv5 32)`, `bvadd`, `concat`, etc.
- ▶ Integers: Integer constants, `+`, `*`, etc.
- ▶ Arrays:
  - ▶ `(store a i e)` puts element `e` into array `a` at index `i` and returns the new array.
  - ▶ `(select a i)` retrieves the element at index `i` from array `a`.

For a full list of all theories and logics, see <http://smtlib.org>

```
void sum(int n) {  
    int i = 0;  
    int r = 0;  
    while(i<=n) {  
        r = r + i;  
        i = i + 1;  
    }  
    assert(r < 100);  
}
```

Step 1: Label program locations

```
void sum(int n) {  
    int i = 0;  
    int r = 0;           //0  
    while(i<=n) {        //1  
        r = r + i;       //2  
        i = i + 1;       //3  
    }  
    assert(r < 100);     //4  
}
```

Step 2: Model the transition relation using a program counter

```
void sum(int n) {  
    int i = 0;  
    int r = 0;           //0  
    while(i<=n) {        //1  
        r = r + i;       //2  
        i = i + 1;       //3  
    }  
    assert(r < 100);     //4  
}
```

$pc = 0 \Rightarrow i' = 0 \wedge r' = 0 \wedge n' = n \wedge pc' = 1$   
 $pc = 1 \wedge i \leq n \Rightarrow r' = r \wedge i' = i \wedge n' = n \wedge pc' = 2$   
 $pc = 1 \wedge i > n \Rightarrow r' = r \wedge n' = n \wedge pc' = 4$   
 $pc = 2 \Rightarrow r' = r + 1 \wedge i' = i \wedge n' = n \wedge pc' = 3$   
 $pc = 3 \Rightarrow r' = r \wedge i' = i + 1 \wedge n' = n \wedge pc' = 1$   
 $pc = 4 \Rightarrow pc' = 5$   
 $pc = 5 \Rightarrow pc' = 5$

See tests/sum.1 for the complete solution.

## Exercise 1

Translate the following program:

```
void gcd(int starta , int startb) {  
    int a = starta , b = startb;  
    int t;  
    while(b!=0) {  
        t = b;  
        b = a % b;  
        a = t;  
    }  
    assert(starta % a == 0);  
    assert(startb % a == 0);  
}
```



We're using *MathSAT 5* as the SMT solver for this exercise:

- ▶ Download from <http://mathsat.fbk.eu/>
- ▶ Available for Linux, Mac and Windows.

The code template is hosted on GitHub:

- ▶ `git clone https://github.com/hguenther/smc.git`
- ▶ This makes it easier to distribute bug fixes
- ▶ If you don't want to install git:
  - ▶ Goto `https://github.com/hguenther/smc`
  - ▶ Click "Download ZIP"

## Compiling the code

Code is compiled using *CMake*.

- ▶ Makes it possible to use Makefiles (Linux) or Visual Studio (Windows) or your favorite (supported) build-tool.
- ▶ Get *CMake* from <http://cmake.org>

For Makefiles:

```
mkdir build
cd build
cmake ..
make
```

See `cmake --help` for a list of other generators available for your platform.

It might be tricky to let CMake know about MathSAT:

- ▶ You can tell CMake where to look for MathSAT by using:

```
cmake . -DMATHSAT_PREFIX=/path/to/mathsat
```

- ▶ Please contact me if there are problems:

```
weissenb@forsyte.at
```

Creates an abstraction of the MathSAT C-interface.

- ▶ Initialize it by giving it the name of the logic we are using:

```
IMathSAT_iface ( " QF_LIA " );
```

- ▶ Provides simplified functions: *check\_sat* checks if the instance is satisfiable etc.

Provides a simple parser for the transition relation format:

- ▶ Creating the parser:

```
int fd = open("myfile",O_RDONLY);  
Program prog(fd);
```

- ▶ *vars*: A mapping of variable names to variable types.
- ▶ *inits*: A vector of initial conditions.
- ▶ *trans*: A vector of expressions forming the transition relation.
- ▶ *props*: A vector of properties that have to be checked.

All types and expressions are unparsed S-expressions.

Provides a way to parse expressions and types into SMTLib format.

- Initialization:

```
IMathSAT iface ("QF_LIA");  
SMTLibParser<IMathSAT> parser (iface);
```

- *add\_named\_term*: Tell the parser that a certain name is bound to a given term:

```
msat_term x = ...;  
parser.add_named_term ("x", x);
```

- *parse\_term*: Parse an S-expression into an SMTLib expression.
- *parse\_type*: Parse an S-expression into an SMTLib type.

This is the template for the “BMC” class that you have to complete.

- ▶ A few helper functions should give you an idea of how to structure your code.
- ▶ To complete the exercise, you have to implement:
  - ▶ *extend()*: This function adds a new state to the vector of states and asserts that it is the successor state of the previous last state.
  - ▶ *check()*: Uses the SMT solver to find out if the BMC instance is solvable.



- ▶ All solutions are due **February 28, 2021**.
- ▶ Please contact me if you need more time.
- ▶ Submit using the TUWEL system.
- ▶ Plan enough time for setting up everything.