## $\mathbb{Z}$ - $\Delta_2$ -FLATNESS CONSTANT CASE 1 OF TRIANGLE COMPUTATIONS

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ABSTRACT. Details for bounding the width of the triangles from case 1 of [1, Table 1]. case\_01.pl contains the polymake code and case\_01.py contains the python code.

Let  $e_1, e_2$  be the standard basis of  $\mathbb{Z}^2$ . Consider the locking points  $(A, B, C) = (e_1 + e_2, -e_2, e_2)$  and let  $P' \subset \mathbb{R}^2$  be their convex hull. Let  $P \subset \mathbb{R}^2$  be a triangle circumscribed around those three locking points. Recall that  $\mathbf{0}$  and  $e_1$  are assumed to be contained in the interior of P. We consider the width directions  $e_1^*, e_2^*, e_1^* - e_2^*$ . The slopes  $m_{XY}, m_{YZ}$  and  $m_{ZX}$  of the facets of P through  $\{X, Y\}, \{Y, Z\}$  and

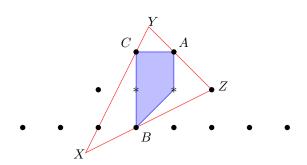


FIGURE 1. A triangle P (in red) with locking points  $(A, B, C) = (e_1 + e_2, -e_2, e_2)$ .

 $\{Z,X\}$  respectively can be expressed in terms of  $\lambda$ ,  $\mu$ , and  $\nu$  (use case\_01.py):

$$m_{XY} = \frac{2-2\lambda}{\mu}, \qquad m_{YZ} = \frac{2\nu}{-1+\mu+\nu}, \qquad m_{ZX} = \frac{2-2\lambda-2\nu}{1-\nu}.$$

Since  $P' \subset P$ , we have  $m_{XY} \geq 0$ ,  $m_{YZ} \leq 0$ , and  $m_{ZX} \leq 1$ . Since  $e_1 - e_2$  is not in the interior of P, we have  $m_{ZX} \geq 0$ . Similarly, since  $-e_1 - e_2$  is not in the interior of P, we have  $m_{XY} \geq 2$ . Hence the slopes of P satisfy:

$$m_{XY} \ge 2, \qquad m_{YZ} \le 0, \qquad 0 \le m_{ZX} \le 1.$$

Arithmetic manipulations of these inequalities yield constraints on the parameters  $\lambda$ ,  $\mu$ , and  $\nu$  which define the polytope of admissible parameters

$$Q = \{(\lambda, \mu, \nu) \in [0, 1]^3 : 1 - \lambda - \mu \ge 0, 1 - \mu - \nu \ge 0, 1 - \lambda - \nu \ge 0, -1 + 2\lambda + \nu \ge 0\}.$$

We now determine the widths of P in the directions  $e_1^*$ ,  $e_2^*$ , and  $e_1^* - e_2^*$  (use case\_01.py). On Q, these are achieved at Z - X, Y - X, and Z - Y respectively:

$$\begin{aligned} \text{width}_{e_1^*}(P) &= e_1^*(Z - X) = \frac{1 - \nu}{\delta} \\ \text{width}_{e_2^*}(P) &= e_2^*(Y - X) = \frac{2 - 2\lambda}{\delta} \\ \text{width}_{e_1^* - e_2^*}(P) &= (e_1^* - e_2^*)(Z - Y) = \frac{1 - \mu + \nu}{\delta}. \end{aligned}$$

We thus obtain

$$\begin{aligned} \operatorname{width}(P) &\leq \min \left\{ \operatorname{width}_{e_1^*}(P), \operatorname{width}_{e_2^*}(P), \operatorname{width}_{e_1^* - e_2^*}(P) \right\} \\ &= \frac{\min \left\{ 1 - \nu, 2 - 2\lambda, 1 - \mu + \nu \right\}}{\delta} =: \frac{f(\lambda, \mu, \nu)}{\delta}. \end{aligned}$$

By using case\_01.pl and Mathematica, we get that

$$\max_{(\lambda,\mu,\nu)\in Q}\frac{f(\lambda,\mu,\nu)}{\delta}=\frac{10}{3},$$

and the maximum is achieved exactly at  $(\lambda, \mu, \nu) = (\frac{3}{5}, \frac{2}{5}, \frac{1}{5})$ . Hence there is a unique maximiser, namely the triangle with vertices given by

$$\frac{1}{3} \begin{pmatrix} -4 & 1 & 6 \\ -5 & 5 & 0 \end{pmatrix}.$$

Notice that the maximiser is an admissible triangle with locking points (A, B, C)

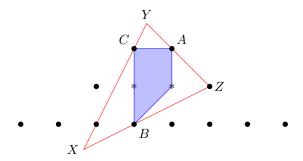


FIGURE 2. The width maximiser (in red) with locking points  $(A,B,C)=(e_1+e_2,-e_2,e_2).$ 

and interior lattice points  $\mathbf{0}$  and  $e_1$ .

## REFERENCES

[1] G. Codenotti, T. Hall, J. Hofscheier, Generalised flatness constants: a framework applied in dimension 2, preprint, arxiv.